## PCFGs: Parsing & Evaluation

LING 571 — Deep Processing Techniques for NLP
October 14, 2020
Shane Steinert-Threlkeld

## Roadmap

- CKY + back-pointers
- PCFGs
- PCFG Parsing (PCKY)
- Inducing a PCFG
- Evaluation
- [Earley parsing]
- HW3 + collaboration

## CKY Parsing: Backpointers

### Backpointers

- Instead of list of possible nonterminals for that node, each cell should have:
  - Nonterminal for the node
  - Pointer to left and right children cells
    - Either direct pointer to cell, or indices

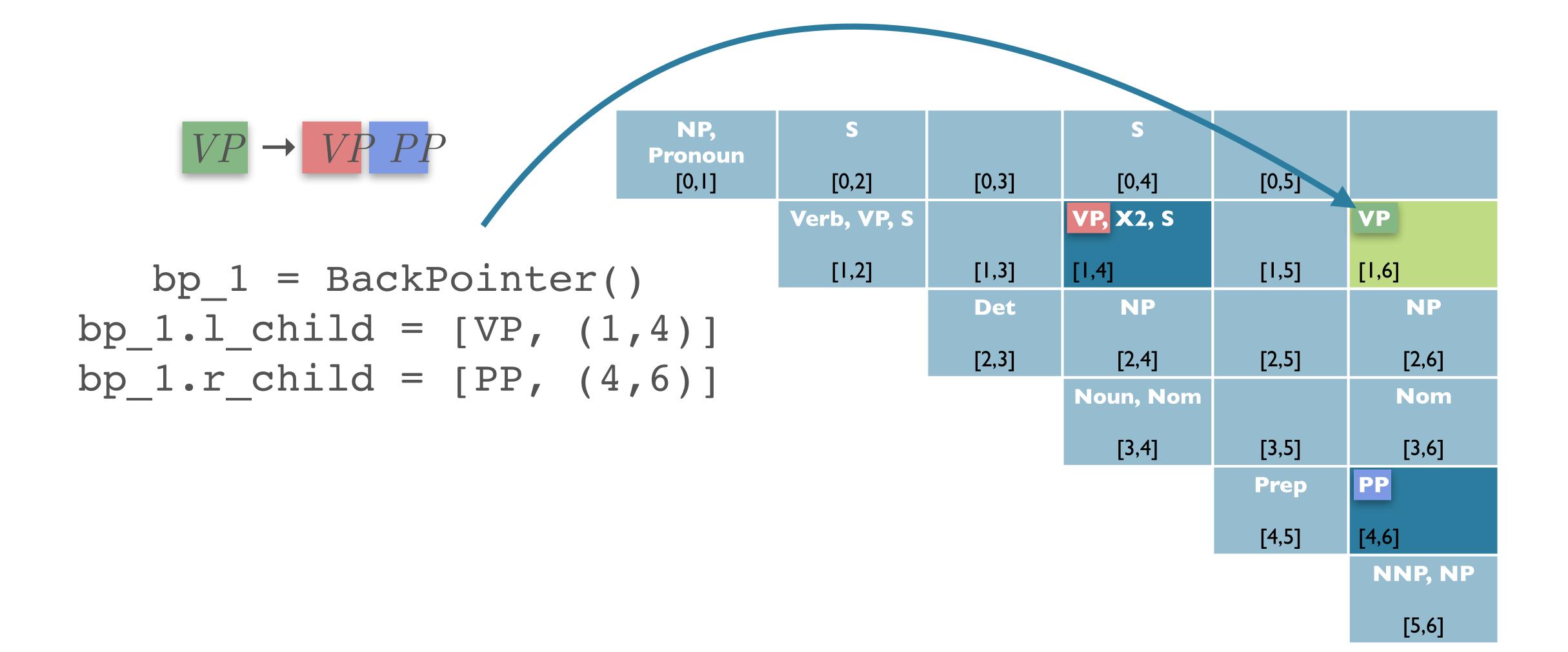
#### For example:

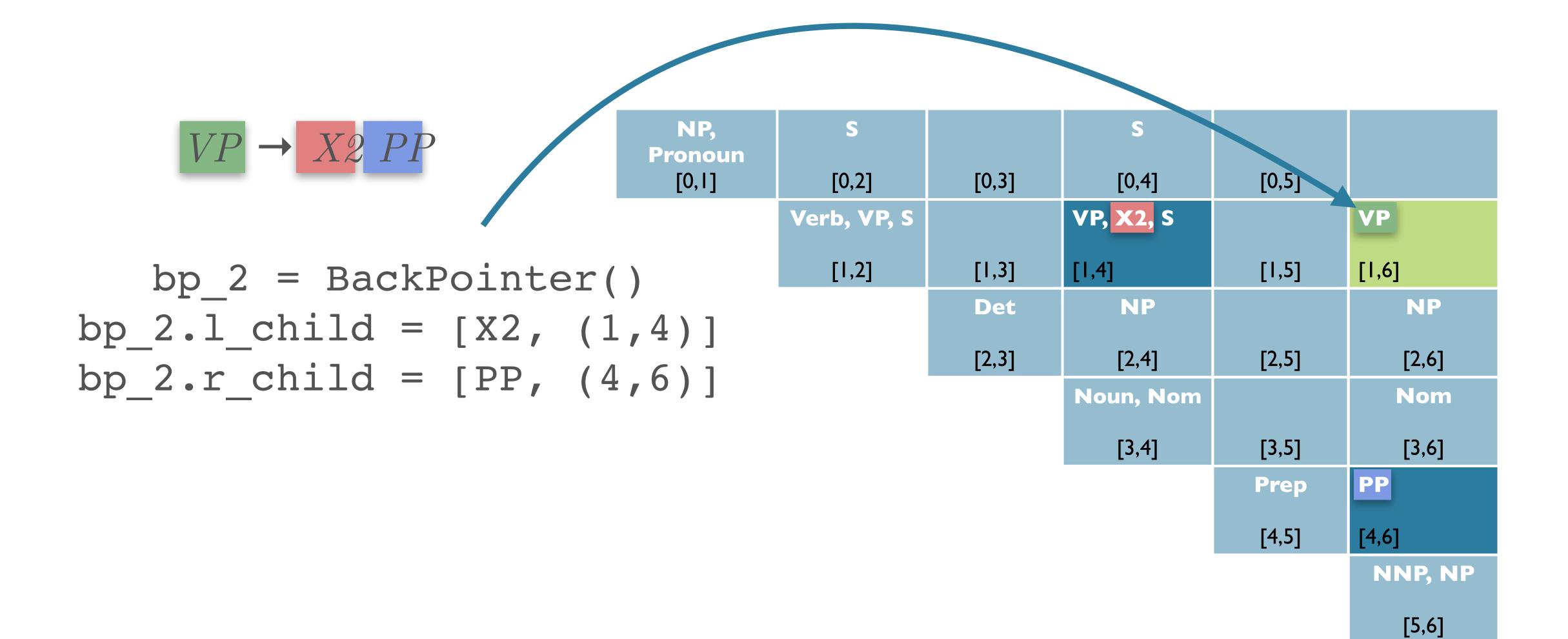
```
bp_2 = BackPointer()
bp_2.l_child = [X2, (1,4)]
bp_2.r_child = [PP, (4,6)]
```

#### CKY Parser

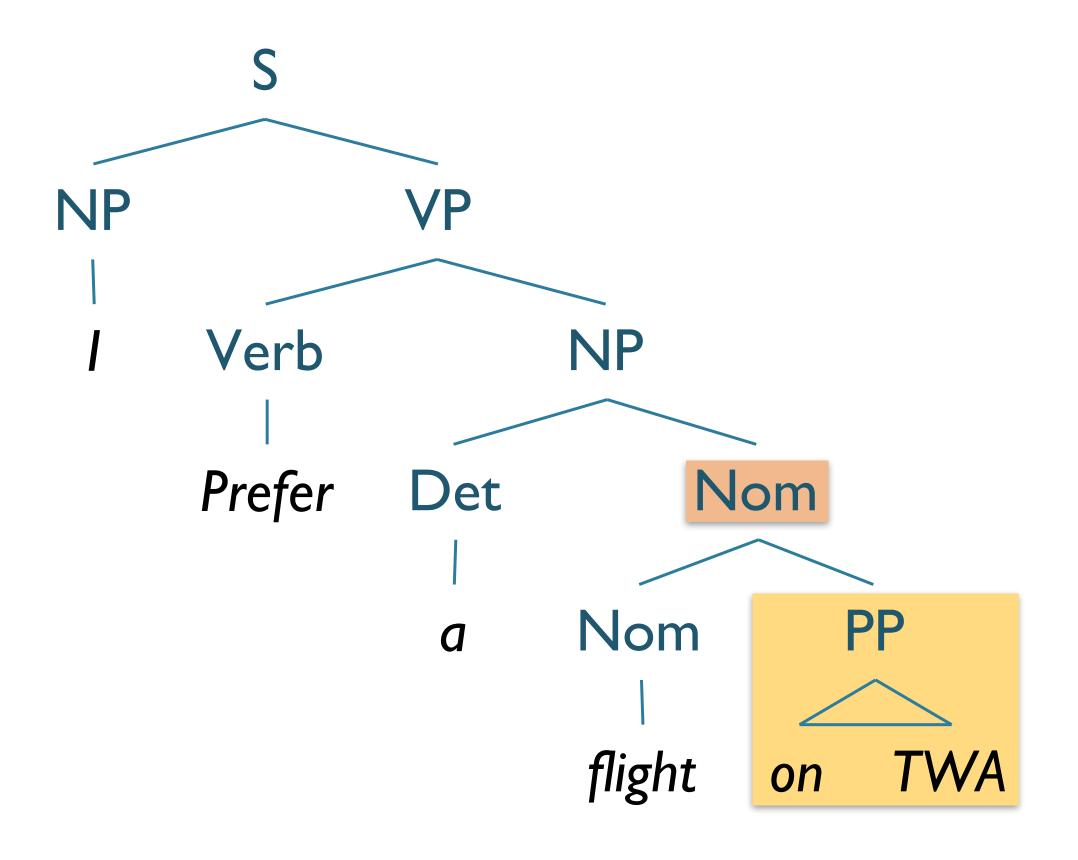
- Pair each nonterminal with back-pointer to cells from which it was derived
- Last step:
  - construct trees from back-pointers in [0, n]

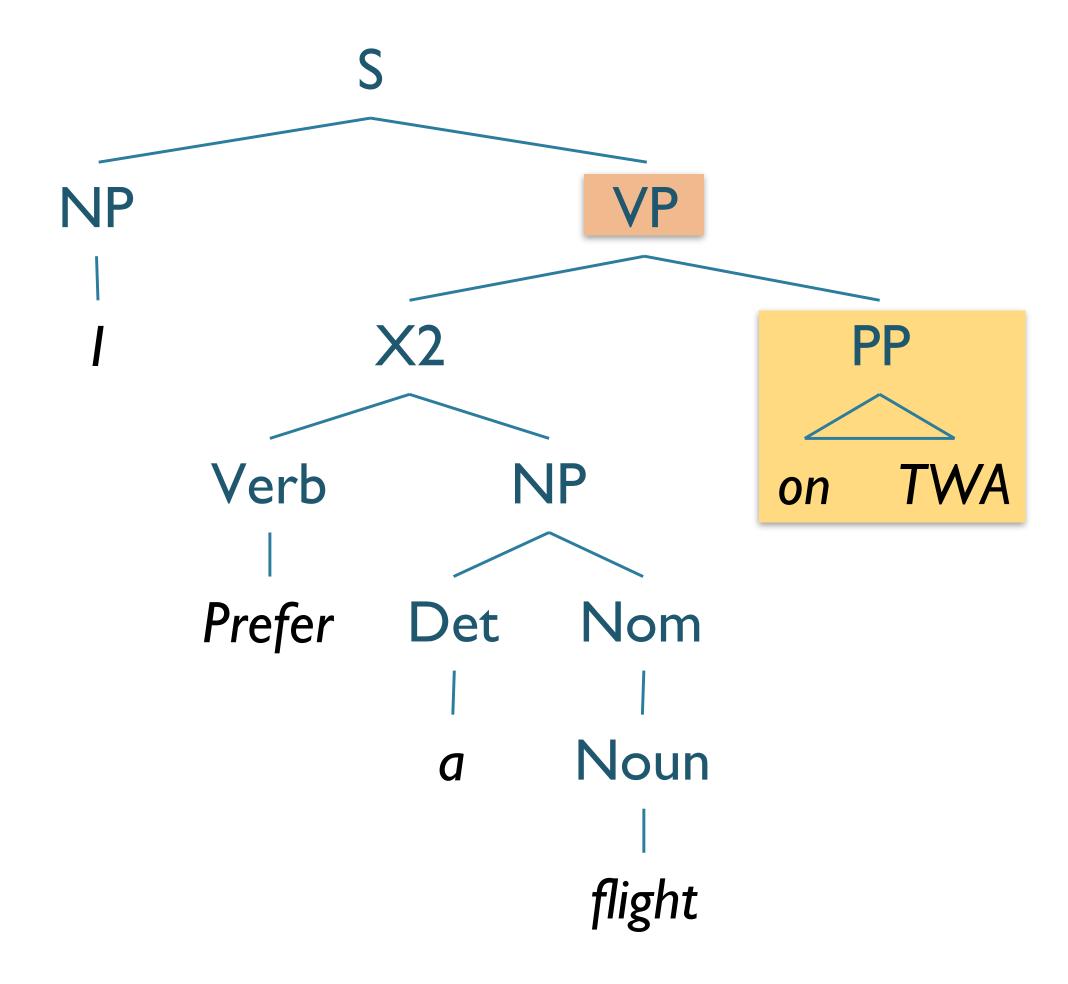
NP, Pronoun	S		S		
[0,1]	[0,2]	[0,3]	[0,4]	[0,5]	
	Verb, VP, S		VP, X2, S		VP
	[1,2]	[1,3]	[1,4]	[1,5]	[1,6]
		Det	NP		NP
		[2,3]	[2,4]	[2,5]	[2,6]
			Noun, Nom		Nom
			[3,4]	[3,5]	[3,6]
				Prep	PP
				[4,5]	[4,6]
					NNP, NP
					[5,6]





