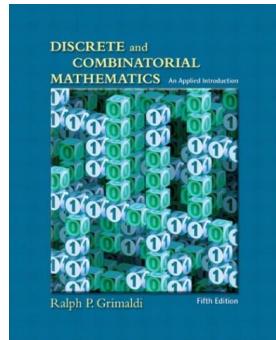
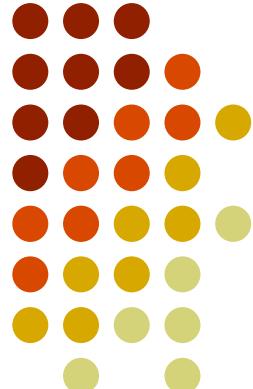


Discrete Mathematics

-- *Chapter 10: Recurrence Relations*



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First glance at “recurrence”

$$F_{n+2} = F_{n+1} + F_n$$

$$\underline{a_{n+1}} = 3\underline{a_n}$$

$$a_n = A^* 3^n$$



Outline

- The first-order **linear** recurrence relation
- The second-order **linear homogeneous** recurrence relation with constant coefficients
- The **nonhomogeneous** recurrence relation
- The method of generating functions

Key1: solve simple recurrence relation

Key2: model recurrence relation



The First-Order Linear Recurrence Relation

- The equation $a_{n+1} = 3a_n$ is a recurrence relation with constant coefficients. Since a_{n+1} *only depends on its immediate predecessor*, the relation is said to be first order.
- The expression $a_0 = A$, where A is a constant, is referred to as an initial condition.
- The **unique** solution of the recurrence relation $a_{n+1} = da_n$, where $n \geq 0$, d is a constant, and $a_0 = A$, is given by $a_n = Ad^n$.



The First-Order Linear Recurrence Relation

- **Ex 10.1** : Solve the recurrence relation $a_n = 7a_{n-1}$, where $n \geq 1$ and $a_2 = 98$.
 - $a_n = a_0(7^n)$, $a_2 = 98 = a_0(7^2) \Rightarrow a_0 = 2$, $a_n = 2(7^n)$.
- **Ex 10.2** : A bank pay 6% annual interest on savings, compounding the interest monthly. If we deposit \$1000, how much will this deposit be worth a year later?
 - $p_{n+1} = (1.005)p_n$, $p_0 = 1000 \Rightarrow p_n = p_0(1.005)^n$
 - $p_{12} = 1000(1.005)^{12} = \1061.68



The First-Order Linear Recurrence Relation

- Refer to examples 1.37, 3.11, 4.12, and 9.12.
- Ex 10.3 : Let a_n count the number of compositions of n , we find that

$$a_{n+1} = 2a_n, n \geq 1, a_1 = 1 \Rightarrow a_n = 2^{n-1}$$

		(1')	4
		(2')	1 + 3
(1)	3	(3')	2 + 2
(2)	1 + 2	(4')	1 + 1 + 2
(3)	2 + 1		
(4)	1 + 1 + 1	(1'')	3 + 1
		(2'')	1 + 2 + 1
		(3'')	2 + 1 + 1
		(4'')	1 + 1 + 1 + 1

Figure 10.1



The First-Order Linear Recurrence Relation

- The recurrence relation $a_{n+1} - da_n = 0$ is called linear because each term appears to the first power.
- Sometimes a nonlinear recurrence (e.g., $a_{n+1} - 3a_n a_{n-1} = 0$) relation can be transformed to a linear one by a suitable algebraic substitution.
- Ex 10.4 : Find a_{12} if $a_{n+1}^2 = 5a_n^2$ where $a_n > 0$ for $n \geq 0$ and $a_0 = 2$.
 - Let $b_n = a_n^2$. Then $b_{n+1} = 5b_n$ (linear) for $n \geq 0$ and $b_0 = 4 \Rightarrow b_n = 4 \cdot 5^n$



Homogeneous and Nonhomogeneous

- The general first-order linear recurrence relation with constant coefficients has the form

$$a_{n+1} + ca_n = f(n).$$

- $f(n) = 0$, the relation is called homogeneous.
 - Otherwise, it is called nonhomogeneous.
- **Ex 10.5** : Let a_n denote the number of comparisons needed to sort n numbers in bubble sort, we find the recurrence relation
 - $a_n = a_{n-1} + (n - 1)$, $n \geq 2$, $a_1 = 0$



$i = 1$	x_1	7	7	7	7	2
	x_2	9	9	9	2	7
	x_3	2	2	2	9	9
	x_4	5	5	5	5	5
	x_5	8	8	8	8	8

Four comparisons and two interchanges.

$i = 2$	x_1	2	2	2	2	
	x_2	7	7	7	5	
	x_3	9	9	5	7	
	x_4	5	5	9	9	
	x_5	8	8	8	8	

Three comparisons and two interchanges.

$i = 3$	x_1	2	2	2		
	x_2	5	5	5		
	x_3	7	7	7		
	x_4	9	8	8		
	x_5	8	9	9		

Two comparisons and one interchange.

$i = 4$	x_1	2				
	x_2	5				
	x_3	7				
	x_4	8				
	x_5	9				

One comparison but no interchanges.

Figure 10.3



The First-Order Linear Recurrence Relation

- **Ex 10.6 :** In Example 9.6 we sought the generating function for the sequence 0, 2, 6, 12, 20, 30, 42,..., due to the observation $a_n = n^2 + n$. If we fail to see this, alternatively

$$\begin{aligned}a_1 - a_0 &= 2 \\a_2 - a_1 &= 4 \\a_3 - a_2 &= 6 \\\vdots &\quad \vdots \quad \vdots \\a_n - a_{n-1} &= 2n.\end{aligned}$$

$$\begin{aligned}a_n - a_0 &= 2 + 4 + 6 + \cdots + 2n = 2(1 + 2 + 3 + \cdots + n) \\&= 2[n(n+1)/2] = n^2 + n.\end{aligned}$$



The First-Order Linear Recurrence Relation

- **Ex 10.7** : Solve the relation $a_n = n \cdot a_{n-1}$, $n \geq 1$, $a_0 = 1$.

$$\begin{matrix} & 1 & 2 \\ 2 & & 1 \end{matrix}$$

x 2

$$\begin{matrix} & 1 & 2 & 3 \\ 1 & & 3 & 2 \\ 3 & 1 & & 2 \\ 3 & 2 & & 1 \\ 2 & & 3 & 1 \\ 2 & & 1 & 3 \end{matrix}$$

x 3

The Second-Order Linear Homogeneous Recurrence Relation with Constant Coefficients



- Linear recurrence relation of order k :
 - $C_0a_n + C_1a_{n-1} + C_2a_{n-2} + \dots + C_ka_{n-k} = f(n), n \geq 0.$
- Homogeneous relation of order 2:
 - $C_0a_n + C_1a_{n-1} + C_2a_{n-2} = 0, n \geq 2.$
- Substituting $a_n = cr^n$ into the equation, we have
 - $C_0cr^n + C_1cr^{n-1} + C_2cr^{n-2} = 0, n \geq 2.$
 - Characteristic equation: $C_0r^2 + C_1r + C_2 = 0, n \geq 2.$
- The roots r_1, r_2 of this equation are called characteristic roots.
- Three cases for the roots:
 - (A) distinct real roots
 - (B) complex conjugate roots
 - (C) equivalent real roots



Case (A): Distinct Real Roots

- **Ex 10.9** : Solve the recurrence relation $a_n + a_{n-1} - 6a_{n-2} = 0$, $n \geq 2$, and $a_0 = -1$ and $a_1 = 8$.

- **Solution**

Let $a_n = cr^n$

$$r^2 + r - 6 = 0 \Rightarrow r = 2, -3$$

$$a_n = c_1(2)^n + c_2(-3)^n$$

$$-1 = a_0 = c_1 + c_2$$

$$8 = a_1 = 2c_1 - 3c_2 \Rightarrow c_1 = 1, c_2 = -2$$

$$\Rightarrow a_n = (2)^n - 2(-3)^n$$

$a_n = 2^n$ and $a_n = (-3)^n$ are both solutions

Linearly independent solutions

$$2^n + 2^{n-1} - 6 * 2^{n-2} = 2^{n-2}(2^2 + 2 - 6) = 0$$



Distinct Real Roots

- **Ex 10.10** : Solve Fibonacci relation, $F_{n+2} = F_{n+1} + F_n$, $n \geq 0$, $F_0 = 0$, $F_1 = 1$.

