

# Computer Organization & Assembly Languages

Midterm Exam - 2005/11/30

Dept. of Engineering Science, National Cheng Kung University

## 1. [24 points] General Knowledge

(A) Short Answer: Assume using the MIPS architecture.

- (1) How many bits are in a word? 32 8
- (2) How many bits are used to represent a "char" in C? 8
- (3) How many bits are used to represent a "double" in C? 64
- (4) How many words of memory are addressable by the MIPS?  $2^{32}$
- (5) One of the most important examples of an abstraction is the interface between hardware and the lowest-level software. What is this abstraction called? 機器語言
- ✓ (6) List the five class components of a computer as defined by the textbook authors.

(B) True/False: Write your response. Ambiguous responses will be marked wrong.

- (7) True/False Assembly language pseudo-operations are executable.
- (8) True/False The stack in MIPS grows from lower address to higher address.
- (9) True/False Conditional branches in MIPS employ PC-relative addressing.
- (10) True/False J-format instructions support conditional branching.
- (11) True/False The lw instruction assembles into two machine code instructions.
- (12) True/False The jump-and-link instruction (jal) in MIPS forms a link to the calling site by storing the address of next instruction in the register \$ra.

## 2. [6 points] Fill in the following blanks.

R-format:

op	rs	rt	rd	shamt	funct
6 bits	5 bits	5 bits	5 bits	5 bits	6 bits

I-format:

op	rs	rt	Addr/immediate
6 bits	5 bits	5 bits	16 bits

J-format:

op	address
6 bits	26 bits

3. [5 points] Suppose die area is  $0.4 \text{ cm}^2$  and there are 4 defects per  $\text{cm}^2$ . Calculate the yield. Then calculate the yield if defects per area can be cut in half. Note that the answers should be in the form of xx.xx%.

$$\text{Yield} = \frac{1}{\left(1 + \text{DefectsPerArea} * \frac{\text{DieArea}}{2}\right)^2}$$

4. [15 points] Translating MIPS code

Here is a mystery function. It expects one non-negative integer argument and returns one integer result.

yuck:

```

    li      $t0, 1
    li      $t1, 1
loop:
    blt    $a0, 2, exit
    add    $t2, $t0, $t1
    move   $t0, $t1
    move   $t1, $t2
    addi   $a0, $a0, -1
    j      loop
exit:
    move   $v0, $t1
    jr    $ra

```

- (1) Translate this yuck function into C. Assume that the usage of registers is known as (n: \$a0 & i, j, k: \$t0~\$t2) where n, i, j, k are all C variables. (12 points)
- (2) What will this function return if it is called with an argument of 4 (i.e; \$a0=4)? (3 points)

5. [15 points] Arithmetic

- (1) Convert -28410 to two's complement. Show the 32-bit binary representation.
- (2) Convert 0xCA50 to binary.
- (3) Convert 0xFFFF to decimal, interpreting it as a signed 16-bit integer.
- (4) Subtract 0x001F from 0xFFFF, interpreting each as 2's complement 16-bit integer. Express the answer both in hex and decimal.
- (5) Show the IEEE 754 representation of  $-1.25_{10}$  in single precision.

6. [8 points] Execute the following MIPS code fragments, showing the changes that occur in the register file and in memory. You only need to show the changes.

(1)

```

addi $19, $0, 0x20
lw    $17, 0x04($19)
add  $20, $19, $16
sw    $20, 0x08($19)

```

Registers	BEFORE	AFTER	Memory	BEFORE	AFTER
S16	0x10	<del>0x10</del>	0x20	0x22	x
S17	0x14	<del>0x14</del>	0x24	0x30	x
S18	0x16	x	0x28	0x40	20
S19	0x28	<del>0x28</del>	0x2C	0x50	x
S20	0x1234	<del>0x1234</del>	0x30	0x60	x
Registers	BEFORE	AFTER	Memory	BEFORE	AFTER
S0	0x00	<del>0x00</del>	0x20	0x10	x
S1	0x14	<del>0x14</del>	0x24	0x30	x
S2	0x16	<del>0x16</del>	0x28	0x40	x
S3	0x28	x	0x2C	0x50	x
S4	0x1234	x	0x30	0x60	x

(2)

```

addi $1, $0, 0x20
lw    $2, 0x04($3)
add  $0, $3, $1
bne  $0, $1, loop

```

(3) Is the branch taken? (answer YES or NO)

7. [5 points] Branch offset computation problem

Fill in the missing line of MIPS assembly code so that it branches to the instruction labeled loop when the contents of the \$t0 register are nonzero.

loop:

```

lbu   $t0, 0($s0)    # get byte pointed to by $s0
addi $s0, $s0, 1       # increment pointer to next byte
andi $t0, $t0, 127     # mask off most significant bit
  

sub   $s2, $s0, $s1     # subtract end pointer from start

```

8. [12 points] Compile the following C program into MIPS instructions. Assume that the usage of registers is specified as (g, h, i, j: \$a0~\$a3 and f: \$s0).

```

int leaf_example (int g, int h, int i, int j){
    int f;
    f = (g+h)-(i+j);
    return f;
}

```

leaf example:

```

addi $sp, $sp, -4
sw    $s0, 0($sp)
add  $t0, $a0, $a1
add  $t1, $a0, $a3
add  $s0, $t0, $t1
move $v0, $s0
lw    $s0, 0($sp)
jr  $ra

```

```

addi $sp, $sp, -4
sw    $s0, 0($sp)
add  $t0, $a0, $a0
add  $t1, $a2, $a3
add  $s0, $t0, $t1
move $v0, $s0

```

9. [10 points] Here is a C fragment that repeatedly calls sequence:

```
void main() {
    int t;
    int sum = 0;
    for(t = 0; t < 5; t++)
        sum = sum + sequence(t);
}
```

This might translate into the following assembly code:

```
main:
    sw    $ra, 0($sp)
    addi $sp, $sp, -4
    # Initialize sum = 0 and t = 0
    add   $s0, $0, $0  # $s0 = sum
    addi $t0, $0, 0    # $t0 = 0
    # Call sequence(t)

loop:
    move  $a0, $t0
    jal   sequence
    # Increment loop index t
    addi $t0, $t0, 1
    # Compute sum = sum + sequence(t)
    add   $s0, $s0, $v0
    # Repeat while (t < 5)
    slti  $t1, $t0, 5
    bne   $t1, $0, loop
    addi $sp, $sp, 4
    lw    $ra, 0($sp)
    jr   $ra
```

This code fragment does not correctly follow the MIPS *procedure calling convention*. What is wrong? Show the changes required to make this code observe the calling convention.

add 580  
sub 0  
lw 59  
sw  
jr