## Resulting Parses

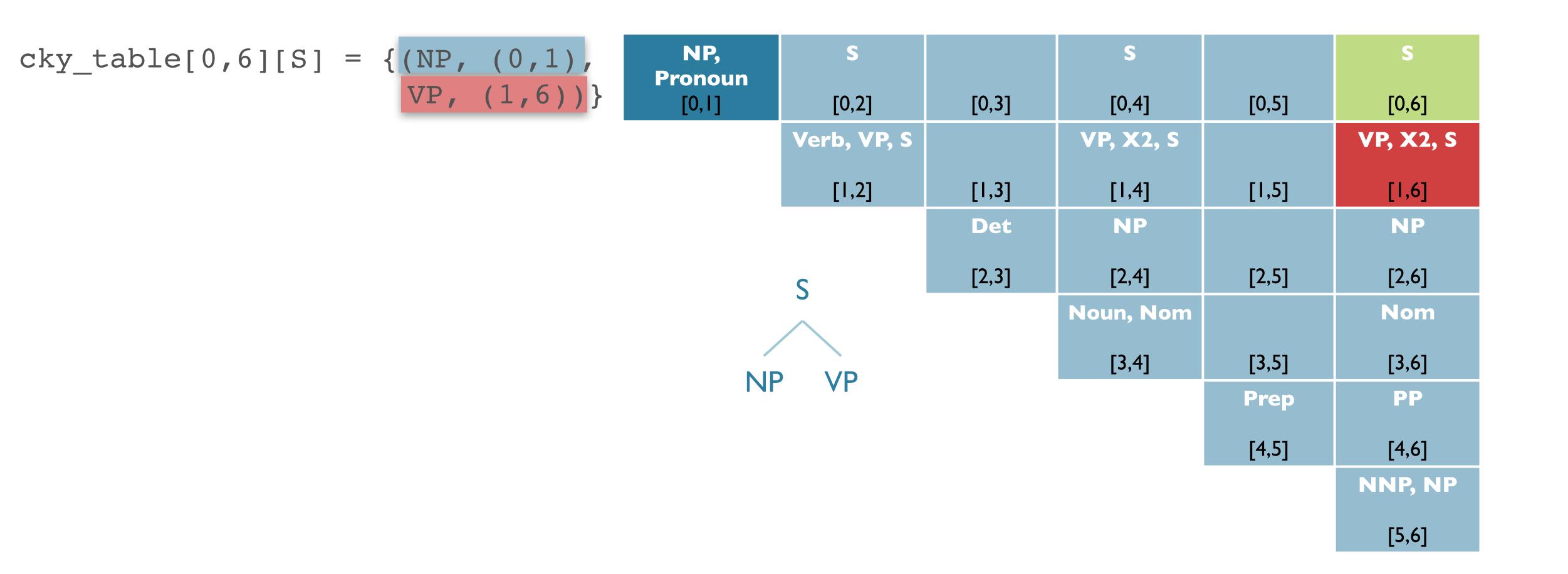


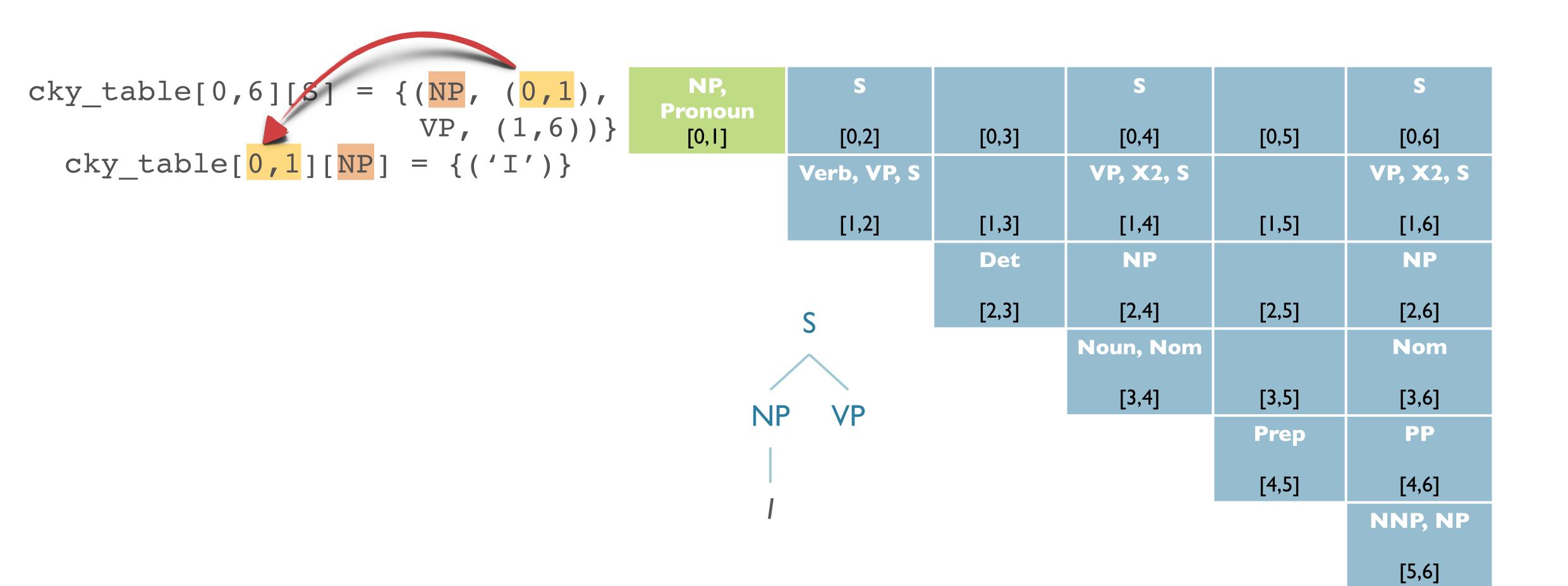


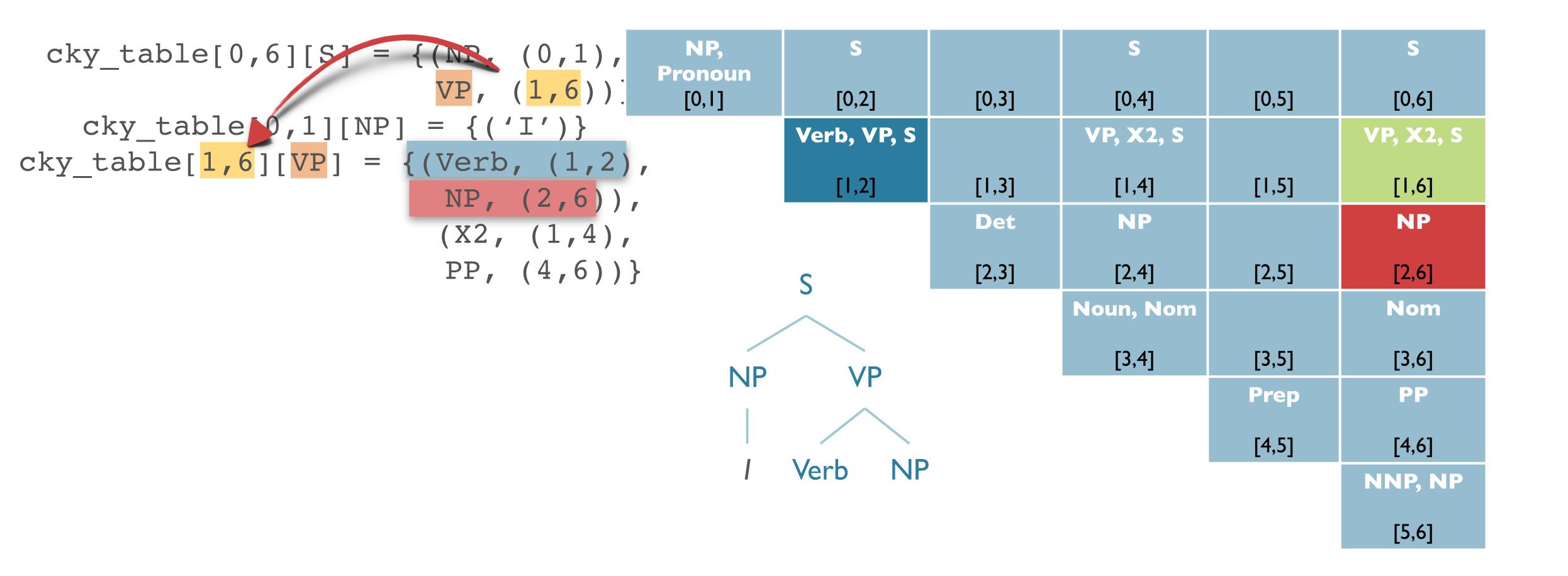
#### CKY Discussion

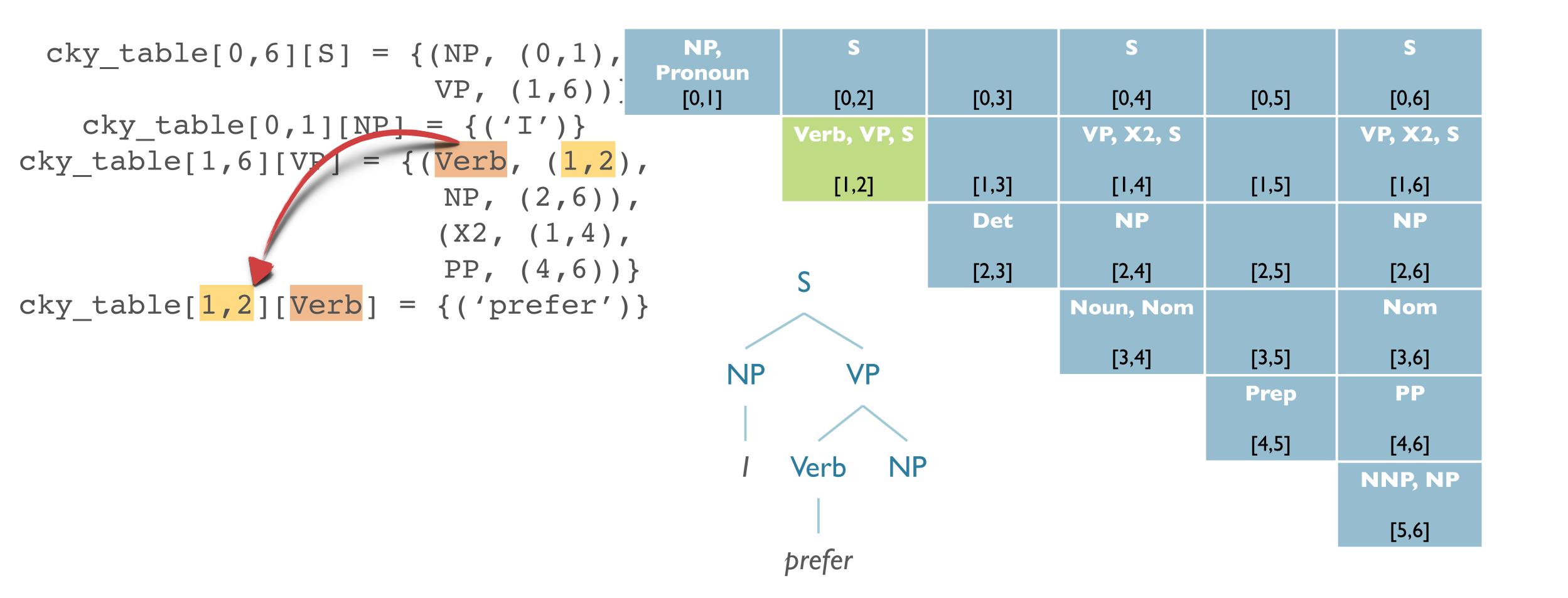
- Running time:
  - ullet  $O(n^3)$  where n is the length of the input string
  - Inner loop grows as square of # of non-terminals
- Expressiveness:
  - As implemented, requires CNF
    - Weak equivalence to original grammar
    - Doesn't capture full original structure
    - Back-conversion?
      - Can do binarization, terminal conversion
    - Unit productions requires change in CKY