• Solution

Let $F_n = cr^n$,

$$r^2 - r - 1 = 0 \Rightarrow r = (1 \pm \sqrt{5})/2$$

$$F_n = c_1 \left(\frac{1 + \sqrt{5}}{2} \right)^n + c_2 \left(\frac{1 - \sqrt{5}}{2} \right)^n.$$

$$F_n = \frac{1}{\sqrt{5}} \left[\left(\frac{1 + \sqrt{5}}{2} \right)^n - \left(\frac{1 - \sqrt{5}}{2} \right)^n \right], \quad n \geq 0.$$

a3=5 \rightarrow $\Phi, \{1\}, \{2\}, \{3\}, \{1,3\}$

a4=8 \rightarrow $\Phi, \{1\}, \{2\}, \{3\}, \{4\}, \{1,3\}, \{2,4\}, \{1,4\}$

a5=13 \rightarrow $\Phi, \{1\}, \{2\}, \{3\}, \{4\}, \{1,3\}, \{2,4\}, \{1,4\}, \underline{\{5\}}, \underline{\{1,5\}}, \underline{\{2,5\}}, \underline{\{3,5\}}, \underline{\{1,3,5\}}$



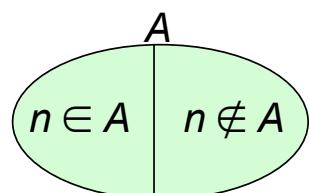
Distinct Real Roots

- **Ex 10.11** : For $n \geq 0$, let $S = \{1, 2, \dots, n\}$, and let a_n denote the number of subsets of S that contains no consecutive integers. Find and solve a recurrence relation for a_n .

- **Solution**

- $a_0 = 1$ and $a_1 = 2$ and $a_2 = 3$ and $a_3 = 5$ (Fibonacci?)
- If $A \subseteq S$ and A is to be counted in a_n , there are two cases
 - (1) $n \in A$, then $A - \{n\}$ would be counted in a_{n-2}
 - (2) $n \notin A$, then A would be counted in a_{n-1}
- $a_n = a_{n-1} + a_{n-2}$, $n \geq 2$

$$a_n = \frac{1}{\sqrt{5}} \left[\left(\frac{1+\sqrt{5}}{2} \right)^{n+2} - \left(\frac{1-\sqrt{5}}{2} \right)^{n+2} \right], \quad n \geq 0.$$



*Exhaustive and
mutually disjoint*



Distinct Real Roots

- **Ex 10.12 :** Suppose we have a $2 \times n$ chessboard. We wish to cover such a chessboard using 2×1 vertical dominoes or 1×2 horizontal dominoes.

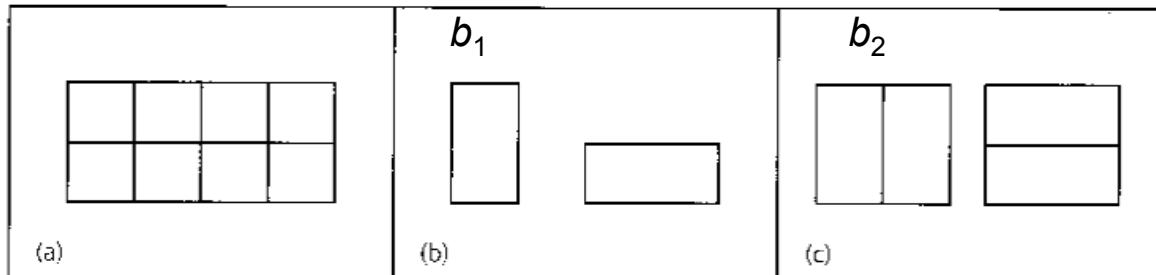


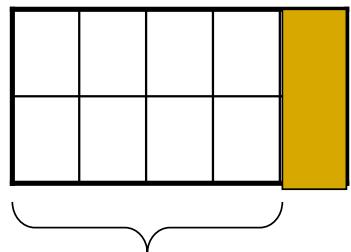
Figure 10.5



Distinct Real Roots

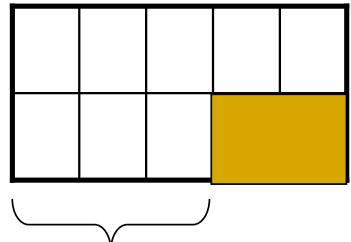
- Let b_n count the number of ways we can cover a $2 \times n$ chessboard by using 2×1 vertical dominoes or 1×2 horizontal dominoes.
- $b_1 = 1$ and $b_2 = 2$
- For $n \geq 3$, consider the last (n -th) column of the chessboard
 - By one 2×1 vertical domino: Now the remaining $2 \times (n - 1)$ subboard can be covered in b_{n-1} ways.
 - By two 1×2 horizontal dominoes, place one above the other: Now the remaining $2 \times (n - 2)$ subboard can be covered in b_{n-2} ways.
- $b_n = b_{n-1} + b_{n-2}$, $n \geq 3$, $b_1 = 1$ and $b_2 = 2$
 $\Rightarrow b_n = F_{n+1}$

$$2n = 2 \times (n - 1) + 2 \times 1$$



b_{n-1} *Disjoint?*

$$2n = 2 \times (n - 2) + 2(1 \times 2)$$



b_{n-2}



Distinct Real Roots

- **Ex 10.14** : Suppose the symbols of legal arithmetic expressions include 0, 1,..., 9, +, *, /.
- Let a_n be the number of legal arithmetic expressions that are made up of n symbols. Solve $a_1=10$ and $a_2=100$

Solution:

- For $n \geq 3$, consider the two cases of length $n - 1$:
 - 1) If x is an expression of $n - 1$ symbols, add a digit to the right of x . $\Rightarrow 10a_{n-1}$ 補上 0, 1, 2, ..., 9
 - 2) If x is an expression of $n - 2$ symbols, we adjoin to the right of x one of the 29 two-symbol expressions: +0, ..., +9, *0, ..., *9, /1, /2, ..., /9. $\Rightarrow 29a_{n-2}$ 不能補2個數字，會與 $n - 1$ 相同
- $a_n = 10a_{n-1} + 29a_{n-2}, n \geq 3$

Idea: use a_{n-1} (or more) to represent a_n



Distinct Real Roots

- **Ex 10.15** (9.13): Palindromes are the compositions of numbers that read the same left to right as right to left.
- Let p_n count the number of palindromes of n .
- $p_n = 2p_{n-2}$, $n \geq 3$, $p_1=1$, $p_2=2$

p_3	p_5	p_4	p_6
(1) 3 (2) $1 + 1 + 1$	(1') 5 (2') $2 + 1 + 2$ (1'') $1 + 3 + 1$ (2'') $1 + 1 + 1 + 1 + 1$	(1) 4 (2) $1 + 2 + 1$ (3) $2 + 2$ (4) $1 + 1 + 1 + 1$	(1') 6 (2') $2 + 2 + 2$ (3') $3 + 3$ (4') $2 + 1 + 1 + 2$ (1'') $1 + 4 + 1$ (2'') $1 + 1 + 2 + 1 + 1$ (3'') $1 + 2 + 2 + 1$ (4'') $1 + 1 + 1 + 1 + 1 + 1$
(') Add 1 to the first and last summands			
(") Append "1+" to the start and "+1" to the end			

Figure 10.6



Distinct Real Roots

$$p_n = 2p_{n-2}, \quad n \geq 3, \quad p_1 = 1, \quad p_2 = 2.$$

Substituting $p_n = cr^n$, for $c, r \neq 0, n \geq 1$, into this recurrence relation, the resulting characteristic equation is $r^2 - 2 = 0$. The characteristic roots are $r = \pm\sqrt{2}$, so $p_n = c_1(\sqrt{2})^n + c_2(-\sqrt{2})^n$. From

$$1 = p_1 = c_1(\sqrt{2}) + c_2(-\sqrt{2})$$

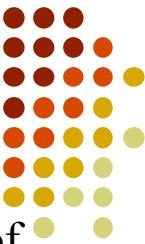
$$2 = p_2 = c_1(\sqrt{2})^2 + c_2(-\sqrt{2})^2$$

we find that $c_1 = \left(\frac{1}{2} + \frac{1}{2\sqrt{2}}\right)$, $c_2 = \left(\frac{1}{2} - \frac{1}{2\sqrt{2}}\right)$, so

$$p_n = \left(\frac{1}{2} + \frac{1}{2\sqrt{2}}\right)(\sqrt{2})^n + \left(\frac{1}{2} - \frac{1}{2\sqrt{2}}\right)(-\sqrt{2})^n, \quad n \geq 1.$$

we consider n even, say $n = 2k$

$$\begin{aligned} p_n &= \left(\frac{1}{2} + \frac{1}{2\sqrt{2}}\right)(\sqrt{2})^{2k} + \left(\frac{1}{2} - \frac{1}{2\sqrt{2}}\right)(-\sqrt{2})^{2k} \\ &= \left(\frac{1}{2} + \frac{1}{2\sqrt{2}}\right)2^k + \left(\frac{1}{2} - \frac{1}{2\sqrt{2}}\right)2^k = 2^k = 2^{n/2} \end{aligned}$$



