

# GraphNet: A Large-Scale Computational Graph Dataset for Tensor Compiler Research

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## Abstract

We introduce GraphNet, a dataset of 2.7K real-world deep learning computational graphs with rich meta-data, spanning six major task categories across multiple frameworks. To evaluate compiler performance on these samples, we define the benchmark metric Speedup Score  $S_t$  that combines runtime speedup and execution correctness under tunable tolerance levels.  $S_t$  reflects the general optimization capability of existing tensor compilers across task categories. Furthermore, we extend  $S_t$  to the Error-aware Speedup Score  $ES_t$ , which incorporates error information and helps compiler developers identify key performance bottlenecks. In this paper, we benchmark the default tensor compilers of PaddlePaddle (CINN) and PyTorch (TorchInductor) on computer vision (CV) and natural language processing (NLP) samples to demonstrate the practicality of GraphNet. The full construction pipeline with graph extraction and compiler evaluation tools is available at <https://github.com/PaddlePaddle/GraphNet>.

## 1 Introduction

The development of high-performance GPU kernels has become critical for computational efficiency in modern deep learning workloads. A widely adopted approach integrates deep learning frameworks, such as PaddlePaddle and PyTorch, with vendor-specific operator libraries (e.g., cuDNN [2], oneDNN [9]). However, emerging hardware with low-precision formats (e.g., BF16, FP8) and advanced memory access patterns create a growing demand for custom operators, which pre-defined vendor libraries often fail to keep up with. To meet these demands, modern deep learning systems increasingly rely on tensor compilers (e.g., TVM [1], XLA [6], BladeDISC [15]) to transform high-level computational graphs into efficient programs for heterogeneous hardware platforms.

Despite these advances, deep learning engineers still rely heavily on manual tuning, as existing tensor compilation frameworks often lack fine-grained hardware control and offer limited support for cross-platform complex algorithms. Meanwhile, researchers are increasingly exploring the use of large language models (LLMs) [10] and AI coding agents [13] to automatically generate efficient operators and kernels. In this challenging context, systematic evaluation of existing tensor compilers is essential to identify performance bottlenecks and guide the evolution of next-generation compilers.

Existing benchmarks are often ad hoc, relying on a small set of hand-picked samples, with limited coverage of up-to-date, real-world models from the community and lacking sufficient support for cross-framework evaluation. To address these limitations, we introduce **GraphNet**, a large-scale dataset of deep learning computational graphs designed to enable systematic evaluation of tensor compilers across tasks and frameworks. We further propose a unified metric to benchmark compiler performance by jointly accounting for runtime speedup, numerical correctness, and compilation failures. Our main contributions are as follows:

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- We collect more than 2.7k computational graphs from real-world models, covering diverse task categories and mainstream frameworks (e.g., PaddlePaddle, PyTorch). All samples are stored in a unified format and are compatible with multiple tensor compiler backends, such as CINN, XLA, and TVM.
- We evaluate the performance of CINN and TorchInductor on CV and NLP samples from GraphNet, and define the **Speedup Score** ( $S_t$ ) to measure the general optimization capability of tensor compilers, as well as the **Error-aware Speedup Score** ( $ES_t$ ), which encodes all types of execution failures.
- We present a detailed study of the dataset construction pipeline and sample constraints, and release the GraphNet dataset together with open-source extraction and evaluation tools.

The rest of this paper is organized as follows. Section 2 introduces dataset properties and its distribution. Section 3 presents the design of our evaluation metric and discusses experimental results. Section 4 describes the dataset construction methodology. Sections 5 and 6 review related works and outline future directions.

## 2 Dataset Properties

**Authenticity:** All samples are extracted from deep learning models in mainstream ecosystems, with a focus on PaddlePaddle libraries (e.g., PaddleNLP, PaddleX, PaddleScience) and PyTorch libraries (e.g., TorchVision, timm, mmseg). These libraries are actively maintained and widely adopted by the community, ensuring that GraphNet reflects real-world workloads rather than synthetic benchmarks.

**Compatibility:** GraphNet samples adopt a standardized format that integrates computational graphs with inputs, weights, and custom operators. They preserve complete computation semantics without information loss, ensuring compatibility with diverse compiler backends, including CINN, BladeDISC, TVM, TorchInductor, XLA, and TensorRT. This enables fair and reproducible evaluation, as well as seamless extensions to new compilers.

**Diversity:** The dataset spans six major task categories, namely computer vision (CV), natural language processing (NLP), audio, multimodal learning, scientific computing, and others. Models vary in scale from a few thousand to 10B parameters, with the most complex sample containing up to 3,686 operators. GraphNet also supports multiple data types, such as BF16, FP16, and FP32, which are essential for evaluating compiler optimizations across different numerical precisions and hardware backends.

## 3 Benchmark of Tensor Compilers

### 3.1 Compiler Evaluation

We introduce an automated compiler evaluation workflow that establishes a comparable performance evaluation platform through unified and repeatable testing of multiple tensor compilers. The evaluation takes GraphNet samples as input and consists of the following key steps:

1. **Baseline Execution:** Execute the original model in Eager mode on a given framework and record its output and execution time  $T_{eager}$  as the baseline for subsequent comparisons.
2. **Compiler Configuration:** Compile the original model into an optimized executable version by specifying the target compiler through the interface `torch.compile(..., backend="tvm")`. Currently supported compilers include CINN, TorchInductor, TensorRT, TVM, BladeDISC, and XLA.
3. **Compiled Execution:** Execute the compiled model after a warmup phase to eliminate cold-start overhead and obtain the pure execution time  $T_{compiled}$ . Record its output to verify accuracy.