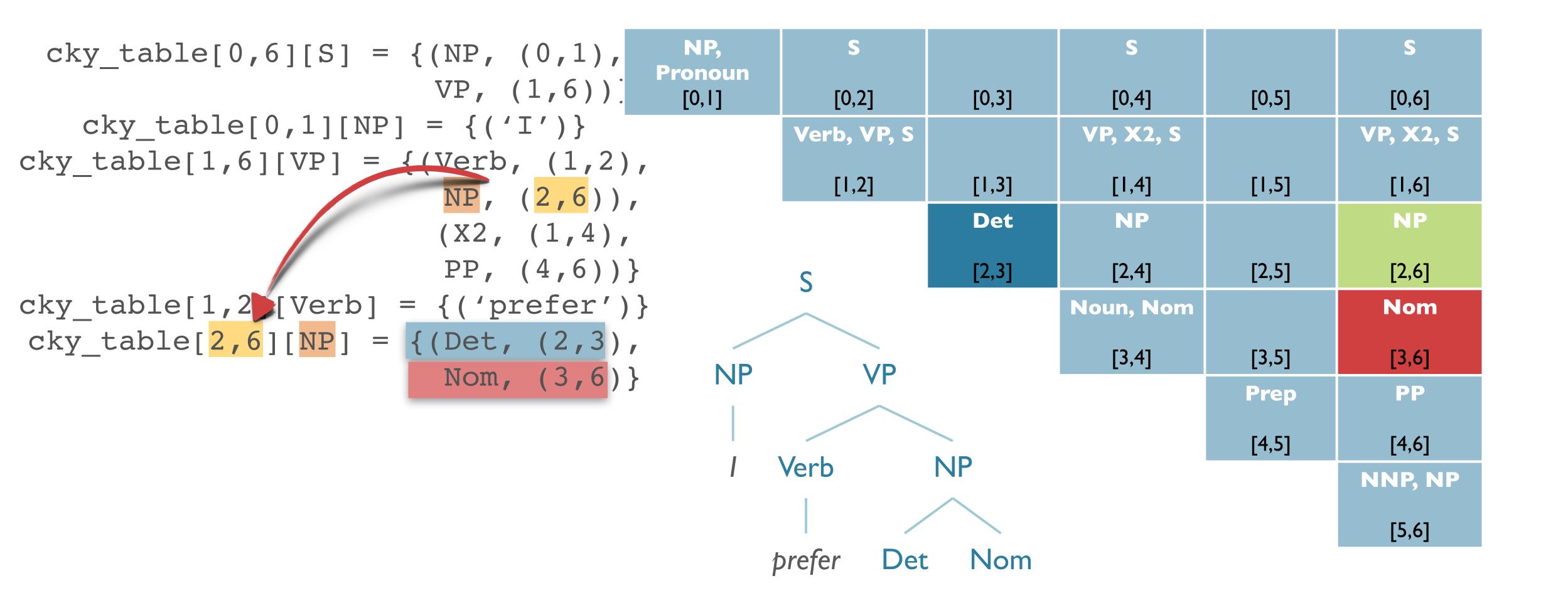
## CKY + Back-pointers Example

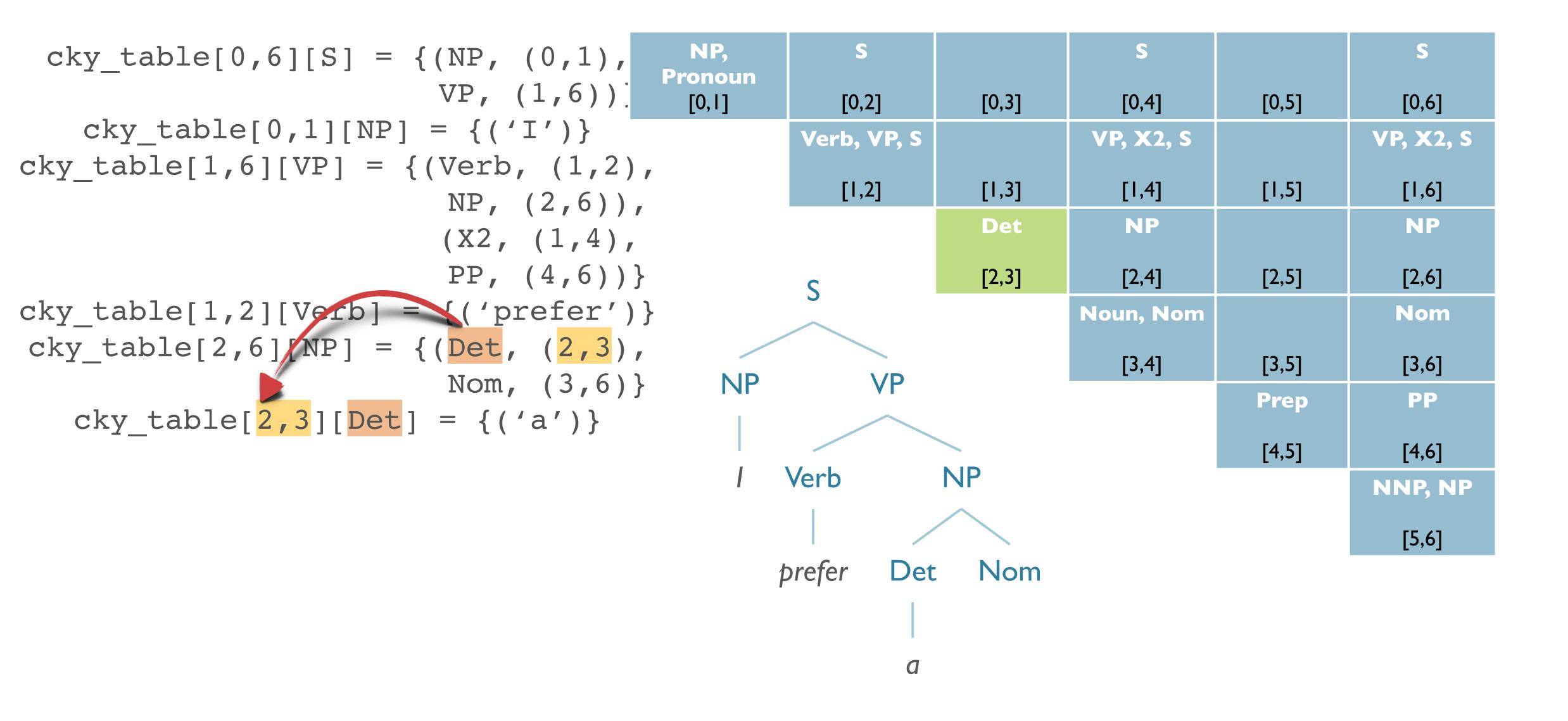


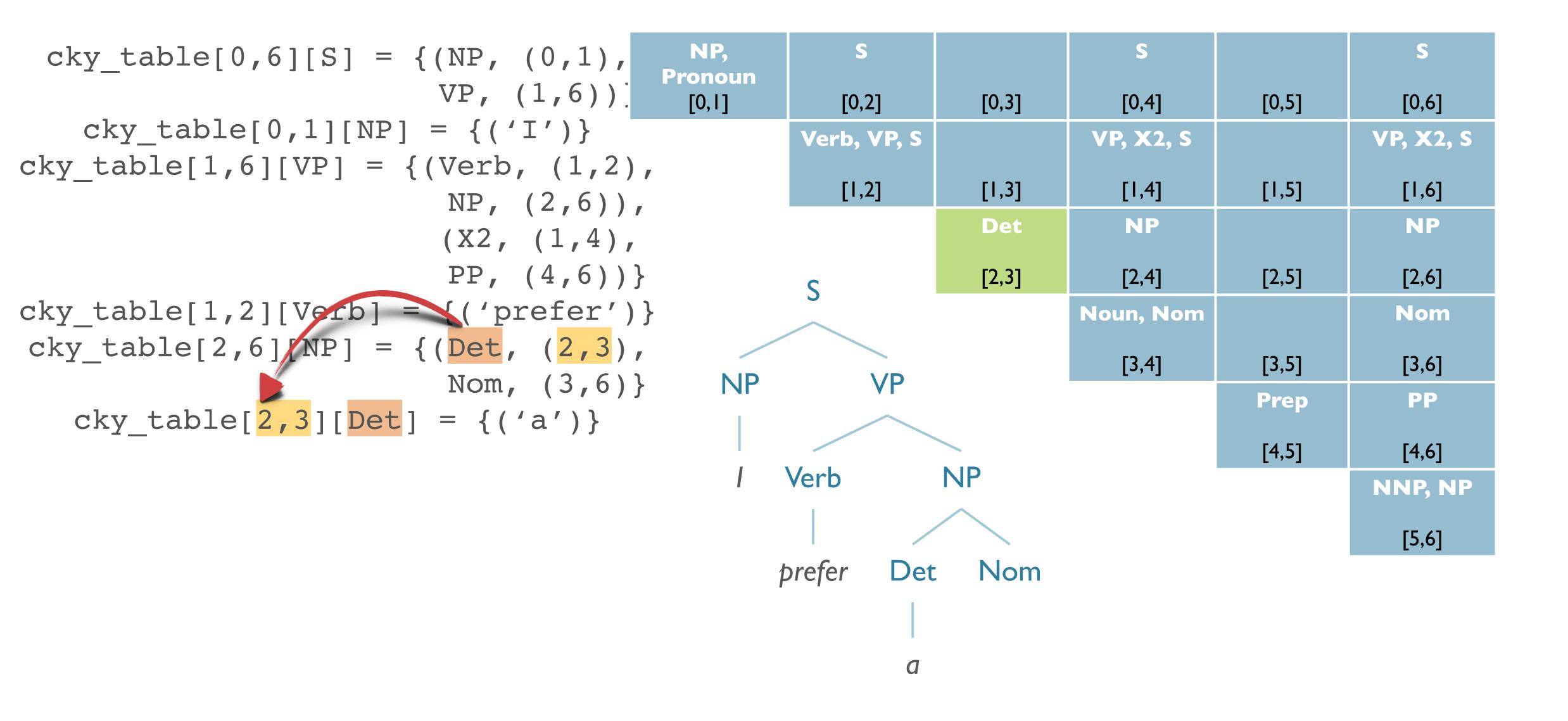


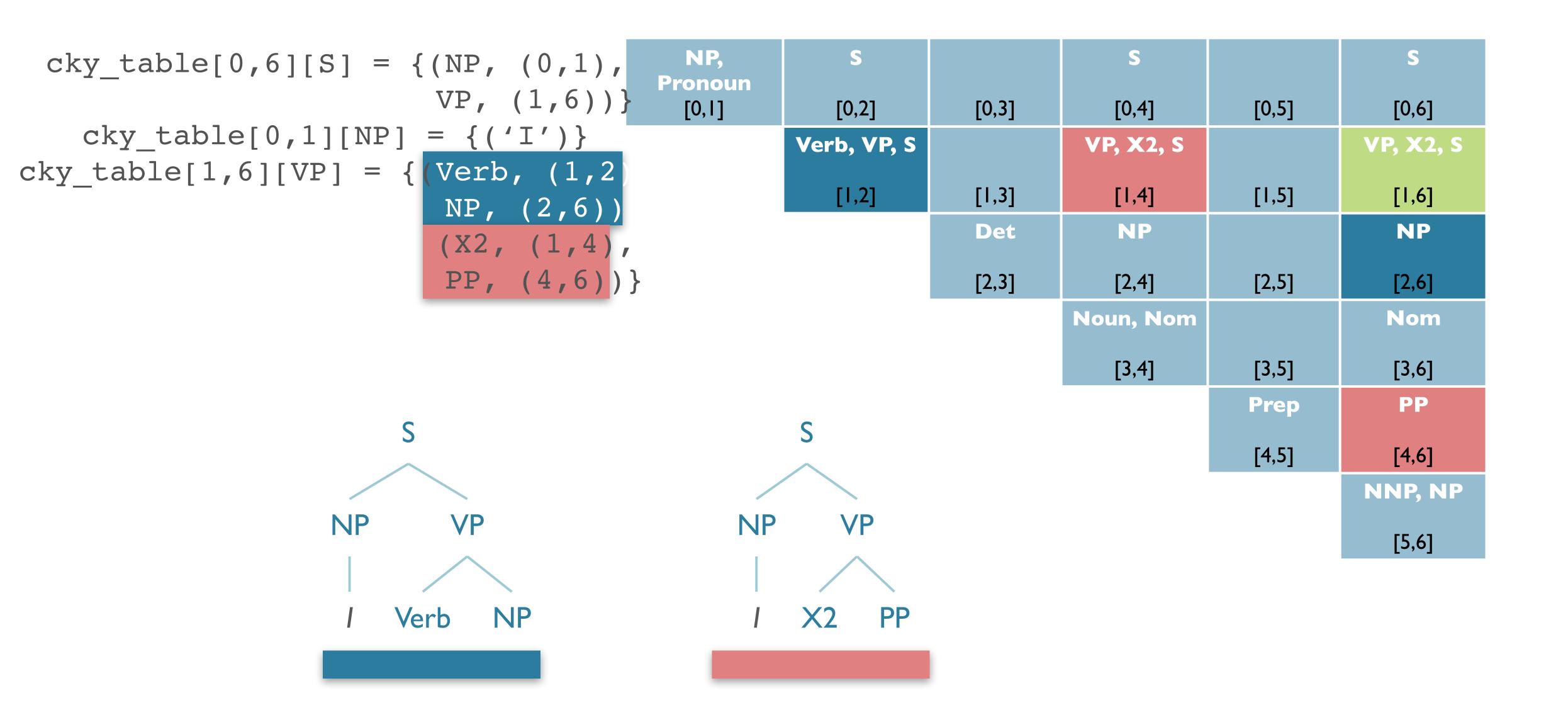












flight

prefer a

TWA

on

19

### Probabilistic Context-Free Grammars

# Probabilistic Context-free Grammars: Roadmap

Motivation: Ambiguity

Approach:

Definition

Disambiguation

Parsing

Evaluation

Enhancements

#### Motivation

What about ambiguity?

Current algorithm can *represent* it

...can't resolve it.

## Probabilistic Parsing

- Provides strategy for solving disambiguation problem
  - Compute the probability of all analyses
  - Select the most probable

- Employed in language modeling for speech recognition
  - N-gram grammars predict words, constrain search
  - Also, constrain generation, translation

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	is a number between $0$ and $1$ expressing $P(oldsymbol{eta} oldsymbol{A})$

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$\overline{S}$	a designated <b>start symbol</b>

#### PCFGs

- Augment each production with probability that LHS will be expanded as RHS
  - $\bullet P(A \rightarrow \beta)$
  - $\bullet P(A \rightarrow \beta | A)$
  - $\bullet P(\beta|A)$
  - $\bullet$   $P(RHS \mid LHS)$
- NB: the first is often used; but the latter are what's really meant.

#### PCFGs

Sum over all possible expansions is 1

$$\sum_{\beta} P(A \to \beta) = 1$$

- A PCFG is consistent if sum of probabilities of all sentences in language is 1
  - Recursive rules often yield inconsistent grammars (Booth & Thompson, 1973)

## Example PCFG: Augmented $\mathcal{L}_1$

Grammar		Lexicon
$S \rightarrow NP VP$	[.80]	$Det \rightarrow that [.10] \mid a [.30] \mid the [.60]$
$S \rightarrow Aux NP VP$	[.15]	$Noun \rightarrow book [.10] \mid flight [.30] \mid meal [.15] \mid money [0.5]$
$S \rightarrow VP$	[.05]	$\mid flights \; [0.40] \mid dinner \; [.10]$
$NP \rightarrow Pronoun$	[.35]	$Verb \rightarrow book [.30] \mid include [.30] \mid prefer [.40]$
$NP \rightarrow Proper-Noun$	[.30]	$Pronoun \rightarrow I[.40] \mid she[.05] \mid me[.15] \mid you[.40]$
$NP \rightarrow Det\ Nominal$	[.20]	$Proper-Noun \rightarrow Houston [.60] \mid NWA [.40]$
$NP \rightarrow Nominal$	[.15]	$Aux \rightarrow does [.60] \mid can [.40]$
$Nominal \rightarrow Noun$	[.75]	$Preposition \rightarrow from [.30] \mid to [.30] \mid on [.20] \mid near [.15]$
$Nominal \rightarrow Nominal \ Noun$	[.20]	$\mid through \; [.05]$
$Nominal \rightarrow Nominal PP$	[.05]	
$VP \rightarrow Verb$	[.35]	
$VP \rightarrow Verb NP$	[.20]	
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$VP \rightarrow VP PP$	[.15]	
$PP \rightarrow Preposition NP$	[1.0]	

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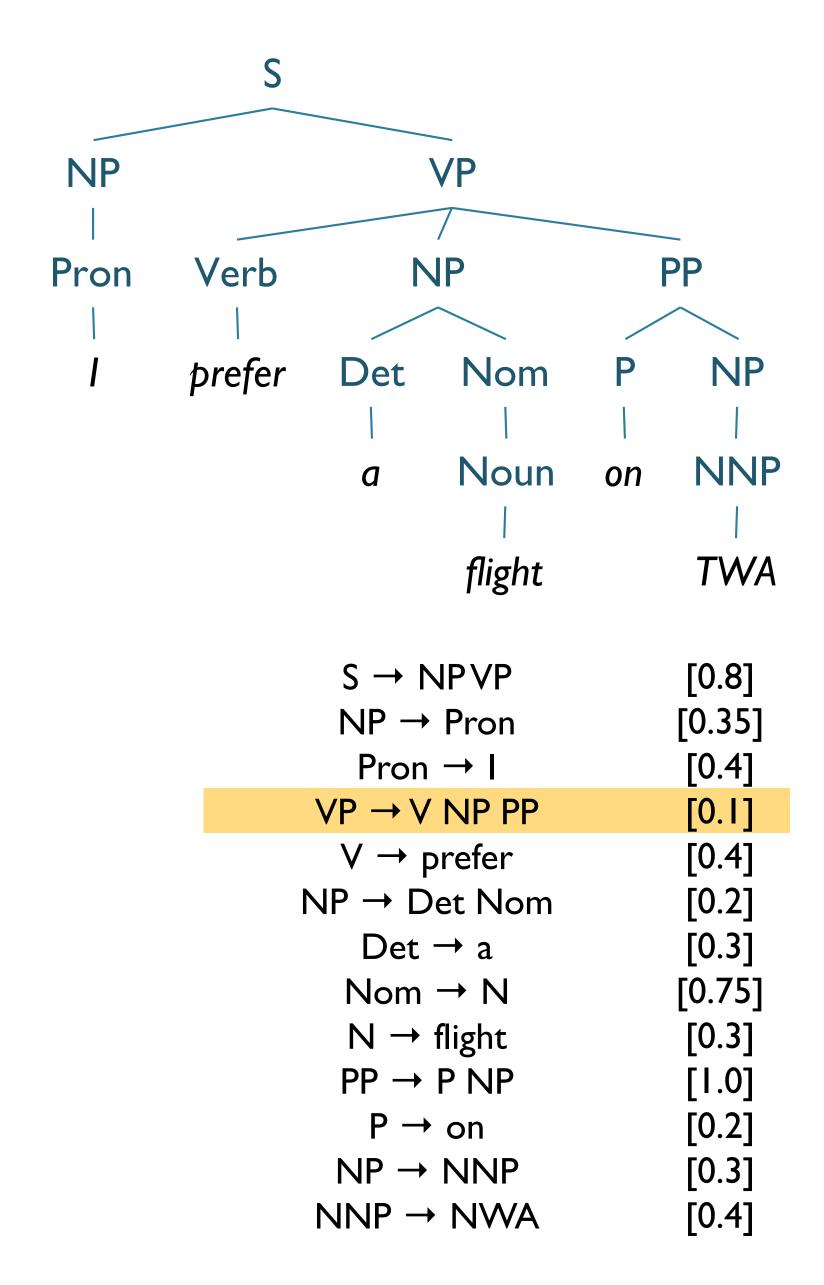
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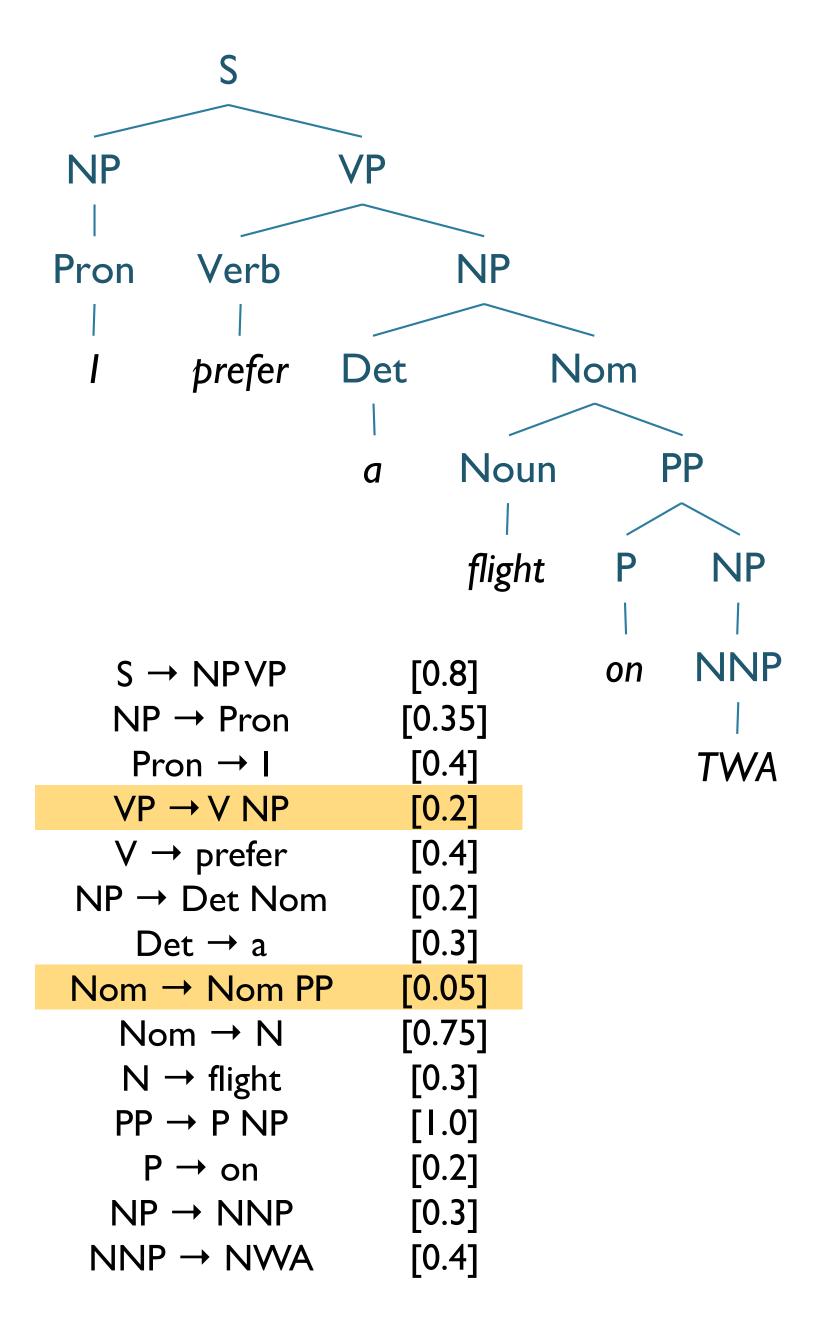
## Disambiguation

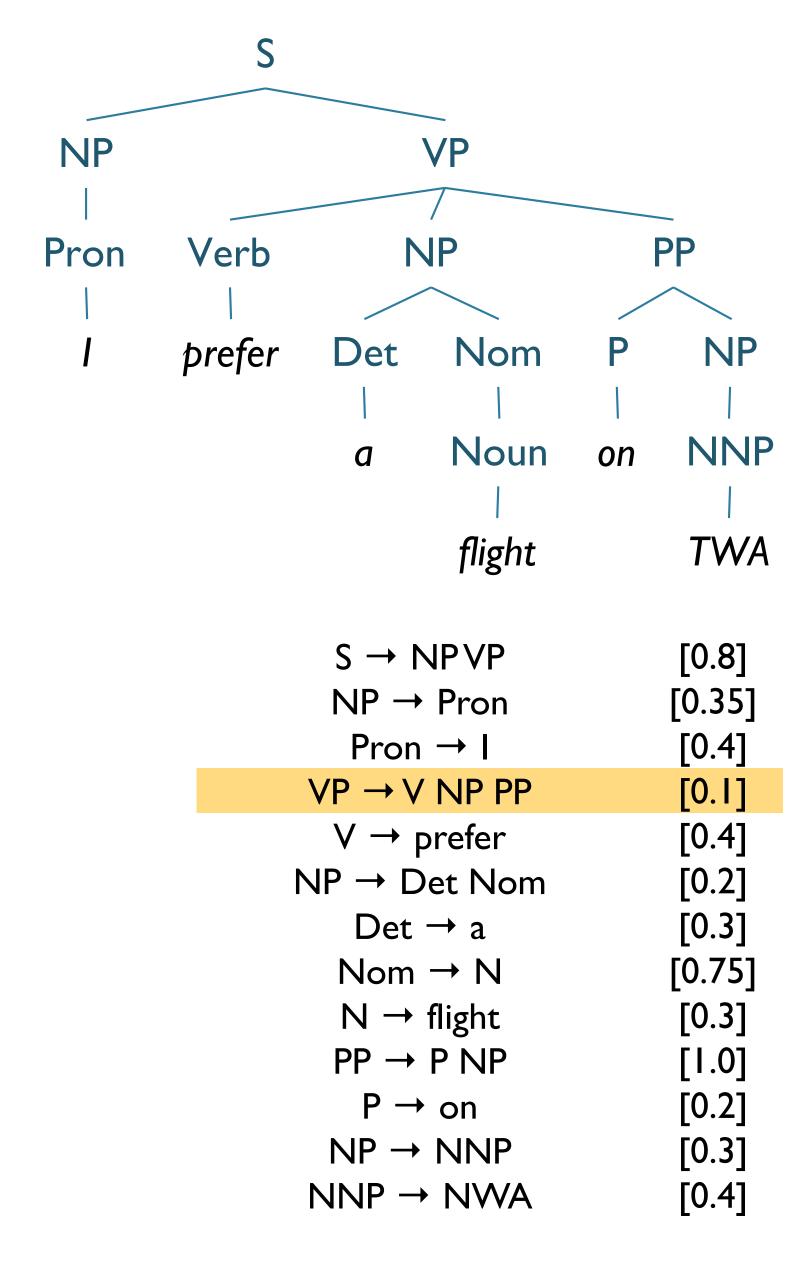
- ullet A PCFG assigns probability to each parse tree T for input S
- ullet Probability of T: product of all rules used to derive T

$$P(T,S) = \prod_{i=1}^{n} P(RHS_i | LHS_i)$$

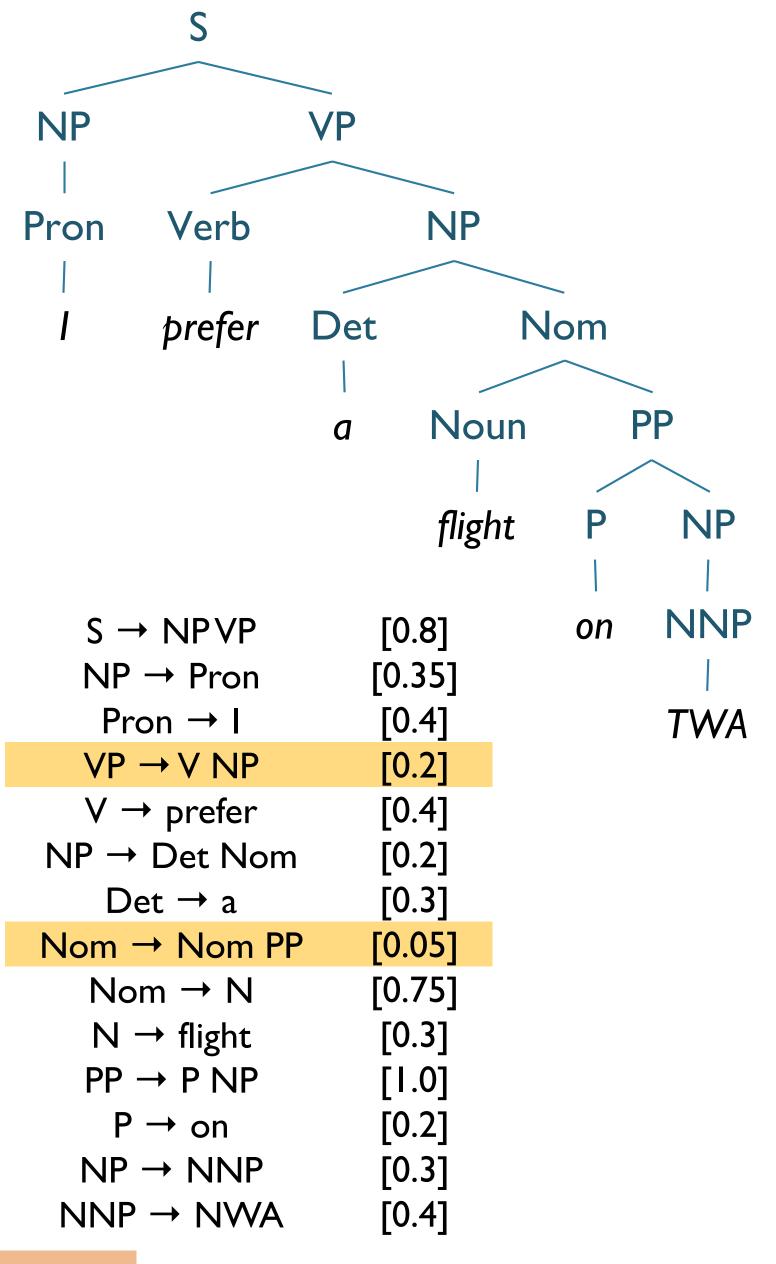
$$P(T,S) = P(T)P(S \mid T) = P(T)$$







 $\sim 1.452 \times 10^{-6}$ 



 $\sim 1.452 \times 10^{-7}$ 

## Parsing Problem for PCFGs

• Select T such that (s.t.)