Distinct Real Roots

- **Ex 10.16** : Find the number of recurrence relation for the number of binary sequences of length n that have no consecutive 0's.
 - Let a_n be the number of such sequences of length n .
 - Let $a_n^{(0)}$ count those end in 0, and $a_n^{(1)}$ count those end in 1
 $\Rightarrow a_n = a_n^{(1)} + a_n^{(0)}$
結尾為1，能為1、0
結尾為0，只能為1
 - Consider x of length $n - 1$
 - If x ends in 1, we can append a 0 or a 1 to it ($2a_{n-1}^{(1)}$).
 - If x ends in 0, we can append a 1 to it ($a_{n-1}^{(0)}$).
 - $a_n = 2a_{n-1}^{(1)} + a_{n-1}^{(0)} = a_{n-1}^{(1)} + \underline{a_{n-1}^{(1)} + a_{n-1}^{(0)}} \rightarrow a_{n-1}$
 - If y is counted in a_{n-2} \Leftrightarrow sequence $y1$ is counted in $a_{n-1}^{(1)}$
 - So, $a_{n-2} = a_{n-1}^{(1)}$.
 - $a_n = a_{n-1} + a_{n-2}, n \geq 3, a_1 = 2, a_2 = 3$

- Try **Ex10.17**



Second- or Higher-Order Recurrence Relation

- **Ex 10.18 :** Solve $2a_{n+3} = a_{n+2} + 2a_{n+1} - a_n$, $n \geq 0$, $a_0 = 0$, $a_1 = 1$, $a_2 = 2$
 - Let $a_n = cr^n$
 - Characteristic equation: $2r^3 - r^2 - 2r + 1 = 0 \Rightarrow r = 1, 1/2, -1$
 - $a_n = c_1(1)^n + c_2(-1)^n + c_3(1/2)^n$
- From $a_0 = 0$, $a_1 = 1$, $a_2 = 2$, derive $c_1 = 5/2$, $c_2 = 1/6$, $c_3 = -8/3$
- $a_n = (5/2) + (1/6)(-1)^n + (-8/3)(1/2)^n$



Second- or Higher-Order Recurrence Relation

- **Ex 10.19 :** We want to tile a $2 \times n$ chessboard using two types of tiles shown in Figure 10.8.

a_2 : 2×2 chessboard

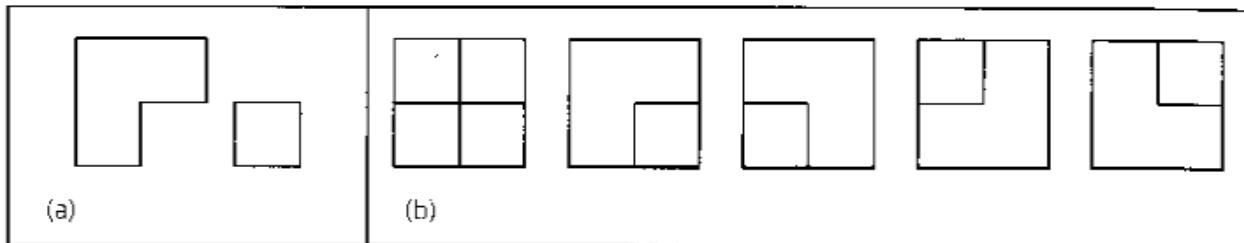


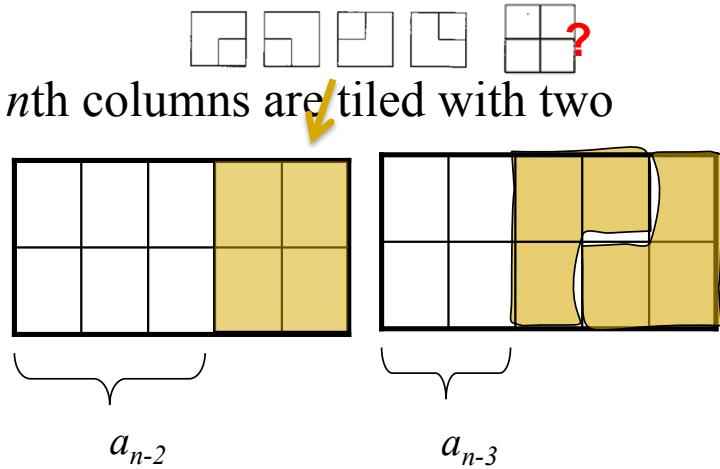
Figure 10.8

a_3 : 2×3 chessboard



Second- or Higher-Order Recurrence Relation

- Let a_n count the number of such tilings.
- $a_1=1$ and $a_2=5$ and $a_3=11$ (try it)
- For $n \geq 4$, consider the last column of the chessboard
 - the n th column is covered by two 1×1 tiles $\Rightarrow a_{n-1}$
 - the $(n - 1)$ st and the n th column are tiled with one 1×1 tile and a larger tile $\Rightarrow 4a_{n-2}$
 - the $(n - 2)$ nd, $(n - 1)$ st and the n th columns are tiled with two large tiles $\Rightarrow 2a_{n-3}$
- $a_n = a_{n-1} + 4a_{n-2} + 2a_{n-3}$, $n \geq 4$





Case (B): Complex Roots

- DeMoivre's Theorem:

$$(\cos\theta + i \sin\theta)^n = \cos n\theta + i \sin n\theta$$

- If $z = x + iy \in \mathbf{C} \Rightarrow z = r(\cos\theta + i \sin\theta)$, $r = \sqrt{x^2 + y^2}$, $\frac{y}{x} = \tan\theta$, for $x \neq 0$

If $x = 0$,

$$\begin{cases} y > 0, z = yi = y \sin(\pi/2) + iy \cos(\pi/2) = y(\cos(\pi/2) + i \sin(\pi/2)) \\ y < 0, z = -|y| + i|y| = |y|(\cos(3\pi/2) + i \sin(3\pi/2)) \end{cases}$$

By DeMoivre's Theorem, $z^n = r^n(\cos n\theta + i \sin n\theta)$



Complex Roots

- **Ex 10.20 :** Determine $(1 + \sqrt{3}i)^{10}$
 - Solution

Complex number $1 + \sqrt{3}i$ is represented as the point $(1, \sqrt{3})$ in the xy -plane

$$r = \sqrt{1^2 + (\sqrt{3})^2} = 2, \theta = \pi/3$$

$$1 + \sqrt{3}i = 2(\cos(\pi/3) + i \sin(\pi/3))$$

$$\begin{aligned}(1 + \sqrt{3}i)^{10} &= 2^{10}(\cos(10\pi/3) + i \sin(10\pi/3)) = 2^{10}(\cos(4\pi/3) + i \sin(4\pi/3)) \\ &= 2^{10}((-1/2) - (\sqrt{3}/2)i) = (-2^9)(1 + \sqrt{3}i)\end{aligned}$$

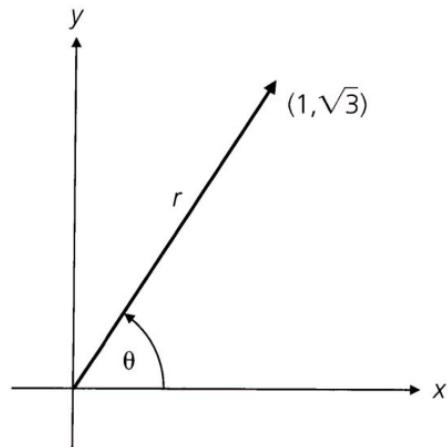


Figure 10.9



Complex Roots

- Ex 10.21 : Solve the recurrence relation $a_n = 2(a_{n-1} - a_{n-2})$ where $n \geq 2$, $a_0 = 1$, $a_1 = 2$.
 - Let $a_n = cr^n$
 - $r^2 - 2r + 2 = 0 \Rightarrow r = 1 \pm i$
 - $1 + i = \sqrt{2}(\cos(\pi/4) + i \sin(\pi/4))$
 - $1 - i = \sqrt{2}(\cos(\pi/4) - i \sin(\pi/4))$



$$\begin{aligned}a_n &= c_1(1+i)^n + c_2(1-i)^n \\&= c_1[\sqrt{2}(\cos(\pi/4) + i \sin(\pi/4))]^n + c_2[\sqrt{2}(\cos(-\pi/4) + i \sin(-\pi/4))]^n \\&= c_1(\sqrt{2})^n(\cos(n\pi/4) + i \sin(n\pi/4)) + c_2(\sqrt{2})^n(\cos(-n\pi/4) + i \sin(-n\pi/4)) \\&= c_1(\sqrt{2})^n(\cos(n\pi/4) + i \sin(n\pi/4)) + c_2(\sqrt{2})^n(\cos(n\pi/4) - i \sin(n\pi/4)) \\&= (\sqrt{2})^n[k_1 \cos(n\pi/4) + k_2 \sin(n\pi/4)],\end{aligned}$$

where $k_1 = c_1 + c_2$ and $k_2 = (c_1 - c_2)i$.

