

SRM INSTITUTE OF SCIENCE & TECHNOLOGY FACULTY OF SCIENCE AND HUMANITIES



Department of Computer Science and Applications (B.Sc. Computer Science)

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SRM INSTITUTE OF SCIENCE AND TECHNOLOGY Ramapuram, Chennai FACULTY OF SCIENCE AND HUMANITIES

Department of Computer Science & Applications (B.Sc Computer Science)				
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INTERNAL EXAMINER 1	INTERNAL EXAMINER 2			

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COMPARISON BETWEEN VARIOUS OPERATING SYSTEMS

Different Operating Systems

1.MS-DOS:

MS-DOS which is short for Microsoft Disk Operating System is a non-graphical command line operating system developed for IBM compatible computers with x86 microprocessor. The operating system used a command line interface for the user to input commands to navigate, open and manipulate files on their computer.



Features:

- It is a single user operating system meaning only one user can operate at a time.
- It is a light weight operating system allowing users to have direct access to the BIOS and its underlying hardware.
- Loads data and programs from external sources and bring them into the internal memory so they can be used on the computer.
- Enables the computer to perform input and output operations such as taking commands from keyboard, printing information on the screen.
- It is very helpful in making file management like creating, editing, deleting files, etc.
- It also controls and manages other external devices such as the printer, keyboard or external hard drive using various drive utilities.

Drawbacks:

- It does not allow multiple users to operate on the system.
- It does not support graphical interface hence mouse cannot be used to operate it.
- It does not support multiprogramming meaning it can only have one process in the ram
- It lacked memory protection which meant no security, and less stability.
- It has difficulty in memory access when addressing more than 640 MB of RAM.

2. Windows Operating System:

Windows is an operating system designed by Microsoft to be used on a standard x86 Intel and AMD processors. It provides an interface, known as a graphical user interface(GUI) which eliminates the need to memorize commands for the command line by using a mouse to navigate through menus, dialog boxes, buttons, tabs, and icons. The operating system was named windows since the programs are displayed in the shape of a square. This Windows operating system has been designed for both a novice user just using at home as well as for professionals who are into development.



Features:

• It is designed to run on any standard x86 Intel and AMD hence most of

the hardware vendors make drivers for windows like Dell, HP, etc.

• It supports enhanced performance by utilizing multi-core processors.

It comes preloaded with many productivity tools which helps to complete all types of everyday tasks on your computer.

- Windows has a very large user base so there is a much larger selection of available software programs, utilities.
- Windows is backward compatible meaning old programs can run on newer versions.
- Hardware is automatically detected eliminating need of manually installing any device drivers.

Drawbacks:

- Windows can be expensive since the OS is paid license and majority of its applications are paid products.
- Windows has high computer resource requirement like it should have high ram capacity, a lot of hard drive space and good graphics card.
- Windows slows and hangs up if the user loads up many programs at the same time.
- Windows includes network sharing that can be useful if user has a network with many PCs.
- Windows is vulnerable to virus attacks since it has a huge user base and users have to update OS to keep up-to-date with security patches.

3.LINUX:

The Linux OS is an open source operating system project that is a freely distributed, cross- platform operating system developed based on UNIX. This operating system is developed by Linus Torvalds. The name Linux comes from the Linux kernel. It is basically the system software on a computer that allows apps and users to perform some specific task on the computer. The development of Linux operating system pioneered the open source development and became the symbol of software collaboration.



Features:

- Linux is free can be downloaded from the Internet or redistribute it under GNU licenses and has the best community support.
- Linux OS is easily portable which means it can be installed on various types of devices like mobile, tablet computers.
- It is a multi-user, multitasking operating system.
- BASH is the Linux interpreter program which can be used to execute commands.
- Linux provides multiple levels of file structures i.e. hierarchical structure in which all the files required by the system and those that are created by the user are arranged.
- Linux provides user security using authentication features and also threat detection and solution is very fast because Linux is mainly community driven.

Drawbacks:

• There's no standard edition of Linux hence confusing for users and also

becoming familiar with the Linux may be a problem for new users.

- More difficult to find applications to support user needs since Linux does not dominate the market.
- Since some applications are developed specifically for Windows and Mac, those might not be compatible with linux and sometimes users might not have much of a choice to choose between different applications like in Windows or Mac since most apps are developed for operating systems that have a huge user base.
- Some hardware may not be incompatible with Linux since it has patchier support for drivers which may result in malfunction.
- There are plenty of forums to resolve Linux issues, but it may not always match the user's own level of technical understanding.

4. Solaris Operating System:

Solaris or SunOS is the name of the Sun company's Unix variant operating system that was originally developed for its family of Scalable Processor Architecture-based processors (SPARC) as well as for Intel-based processors. The UNIX workstation market had been largely dominated by this operating system during its time. As the Internet grew Sun's Solaris systems became the most widely installed servers for Web sites. Oracle purchased Sun and later renamed to Oracle Solaris.