$$\hat{T}(S) = \underset{T \text{ s.t. } S = yield(T)}{\operatorname{argmax}} P(T)$$

- String of words S is yield of parse tree
- Select the tree T that maximizes the probability of the parse

# Application: Language Modeling

• *n*-grams helpful for modeling the probability of a string

32

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- PCFGs are able to give probability of entire string without as bad sparsity
- Model probability of syntactically valid sentences
  - Not just probability of sequence of words

## PCFGs: Parsing

## Probabilistic CKY (PCKY)

- Like regular CKY
  - Assumes grammar in Chomsky Normal Form (CNF)
    - $\bullet$   $A \rightarrow B C$
    - $\bullet$   $A \rightarrow w$
  - Represent input with indices b/t words:
    - Book 1 that 2 flight 3 through 4 Houston 5

### Probabilistic CKY (PCKY)

- For input string length n and non-terminals V
  - Cell [ i, j, A ] in ( n+1 ) × ( n+1 ) × V matrix
  - Contains probability that A spans [i, j]

```
function Probabilistic-CKY-PARSE(words, grammar) returns most probable parse and its probability
for j ← from 1 to LENGTH(words) do
for all \{A \mid A \rightarrow words[j] \in grammar\}
      table[j-1, j, A] \leftarrow P(A \rightarrow words[j])
 for i \leftarrow \text{from } j-2 \text{ downto } 0 \text{ do}
  for k \leftarrow i + 1 to j-1 do
  for all \{A \mid A \rightarrow B \ C \in grammar,
        and table[i, k, B] > 0 and table[k, j, C] > 0
  if (table[i, j, A] < P(A \rightarrow BC) \times table[i, k, B] \times table[k, j, C]) then
      table[i, j, A] \leftarrow P(A \rightarrow BC) \times table[i, k, B] \times table[k, j, C]
      back[i, j, A] \leftarrow \{k, B, C\}
  return Build_Tree(back[1, Length(words), S]), table[1, Length(words), S]
```

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function Probabilistic-Cky-Parse (words, grammar) returns most probable parse and its probability
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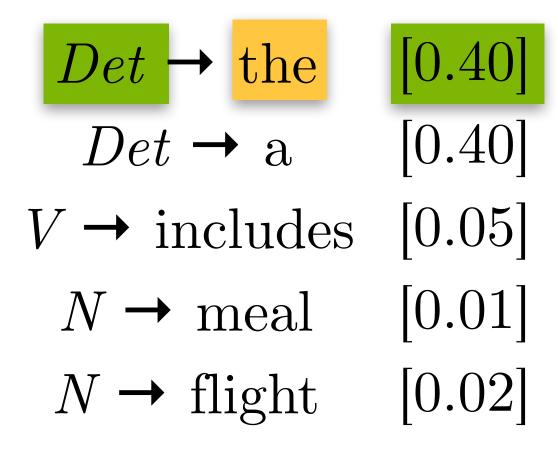
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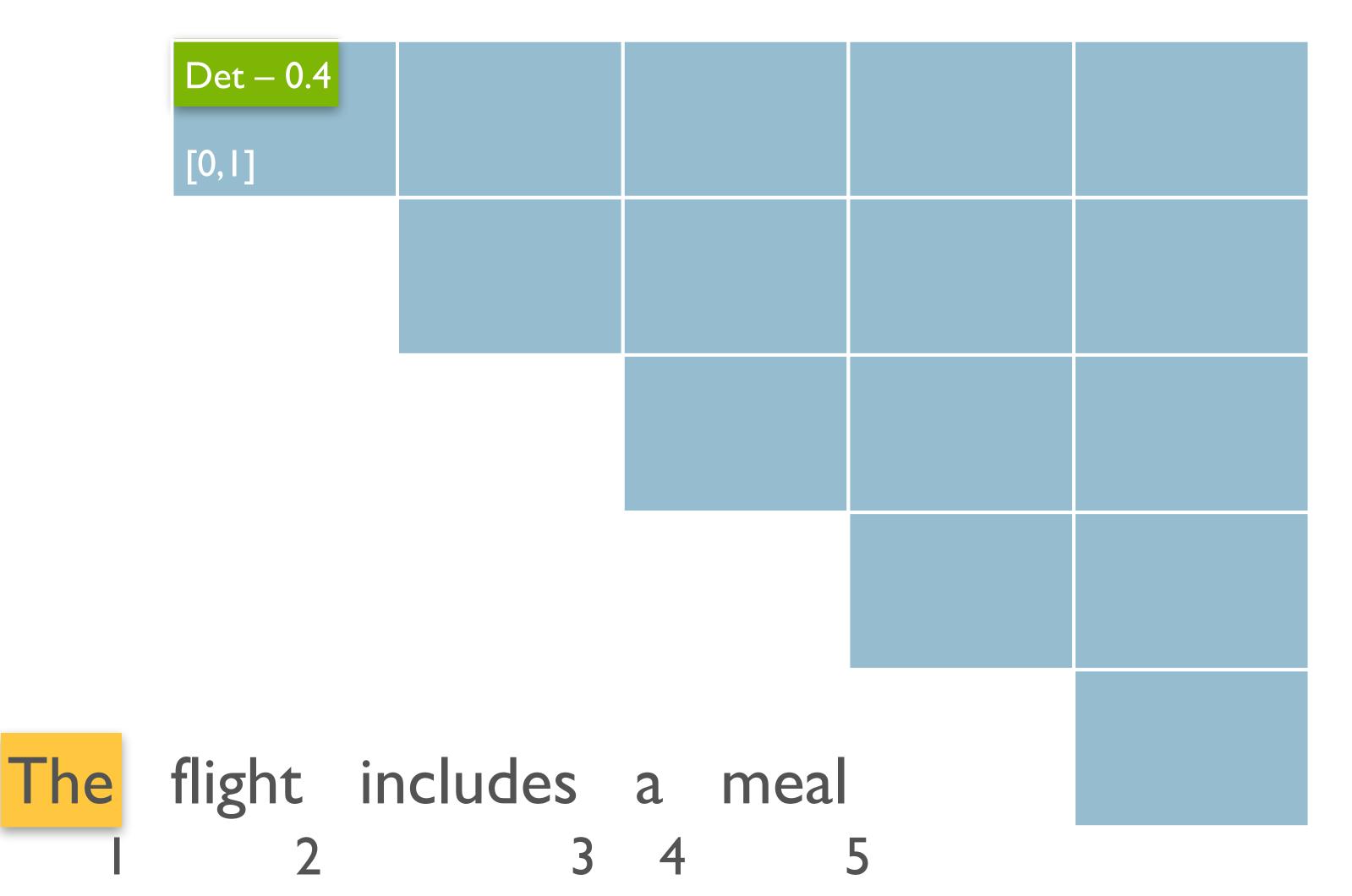
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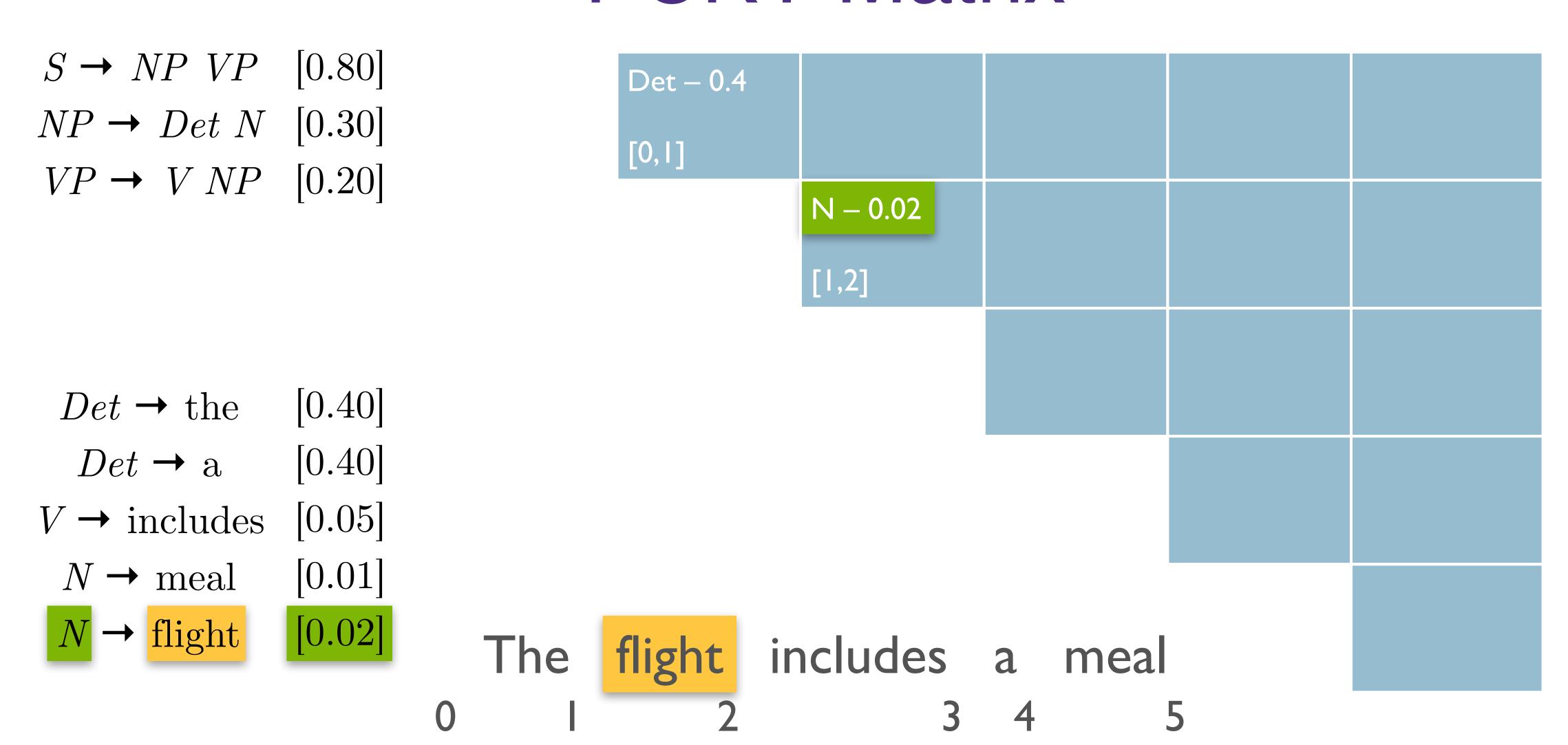
## PCKY Grammar Segment

```
S \rightarrow NP \ VP \quad [0.80] Det \rightarrow the \quad [0.40] NP \rightarrow Det \ N \quad [0.30] Det \rightarrow a \quad [0.40] VP \rightarrow V \ NP \quad [0.20] V \rightarrow includes \quad [0.05] N \rightarrow meal \quad [0.01] N \rightarrow flight \quad [0.02]
```

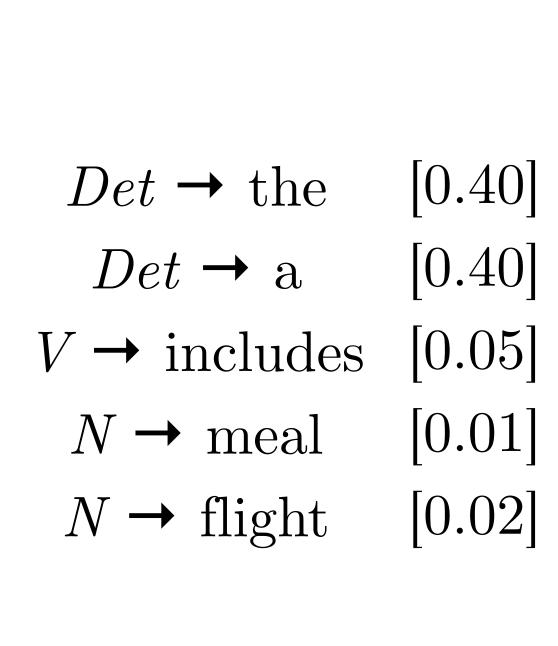
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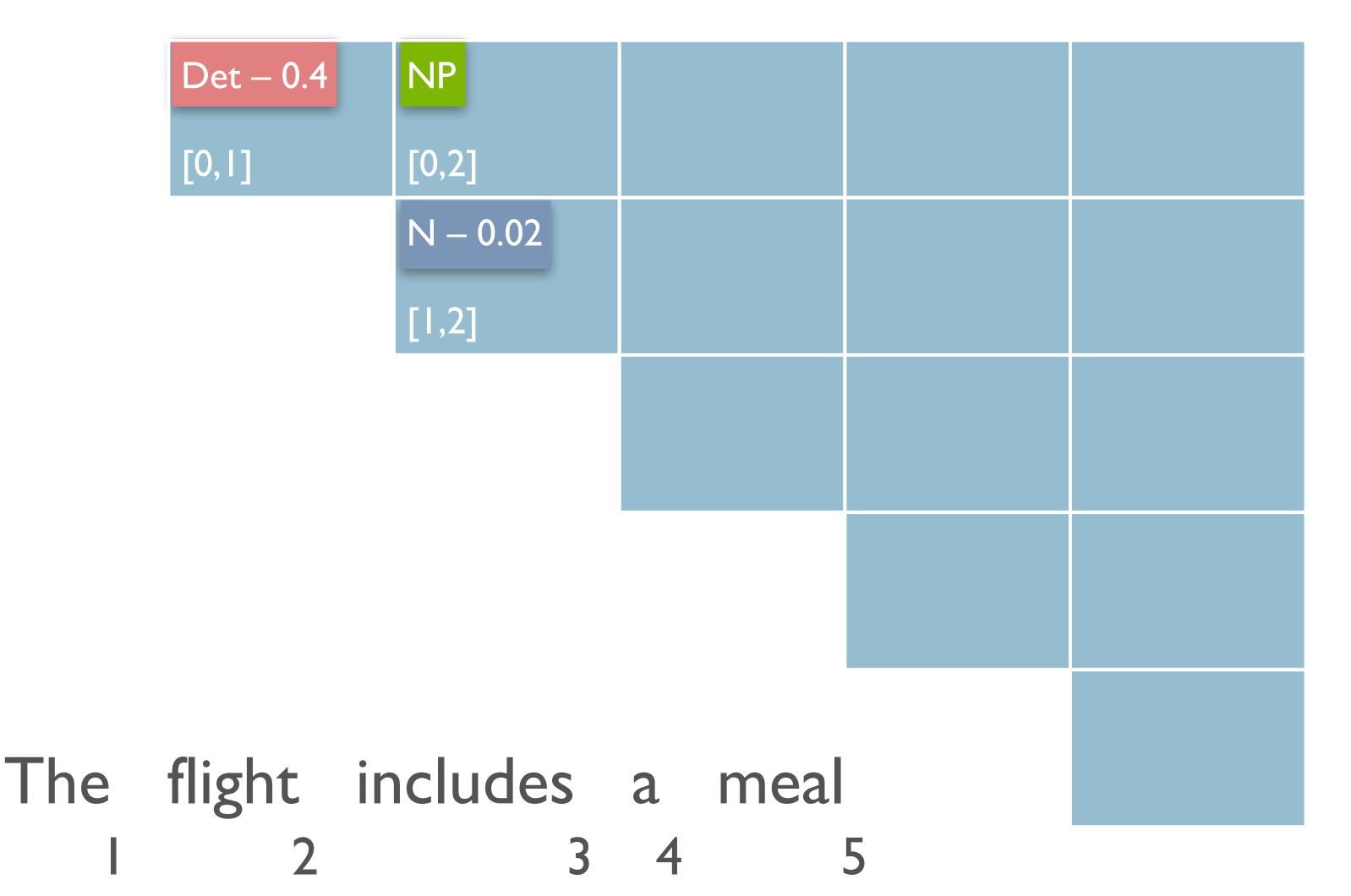






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Det \rightarrow a \qquad [0.40]
V \rightarrow includes \qquad [0.05]
N \rightarrow meal \qquad [0.01]
N \rightarrow flight \qquad [0.02]
```

```
Det – 0.4
                 NP
       [0,1]
                 [0,2]
                 N - 0.02
                 [1,2]
The flight includes a meal
```

$$S \rightarrow NP \ VP \quad [0.80]$$
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$$Det \rightarrow \text{the}$$
 [0.40]  
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[0.02]

 $N \rightarrow \text{flight}$ 