$$1 = a_0 = [k_1 \cos 0 + k_2 \sin 0] = k_1$$

$$2 = a_1 = \sqrt{2}[1 \cdot \cos(\pi/4) + k_2 \sin(\pi/4)], \text{ or } 2 = 1 + k_2, \text{ and } k_2 = 1.$$

The solution for the given initial conditions is then given by

$$a_n = (\sqrt{2})^n[\cos(n\pi/4) + \sin(n\pi/4)], \quad n \geq 0.$$



Complex Roots

- **Ex 10.22** : Let a_n denote the value of the $n \times n$ determinant

D_n

- $a_1 = b$, $a_2 = 0$ and $a_3 = -b^3$
- $D_n = bD_{n-1} - b^2D_{n-2}$
- $a_n = ba_{n-1} - b^2a_{n-2}$

$$a_1 = |b| = b \quad \text{and} \quad a_2 = \begin{vmatrix} b & b \\ b & b \end{vmatrix} = 0 \quad (\text{and} \quad a_3 = \begin{vmatrix} b & b & 0 \\ b & b & b \\ 0 & b & b \end{vmatrix} = -b^3).$$

$$\begin{vmatrix} b & b & 0 & 0 & 0 & \dots & 0 & 0 & 0 & 0 & 0 \\ b & b & b & 0 & 0 & \dots & 0 & 0 & 0 & 0 & 0 \\ 0 & b & b & b & 0 & \dots & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & b & b & b & \dots & 0 & 0 & 0 & 0 & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots & \ddots & \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 0 & 0 & 0 & \dots & b & b & b & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & \dots & 0 & b & b & b & 0 \\ 0 & 0 & 0 & 0 & 0 & \dots & 0 & 0 & b & b & b \\ 0 & 0 & 0 & 0 & 0 & \dots & 0 & 0 & 0 & b & b \end{vmatrix}$$

$$= b \underbrace{\begin{vmatrix} b & b & 0 & 0 & \cdots & 0 & 0 & 0 & 0 \\ b & b & b & 0 & \cdots & 0 & 0 & 0 & 0 \\ 0 & b & b & b & \cdots & 0 & 0 & 0 & 0 \\ \vdots & \vdots & \vdots & \vdots & \ddots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 0 & 0 & \cdots & b & b & b & 0 \\ 0 & 0 & 0 & 0 & \cdots & 0 & b & b & b \\ 0 & 0 & 0 & 0 & \cdots & 0 & 0 & b & b \end{vmatrix}}_{\text{(This is } D_{n-1}\text{.)}} - b \begin{vmatrix} b & b & 0 & 0 & \cdots & 0 & 0 & 0 & 0 \\ 0 & b & b & 0 & \cdots & 0 & 0 & 0 & 0 \\ 0 & b & b & b & \cdots & 0 & 0 & 0 & 0 \\ \vdots & \vdots & \vdots & \vdots & \ddots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 0 & 0 & \cdots & b & b & b & 0 \\ 0 & 0 & 0 & 0 & \cdots & 0 & b & b & b \\ 0 & 0 & 0 & 0 & \cdots & 0 & 0 & b & b \end{vmatrix}$$



If we let $a_n = cr^n$ for $c, r \neq 0$ and $n \geq 1$, the characteristic equation produces the roots $b[(1/2) \pm i\sqrt{3}/2]$.

Hence

$$\begin{aligned}a_n &= c_1[b((1/2) + i\sqrt{3}/2)]^n + c_2[b((1/2) - i\sqrt{3}/2)]^n \\&= b^n[c_1(\cos(\pi/3) + i\sin(\pi/3))^n + c_2(\cos(\pi/3) - i\sin(\pi/3))^n] \\&= b^n[k_1 \cos(n\pi/3) + k_2 \sin(n\pi/3)].\end{aligned}$$

$b = a_1 = b[k_1 \cos(\pi/3) + k_2 \sin(\pi/3)]$, so $1 = k_1(1/2) + k_2(\sqrt{3}/2)$, or $k_1 + \sqrt{3}k_2 = 2$.

$0 = a_2 = b^2[k_1 \cos(2\pi/3) + k_2 \sin(2\pi/3)]$, so $0 = (k_1)(-1/2) + k_2(\sqrt{3}/2)$, or
 $k_1 = \sqrt{3}k_2$.

Hence $k_1 = 1$, $k_2 = 1/\sqrt{3}$ and the value of D_n is

$$b^n[\cos(n\pi/3) + (1/\sqrt{3}) \sin(n\pi/3)].$$



Case (C): Repeated Real Roots

- **Ex 10.23** : Solve the recurrence relation

$$a_{n+2} = 4a_{n+1} - 4a_n \text{ where } n \geq 0, a_0 = 1, a_1 = 3$$

- **Solution**

Let $a_n = cr^n$

$$a_n = c_1 2^n + c_2 2^n ?$$



$r^2 - 4r + 4 = 0 \Rightarrow r = 2 \Rightarrow 2^n$ and 2^n are not independent solutions, need another independent solution

So, try $g(n)2^n$, where $g(n)$ is not a constant

$$\Rightarrow g(n+2)2^{n+2} = 4g(n+1)2^{n+1} - 4g(n)2^n$$

$$\Rightarrow g(n+2) = 2g(n+1) - g(n) \Rightarrow g(n) = n, \therefore n2^n \text{ is a solution}$$

$$a_n = c_1(2^n) + c_2(n2^n) \text{ with } a_0 = 1, a_1 = 3$$

$$a_n = 1(2^n) + (1/2)(n2^n)$$



Repeated Real Roots

- In general, if

$$C_0 a_n + C_1 a_{n-1} + C_2 a_{n-2} + \dots + C_k a_{n-k} = 0$$

with r , a characteristic root of multiplicity m , then the part of the general solution that involves the root r has the form

$$\begin{aligned} & A_0 r^n + A_1 n r^n + A_2 n^2 r^n + A_3 n^3 r^n + \dots + A_{m-1} n^{m-1} r^n \\ &= (A_0 + A_1 n + A_2 n^2 + A_3 n^3 + \dots + A_{m-1} n^{m-1}) r^n \end{aligned}$$



Repeated Real Roots

- **Ex 10.24 :** Let p_n denote the probability that at least one case of measles is reported during the n th week after the first recorded case. School records provide evidence that $p_n = p_{n-1} - (0.25)p_{n-2}$, for $n \geq 2$. Since $p_0 = 0$ and $p_1 = 1$, if the first case is recorded on Monday, March 3, 2003, when did the probability for the occurrence of a new case decrease to less than 0.01 for the first time?

- **Solution**

$$\text{Let } p_n = cr^n$$

$$r^2 - r + (1/4) = 0 \Rightarrow r = 1/2$$

$$p_n = (c_1 + c_2 n)(1/2)^n \Rightarrow c_1 = 0 \text{ and } c_2 = 2 \Rightarrow p_n = n2^{-n+1}$$

$$p_n < 0.01 \Rightarrow \text{the first } n \text{ is 12, the week of May 19, 2003.}$$



10.3 The Nonhomogeneous Recurrence Relation

- $a_n + C_1 a_{n-1} = f(n), n \geq 1,$
- $a_n + C_1 a_{n-1} + C_2 a_{n-2} = f(n), n \geq 2$
- Let $a_n^{(h)}$ denote the general solution of the associated homogeneous relation.
- Let $a_n^{(p)}$ denote a solution of the given nonhomogeneous relation. (particular solution)
- Then $a_n = a_n^{(h)} + a_n^{(p)}$ is the general solution of the recurrence relation.



The Nonhomogeneous Recurrence Relation

$a_n - a_{n-1} = f(n)$, we have

$$a_1 = a_0 + f(1)$$

$$a_2 = a_1 + f(2) = a_0 + f(1) + f(2)$$

$$a_3 = a_2 + f(3) = a_0 + f(1) + f(2) + f(3)$$

⋮

$$a_n = a_{n-1} + f(n) = a_0 + f(1) + \cdots + f(n) = a_0 + \sum_{i=1}^n f(i).$$

• Ex 10.25

Solve the recurrence relation $a_n - a_{n-1} = 3n^2$, where $n \geq 1$ and $a_0 = 7$.