Features:

- Solaris is known for its scalability. It can handle a large workload and still delivers indisputable performance advantages for database, Web, and Java technology-based services.
- Solaris systems were known to their availability meaning that these operating systems hardly crashes at anytime and because of its internet networking oriented design and broad scope of features it makes the job of adding new features or fixing any problems easy.
- It is built for network computing as it provides optimized network stack and support for advanced network computing protocols that delivers high-performance networking to most applications.
- Solaris has advanced, unique security capabilities which includes some of the world's most advanced security features, such as user rights management, cryptographic Framework and secure by default networking that allows users to safely deliver new solutions.
- Provides tools to enable seamless interoperability, test new software and efficiently consolidate application workloads.

Drawbacks:

- Solaris is quite expensive since it's an enterprise operating system. Also, Solaris doesn't provide updates for free.
- Solaris lacks a good graphical user interface support and is not user friendly.
- Hardware support is not nearly as good as many other operating systems.
- Performance would degrade considerably since Solaris cannot make use of different hardware that efficiently.
- Solaris sometimes becomes unstable and crashes due to total consumption of CPU and memory.

Result:

Based on comparative studies of various operating systems, the result highlights that the choice of an OS depends on its intended application, with Windows excelling in user-friendliness, Linux offering flexibility and performance, and macOS prioritizing seamless integration and design.

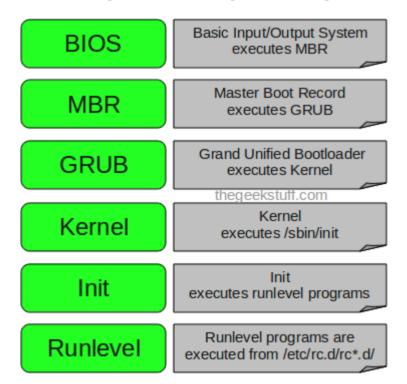
Ex.No:2	Name:
Date :	Reg.No:

BOOTING PROCESS IN LINUX OS

SIX Stages of Linux Boot Process (Startup Sequence)

When you press the power button, a fascinating sequence unfolds: the system's firmware (BIOS or UEFI) initializes hardware, checks for issues, and locates the bootloader. The bootloader then loads the Linux kernel into memory, which initializes drivers and system processes. Finally, the init system (like systemd) sets up user space, launching services and bringing you to the login prompt—all in a matter of seconds! It's a symphony of software and hardware coordination.

The following are the 6 high level stages of a typical Linux boot process.



1. BIOS

- BIOS stands for Basic Input/Output System
- Performs some system integrity checks
- Searches, loads, and executes the boot loader program.
- It looks for boot loader in floppy, cd-rom, or hard drive. You can press a key (typically F12 of F2, but it depends on your system) during the BIOS startup to change the boot sequence.
- Once the boot loader program is detected and loaded into the memory, BIOS gives the control to it.
- So, in simple terms BIOS loads and executes the MBR boot loader.

2. MBR

- MBR stands for Master Boot Record.
- It is located in the 1st sector of the bootable disk. Typically /dev/hda, or /dev/sda

MBR is less than 512 bytes in size. This has three components 1) primary boot loader info in 1st 446 bytes 2) partition table info in next 64 bytes 3) mbr validation check in last 2 bytes.

- It contains information about GRUB (or LILO in old systems).
- So, in simple terms MBR loads and executes the GRUB boot loader.

3. GRUB

- GRUB stands for Grand Unified Bootloader.
- If you have multiple kernel images installed on your system, you can choose which one to be executed.
- GRUB displays a splash screen, waits for few seconds, if you don't enter anything, it loads the default kernel image as specified in the grub configuration file.
- GRUB has the knowledge of the filesystem (the older Linux loader LILO didn't understand filesystem).
- Grub configuration file is /boot/grub/grub.conf (/etc/grub.conf is a link to this). The following is sample grub.conf of CentOS.

```
#boot=/dev/sda

default=0

timeout=5

splashimage=(hd0,0)/boot/grub/splash.xpm.gz

hiddenmenu

title CentOS (2.6.18-194.el5PAE)

root (hd0,0)

kernel /boot/vmlinuz-2.6.18-194.el5PAE ro root=LABEL=/

initrd /boot/initrd-2.6.18-194.el5PAE.img
```

- As you notice from the above info, it contains kernel and initrd image.
- So, in simple terms GRUB just loads and executes Kernel and initrd images.

4. Kernel

- Mounts the root file system as specified in the "root=" in grub.conf
- Kernel executes the /sbin/init program
- Since init was the 1st program to be executed by Linux Kernel, it has the process id (PID) of 1. Do a 'ps -ef | grep init' and check the pid.
- initrd stands for Initial RAM Disk.

initrd is used by kernel as temporary root file system until kernel is booted and the real root file system is mounted. It also contains necessary drivers compiled inside, which helps it to access the hard drive partitions, and other hardware.

- 5. I /etc/inittab file to decide the Linux run level.
 - Following are the available run levels
 - i 0 − halt
 - t 1 Single user mode
 - 2 Multiuser, without NFS
 - 3 Full multiuser mode
- L 4 unused
 - o 5 X11
 - o 6 − reboot
 - Init identifies the default initlevel from /etc/inittab and uses that to load
 - s all appropriate program.
 - Execute 'grep initdefault /etc/inittab' on your system to identify the default
 - t run level
 - If you want to get into trouble, you can set the default run level to 0 or 6.
 - ⁿ Since you know what 0 and 6 means, probably you might not do that.
 - Typically you would set the default run level to either 3 or 5.