```
Det – 0.4
                 NP
       [0,1]
                 [0,2]
                 N - 0.02
                 [1,2]
               = 0.00024
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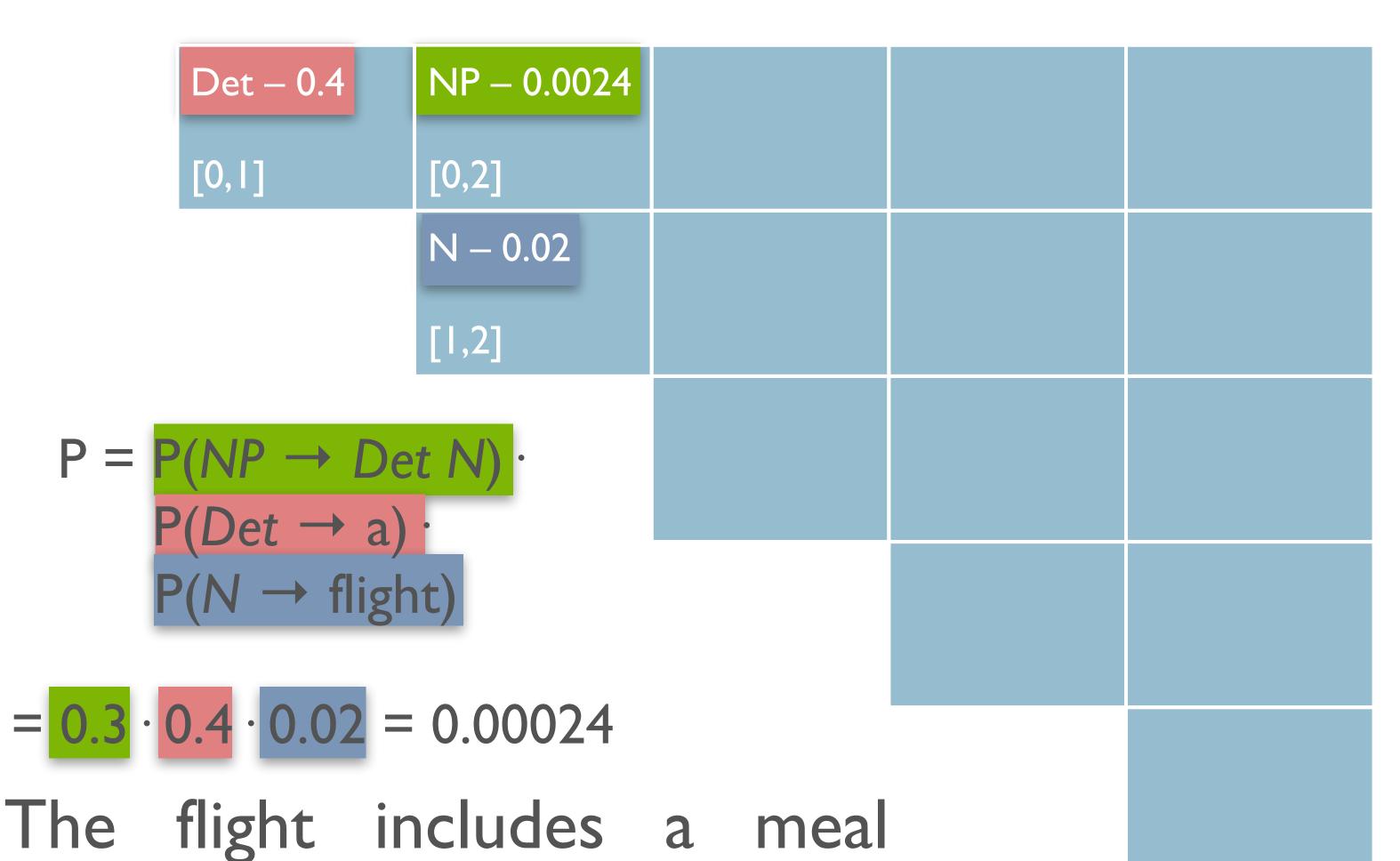
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 [0.01]

$$N \rightarrow \text{flight} \quad [0.02]$$



$S \rightarrow NP VP$	[0.80]
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Det - 0.4	NP - 0.0024			S - 2.304×10-8
[0,1]	[0,2]	[0,3]	[0,4]	[0,5]
	N - 0.02			
	[1,2]	[1,3]	[1,4]	[1,5]
		V — 0.05		VP - 1.2×10-5
		[2,3]	[2,4]	[2,5]
			Det – 0.4	NP - 0.0012
			[3,4]	[3,5]
			N - 0.01	
flight includes a meal			[4,5]	

## Inducing a PCFG

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$$P(\alpha \to \beta \mid \alpha) = \frac{Count(\alpha \to \beta)}{\sum_{\gamma} Count(\alpha \to \gamma)} = \frac{Count(\alpha \to \beta)}{Count(\alpha)}$$

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- Alternative: Learn probabilities by re-estimating
  - (Later)

### Probabilistic Parser Development Paradigm

	Train	Dev	Test
	Large	Small	Small/Med
Size	(eg.WSJ 2-21, 39,830 sentences)	(e.g.WSJ 22)	(e.g. WSJ, 23, 2,416 sentences)
Usage	Estimate rule probabilities	Tuning/Verification, Check for Overfit	Held Out, Final Evaluation

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  - Maximally strict: identical to 'gold standard'

- Assume a 'gold standard' set of parses for test set
- How can we tell how good the parser is?
- How can we tell how good a parse is?
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  - Partial credit:

#### Parser Evaluation

- Assume a 'gold standard' set of parses for test set
- How can we tell how good the parser is?
- How can we tell how good a parse is?
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  - Partial credit:
    - Constituents in output match those in reference

#### Parser Evaluation

- Assume a 'gold standard' set of parses for test set
- How can we tell how good the parser is?
- How can we tell how good a parse is?
  - Maximally strict: identical to 'gold standard'
  - Partial credit:
    - Constituents in output match those in reference
      - Same start point, end point, non-terminal symbol

#### Parseval

- How can we compute parse score from constituents?
- Multiple Measures:

```
Labeled Recall (LR) = 
# of correct constituents in hypothetical parse

# of total constituents in reference parse

Labeled Precision (LP) = 
# of correct constituents in hypothetical parse

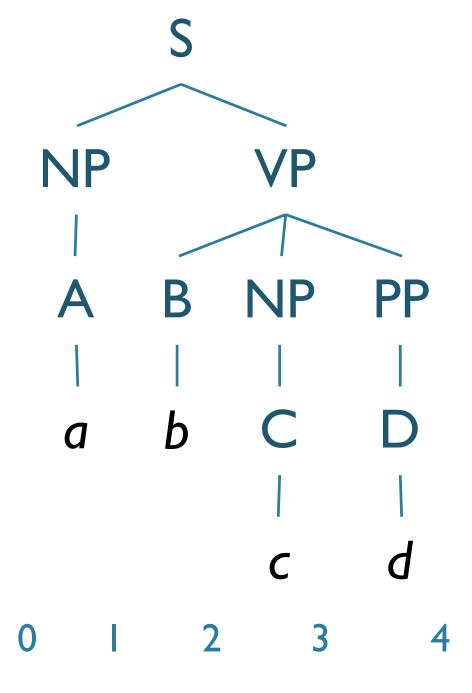
# of total consituents in hypothetical parse
```

#### Parseval

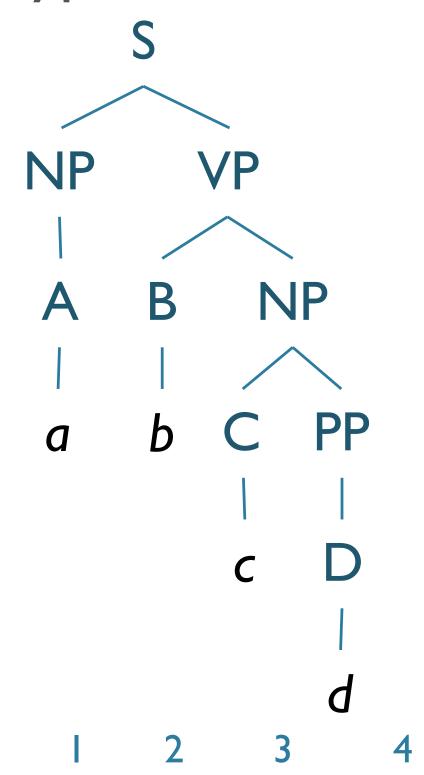
#### • F-measure:

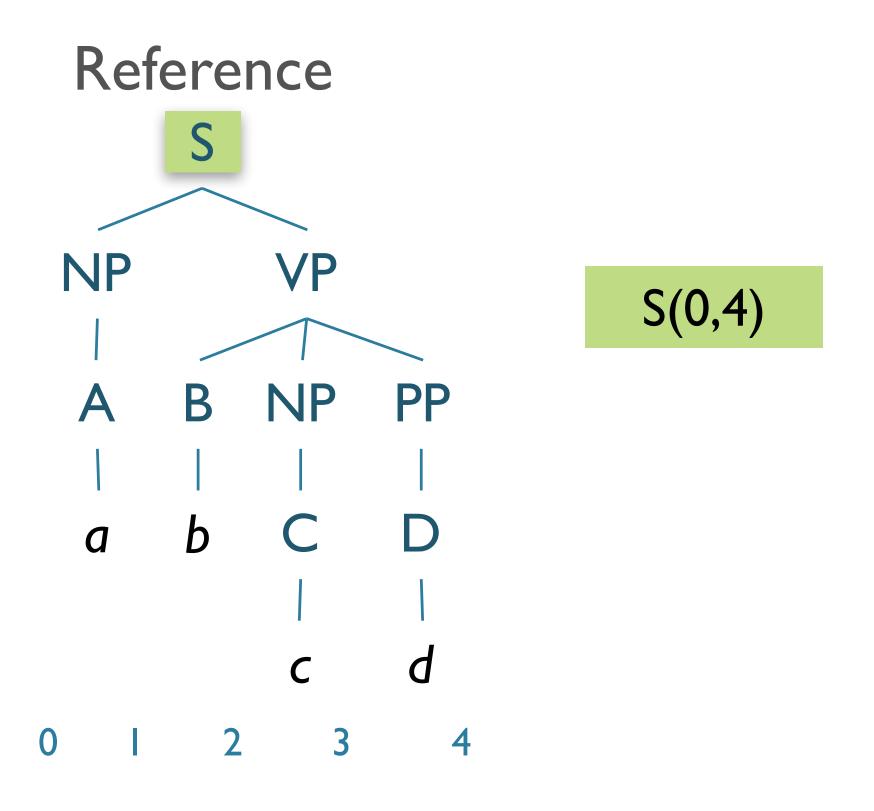
- Combines precision and recall
- Let  $\beta \in \mathbb{R}$ ,  $\beta > 0$  that adjusts P vs. R s.t.  $\beta \propto \frac{R}{P}$
- $F_{\beta}$ -measure is then:  $F_{\beta} = (1 + \beta^2) \cdot \frac{P \cdot R}{\beta^2 \cdot P + R}$
- With F1-measure as  $F_1 = \frac{2PR}{P+R}$