Here $f(n) = 3n^2$, so the unique solution is

$$a_n = a_0 + \sum_{i=1}^n f(i) = 7 + 3 \sum_{i=1}^n i^2 = 7 + \frac{1}{2}(n)(n+1)(2n+1).$$



The Nonhomogeneous Recurrence Relation

- **Ex 10.26** : Solve the recurrence relation

$$a_n - 3a_{n-1} = 5(7^n) \text{ for } n \geq 1 \text{ and } a_0 = 2.$$

- **Solution**

The solution for $a_n - 3a_{n-1} = 0$ is $a_n^{(h)} = c(3^n)$. First order linear recurrence

Since $f(n) = 5(7^n)$, let $a_n^{(p)} = A(7^n)$

$$\Rightarrow A(7^n) - 3A(7^{n-1}) = 5(7^n) \Rightarrow A = 35/4$$

$$a_n^{(p)} = (35/4)7^n = (5/4)7^{n+1}.$$

The general solution is $a_n = a_n^{(h)} + a_n^{(p)} = c(3^n) + (5/4)7^{n+1}$

$$\text{So, } a_n = (-1/4)(3^{n+3}) + (5/4)7^{n+1}$$



The Nonhomogeneous Recurrence Relation

- **Ex 10.27** : Solve the recurrence relation $a_n - 3a_{n-1} = 5(3^n)$ for $n \geq 1$ and $a_0 = 2$.

• Solution

Let $a_n^{(h)} = c(3^n)$.
多乘一個 n

Since $a_n^{(h)}$ and $f(n)$ are not linearly independent, let

$$a_n^{(p)} = Bn(3^n) \Rightarrow Bn(3^n) - 3B(n-1)(3^{n-1}) = 5(3^n). \Rightarrow B=5.$$

The general solution is $a_n = a_n^{(h)} + a_n^{(p)} = c(3^n) + (5)n3^{n+1}$

$$a_n = (2 + 5n)(3^n)$$



Solution for the Nonhomogeneous First-Order Relation

- $a_n + C_1 a_{n-1} = kr^n.$
 - If r^n is not a solution of the homogeneous relation $a_n + C_1 a_{n-1} = 0$, then $a_n^{(p)} = Ar^n.$
 - If r^n is a solution of the homogeneous relation, then $a_n^{(p)} = Bnr^n.$



Solution for the Nonhomogeneous Second-Order Relation

- $a_n + C_1 a_{n-1} + C_2 a_{n-2} = kr^n.$
 - If r^n is not a solution of the homogeneous relation, then $a_n^{(p)} = Ar^n.$
 - If $a_n^{(h)} = c_1 r^n + c_2 r_1^n$, where $r \neq r_1$, then $a_n^{(p)} = Bnr^n.$
 - If $a_n^{(h)} = (c_1 + c_2 n)r^n$, then $a_n^{(p)} = Cn^2r^n.$



The Nonhomogeneous Recurrence Relation

- **Ex 10.28** : The Towers of Hanoi.

- Let count the minimum number of moves it takes to transfer n disks from peg 1 to peg 3.
- $a_{n+1} = 2a_n + 1$
 - Transfer the top n disks from peg 1 to peg 2, need a_n moves.
 - Transfer the largest disk from peg 1 to peg 3, need 1 moves.
 - Transfer the n disks on peg 2 onto the largest disk, need a_n moves.

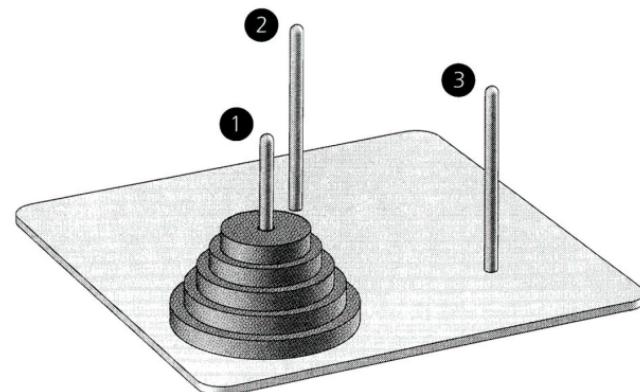


Figure 10.11

For $a_{n+1} - 2a_n = 1$, we know that $a_n^{(h)} = c(2^n)$. Since $f(n) = 1 = (1)^n$ is not a solution of $a_{n+1} - 2a_n = 0$, we set $a_n^{(p)} = A(1)^n = A$ and find from the given relation that $A = 2A + 1$, so $A = -1$ and $a_n = c(2^n) - 1$. From $a_0 = 0 = c - 1$ it then follows that $c = 1$, so $a_n = 2^n - 1$, $n \geq 0$.



The Nonhomogeneous Recurrence Relation

- **Ex 10.29 :** Let a_n denote the amount still owed on the loan at the end of the n th period.
(r is the interest rate, P is payment, S is loan)

$$a_{n+1} = a_n + ra_n - P, \quad 0 \leq n \leq T-1, \quad a_0 = S, \quad a_T = 0$$

$$a_n^{(h)} = c(1+r)^n.$$

$$\text{Let } a_n^{(p)} = A, \quad A - (1+r)A = -P \Rightarrow A = P/r, \quad a_n^{(p)} = P/r.$$

$$a_n = a_n^{(h)} + a_n^{(p)} = c(1+r)^n + P/r \Rightarrow c = S - (P/r)$$

$$a_n = (S - (P/r))(1+r)^n + (P/r)$$

$$\text{Since } 0 = a_T, \text{ we have } P = (Sr)[1 - (1 + r)^{-T}]^{-1}$$



The Nonhomogeneous Recurrence Relation

- **Ex 10.30 :** Let S be a set containing 2^n real numbers. Find the maximum and minimum in S . We wish to determine the number of comparisons made between pairs of elements in S .
 - Let a_n denote the number of needed comparisons.

$$n = 2, |S| = 2^2 = 4, S = \{x_1, x_2, y_1, y_2\} = S_1 \cup S_2,$$

$$S_1 = \{x_1, x_2\}, S_2 = \{y_1, y_2\}$$

$$a_{n+1} = 2a_n + 2, n \geq 1.$$

$$a_n^{(h)} = c(2^n), a_n^{(p)} = A$$

$$a_1 = 1 \rightarrow a_n = (3/2)(2^n) - 2$$



The Nonhomogeneous Recurrence Relation

- **Ex 10.31** : For the alphabet = {0,1,2,3}, how many strings of length n contains an even number of 1's.
 - Let a_n count those strings among the 4^n strings.

Consider the n th symbol of a string of length n

1. The n th symbol is 0, 2, 3 $\Rightarrow 3a_{n-1}$ 前面 $n-1$ 個 symbol 為偶數個 1，最後一個可為 0、2、3
2. The n th symbol is 1 \Rightarrow there must be an odd number of 1's among the first $n-1$ symbols $\Rightarrow 4^{n-1} - a_{n-1}$

$$a_n = 3a_{n-1} + (4^{n-1} - a_{n-1}) = 2a_{n-1} + 4^{n-1}$$

$$a_n^{(h)} = c(2^n), a_n^{(p)} = A(4^{n-1})$$

前面 $n-1$ 個 symbol 為奇數個 1，最後一個只能為 1 (全部 - 偶數個 1)

$$a_1 = 3 \rightarrow a_n = 2^{n-1} + 2(4^{n-1})$$



The Nonhomogeneous Recurrence Relation

$$\begin{aligned}f(x) &= \left(1 + x + \frac{x^2}{2!} + \dots\right) \left(1 + \frac{x^2}{2!} + \frac{x^4}{4!} + \dots\right) \left(1 + x + \frac{x^2}{2!} + \dots\right) \left(1 + x + \frac{x^2}{2!} + \dots\right) \\&= e^x \cdot \left(\frac{e^x + e^{-x}}{2}\right) \cdot e^x \cdot e^x \\&= \left(\frac{1}{2}\right) e^{4x} + \left(\frac{1}{2}\right) e^{2x} \\&= \left(\frac{1}{2}\right) \sum_{n=0}^{\infty} \frac{(4x)^n}{n!} + \left(\frac{1}{2}\right) \sum_{n=0}^{\infty} \frac{(2x)^n}{n!}.\end{aligned}$$

Here $a_n =$ the coefficient of $\frac{x^n}{n!}$ in $f(x) = \left(\frac{1}{2}\right) 4^n + \left(\frac{1}{2}\right) 2^n = 2^{n-1} + 2(4^{n-1})$, as above.