6. Runlevel programs

- When the Linux system is booting up, you might see various services getting started. For example, it might say "starting sendmail OK". Those are the runlevel programs, executed from the run level directory as defined by your run level.
- Depending on your default init level setting, the system will execute the programs from one of the following directories.
 - Run level 0 /etc/rc.d/rc0.d/
 - Run level 1 /etc/rc.d/rc1.d/
 - Run level 2 /etc/rc.d/rc2.d/
 - Run level 3 /etc/rc.d/rc3.d/
 - Run level 4 /etc/rc.d/rc4.d/
 - Run level 5 /etc/rc.d/rc5.d/
 - Run level 6 /etc/rc.d/rc6.d/
- Please note that there are also symbolic links available for these directory under /etc directly. So, /etc/rc0.d is linked to /etc/rc.d/rc0.d.
- Under the /etc/rc.d/rc*.d/ directories, you would see programs that start with S and K.
- Programs starts with S are used during startup. S for startup.
- Programs starts with K are used during shutdown. K for kill.

- There are numbers right next to S and K in the program names. Those are the sequence number in which the programs should be started or killed.
- For example, S12syslog is to start the syslog deamon, which has the sequence number of 12. S80sendmail is to start the sendmail daemon, which has the sequence number of 80. So, syslog program will be started before sendmail.

Result: The Linux booting process involves hardware initialization by BIOS/UEFI, bootloader execution to load the kernel, kernel initialization, and system setup by the init system, culminating in the login prompt.

Ex.No:3 Name:

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MULTI-THREAD PROGRAMMING

```
#include<stdi.h>
#include<conio.h>
#include <dos.h>
void threadFunction(int id) {
  printf("Thread %d: Hello from simulated thread!\n",
  id); delay(500); // Simulate thread execution delay
}
int main() {
  for (int i = 1; i \le 3; i++) {
    threadFunction(i); // Simulate thread-like function calls
  }
  printf("Main:
                          simulated
                   All
                                       threads
finished.\n"); getch();
  return 0;
}
```

```
C:\TURBOC3\BIN>TC
Thread 1: Hello from simulated thread!
Thread 2: Hello from simulated thread!
Thread 3: Hello from simulated thread!
Main: All simulated threads finished.
```

Result: Multi-thread programming enables concurrent execution of multiple threads within a process, thus enhancing performance and efficiency by utilizing system resources effectively.

Ex.No:4 Name:

Date: Reg.No:

Simulation of FCFS CPU scheduling algorithm

```
#include<stdio.h>
#include<conio.h>
int main() {
  int n, bt[10], wt[10] = \{0\}, tat[10], i;
  float total_wt = 0, total_tat = 0;
  printf("Enter number of processes: ");
scanf("%d", &n);
  printf("Enter burst times:\n");
for (i = 0; i < n; i++) scanf("%d",
  &bt[i]); for (i = 1; i < n; i++)
wt[i] = wt[i - 1] + bt[i - 1];
  for (i = 0; i < n; i++)
tat[i] = wt[i] +
bt[i]; total_wt +=
wt[i]; total_tat +=
tat[i];
}
   printf("P\tBT\tWT\tTAT
  n''; for (i = 0; i < n; i++)
printf("%d\t\% d\t\% d\t\% d\t\% d\t\% i, i + 1, bt[i], wt[i], tat[i]);
  printf("Avg WT = \%.2f, Avg TAT = \%.2f\n", total_wt / n,
  total_tat / n); getch(); return 0;
}
```

```
C:\TURBOC3\BIN\TC
Enter number of processes: 5
Enter burst times:
2 5 8 10 14
P BT WT TAT
1 2 0 2
2 5 2 7
3 8 7 15
4 10 15 25
5 14 25 39
Aug WT = 9.80, Aug TAT = 17.60
```

Result: Thus, the simulation of the FCFS (First-Come, First-Served) CPU scheduling algorithm processes jobs sequentially based on their arrival time, ensuring simplicity while risking inefficiencies like the convoy effect.

Ex.No:5 Name:

Date: Reg.No:

Priority CPU scheduling algorithm

```
#include <stdio.h>
#include<conio.h>
int main() {
  int n, i, j, temp;
  int p[10], bt[10], pr[10], wt[10] = \{0\}, tat[10]; float
  total_wt = 0, total_tat = 0;
  printf("Enter number of processes: ");
scanf("%d", &n);
  printf("Enter burst time and priority for each
  process:\n"); for (i = 0; i < n; i++) \{ p[i] = i + 1;
scanf("%d%d", &bt[i],
  &pr[i]); } for (i = 0; i < n)
  - 1; i++)
for (j = i + 1; j < n;
j++) if (pr[i] >
pr[j]) {
    temp = pr[i]; pr[i] =
pr[j]; pr[j] = temp;
    temp = bt[i];
bt[i] = bt[i]; bt[i] =
    temp; temp =
    p[i];
p[i] = p[j]; p[j] = temp;
  }
  for (i = 1; i < n;
i++) wt[i] = wt[i-1]
```

```
+ bt[i - 1];
  for (i = 0; i < n; i++)
 {
tat[i] = wt[i] +
bt[i]; total_wt +=
wt[i]; total_tat +=
tat[i]; }
  printf("P\tBT\tPr\tWT\tTAT\n");
             for (i = 0; i < n; i++)
printf("\%d\t\%d\t\%d\t\%d\t\%d\t\%d\n", p[i], bt[i], pr[i], wt[i], tat[i]);
  printf("Avg WT = \%.2f, Avg TAT = \%.2f\n", total_wt / n,
  total_tat / n);
getch();
  return
  0;
}
```

```
C:\TURBOC3\BIN>TC
Enter number of processes: 3
Enter burst time and priority for each process:
4 6 9
2 1 3
P BT Pr WT TAT
2 9 2 0 9
3 1 3 9 10
1 4 6 10 14
Avg WT = 6.33, Avg TAT = 11.00
```