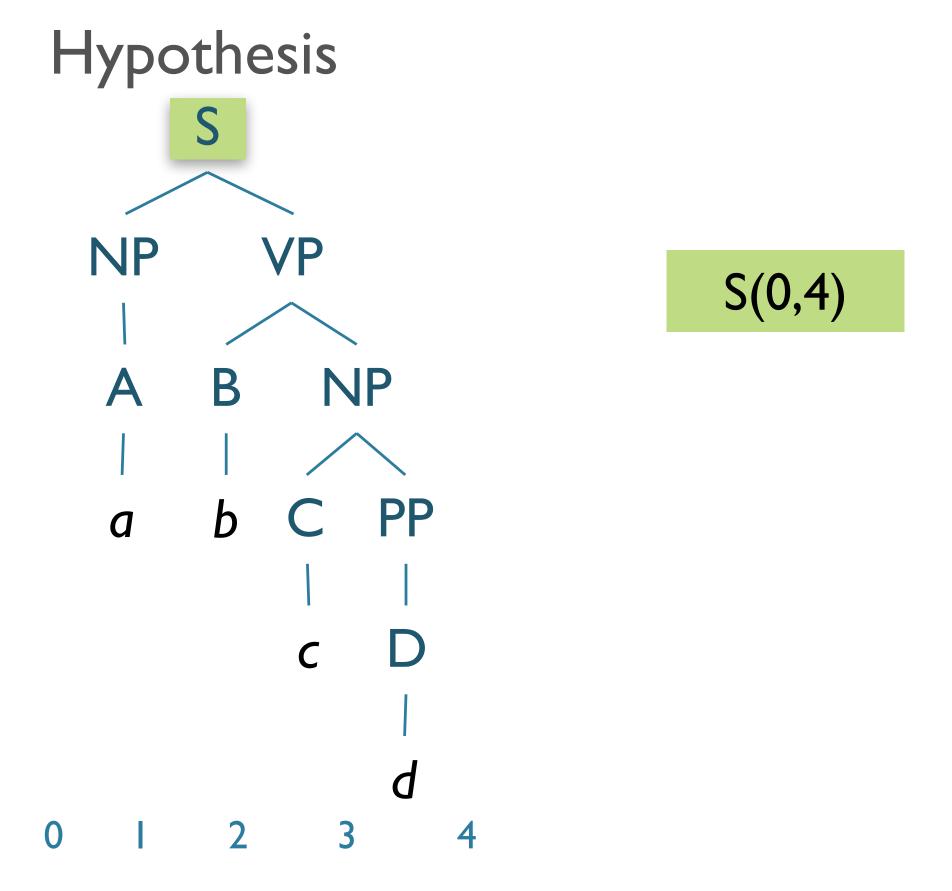
#### Reference

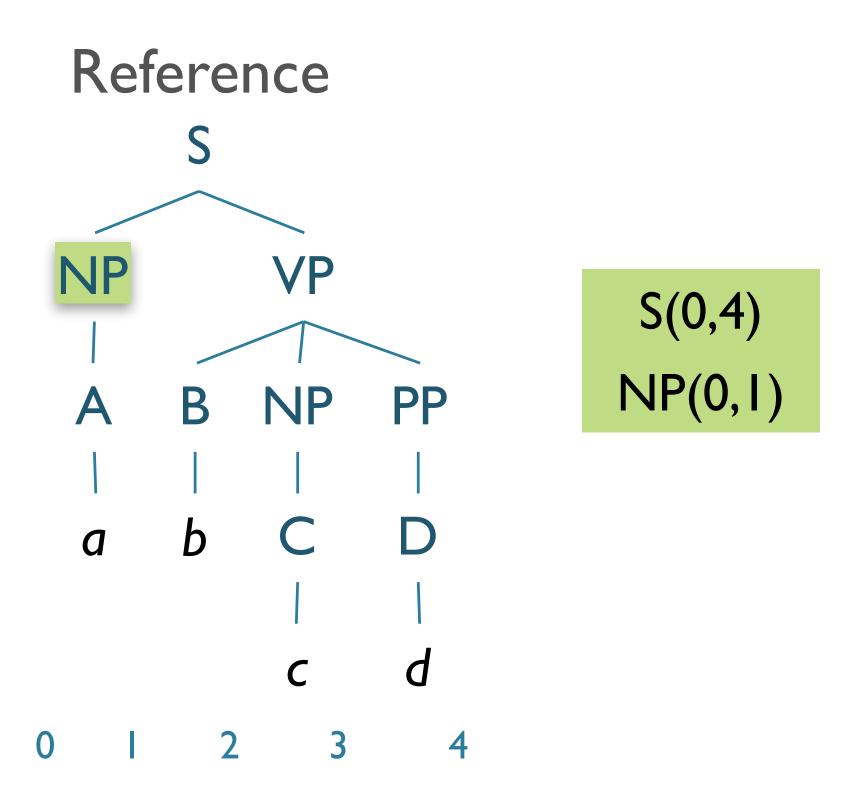


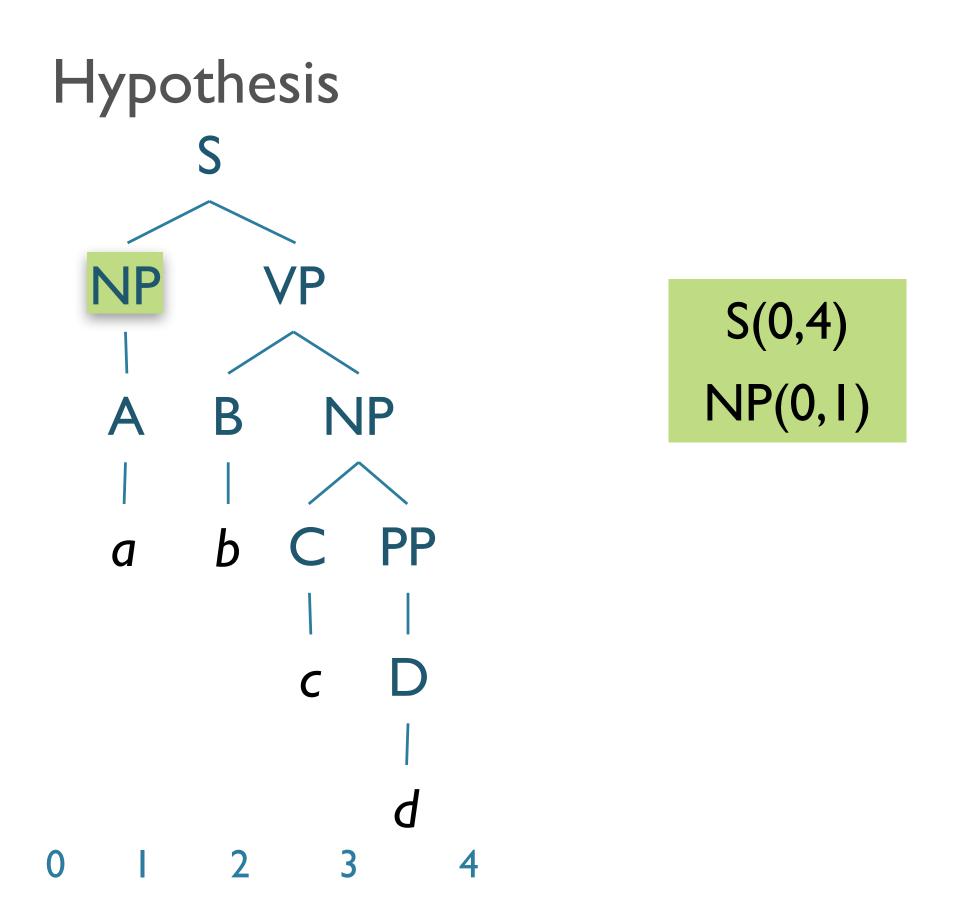
#### Hypothesis

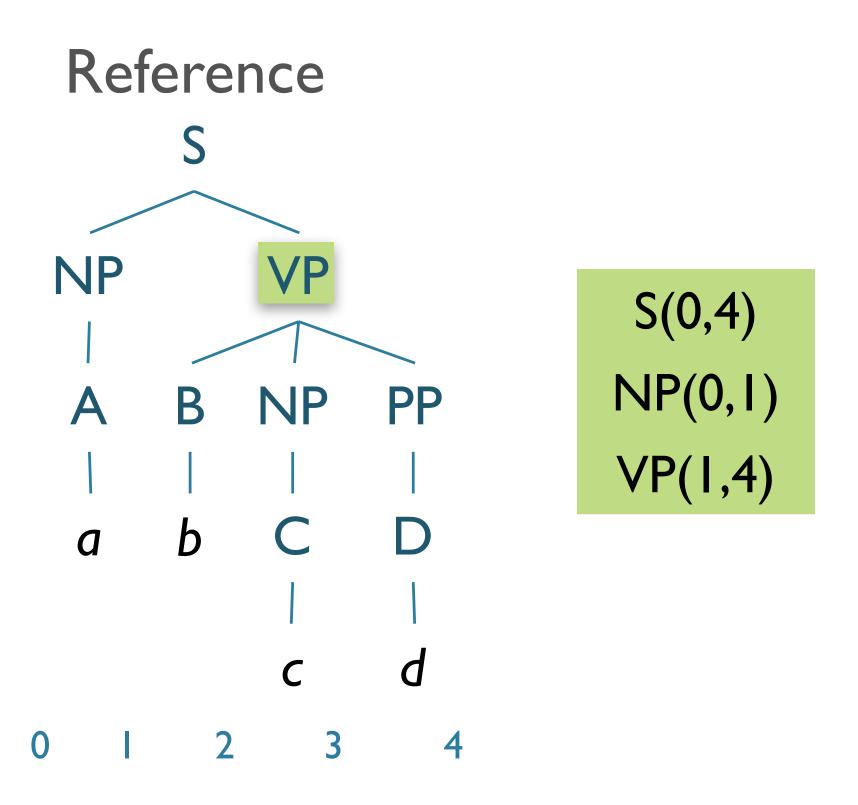


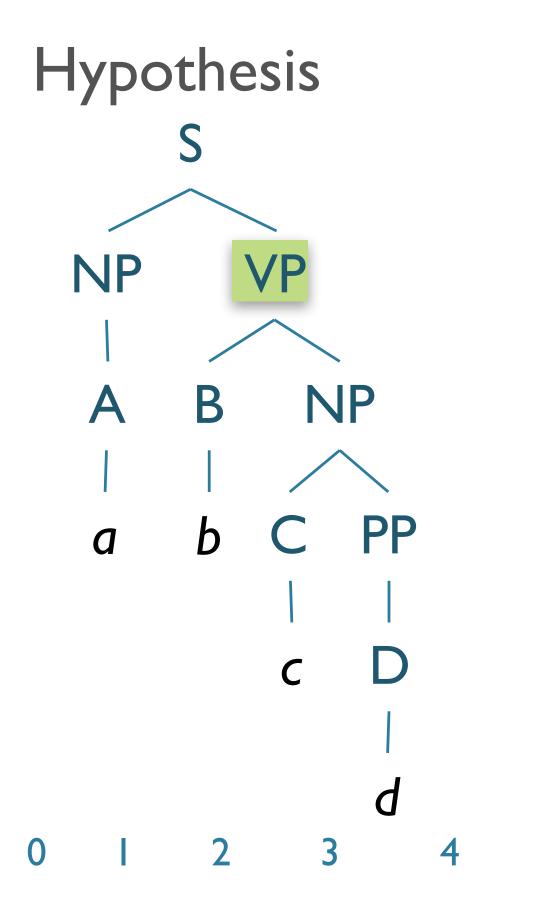




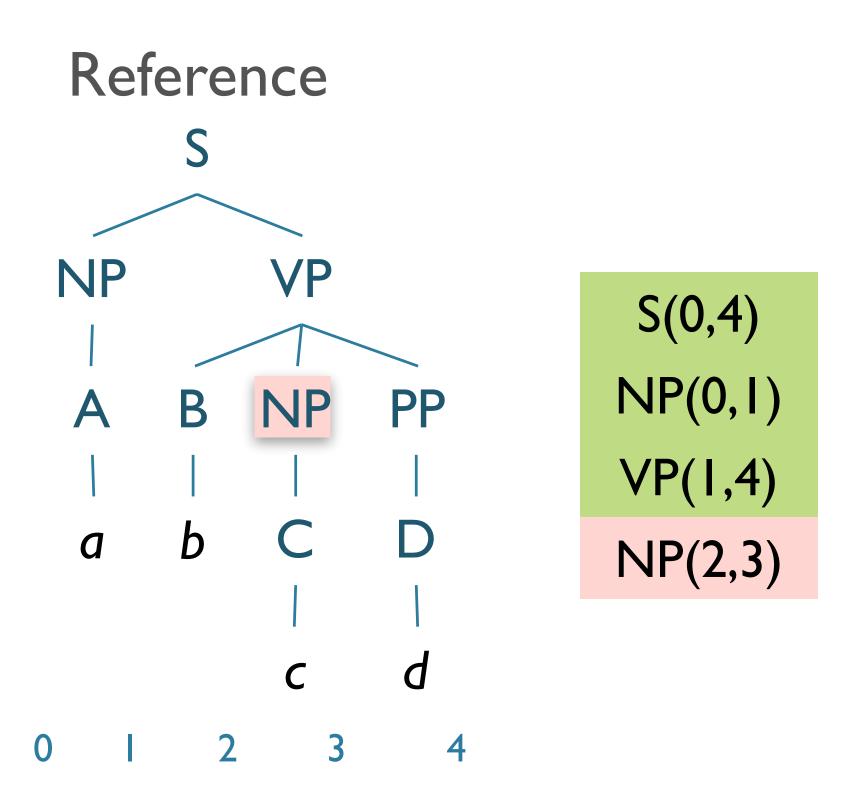


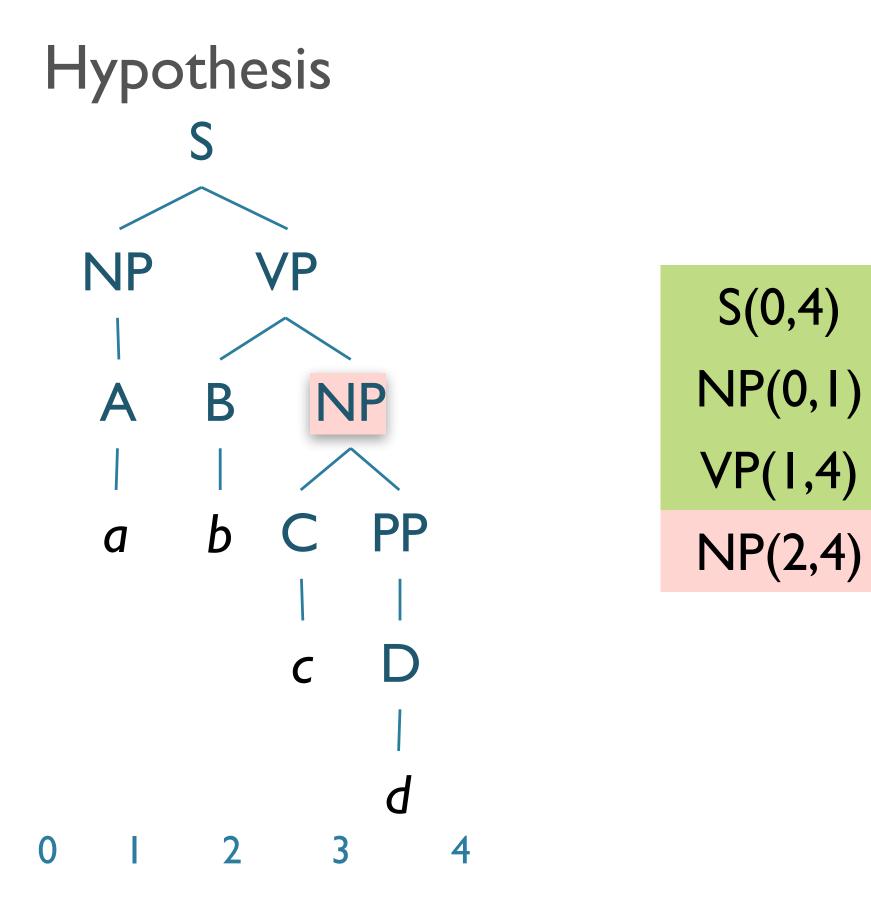


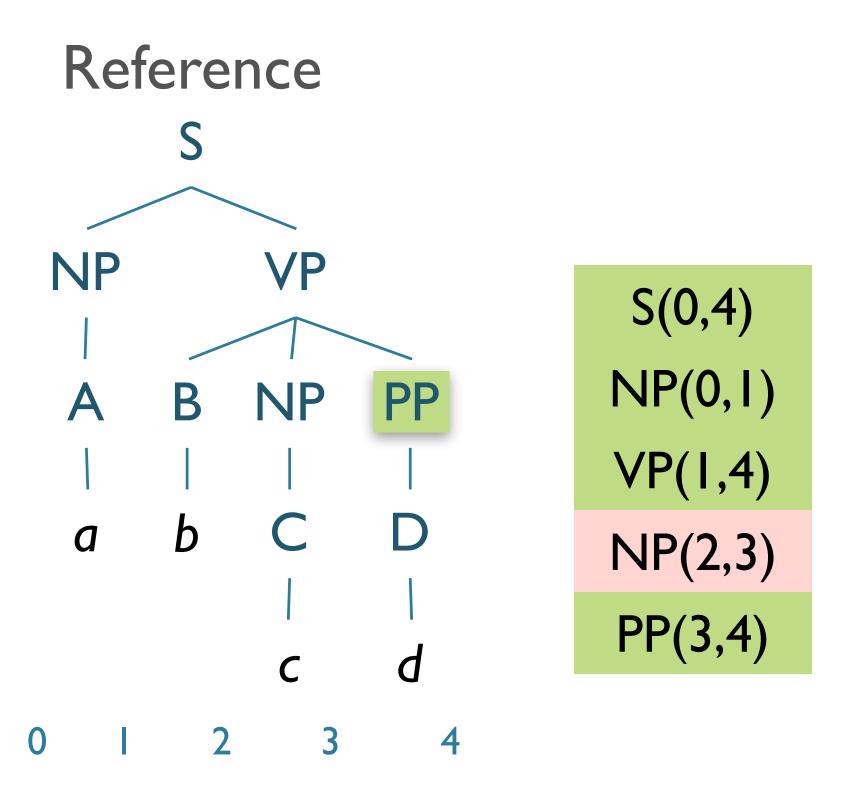


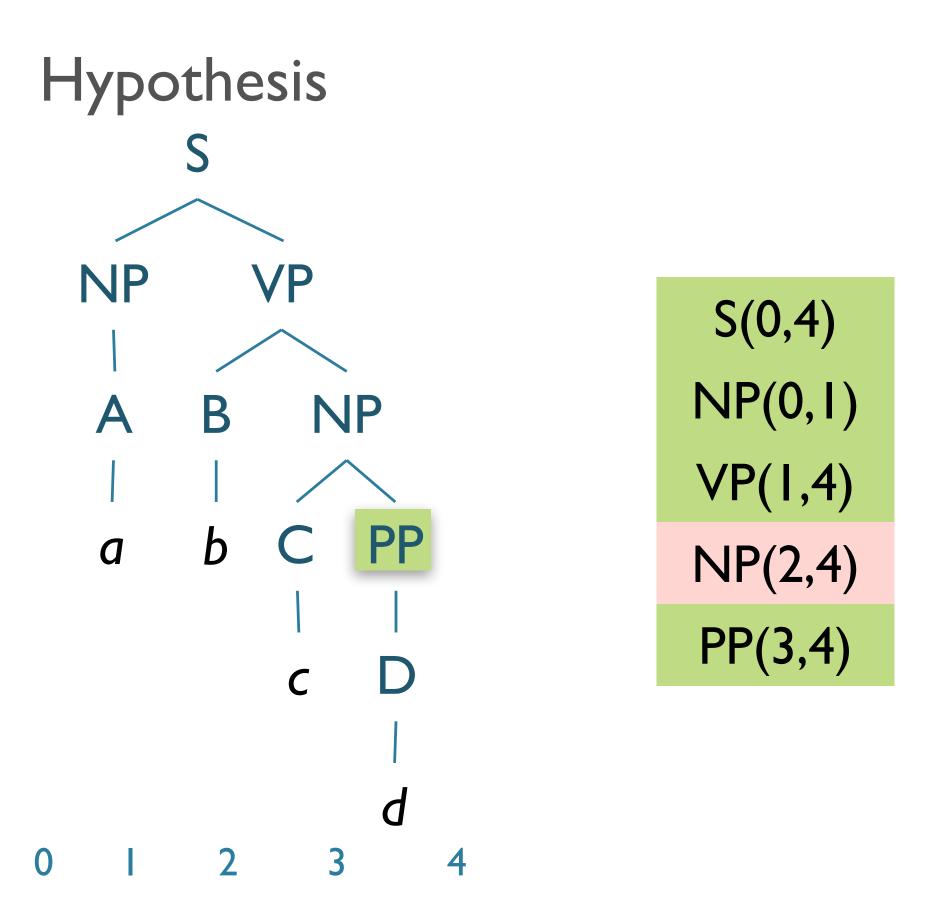


S(0,4) NP(0,1) VP(1,4)

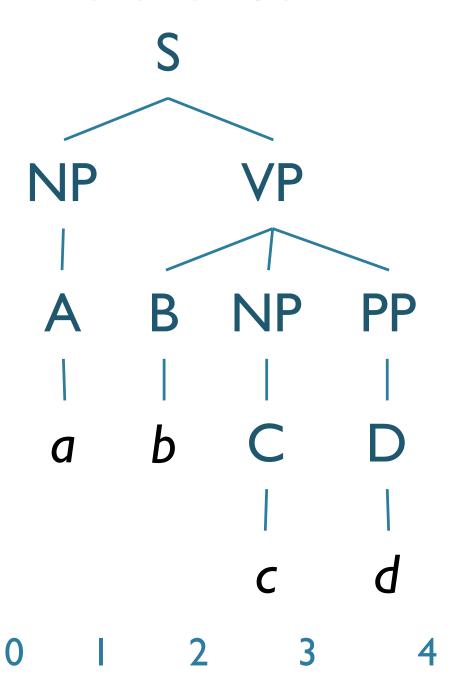








#### Reference



S(0,4) NP(0,1) VP(1,4) NP(2,3) PP(3,4)

> LP: 4/5 LR: 4/5

# Hypothesis **VP** NP NP a

S(0,4) NP(0,1) VP(1,4) NP(2,4) PP(3,4)

#### Parser Evaluation

- Crossing Brackets:
  - # of constituents where produced parse has bracketings that overlap for the siblings:
  - ((A B) C) { (0,2), (2,3) }
     and hyp. has
     (A (B C)) { (0,1), (1,3) }

```
TOP

A B A C

B B
```

```
/* crossing is counted based on the brackets */
/* in test rather than gold file (by Mike) */
for(j=0;j<bn2;j++){
  for(i=0;i<bn1;i++){
    if(bracket1[i].result != 5 &&
        bracket2[j].result != 5 &&
        ((bracket1[i].start < bracket2[j].start &&
        bracket1[i].end > bracket2[j].start &&
        bracket1[i].end < bracket2[j].end) ||
        (bracket1[i].start > bracket2[j].start &&
        bracket1[i].start < bracket2[j].end &&
        bracket1[i].end > bracket2[j].end))){
```

## State-of-the-Art Parsing

- Parsers trained/tested on Wall Street Journal PTB
  - LR: 90%+;
  - LP: 90%+;
  - Crossing brackets: 1%

- Standard implementation of Parseval:
  - evalb

#### Evaluation Issues

- Only evaluating constituency
- There are other grammar formalisms:
  - LFG (Constraint-based)
  - Dependency Structure
- Extrinsic evaluation
  - How well does getting the correct parse match the semantics, etc?

## Earley Parsing

### Earley vs. CKY

- CKY doesn't capture full original structure
  - Can back-convert binarization, terminal conversion
  - Unit non-terminals require change in CKY

### Earley vs. CKY

- CKY doesn't capture full original structure
  - Can back-convert binarization, terminal conversion
  - Unit non-terminals require change in CKY
- Earley algorithm
  - Supports parsing efficiently with arbitrary grammars
  - Top-down search
  - Dynamic programming
    - Tabulated partial solutions
  - Some bottom-up constraints

## Earley Algorithm

- Another dynamic programming solution
  - Partial parses stored in "chart"
  - Compactly encodes ambiguity
  - O(N<sup>3</sup>)
- Chart entries contain:
  - Subtree for a single grammar rule
  - Progress in completing subtree
  - Position of subtree w.r.t. input

## Earley Algorithm

- First, left-to-right pass fills out a chart with *N+1* states
  - Chart entries sit between words in the input string
  - Keep track of states of the parse at those positions
  - For each word position, chart contains set of states representing all partial parse trees generate so far
    - e.g. chart[0] contains all partial parse trees generated at the beginning of sentence

#### Chart Entries

- Three types of constituents:
  - Predicted constituents
  - In-progress constituents
  - Completed constituents

## Parse Progress

- Represented by Dotted Rules
  - Position of indicates type of constituent
- 0 Book 1 that 2 flight 3
  - $S \rightarrow VP$  [0,0] (predicted)
  - $NP \rightarrow Det \cdot Nom$  [1,2] (in progress)
  - $VP \rightarrow VNP$  [0,3] (completed)
- [x,y] tells us what portion of the input is spanned so far by rule
- Each state *s<sub>i</sub>*: <*dotted rule*>, [<*back pointer*>, <*current position*>]