The Nonhomogeneous Recurrence Relation

- **Ex 10.32** : Snowflake curve shown in Figure 10.12.
- Let a_n denote the area of the polygon P_n obtained from the original equilateral triangle after we apply n transformations.

$$a_0 = \sqrt{3}/4$$

$$a_1 = (\sqrt{3}/4) + (3)(\sqrt{3}/4)(1/3)^2 = \sqrt{3}/3$$

$$a_2 = a_1 + (4)(3)(\sqrt{3}/4)[(1/3)^2]^2 = 10\sqrt{3}/27$$

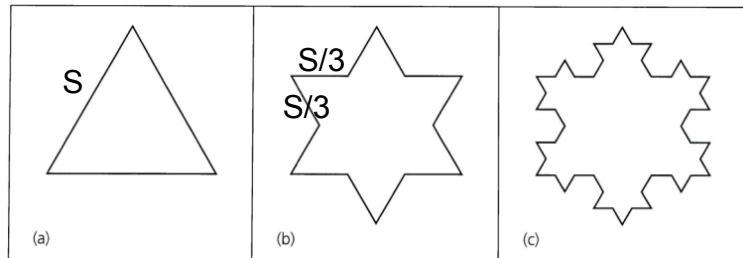


Figure 10.12

A special kind of fractal curves 1904, Helge von Koch

http://en.wikipedia.org/wiki/Koch_snowflake

#segment in each side



$$a_{n+1} = a_n + (4^n(3))(\sqrt{3}/4)(1/3^{n+1})^2 = a_n + (1/(4\sqrt{3}))(4/9)^n$$

$$a_n = a_n^{(h)} + a_n^{(p)} = A(1)^n + B(4/9)^n = (6/(5\sqrt{3})) - (1/(5\sqrt{3}))(4/9)^{n-1} \approx 6/(5\sqrt{3})$$



The Nonhomogeneous Recurrence Relation

- **Ex 10.34 :** Solve the recurrence relation $a_{n+2} - 4a_n + 3a_n = -200$ for $n \geq 0$ and $a_0 = 3000$ and $a_1 = 3300$.

- **Solution**

$$a_n^{(h)} = c_1(3^n) + c_2(1^n).$$

$$\text{Let } a_n^{(p)} = An \Rightarrow A(n+2) - 4A(n+1) + 3An = -200$$

$$\Rightarrow a_n^{(p)} = 100n.$$

$$a_n = a_n^{(h)} + a_n^{(p)} = c_1(3^n) + c_2(1^n) + 100n$$

$$\Rightarrow a_n = 100(3^n) + 2900 + 100n$$



The Nonhomogeneous Recurrence Relation

- Two procedures of computing the n th Fibonacci number in Figure 10.15. Which one is more efficient?
- $a_n = a_{n-1} + a_{n-2} + 1$

$$\begin{aligned}a_n &= \left(\frac{1+\sqrt{5}}{2\sqrt{5}}\right) \left(\frac{1+\sqrt{5}}{2}\right)^n - \left(\frac{1-\sqrt{5}}{2\sqrt{5}}\right) \left(\frac{1-\sqrt{5}}{2}\right)^n - 1 \\&= \frac{1}{\sqrt{5}} \left(\frac{1+\sqrt{5}}{2}\right)^{n+1} - \frac{1}{\sqrt{5}} \left(\frac{1-\sqrt{5}}{2}\right)^{n+1} - 1.\end{aligned}$$

```
procedure FibNum2(n: nonnegative integer)
begin
  if n = 0 then
    fib := 0
  else if n = 1 then
    fib := 1
  else
    fib := FibNum2(n - 1) + FibNum2(n - 2)
  end
(b)
```

recursive

$$F_n = \frac{1}{\sqrt{5}} \left[\left(\frac{1+\sqrt{5}}{2} \right)^n - \left(\frac{1-\sqrt{5}}{2} \right)^n \right], \quad n \geq 0.$$

```
procedure FibNum1(n: nonnegative integer)
begin
  if n = 0 then
    fib := 0
  else if n = 1 then
    fib := 1
  else
    begin
      last := 1
      next_to_last := 0
      for i := 2 to n do
        begin
          temp := last
          last := last + next_to_last
          next_to_last := temp
        end
        fib := last
      end
    end
(a)
```

iterative

Figure 10.15



Particular Solutions to Nonhomogeneous Recurrence Relation

- $C_0a_n + C_1a_{n-1} + C_2a_{n-2} + \dots + C_ka_{n-k} = f(n)$
- (1) If $f(n)$ is a constant multiple of one of the forms in the first column of Table 10.2 $\Rightarrow a_n^{(p)}$ in the second column .
 - (2) When $f(n)$ comprises a sum of constant multiples of terms.
 - E.g., $f(n) = n^2 + 3\sin 2n \Rightarrow a_n^{(p)} = (A_2n^2 + A_1n + A_0) + (A \sin 2n + B \cos 2n)$
 - (3) If a summand $f_1(n)$ of $f(n)$ is a solution of the associated homogeneous relation.
 - If $f_1(n)$ causes this problem, we multiply the trial solution $(a_n^{(p)})_1$ corresponding to $f_1(n)$ by the smallest power of n , say n^s , for which no summand of $n^s f_1(n)$ is a solution of the associated homogeneous relation. Thus, $n^s(a_n^{(p)})_1$ is the corresponding part of $a_n^{(p)}$.



Table 10.2

$f(n)$	$a_n^{(p)}$
c , a constant	A , a constant
n	$A_1n + A_0$
n^2	$A_2n^2 + A_1n + A_0$
n^t , $t \in \mathbf{Z}^+$	$A_t n^t + A_{t-1} n^{t-1} + \cdots + A_1 n + A_0$
r^n , $r \in \mathbf{R}$	Ar^n
$\sin \theta n$	$A \sin \theta n + B \cos \theta n$
$\cos \theta n$	$A \sin \theta n + B \cos \theta n$
$n^t r^n$	$r^n (A_t n^t + A_{t-1} n^{t-1} + \cdots + A_1 n + A_0)$
$r^n \sin \theta n$	$Ar^n \sin \theta n + Br^n \cos \theta n$
$r^n \cos \theta n$	$Ar^n \sin \theta n + Br^n \cos \theta n$



Particular Solutions to Nonhomogeneous Recurrence Relation

- **Ex 10.36 :** For n people at a party, each of them shakes hands with others. $C(n, 2)$

- a_n counts the total number of handshakes:

$$a_{n+1} = a_n + n, n \geq 2, a_2 = 1$$

- $a_n^{(h)} = c(1^n) = c.$

- Let $a_n^{(p)} = A_1 n + A_0$

- By the third remark stated above, multiplying $a_n^{(p)}$ by n^1 , then
 $a_n^{(p)} = A_1 n^2 + A_0 n$

- $A_1 = \frac{1}{2}, A_0 = -\frac{1}{2} \Rightarrow a_n^{(p)} = (\frac{1}{2})n^2 + (-\frac{1}{2})n.$

- $a_n = a_n^{(h)} + a_n^{(p)} = c + (\frac{1}{2})n^2 + (-\frac{1}{2})n \Rightarrow c = 0$

- $a_n = (\frac{1}{2})n(n-1)$



Particular Solutions to Nonhomogeneous Recurrence Relation

- **Ex 10.37** : $a_{n+2} - 10a_{n+1} + 21a_n = f(n)$, $n \geq 0$
- $a_n^{(h)} = c_1(3^n) + c_2(7^n)$.

Table 10.3

$f(n)$	$a_n^{(p)}$
5	A_0
$3n^2 - 2$	$A_3n^2 + A_2n + A_1$
$7(11^n)$	$A_4(11^n)$
$31(r^n)$, $r \neq 3, 7$	$A_5(r^n)$
$6(3^n)$	$\underline{A_6n3^n}$
$2(3^n) - 8(9^n)$	$\underline{A_7n3^n} + A_8(9^n)$
$4(3^n) + 3(7^n)$	$\underline{A_9n3^n} + \underline{A_{10}n7^n}$



10.4 The Method of Generating Functions

- **Ex 10.38 :** Solve the relation $a_n - 3a_{n-1} = n$, $n \geq 1$, $a_0 = 1$.

- Let $f(x) = \sum_{n=0}^{\infty} a_n x^n$ be the generating function for a_0, a_1, \dots, a_n .

$$\sum_{n=1}^{\infty} a_n x^n - 3 \sum_{n=1}^{\infty} a_{n-1} x^n = \sum_{n=1}^{\infty} n x^n.$$

$$(f(x) - a_0) - 3x \sum_{n=1}^{\infty} a_{n-1} x^{n-1} = \sum_{n=1}^{\infty} n x^n \left(= \sum_{n=0}^{\infty} n x^n \right).$$

$$\frac{x}{(1-x)^2} = x + 2x^2 + 3x^3 + \dots,$$

$$(f(x) - 1) - 3x f(x) = \frac{x}{(1-x)^2}, \quad \text{and} \quad f(x) = \frac{1}{(1-3x)} + \frac{x}{(1-x)^2(1-3x)}.$$



$$\begin{aligned}f(x) &= \frac{1}{1-3x} + \frac{(-1/4)}{(1-x)} + \frac{(-1/2)}{(1-x)^2} + \frac{(3/4)}{(1-3x)} \\&= \frac{(7/4)}{(1-3x)} + \frac{(-1/4)}{(1-x)} + \frac{(-1/2)}{(1-x)^2}.\end{aligned}$$

We find a_n by determining the coefficient of x^n in each of the three summands.