Result: Thus, the simulation of the Priority CPU scheduling algorithm executes processes based on their priority levels, ensuring critical tasks are handled first, but may lead to starvation of lower-priority processes.

Ex.No:6 Name:

Date: Reg.No:

Simulation of Round Robin CPU scheduling algorithm

```
#include<stdio.h>
#include<conio.h>
int main() {
  int n, tq, bt[10], wt[10] = \{0\}, tat[10], rem_bt[10], time = 0,
  done; float total_wt = 0, total_tat = 0;
  printf("Enter number of processes: ");
scanf("%d", &n);
  printf("Enter
                          burst
times:\n"); for (int i = 0; i < n;
i++)
{
scanf("%d", &bt[i]); rem_bt[i] = bt[i];
}
  printf("Enter time quantum: "); scanf("%d", &tq);
  do {
    done = 1;
    for (int i = 0; i < n; i++)
{
      if (rem_bt[i] >
        0) { done = 0;
        if (rem_bt[i] > tq)
{
time += tq; rem_bt[i] -= tq;
}
        else
```

{

```
time+=
rem_bt[i]; wt[i]
= time - bt[i];
rem_bt[i] = 0;}
      }
    }
  } while (!done);
  for (int i = 0; i < n; i++)
{
tat[i] = wt[i] + bt[i];
total_wt += wt[i];
total_tat += tat[i];
}
  printf("P\tBT\tWT\tTAT\
  n"); for (int i = 0; i < n;
  i++)
printf("\%d\t\%d\t\%d\t\%', i+1, bt[i], wt[i], tat[i]);
  printf("Avg WT = \%.2f, Avg TAT = \%.2f\n", total_wt / n,
total_tat / n); getch();
  return 0;
}
```

```
C:\TURBOC3\BIN>TC
Enter number of processes: 3
Enter burst times:
5 9 12
Enter time quantum: 3
P BT WT TAT
1 5 6 11
2 9 11 20
3 12 14 26
Avg WT = 10.33, Avg TAT = 19.00
```

Result: Thus, the simulation of the Round Robin CPU scheduling algorithm processes tasks in a cyclic order with a fixed time quantum, ensuring fairness but possibly increasing context switching overhead.

Ex.No:7 Name:

Date: Reg.No:

WRITE A PROCEDURE FOR TIMER INTERRUPT HANDLER

```
#include <dos.h>
#include <stdio.h>
#include <conio.h>
void interrupt (*oldTimerISR)(); // Pointer to store the old interrupt service routine
void
                 interrupt
  newTimerISR() { static
  unsigned int ticks = 0;
  ticks++;
  if (ticks % 18 == 0) { // Approximately 1-second interval (18.2 ticks/sec)
    printf("Timer Interrupt Triggered! %u seconds elapsed.\n", ticks / 18);
  }
  (*oldTimerISR)(); // Call the original ISR to maintain timer functionality
void setupTimerInterrupt() {
  oldTimerISR = getvect(0x08); // Save the old ISR for interrupt 0x08 (timer
  interrupt) setvect(0x08, newTimerISR); // Set our custom ISR
void restoreTimerInterrupt() {
  setvect(0x08, oldTimerISR); // Restore the original ISR
int main() {
  printf("Installing custom timer interrupt handler...\n");
  setupTimerInterrupt();
```

```
printf("Custom timer handler installed. Press any key to
stop...\n"); while (!kbhit()); // Wait for key press
restoreTimerInterrupt();
printf("Original timer interrupt handler restored.
Exiting.\n"); getch();
return 0;
}
```

```
C:\TURBOC3\BIN>TC
Installing custom timer interrupt handler...
Custom timer handler installed. Press any key to stop...
Timer Interrupt Triggered! 1 seconds elapsed.
Timer Interrupt Triggered! 2 seconds elapsed.
Timer Interrupt Triggered! 3 seconds elapsed.
Timer Interrupt Triggered! 4 seconds elapsed.
Timer Interrupt Triggered! 5 seconds elapsed.
Timer Interrupt Triggered! 6 seconds elapsed.
Timer Interrupt Triggered! 7 seconds elapsed.
Timer Interrupt Triggered! 8 seconds elapsed.
Timer Interrupt Triggered! 9 seconds elapsed.
Timer Interrupt Triggered! 10 seconds elapsed.
Timer Interrupt Triggered! 11 seconds elapsed.
Timer Interrupt Triggered! 12 seconds elapsed.
Original timer interrupt handler restored. Exiting.
Installing custom timer interrupt handler...
Custom timer handler installed. Press any key to stop...
Timer Interrupt Triggered! 1 seconds elapsed.
Timer Interrupt Triggered! 2 seconds elapsed.
Timer Interrupt Triggered! 3 seconds elapsed.
Timer Interrupt Triggered! 3 seconds elapsed.
Timer Interrupt Triggered! 4 seconds elapsed.
Timer Interrupt Triggered! 4 seconds elapsed.
```

Result: Successfully implemented a timer interrupt handler that executes a predefined task upon timer expiration.