### o Book 1 that 2 flight 3

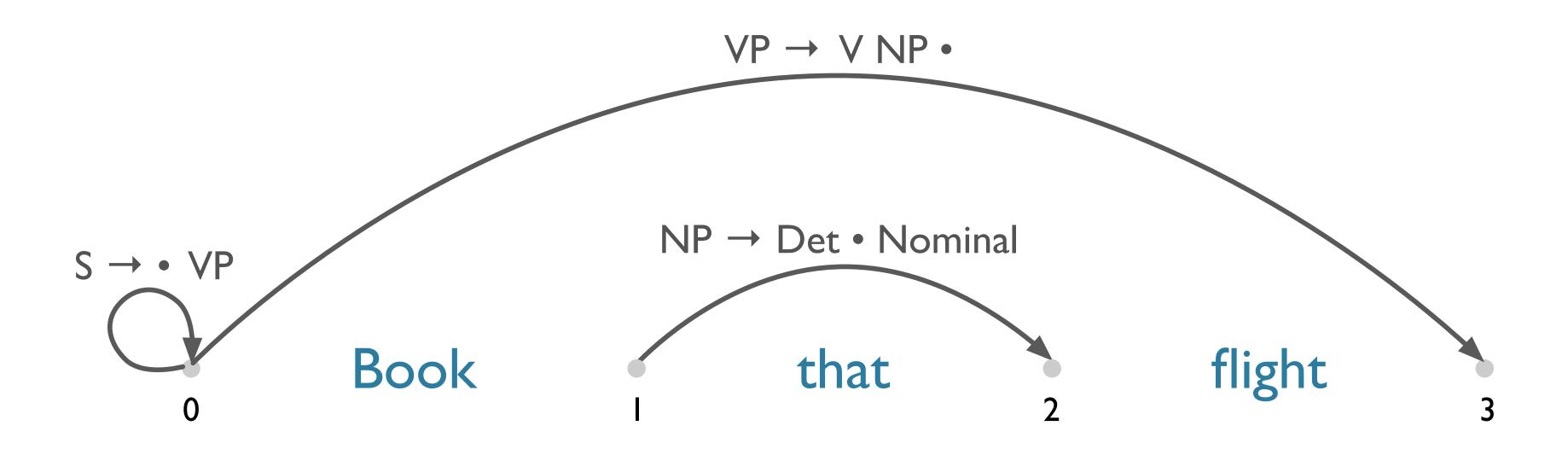
- $S \to VP, [0,0]$ 
  - First 0 means S constituent begins at the start of input
  - Second 0 means the dot is here too
  - So, this is a top-down prediction

### o Book 1 that 2 flight 3

- $S \rightarrow VP$ , [0,0]
  - First 0 means S constituent begins at the start of input
  - Second 0 means the dot is here too
  - So, this is a top-down prediction
- *NP* → *Det Nom*, [1,2]
  - the NP begins at position 1
  - the dot is at position 2
  - so, Det has been successfully parsed
  - Nom predicted next

## <sub>0</sub> Book <sub>1</sub> that <sub>2</sub> flight <sub>3</sub> (continued)

- $V \rightarrow V NP \cdot [0,3]$ 
  - Successful VP parse of entire input



#### Successful Parse

- Final answer found by looking at last entry in chart
- If entry resembles  $S \rightarrow \alpha \cdot [0,N]$  then input parsed successfully
- Chart will also contain record of all possible parses of input string, given the grammar

#### Parsing Procedure for the Earley Algorithm

- Move through each set of states in order, applying one of three operations:
  - predictor: add predictions to the chart
  - scanner: read input and add corresponding state to chart
  - completer: move dot to right when new constituent found
- Results (new states) added to current or next set of states in chart
- No backtracking and no states removed: keep complete history of parse

### Earley Algorithm

```
function Earley-Parse(words, grammar) returns chart
 ENQUEUE((\gamma \longrightarrow \bullet S, [0,0]), chart[\theta])
 for i \leftarrow from 0 to LENGTH(words) do
   for each state in chart[i] do
     if Incomplete?(state) and
         Next-Cat(state) is not a part of speech then
       \mathbf{PREDICTOR}(state)
     elseif Incomplete?(state) and
         NEXT-CAT(state) is a part of speech then
       SCANNER(state)
     else
       Complete Eta(state)
     end
   end
 return(chart)
```

#### Earley Algorithm

```
procedure PREDICTOR((A \rightarrow \alpha \bullet B \beta, (i,j)))
  for each (B \rightarrow \gamma) in Grammar-Rules-For(B,grammar) do
    ENQUEUE((B \rightarrow \bullet \gamma, [j,j]), chart[j])
  end
procedure SCANNER((A \rightarrow a \bullet B \beta, (i,j)))
  if B c Parts-of-Speech(word/j/) then
    ENQUEUE((B \rightarrow word[j] \bullet, [j,j+1]), chart[j+1])
procedure Completer (B \rightarrow \gamma \bullet, [j,k])
 for each (A \rightarrow \alpha \bullet B \beta, [i,j]) in chart[j] do
    ENQUEUE((A \rightarrow a B \bullet \beta, [i,k]), chart[k])
  end
```

#### 3 Main Subroutines of Earley

- Predictor
  - Adds predictions into the chart
- Scanner
  - Reads the input words and enters states representing those words into the chart
- Completer
  - Moves the dot to the right when new constituents are found

#### Predictor

- Intuition:
  - Create new state for top-down prediction of new phrase
- Applied when non part-of-speech non-terminals are to the right of a dot:
  - $S \rightarrow VP[0,0]$
- Adds new states to current chart
  - One new state for each expansion of the non-terminal in the grammar

$$VP \rightarrow \cdot V$$
 [0,0]  
 $VP \rightarrow \cdot V NP$  [0,0]

## Chart[0]

S0	$\gamma \rightarrow \cdot S$	[0,0]	Dummy start state
S1	$S \rightarrow \cdot NP VP$	[0.0]	Predictor
S2	$S \rightarrow \cdot Aux NP VP$	[0,0]	Predictor
S3	$S \rightarrow \cdot VP$	[0,0]	Predictor
S4	$NP \rightarrow \cdot Pronoun$ $NP \rightarrow \cdot Proper-Noun$ $NP \rightarrow \cdot Det Nominal$	[0,0]	Predictor
S5		[0,0]	Predictor
S6		[0,0]	Predictor
S7 S8 S9 S10 S11	$VP \rightarrow \cdot Verb$ $VP \rightarrow \cdot Verb NP$ $VP \rightarrow \cdot Verb NP PP$ $VP \rightarrow \cdot Verb PP$ $VP \rightarrow \cdot VP PP$	[0,0] [0,0] [0,0] [0,0]	Predictor Predictor Predictor Predictor Predictor

## Chart[1]

S12	Verb → book •	[0,1]	Scanner
S13 S14 S15 S16	$VP \rightarrow Verb \cdot VP \rightarrow Verb \cdot NP \ VP \rightarrow Verb \cdot NP PP \ VP \rightarrow Verb \cdot PP$	[0,1] [0,1] [0,1]	Completer Completer Completer Completer
S17	$S \rightarrow VP$ .	[0,1]	Completer
S18	$VP \rightarrow VP \cdot PP$	[0,1]	Completer
S19 S20 S21 S22	$NP \rightarrow \cdot Pronoun$ $NP \rightarrow \cdot Proper-Noun$ $NP \rightarrow \cdot Det Nominal$ $PP \rightarrow \cdot Prep NP$	[1,1] [1,1] [1,1] [1,1]	Predictor Predictor Predictor Predictor

S0:  $\gamma \to S[0,0]$ 

S0: 
$$\gamma \rightarrow \cdot S[0,0]$$

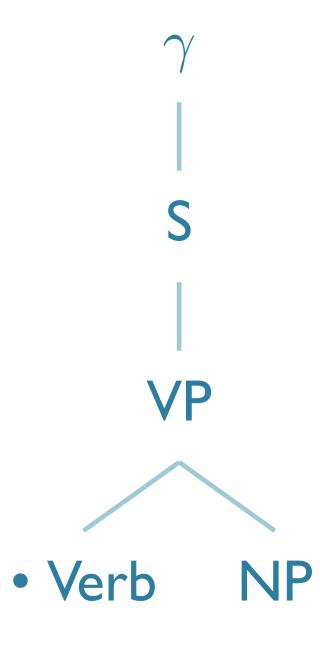
S3: 
$$S \rightarrow VP[0,0]$$



S0:  $\gamma \to S[0,0]$ 

S3:  $S \rightarrow VP[0,0]$ 

S8: *VP* → • *Verb NP* [0,0]

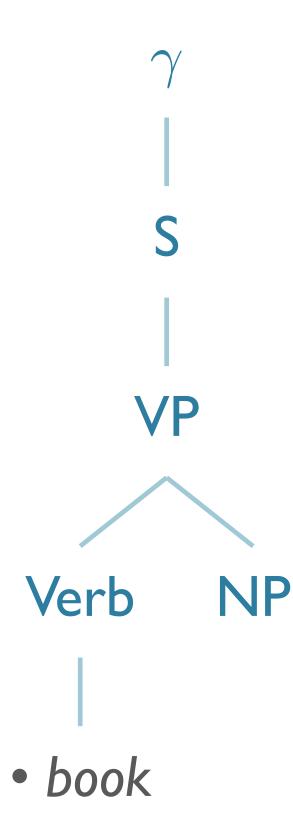


```
S0: \gamma \to S[0,0]
```

S3:  $S \rightarrow VP[0,0]$ 

S8: *VP* → • *Verb NP* [0,0]

S12: *Verb* → • *book* [0,0]

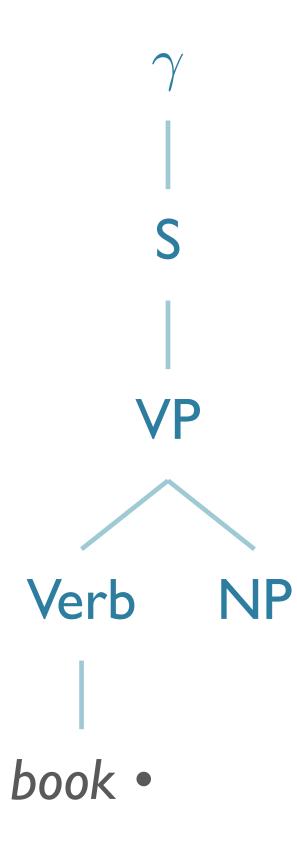


```
S0: \gamma \to S[0,0]
```

S3:  $S \rightarrow VP[0,0]$ 

S8: *VP* → • *Verb NP* [0,0]

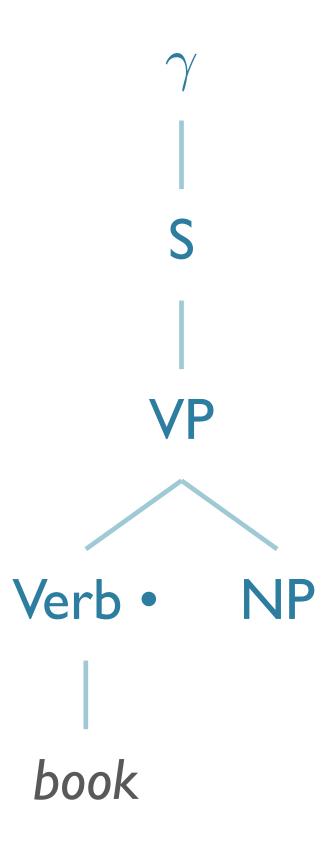
S12: *Verb* → *book* • [0,1]



```
S0: \gamma \to S[0,0]
```

S3:  $S \rightarrow VP[0,0]$ 

S8: *VP* → *Verb* • *NP* [0,1]



```
S0: \gamma \to S[0,0]
```

S3:  $S \rightarrow VP \cdot [0,1]$ 

S8: *VP* → *Verb* • *NP* [0,1]



Det

book

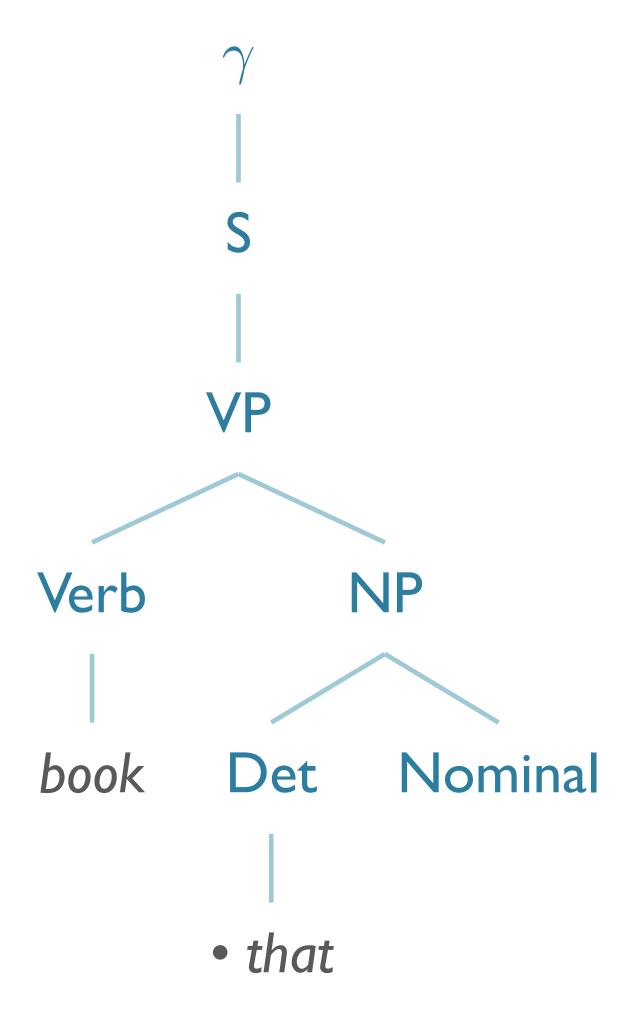
Nominal

```
S0: \gamma \rightarrow \cdot S[0,0] \gamma

S3: S \rightarrow VP \cdot [0,1] | S8: VP \rightarrow Verb \cdot NP[0,1] S S21: NP \rightarrow \cdot Det Nominal[1,1] VP
```

```
S0: \gamma \to S[0,0]
S3: S \rightarrow VP \cdot [0,1]
S8: VP → Verb • NP [0,1]
S21: NP → • Det Nominal [1,1]
```

S23: *Det* → • "that" [1,1]

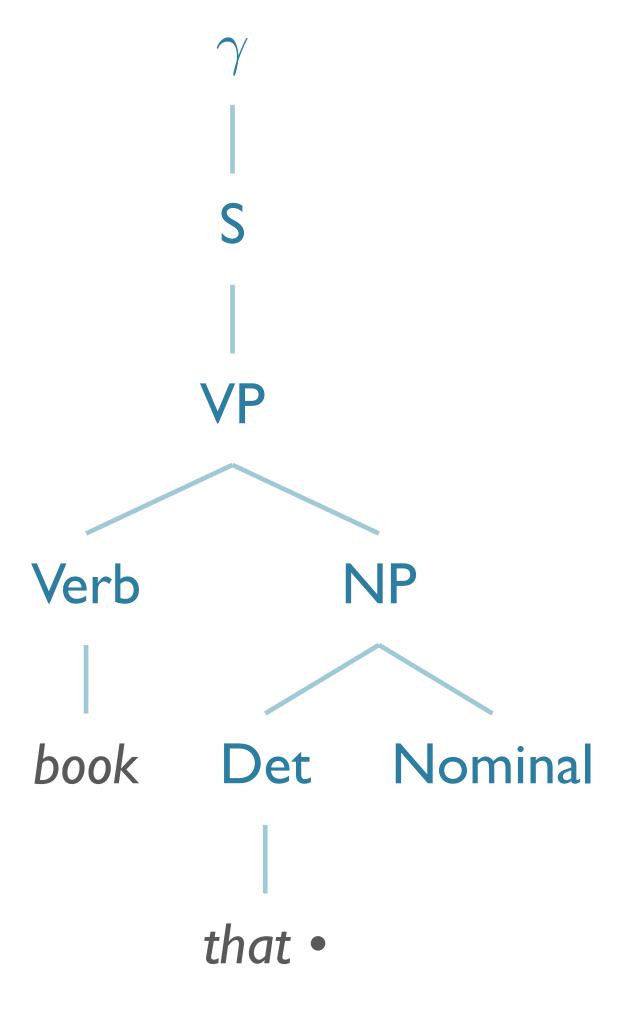


```
S0: \gamma \to \cdot S[0,0]
S3: S \to VP \cdot [0,1]
```

S8: *VP* → *Verb* • *NP* [0,1]

S21: *NP* → • *Det Nominal* [1,1]

S23: *Det* → "that" • [1,2]



that

```
S0: \gamma \to S[0,0]
S3: S \rightarrow VP \cdot [0,1]
S8: VP → Verb • NP [0,1]
S21: NP → Det • Nominal [1,2]
                                              VP
                                      Verb
                                                    NP
                                                      Nominal
                                      book
                                              Det •
```

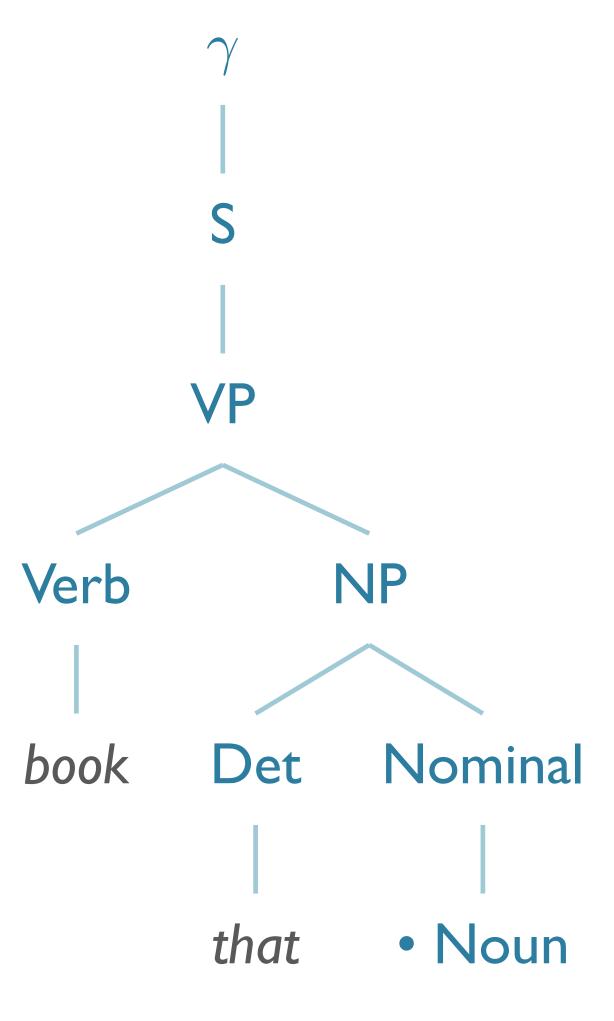
```
S0: \gamma \rightarrow \cdot S[0,0]
```

S3:  $S \rightarrow VP \cdot [0,1]$ 

S8: *VP* → *Verb* • *NP* [0,1]

S21: *NP* → *Det* • *Nominal* [1,2]