- a) $(7/4)/(1-3x) = (7/4)[1/(1-3x)]$
 $= (7/4)[1 + (3x) + (3x)^2 + (3x)^3 + \dots]$, and the coefficient of
 x^n is $(7/4)3^n$.
- b) $(-1/4)/(1-x) = (-1/4)[1 + x + x^2 + \dots]$, and the coefficient of x^n here is
 $(-1/4)$.
- c) $(-1/2)/(1-x)^2 = (-1/2)(1-x)^{-2}$
 $= (-1/2) \left[\binom{-2}{0} + \binom{-2}{1}(-x) + \binom{-2}{2}(-x)^2 + \binom{-2}{3}(-x)^3 + \dots \right]$
and the coefficient of x^n is given by $(-1/2)\binom{-2}{n}(-1)^n = (-1/2)(-1)^n\binom{2+n-1}{n} \cdot (-1)^n = (-1/2)(n+1)$.

Therefore $a_n = \underline{\hspace{2cm}}(7/4)3^n - (1/2)n - (3/4)\underline{\hspace{2cm}}, n \geq 0$.

$a_n^{(h)}$

$a_n^{(p)}$



The Method of Generating Functions

- **Ex 10.39** : Solve the relation

$$a_{n+2} - 5a_{n+1} + 6a_n = 2, \quad n \geq 0, \quad a_0 = 3, \quad a_1 = 7.$$

- Let $f(x) = \sum a_n x^n$ be the generating function for a_0, a_1, \dots, a_n

- 1) We first multiply this given relation by x^{n+2} because $n + 2$ is the largest subscript that appears. This gives us

$$a_{n+2}x^{n+2} - 5a_{n+1}x^{n+2} + 6a_nx^{n+2} = 2x^{n+2}.$$

- 2) Then we sum all of the equations represented by the result in step (1) and obtain

$$\sum_{n=0}^{\infty} a_{n+2}x^{n+2} - 5 \sum_{n=0}^{\infty} a_{n+1}x^{n+2} + 6 \sum_{n=0}^{\infty} a_nx^{n+2} = 2 \sum_{n=0}^{\infty} x^{n+2}.$$

- 4) Let $f(x) = \sum_{n=0}^{\infty} a_n x^n$ be the generating function for the solution. The equation in step (3) now takes the form

$$(f(x) - a_0 - a_1x) - 5x(f(x) - a_0) + 6x^2 f(x) = \frac{2x^2}{1-x},$$



5) Solving for $f(x)$ we have

$$(1 - 5x + 6x^2)f(x) = 3 - 8x + \frac{2x^2}{1-x} = \frac{3 - 11x + 10x^2}{1-x},$$

from which it follows that

$$f(x) = \frac{3 - 11x + 10x^2}{(1 - 5x + 6x^2)(1 - x)} = \frac{(3 - 5x)(1 - 2x)}{(1 - 3x)(1 - 2x)(1 - x)} = \frac{3 - 5x}{(1 - 3x)(1 - x)}.$$

A partial-fraction decomposition (by hand, or via a computer algebra system) gives us

$$f(x) = \frac{2}{1 - 3x} + \frac{1}{1 - x} = 2 \sum_{n=0}^{\infty} (3x)^n + \sum_{n=0}^{\infty} x^n.$$

Consequently, $a_n = 2(3^n) + 1$, $n \geq 0$.



The Method of Generating Functions

- **Ex 10.40 :** Let $a(n, r)$ = the number of ways we can select, with repetitions allowed, r objects from a set of n distinct objects.
- Let $\{b_1, b_2, \dots, b_n\}$ be the set, consider b_1
 - b_1 is never selected: the r objects from $\{b_2, \dots, b_n\} \Rightarrow a(n - 1, r)$
 - b_1 is selected at least once: must select $r - 1$ objects from $\{b_1, b_2, \dots, b_n\} \Rightarrow a(n, r - 1)$
- Then $\textcolor{red}{a(n, r) = a(n-1, r) + a(n, r-1)}$.
- Let $f_n = \sum_{r=0}^{\infty} a(n, r)x^r$ be the generating function for $a(n, 0), a(n, 1), a(n, 2), \dots,$



$$a(n, r)x^r = a(n - 1, r)x^r + a(n, r - 1)x^r \quad \text{and}$$

$$\sum_{r=1}^{\infty} a(n, r)x^r = \sum_{r=1}^{\infty} a(n - 1, r)x^r + \sum_{r=1}^{\infty} a(n, r - 1)x^r.$$

Realizing that $a(n, 0) = 1$ for $n \geq 0$ and $a(0, r) = 0$ for $r > 0$, we write

$$f_n - a(n, 0) = f_{n-1} - a(n - 1, 0) + x \sum_{r=1}^{\infty} a(n, r - 1)x^{r-1}.$$

so $f_n - 1 = f_{n-1} - 1 + xf_n$. Therefore, $f_n - xf_n = f_{n-1}$, or $f_n = f_{n-1}/(1 - x)$.

If $n = 5$, for example, then

$$\begin{aligned} f_5 &= \frac{f_4}{(1 - x)} = \frac{1}{(1 - x)} \cdot \frac{f_3}{(1 - x)} = \frac{f_3}{(1 - x)^2} = \frac{f_2}{(1 - x)^3} = \frac{f_1}{(1 - x)^4} \\ &= \frac{f_0}{(1 - x)^5} = \frac{1}{(1 - x)^5}, \end{aligned}$$

since $f_0 = a(0, 0) + a(0, 1)x + a(0, 2)x^2 + \dots = 1 + 0 + 0 + \dots$.

In general, $f_n = 1/(1 - x)^n = (1 - x)^{-n}$, so $a(n, r)$ is the coefficient of x^r in $(1 - x)^{-n}$, which is $\binom{-n}{r}(-1)^r = \binom{n+r-1}{r}$.



10.5 A Special Kind of Nonlinear Recurrence Relation

- **Ex 10.42 :** Let b_n denote the number of rooted ordered binary trees on n vertices.
- $b_3 = 5$ is shown in Figure 10.18.
- $b_{n+1} = b_0 b_n + b_1 b_{n-1} + \dots + b_{n-1} b_1 + b_n b_0$
- Let $f(x) = \sum_{n=0}^{\infty} b_n x^n$ be the generating function for b_0, b_1, \dots, b_n .

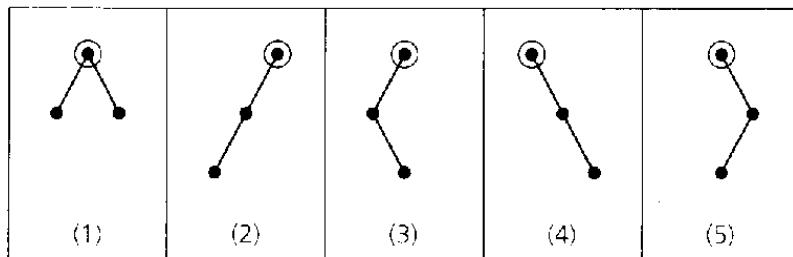


Figure 10.18



A Special Kind of Nonlinear Recurrence Relation

- $b_{n+1} = b_0 b_n + b_1 b_{n-1} + \dots + b_{n-1} b_1 + b_n b_0$
 1. 0 vertices on the left, n vertices on the right
 $\Rightarrow b_0 b_n$
 2. 1 vertices on the left, $n - 1$ vertices on the right $\Rightarrow b_1 b_{n-1}$
 3. i vertices on the left, $n - i$ vertices on the right $\Rightarrow b_i b_{n-i}$
 4. n vertices on the left, *none* on the right $\Rightarrow b_n b_0$

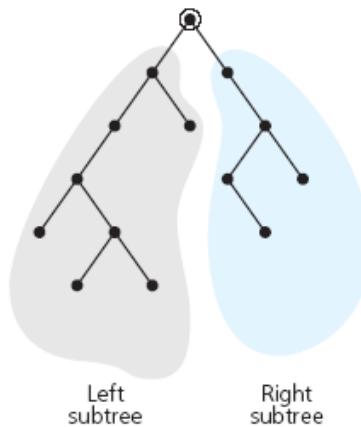


Figure 10.19



$$b_{n+1} = b_0 b_n + b_1 b_{n-1} + b_2 b_{n-2} + \cdots + b_{n-1} b_1 + b_n b_0,$$

$$\sum_{n=0}^{\infty} b_{n+1} x^{n+1} = \sum_{n=0}^{\infty} (b_0 b_n + b_1 b_{n-1} + \cdots + b_{n-1} b_1 + b_n b_0) x^{n+1}.$$

$$(f(x) - b_0) = x \sum_{n=0}^{\infty} (b_0 b_n + b_1 b_{n-1} + \cdots + b_n b_0) x^n = x [f(x)]^2.$$

$$x [f(x)]^2 - f(x) + 1 = 0, \quad \text{so } \underline{f(x) = [1 \pm \sqrt{1 - 4x}] / (2x)}.$$