Ex.No:8 Name:

Date: Reg.No:

CLASSICAL INTER PROCESS COMMUNICATION PROBLEM (PRODUCER CONSUMER)

```
#include <stdio.h>
#include <conio.h>
#define
BUFFER_SIZE 5
int buffer[BUFFER_SIZE], item; // Shared buffer
int in = 0, out = 0, count = 0; // Indices and
counter void producer() {
  if (count == BUFFER_SIZE) {
    printf("Buffer full! Producer
    waiting...\n"); return;
  }
  item = (in + 1) * 10; // Produce an item
  buffer[in] = item;
  in = (in + 1) \% BUFFER_SIZE;
  count++;
  printf("Produced: %d\n", item);
}
void
  consumer() {
  if (count ==
  0) {
    printf("Buffer empty! Consumer waiting...\n");
    return;
  item = buffer[out];
  out = (out + 1) \% BUFFER\_SIZE;
  count--;
```

```
printf("Consumed: %d\n", item);
  }
 int main()
    { int
   choice;
   do {
      printf("\n1. Produce\n2. Consume\n3. Exit\n");
printf("Enter your choice: "); scanf("%d", &choice); switch (choice) {
      case 1:
        producer();
        break;
      case 2:
        consumer
        (); break;
      case 3:
        printf("Exiting...\n");
        break;
      default:
        printf("Invalid choice!\n");
      }
    } while (choice !=
   3); getch();
   return 0;
  }
```

```
1. Produce
2. Consume
3. Exit
Enter your choice: 1
Produced: 20
1. Produce
2. Consume
3. Exit
Enter your choice: 1
Produced: 30
1. Produce
2. Consume
3. Exit
Enter your choice: 2
Consumed: 10

    Produce
    Consume

3. Exit
Enter your choice: 3
Exiting...
```

Result: This Program implemented the Producer-Consumer problem using inter-process communication, ensuring synchronized data exchange between processes.

Ex.No:9	Name:	

Reg.No:

Date:

WRITE A PROCEDURE TO MAKE MESSAGE PASSING IN INTER PROCESS COMMUNICATION

```
#include<stdio.h>
#include<conio.h>
#include <string.h>
#define BUFFER_SIZE 100
char buffer[BUFFER_SIZE]; // Shared memory buffer
int message_flag = 0; // Message availability flag (0 = empty, 1 = full)
// Function to simulate the sender
process void sender() {
  if (message\_flag == 1) {
    printf("Buffer full! Sender
    waiting...\n"); return;
  }
  printf("Enter a message to send: ");
  fflush(stdin); // Clear input buffer
  gets(buffer); // Read message into
  buffer message_flag = 1; // Mark buffer
  as full printf("Message sent
  successfully!\n");
}
// Function to simulate the receiver
process void receiver() {
  if (message\_flag == 0) {
    printf("Buffer empty! Receiver
    waiting...\n"); return;
```

```
printf("Received message: %s\n", buffer);
  strcpy(buffer, ""); // Clear the buffer
  message_flag = 0; // Mark buffer as empty
  printf("Message received successfully!\n");
}
int main()
  { int
  choice;
  do {
    printf("\n1. Send Message\n2. Receive Message\n3. Exit\n");
    printf("Enter your choice: ");
    scanf("%d", &choice);
    switch (choice) { case 1:
      sender
      ();
      break;
    case 2:
      receiver();
      break;
    case 3:
      printf("Exiting...\n");
      break;
    default:
      printf("Invalid choice!\n");
  } while (choice != 3);
  getch(); // Wait for key press before exiting
  return 0;}
```

```
C:\TURBOC3\BIN>TC
1. Send Message
2. Receive Message
3. Exit
Enter your choice: 1
Enter a message to send: Hello SRM Message sent successfully!
1. Send Message
2. Receive Message
3. Exit
Enter your choice: 2
Received message: Hello SRM
Message received successfully!
1. Send Message
2. Receive Message
3. Exit
Enter your choice: 3
Exiting...
```

Result: Thus the program, implemented message passing in IPC, enabling efficient and synchronized communication between processes.