S25: *Nominal* → • *Noun* [2,2]



```
S0: \gamma \rightarrow \cdot S[0,0]
```

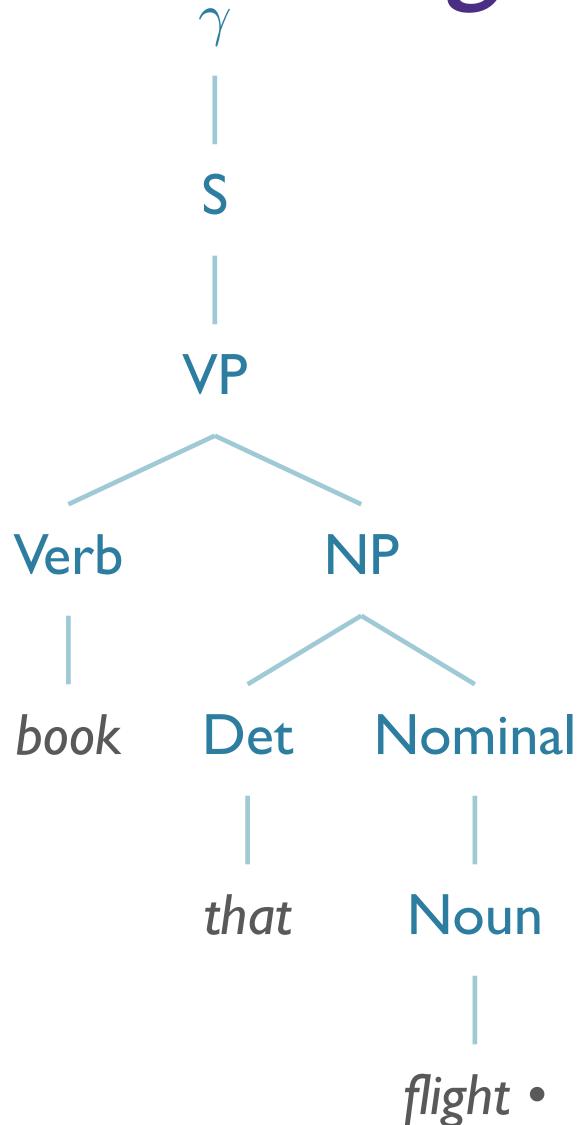
S3: 
$$S \rightarrow VP \cdot [0,1]$$

S8: *VP* → *Verb* • *NP* [0,1]

S21: *NP* → *Det* • *Nominal* [1,2]

S25: *Nominal* → • *Noun* [2,2]

S28: *Noun* → "*flight*" • [2,3]



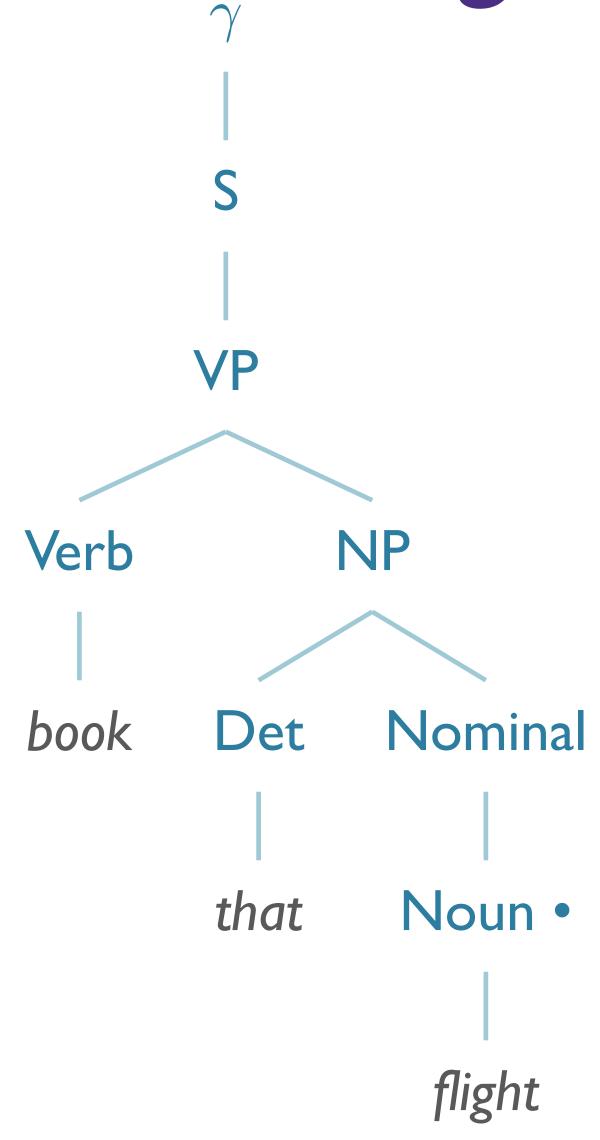
```
S0: \gamma \rightarrow \cdot S[0,0]
```

S3:  $S \rightarrow VP \cdot [0,1]$ 

S8: *VP* → *Verb* • *NP* [0,1]

S21: *NP* → *Det* • *Nominal* [1,2]

S25: *Nominal* → *Noun* • [2,3]

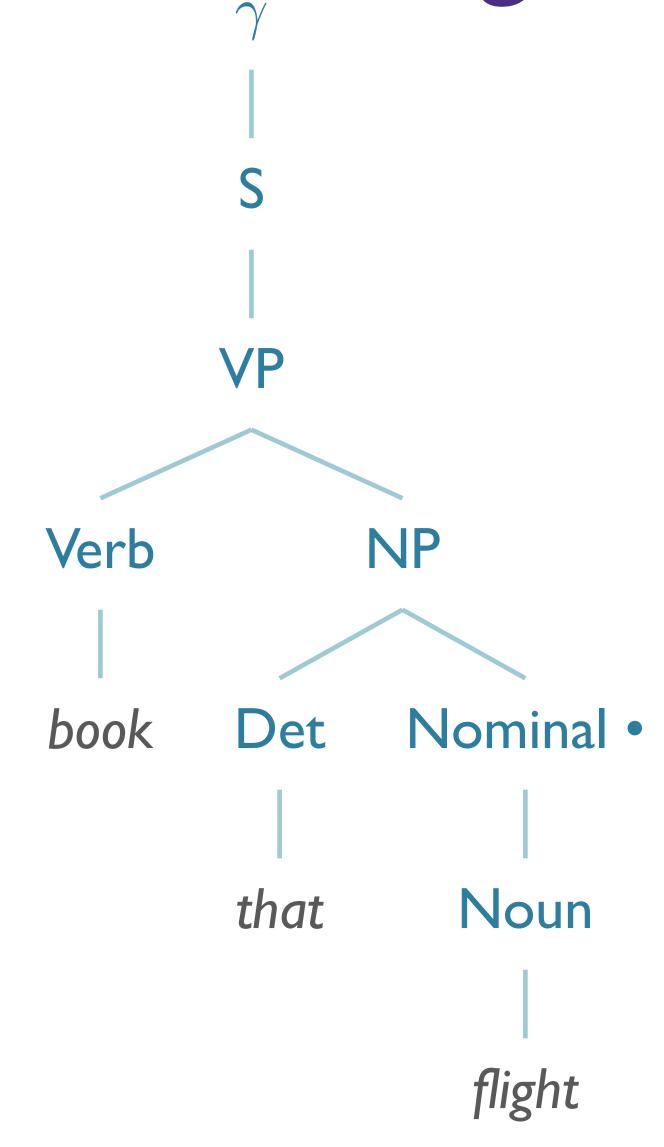


S0:  $\gamma \rightarrow \cdot S[0,0]$ 

S3:  $S \rightarrow VP \cdot [0,1]$ 

S8: *VP* → *Verb* • *NP* [0,1]

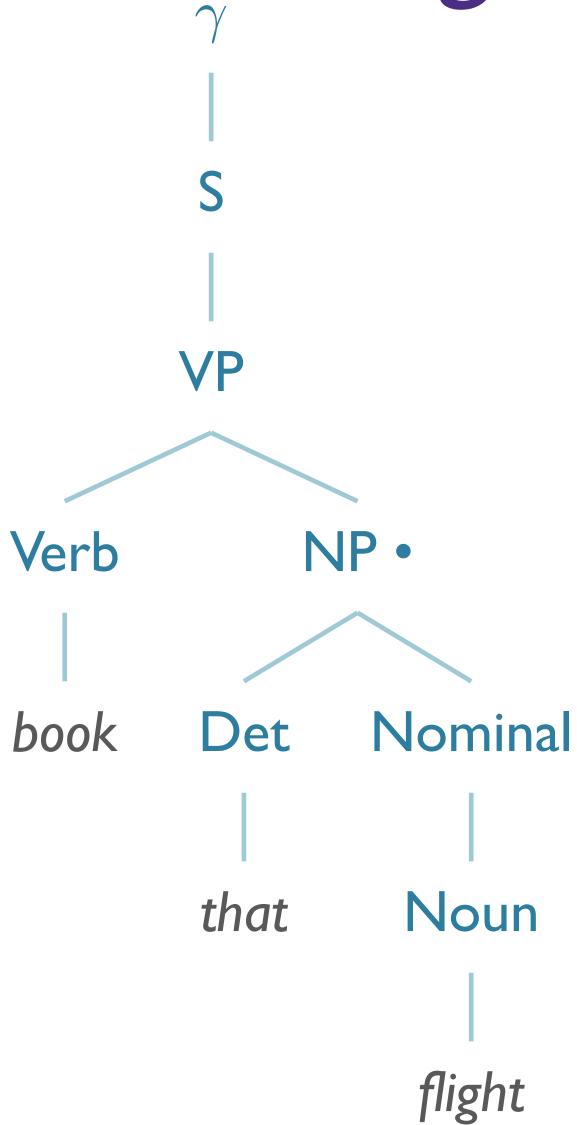
S21: *NP* → *Det Nominal* • [1,3]



S0:  $\gamma \to S[0,0]$ 

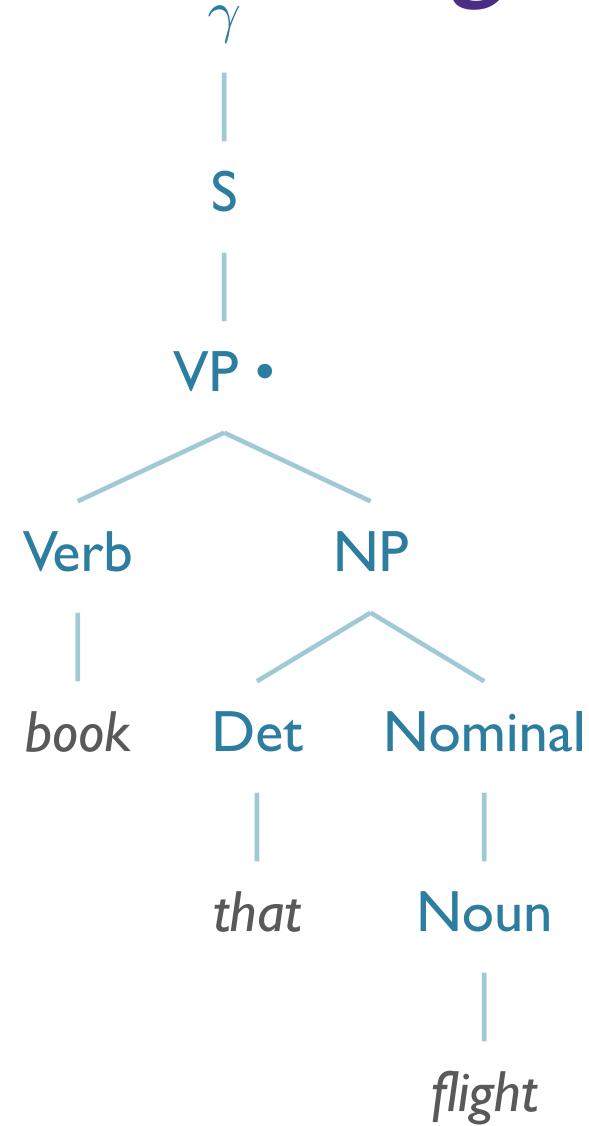
S3:  $S \rightarrow VP \cdot [0,1]$ 

S8: *VP* → *Verb NP* • [0,3]



S0:  $\gamma \to S[0,0]$ 

S3:  $S \rightarrow VP \cdot [0,3]$ 



## What About Dead Ends?

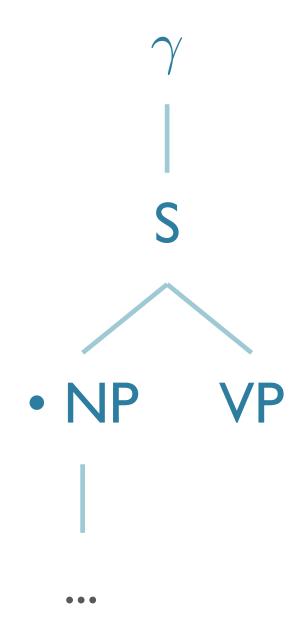
S0:  $\gamma \to S[0,0]$ 

S1:  $S \rightarrow \cdot NP VP$  [0,0]

*NP* → • *Pronoun* 

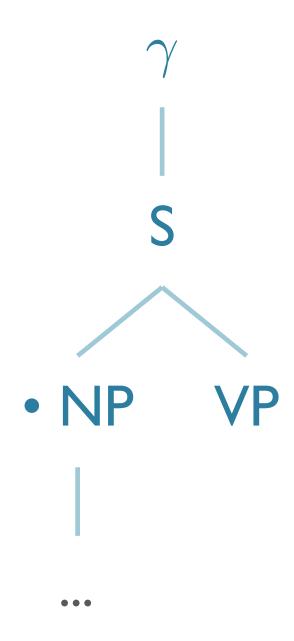
NP → • Proper-Noun

NP → • Det Nominal



book

```
S0: \gamma \to \cdot S[0,0]
S1: S \to \cdot NP VP[0,0]
```



book

 We now have a top-down parser in hand. Does it enter infinite loops on rules like S -> S 'and' S?

- We now have a top-down parser in hand. Does it enter infinite loops on rules like S -> S 'and' S?
- No!

```
procedure ENQUEUE(state, chart-entry)
  if state is not already in chart-entry then
    Push(state, chart-entry)
  end
```

- We now have a top-down parser in hand. Does it enter infinite loops on rules like S -> S 'and' S?
- No!

```
procedure ENQUEUE(state, chart-entry)
if state is not already in chart-entry then
    Push(state, chart-entry)
end
```

**Exercise**: parse 'table and chair' using the very simple grammar Nom -> Nom 'and' Nom | 'table' | 'chair'

# HW#3

# CKY Parsing: Goals

Complete implementation of CKY parser

Implement dynamic programming approach

Incorporate/follow backpointers to recover parse

# Implementation

- Build full parser
- Can use any language, per course policies
- You may use existing data structures for rules, trees
   e.g. NLTK has nice tree data structure
   CKY algorithm must be your own
- Dynamic programming table filling crucial!
- Will use smaller grammar (similar to HW #1)
- Back to ATIS for HW #4

# Implementation

- For CKY Implementation:
  - NLTK's **cfg.productions()** method:
    - optional rhs= argument only looks at first token of RHS

#### Notes

• Teams:

You may work in teams of two on this assignment

Test grammar

Pre-converted to CNF

Start symbol: TOP

Parse should span input and be rooted at: TOP

## Some Collaboration Basics

### Git Branches

Good for semi-isolating your development code from the shared, reviewed

code heart\_glasses branch master branch master branch

#### Recommended Git Flow

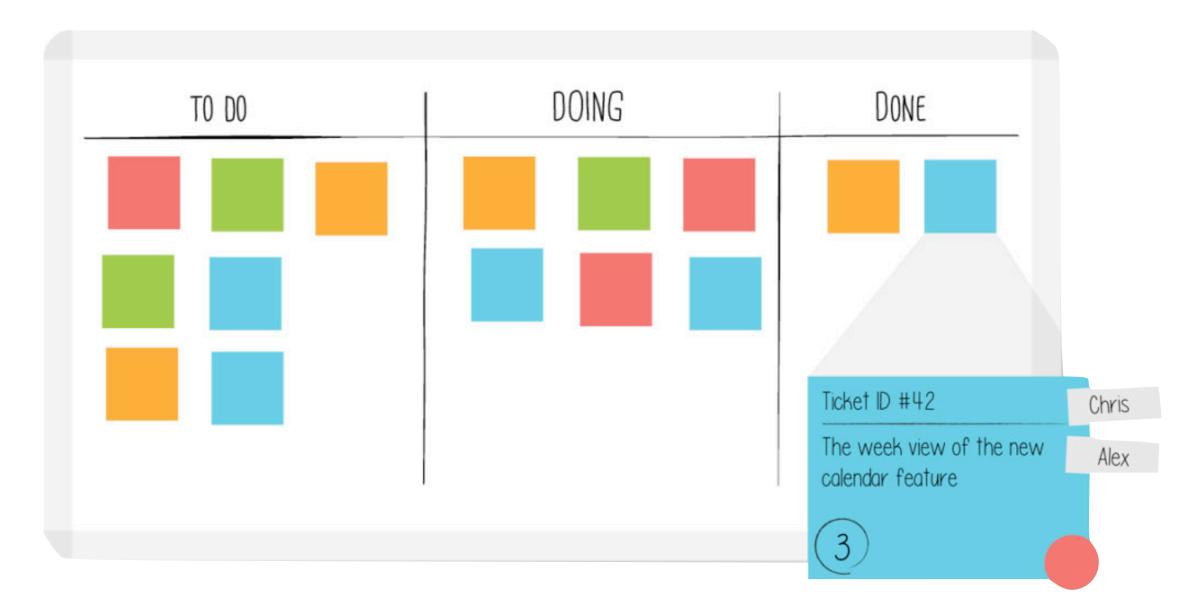
- Initialize a git repository, with a master branch
  - (Create initial checkin, if necessary)
- Create a new branch, maybe "adding\_rule\_objects"
- Make regular checkins on your branch (like saving)
- Switch to master branch, and "pull"
- Merge your branch to master
- ...rinse & repeat
- If using GitHub (or GitLab, etc): MUST BE PRIVATE REPO!

### Communication: Check-ins

- For check-ins, three main points:
  - What have you been working on?
  - What do you plan to work on next?
  - Is there anything "blocking" you?
- In industry, these brief check-ins among small teams are often done daily

# Project Planning: Kanban Boards

- Before you start working:
  - Write out tasks on sticky notes.
  - Place in three columns:
    - To-Do
    - Doing
    - Done



- As you work, you can move them from column to column
- Add tasks as new issues come up
- trello.com has free online implementation of Kanban Boards