But $\sqrt{1 - 4x} = (1 - 4x)^{1/2} = \binom{1/2}{0} + \binom{1/2}{1}(-4x) + \binom{1/2}{2}(-4x)^2 + \dots$, where the coefficient of x^n , $n \geq 1$, is

$$\begin{aligned}\binom{1/2}{n}(-4)^n &= \frac{(1/2)((1/2)-1)((1/2)-2)\cdots((1/2)-n+1)}{n!}(-4)^n \\ &= (-1)^{n-1} \frac{(1/2)(1/2)(3/2)\cdots((2n-3)/2)}{n!}(-4)^n \\ &= \frac{(-1)2^n(1)(3)\cdots(2n-3)}{n!} \\ &= \frac{(-1)2^n(n!)(1)(3)\cdots(2n-3)(2n-1)}{(n!)(n!)(2n-1)} \\ &= \frac{(-1)(2)(4)\cdots(2n)(1)(3)\cdots(2n-1)}{(2n-1)(n!)(n!)} = \frac{(-1)}{(2n-1)} \binom{2n}{n}.\end{aligned}$$



$$f(x) = \frac{1}{2x} \left[\text{X} - \left[1 - \sum_{n=1}^{\infty} \frac{1}{(2n-1)} \binom{2n}{n} x^n \right] \right],$$

b_n , the coefficient of x^n in $f(x)$, is half the coefficient of x^{n+1} in

$$\sum_{n=1}^{\infty} \frac{1}{(2n-1)} \binom{2n}{n} x^n.$$

$$b_n = \frac{1}{2} \left[\frac{1}{2(n+1)-1} \right] \binom{2(n+1)}{n+1} = \frac{(2n)!}{(n+1)!(n!)} = \frac{1}{(n+1)} \binom{2n}{n}$$

- b_n are called *Catalan numbers*



A Special Kind of Nonlinear Recurrence Relation

$n=2, 12, 21$

$n=3, \textcolor{red}{123, 132, 213, 321, 231, 312}$

- **Ex 10.43** : permute $1, 2, 3, \dots, n$, which must be pushed onto the top of the stack in the order given.

- $n = 0 \Rightarrow 1$
- $n = 1 \Rightarrow 1$
- $n = 2 \Rightarrow 2$
- $n = 3 \Rightarrow 5$
- $n = 4 \Rightarrow 14$

Table 10.4

$a_4 = a_0a_3 + a_1a_2 + a_2a_1 + a_3a_0$			
1, 2, 3, 4	2, 1, 3, 4	2, 3, 1, 4	2, 3, 4, 1
1, 2, 4, 3	2, 1, 4, 3	3, 2, 1, 4	2, 4, 3, 1
1, 3, 2, 4			3, 2, 4, 1
1, 3, 4, 2			3, 4, 2, 1
1, 4, 3, 2			4, 3, 2, 1

Output

1, 2, 3, ..., n Input

Stack

- 1) There are five permutations with 1 in the first position, because after 1 is pushed onto and popped from the stack, there are five ways to permute 2, 3, 4 using the stack.
- 2) When 1 is in the second position, 2 must be in the first position. This is because we pushed 1 onto the (empty) stack, then pushed 2 on top of it and then popped 2 and then 1. There are two permutations in column 2, because 3, 4 can be permuted in two ways on the stack.

A Special Kind of Nonlinear Recurrence Relation



- 3) For column 3 we have 1 in position three. We note that the only numbers that can precede it are 2 and 3, which can be permuted on the stack (with 1 on the bottom) in two ways. Then 1 is popped, and we push 4 onto the (empty) stack and then pop it.
- 4) In the last column we obtain five permutations: After we push 1 onto the top of the (empty) stack, there are five ways to permute 2, 3, 4 using the stack (with 1 on the bottom). Then 1 is popped from the stack to complete the permutation.

- $a_4 = a_0a_3 + a_1a_2 + a_2a_1 + a_3a_0$
- $a_{n+1} = a_0a_n + a_1a_{n-1} + \dots + a_{n-1}a_1 + a_na_0$
- $a_n = \frac{1}{n+1} \binom{2n}{n}$
- Push, pop permutation with limitation (Ex 1.43)



10.6 Divide-and-Conquer Algorithms

- In general, solve a given problem of size n by
 - Solving the problem for a small value of n directly.
 - Breaking the problem into a smaller problems of the same type and the same size $\lceil n/b \rceil$ or $\lfloor n/b \rfloor$
- Divide-and-conquer algorithms
 - 1) The time to solve the initial problem of size $n = 1$ is a constant $c \geq 0$, and
 - 2) The time to break the given problem of size n into a smaller (similar) problems, together with the time to combine the solutions of these smaller problems to get a solution for the given problem, is $h(n)$, a function of n .
- Time complexity function $f(n)$
 - $f(1) = c$
 - $f(n) = af(n/b) + h(n)$ for $n = b^k$



Divide-and-Conquer Algorithms

- Theorem 10.1: Let $a, b, c \in \mathbb{Z}^+$ with $b \geq 2$, and let $f: \mathbb{Z}^+ \rightarrow \mathbb{R}$. If

$$f(1) = c, \quad \text{and}$$

$$f(n) = af(n/b) + c, \quad \text{for } n = b^k, \quad k \geq 1,$$

then for all $n = 1, b, b^2, b^3, \dots$,

1) $f(n) = c(\log_b n + 1)$, when $a = 1$, and

2) $f(n) = \frac{c(an^{\log_b a} - 1)}{a - 1}$, when $a \geq 2$.

Proof: For $k \geq 1$ and $n = b^k$, we write the following system of k equations. [Starting with the second equation, we obtain each of these equations from its immediate predecessor by (i) replacing each occurrence of n in the prior equation by n/b and (ii) multiplying the resulting equation in (i) by a .]

$$\begin{aligned} f(n) &= af(n/b) + c \\ af(n/b) &= a^2 f(n/b^2) + ac \\ a^2 f(n/b^2) &= a^3 f(n/b^3) + a^2 c \\ &\vdots \quad \vdots \quad \vdots \\ a^{k-2} f(n/b^{k-2}) &= a^{k-1} f(n/b^{k-1}) + a^{k-2} c \\ a^{k-1} f(n/b^{k-1}) &= a^k f(n/b^k) + a^{k-1} c \end{aligned}$$



$$f(n) = a^k f(n/b^k) + [c + ac + a^2 c + \cdots + a^{k-1} c].$$

Since $n = b^k$ and $f(1) = c$, we have

$$\begin{aligned} f(n) &= a^k f(1) + c[1 + a + a^2 + \cdots + a^{k-1}] \\ &= c[1 + a + a^2 + \cdots + a^{k-1} + a^k]. \end{aligned}$$

1) If $a = 1$, then $f(n) = c(k + 1)$. But $n = b^k \Leftrightarrow \log_b n = k$, so $f(n) = c(\log_b n + 1)$, for $n \in \{b^i | i \in \mathbf{N}\}$.

2) When $a \geq 2$, then $f(n) = \frac{c(1 - a^{k+1})}{1 - a} = \frac{c(a^{k+1} - 1)}{a - 1}$, from identity 4 of Table 9.2.

Now $n = b^k \Leftrightarrow \log_b n = k$, so

$$a^k = a^{\log_b n} = (b^{\log_b n})^{\log_b n} = (b^{\log_b n})^{\log_b a} = n^{\log_b a},$$

and

$$f(n) = \frac{c(an^{\log_b a} - 1)}{(a - 1)}, \quad \text{for } n \in \{b^i | i \in \mathbf{N}\}.$$



Divide-and-Conquer Algorithms

- **Ex 10.45 :**

$$f(n) = \frac{c(an^{\log_b a} - 1)}{(a - 1)}$$

(a) $f(1) = 3$ and $f(n) = f(n/2) + 3$ for $n = 2^k$

$$c = 3, b = 2, a = 1$$

$$f(n) = 3(\log_2 n + 1)$$

(b) $g(1) = 7$ and $g(n) = 4g(n/3) + 7$ for $n = 3^k$

$$c = 7, b = 3, a = 4$$

$$g(n) = (7/3)(4n^{\log_3 4} - 1)$$

(c) $h(1) = 5$ and $h(n) = 7h(n/7) + 5$ for $n = 7^k$

$$c = 5, b = 7, a = 7$$

$$h(n) = (5/6)(7n - 1)$$



Changing variables

$$T(n) = 2T(\sqrt{n}) + \lg n$$

Let $m = \lg n$.

$$T(2^m) = 2T(2^{m/2}) + m$$

$$T(n) = T(2^{\sqrt{\lg n}}) + c ?$$

$$Ans : T(n) = O(\lg \lg \lg n)$$

Suppose $S(m) = T(2^m)$,

Then $S(m) = 2S(m/2) + m$.

$$\Rightarrow S(m) = O(m \lg m)$$

$$\begin{aligned}\Rightarrow T(n) &= T(2^m) = S(m) = O(m \lg m) \\ &= O(\lg n \lg \lg n)\end{aligned}$$