Ex.No:10 Name:

Date: Reg.No:

PROGRAM TO IMPLEMENT BANKERS' ALGORITHM

```
#include <stdio.h>
#include <conio.h>
#define MAX 10
int allocation[MAX][MAX], max[MAX][MAX], need[MAX][MAX],
available[MAX]; int n, m; // Number of processes (n) and resources (m)
void input()
  { int i, j;
  printf("Enter the number of processes: ");
  scanf("%d", &n);
  printf("Enter the number of resources: ");
  scanf("%d", &m);
  printf("Enter
                         allocation
  matrix:\n"); for (i = 0; i < n; i++)
    for (j = 0; j < m; j++)
      scanf("%d",
      &allocation[i][j]);
  printf("Enter
                        maximum
  matrix:\n"); for (i = 0; i < n; i++)
    for (j = 0; j < m; j++)
      scanf("%d",
      &max[i][j]);
 printf("Enteravailable resources:\n"); for (i = 0; i < m; i++)
```

```
scanf("%d", &available[i]);
            // Calculate the need matrix for (i = 0; i < n; i++)
    for (j = 0; j < m; j++)
      need[i][j] = max[i][j] - allocation[i][j];
}
int is_safe() {
  int work[MAX], finish[MAX] = \{0\}, safe_seq[MAX],
  count = 0; int i, j;
  for (i = 0; i < m; i++)
    work[i]
    available[i];
  while (count < n)
    \{ \text{ int found} = 0; 
    for (i = 0; i < n;
      i++) { if
      (!finish[i]) {
        for (j = 0; j < m; j++)
           if (need[i][j]
             work[j]) break;
        if (j == m) {
           for (j = 0; j < m; j++)
             work[j]
           allocation[i][j];
           safe\_seq[count++] = i;
           finish[i] = 1;
           found = 1;
         }
      }
```

```
}
 if (!found) {
      printf("System is in an unsafe state.\n");
      return 0;
    }
  }
  printf("System is in a safe state.\nSafe sequence is:
  "); for (i = 0; i < n; i++)
    printf("P%d
  safe_seq[i]); printf("\n");
  return 1;
int main()
{
input();
is_safe();
getch();
return 0;
}
```

```
C:\TURBOC3\BIN>TC
Enter the number of processes: 5
Enter the number of resources: 3
Enter allocation matrix:
0 1 0
2 0 0
3 0 2
2 1 1
0 0 2
Enter maximum matrix:
7 5 3
3 2 2
9 0 2
2 2 2
4 3 3
Enter available resources:
3 3 2
System is in a safe state.
Safe sequence is: P1 P3 P4 P0 P2
```

Result: Successfully implemented the Banker's Algorithm, ensuring safe resource allocation and deadlock avoidance.

Ex.No:11 Name:

Date: Reg.No:

PROGRAM TO IMPLEMENT MEMORY ALLOCATION WITH PAGES

```
#include <stdio.h>
#include <conio.h>
#define MAX 10 // Maximum number of pages or frames
int main() {
  int page_table[MAX], frames[MAX], n_pages, n_frames, i, page, frame;
  printf("Enter the number of pages (max %d): ",
  MAX); scanf("%d", &n_pages);
  printf("Enter the number of frames (max %d): ",
  MAX); scanf("%d", &n_frames);
  if (n_pages > MAX || n_frames > MAX) {
    printf("Number of pages or frames exceeds
    limit!\n");
    return 1;
  // Initialize all frames to -1
  (empty) for (i = 0; i <
  n_frames; i++)
    frames[i] = -1;
  // Simulate page table mapping
  printf("Enter the frame numbers for each page (0 to %d):\n", n_frames
  - 1); for (i = 0; i < n_pages; i++) {
    printf("Page %d -> Frame: ", i);
    scanf("%d", &frame);
    if (frame < 0 \parallel frame >= n_frames) {
      printf("Invalid frame number!\n");
```

```
return 1;
  }
  page_table[i] = frame;
  frames[frame] = i; // Map frame to page
}
// Display page table
printf("\nPage
Table:\n");
printf("Page\tFrame\n")
; for (i = 0; i < n_pages;
i++) {
  printf("%d\t%d\n", i, page_table[i]);
}
// Simulate memory access
printf("\nEnter a page number to access (-1 to
exit): "); while (1) {
  scanf("%d",
  &page); if (page
  == -1) {
    printf("Exiting...\n");
    break;
  }
  if (page < 0 \parallel page >= n_pages) {
    printf("Invalid page number!\n");
  } else {
    printf("Page %d is in Frame %d.\n", page, page_table[page]);
  }
getch();
return 0;
```

Result: Successfully implemented memory allocation using paging, ensuring efficient memory management and address translation.

Ex.No:12 Name:

Date: Reg.No:

SIMULATION OF FIFO PAGE REPLACEMENT ALGORITHM

```
#include <stdio.h>
#include <conio.h>
#define MAX FRAMES 10
#define MAX_PAGES 50
int main() {
   int frames[MAX_FRAMES],
   pages[MAX_PAGES]; int n_frames,
   n_pages, i, j, k, page_faults = 0;
   int index = 0; // Points to the oldest frame for replacement
   printf("Enter the number of frames (max %d): ",
   MAX_FRAMES); scanf("%d", &n_frames);
   printf("Enter the number of pages (max %d): ",
   MAX_PAGES); scanf("%d", &n_pages);
   printf("Enter the reference string (page
   numbers):\n"); for (i = 0; i < n_pages; i++) \{
     scanf("%d", &pages[i]);
   }
   // Initialize frames to -1 (indicating
   empty) for (i = 0; i < n_{frames}; i++) \{
     frames[i] = -1;
   }
   // Simulate FIFO Page Replacement for (i = 0; i < n_pages; i++) {
     int found = 0;
     // Check if the page is already in one of the
     frames for (j = 0; j < n_frames; j++) {
       if (frames[j] ==
```

```
pages[i]) { found = 1;
      break;
    }
  }
  if (!found) { // Page fault occurs
    frames[index] = pages[i]; // Replace the oldest page
    index = (index + 1) % n_frames; // Update the index to point to the next
    frame page_faults++;
  }
  // Display current frame contents
  printf("Page %d -> Frames: ", pages[i]); for (k = 0; k < n_frames; k++) {
    if (frames[k] != -1)
      printf("%d",
    frames[k]); else
      printf("- ");
  }
  printf("\n");
}
printf("\nTotal Page Faults: %d\n",
page_faults); getch();
return 0;
```

}

```
C:\TURBOC3\BIN>TC
Enter the number of frames (max 10): 3
Enter the number of pages (max 50): 5
Enter the reference string (page numbers):
7 0 1 2 3
Page 7 -> Frames: 7 - -
Page 0 -> Frames: 7 0 -
Page 1 -> Frames: 7 0 1
Page 2 -> Frames: 2 0 1
Page 3 -> Frames: 2 3 1

Total Page Faults: 5
```

Result: Thus, the simulated the FIFO page replacement algorithm, demonstrating page faults and efficient memory management.

Name:

Date: Reg.No:

MULTIPLE PARTITION (DYNAMIC MEMORY ALLOCATION METHOD)

```
#include <stdio.h>
#include <conio.h>
#define MAX 10
typedef struct
{
int size;
  int allocated; int process;
} Partition;
int main() {
  Partition memory[MAX];
  int n_partitions, n_processes, process[MAX], i, j;
  printf("Enter the number of memory partitions (max %d): ",
  MAX); scanf("%d", &n_partitions);
  printf("Enter the sizes of the
  partitions:\n"); for (i = 0; i <
  n_partitions; i++) {
    printf("Partition %d size: ", i +
    1); scanf("%d",
    &memory[i].size);
    memory[i].allocated = 0;
    memory[i].process = -1;
  printf("Enter the number of processes (max %d): ",
  MAX); scanf("%d", &n_processes);
  printf("Enter the sizes of the processes:\n");
```

```
for (i = 0; i < n_processes; i++) {
    printf("Process %d size: ", i +
    1); scanf("%d", &process[i]);
 // Allocate processes using First
 Fit for (i = 0; i < n_processes;
 i++) {
   for (j = 0; j < n_partitions; j++)
      if (!memory[j].allocated && memory[j].size >= process[i])
        { memory[i].allocated = 1;
        memory[j].process = i + 1;
        printf("Process %d allocated to Partition %d\n", i + 1, j + 1);
        break;
      }
    }
    if (j == n_partitions) {
      printf("Process %d cannot be allocated (Insufficient memory).\n", i + 1);
    }
  }
 // Display memory allocation
 printf("\nPartition\tSize\tAllocated\tProcess\
 n"); for (i = 0; i < n_partitions; i++) {
    printf("%d\t\t%d\t%s\t\t", i + 1, memory[i].size, memory[i].allocated? "Yes"
    : "No"); if (memory[i].allocated) {
      printf("P%d\n", memory[i].process);
    } else {
      printf("-\n");
} }
 getch();
 return
 0;
}
```

```
Enter the number of memory partitions (max 10): 5
Enter the sizes of the partitions:
Partition 1 size: 50
Partition 2 size: 100
Partition 3 size: 500
Partition 4 size: 350
Partition 5 size: 400
Enter the number of processes (max 10): 4 Enter the sizes of the processes:
Process 1 size: 75
Process 2 size: 50
Process 3 size: 100
Process 4 size: 112
Process 1 allocated to Partition 2
Process 2 allocated to Partition 1
Process 3 allocated to Partition 3
Process 4 allocated to Partition 4
                              Allocated
                    Size
Partition
                                                   Process
                    50
                              Yes
                                                   PZ
2 3
                     100
                              Yes
                                                   P1
                    500
                                                   Р3
                               Yes
4
                    350
                               Yes
                                                   P4
                     400
                              No
```

Result: Successfully implemented multiple partition dynamic memory allocation, thus optimizing memory utilization and reducing fragmentation.

Ex.No:14 Name:

Date: Reg.No:

SIMULATION OF LRU PAGE REPLACEMENT ALGORITHM

```
#include <stdio.h>
#include <conio.h>
#define MAX_FRAMES 10
#define MAX_PAGES 50
int findLRU(int time[], int
  n) { int i, min = time[0],
  pos = 0; for (i = 1; i < n;
 i++) {
   if (time[i] <
     min) \{ min =
     time[i]; pos
     = i;
    }
 return pos;
}
int main() {
  int frames[MAX_FRAMES], pages[MAX_PAGES],
  time[MAX_FRAMES]; int n_frames, n_pages, i, j, pos,
  page_faults = 0, counter = 0, flag1, flag2;
  printf("Enter the number of frames (max %d): ",
  MAX_FRAMES); scanf("%d", &n_frames);
  printf("Enter the number of pages (max %d): ",
  MAX_PAGES); scanf("%d", &n_pages);
  printf("Enter the reference string (page
```

```
numbers): \label{eq:numbers} $$ numbers): \n"); for (i = 0; i < n\_pages; i++) \{ $$ scanf("%d", &pages[i]); $$
```

```
}
// Initialize frames and time
arrays for (i = 0; i < n_frames;
i++) {
  frames[i] = -1;
  time[i] = 0;
}
// Simulate LRU Page
Replacement for (i = 0; i <
n_pages; i++) {
  flag1 = flag2 = 0;
 // Check if the page is already in a
  frame for (j = 0; j < n_frames; j++) {
    if (frames[j] ==
      pages[i]) { flag1 =
      flag2 = 1;
      time[j] = ++counter; // Update time for the page
      break;
    }
  }
  // If the page is not in a
  frame if (!flag1) {
    for (j = 0; j < n_frames; j++) {
      if (frames[j] == -1) \{ // Check for an empty \}
        frame frames[j] = pages[i];
        time[j] =
        ++counter;
        page_faults++;
        flag2 = 1;
        break;
```

```
}
        }
      }
     // If no empty frame is available, replace the least recently used
      page if (!flag2) {
pos = findLRU(time, n_frames); frames[pos] = pages[i]; time[pos] = ++counter; page_faults++;
      // Display the current state of frames
      printf("Page %d -> Frames: ",
      pages[i]); for (j = 0; j < n_frames;
     j++) {
        if (frames[j] != -1)
          printf("%d",
        frames[j]); else
          printf("- ");
      }
      printf("\n");
   printf("\nTotal Page Faults: %d\n",
    page_faults); getch();
    return 0;
  }
```

```
C:\TURBOC3\BIN\TC
Enter the number of frames (max 10): 3
Enter the number of pages (max 50): 5
Enter the reference string (page numbers):

1
2
0
7
3
Page 1 -> Frames: 1 - -
Page 2 -> Frames: 1 2 -
Page 0 -> Frames: 1 2 0
Page 7 -> Frames: 7 2 0
Page 3 -> Frames: 7 3 0

Total Page Faults: 5
```

Result: Successfully simulated the LRU page replacement algorithm, thus optimizing page fault handling and memory management.

Ex.No:15 Name:

Date: Reg.No:

SIMULATION OF OPTIMAL PAGE REPLACEMENT ALGORITHM

```
#include
            <stdio.h>
#include <conio.h>
#define MAX_FRAMES 10
#define MAX_PAGES 50
// Function to find the page to replace
int findOptimal(int pages[], int frames[], int n_pages, int current, int n_frames) { int farthest
  = current, pos = -1, i, j;
  for (i = 0; i < n_frames;
    i++) { int found = 0;
    for (j = current + 1; j < n_pages;
      j++) { if (frames[i] == pages[j])
      {
        if
              (i >
          farthest) {
          farthest = j;
          pos = i;
        }
        found = 1;
        break;
      }
    if (!found) return i; // If the frame is never used in the future
  }
  return pos == -1 ? 0 : pos; // Replace the least recently needed page
}
int main() {
```

```
int frames[MAX_FRAMES], pages[MAX_PAGES];
int n_frames, n_pages, i, j, page_faults = 0, flag1, flag2;
   printf("Enter
                  the number of frames
                                                (max
                                                        %d): ",
   MAX_FRAMES); scanf("%d", &n_frames);
   printf("Enter the number of pages (max %d): ", MAX_PAGES);
   scanf("%d", &n_pages);
   printf("Enter
                   the
                         reference
                                     string
                                              (page
   numbers):\n"); for (i = 0; i < n_pages; i++) {
     scanf("%d", &pages[i]);
   }
   // Initialize frames
   for (i = 0; i < n_frames; i++)
     { frames[i] = -1;
   }
   //
         Simulate
                     Optimal
                                 Page
   Replacement for (i = 0; i < n_pages;
   i++) {
     flag1 = flag2 = 0;
     // Check if the page is already in a
     frame for (j = 0; j < n_frames; j++) {
       if
              (frames[j]
         pages[i]) \{ flag1 = 1;
         break;
       }
     // If the page is not in a
     frame if (!flag1) {
```

```
for (j = 0; j < n_frames; j++) {
      if (frames[j] == -1) \{ // Check for an empty frame \}
        frames[j] = pages[i];
        page_faults
        ++; flag2 =
        1;
        break;
      }
    }
  }
  // If no empty frame is available, replace the optimal
  page if (!flag2 && !flag1) {
    int pos = findOptimal(pages, frames, n_pages, i, n_frames);
    frames[pos] = pages[i];
    page_faults++;
  }
  // Display the current state of frames
  printf("Page %d -> Frames: ",
  pages[i]); for (j = 0; j < n_frames;
  j++) {
    if (frames[j] != -1)
      printf("%d
    frames[j]); else
      printf("- ");
  }
  printf("\n");
printf("\nTotal
                  Page
                           Faults:
                                     %d\n'',
page_faults); getch();
return 0;
```

}

```
C:\TURBOC3\BIN>TC
Enter the number of frames (max 10): 4
Enter the number of pages (max 50): 3
Enter the reference string (page numbers):
2 5 3
Page 2 -> Frames: 2 - - -
Page 5 -> Frames: 2 5 - -
Page 3 -> Frames: 2 5 3 -

Total Page Faults: 3
```

Result: This program successfully simulated the Optimal Page Replacement Algorithm, achieving minimal page faults and efficient memory management.