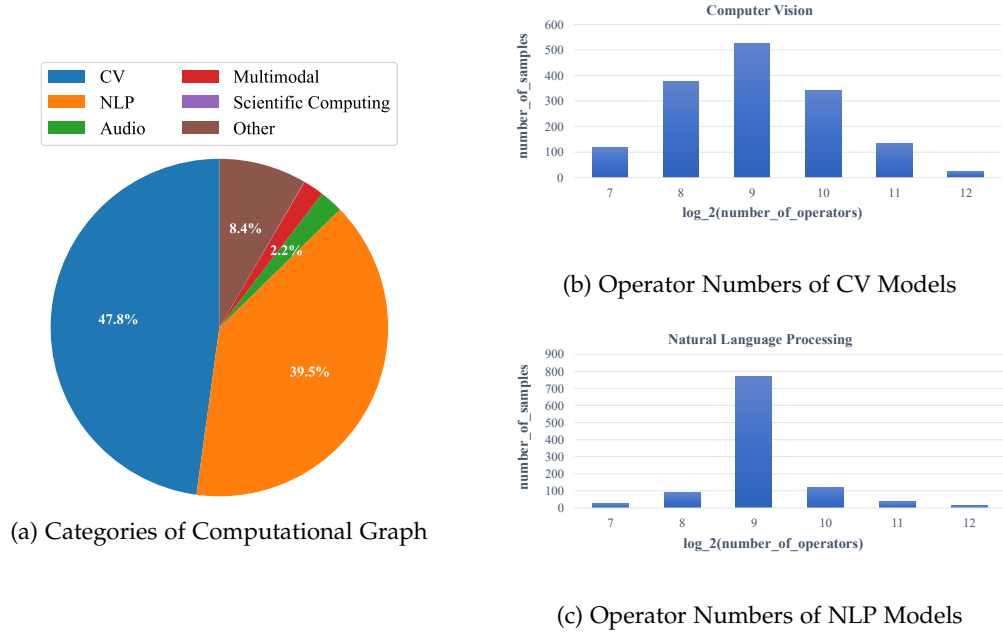


Figure 1: Properties of Computational Graphs Dataset

4. **Performance Analysis:** Compare the outputs collected from baseline and compiled execution to validate correctness. Quantify compiler performance by considering speedup ratio, execution correctness, and performance degradation. The details of the evaluation metrics are discussed in Section 3.2.

## 3.2 Evaluation Metrics

We design two distinct metrics: one for benchmarking tensor compilers and one for compiler development. Section 3.2.1 focuses primarily on acceleration, treating the error penalty as a fixed constant. Section 3.2.2 exposes detailed error penalty information, allowing developers to configure it on demand.

### 3.2.1 Metrics for Compiler Benchmark

The performance evaluation metric is parameterized by a tunable **tolerance**  $t$ , which enables the metric to capture compiler performance under varying correctness criteria, from strict numerical precision to more relaxed settings. Tolerance  $t$  determines whether a sample is considered correctly executed, based on the check `assert_close(x, y, rtol(t), atol(t))`. The values of relative tolerance (rtol) and absolute tolerance (atol) corresponding to different  $t$  settings and data types are summarized in Appendix A.

Based on tolerance  $t$  and its associated correctness criteria, we define the **Speedup Score**  $S_t$  as a unified metric that combines speedup, accuracy, and penalty factors, providing a comprehensive measure of compiler performance:

$$S_t = \alpha^\lambda \cdot \beta^{\lambda\eta p} \cdot b^{1-\lambda} \quad (1)$$

For readability, we omit the explicit ( $t$ ) notation in this expression, while noting that  $\alpha, \beta, \lambda$ , and  $\eta$  are all tolerance( $t$ ) dependent. Eq. (1) can be viewed as the product of three components:

- **Correct executions:**  $\alpha^\lambda$  where  $\alpha$  is the geometric mean speedup of all correctly executed samples, and  $\lambda$  is the fraction of correctly executed samples, acting as the weight.
- **Performance degradation:**  $\beta^{\lambda\eta p}$  where  $\eta$  is the fraction of correctly executed samples that have

speedup  $< 1$  (i.e., slowdowns among correct executions),  $\beta$  is the geometric mean speedup of these slowdown cases, and  $p \in (0, 1)$  is the penalty (default 0.1) applied to emphasize their impact.

- **Failures:**  $b^{1-\lambda}$  where  $1 - \lambda$  denotes the fraction of samples that execute incorrectly (tolerance violations, compilation failures, or runtime crashes), and  $b \in (0, 1)$  is the penalty factor (default 0.1).

The equivalence between the macro formulation and its sample-level interpretation is formally proven in Proposition B.1 (Appendix B).

**Experiment Result:** We conducted experiments on the NVIDIA H20 GPU using CV and NLP samples from GraphNet, evaluating **CINN** (the default compiler backend of PaddlePaddle) and **TorchInductor** (the default compiler backend of PyTorch) on their respective frameworks. As the workloads originate from their native ecosystems, we refer to the two settings as **PaddlePaddle** and **PyTorch** in the following discussion. Detailed experimental settings, including framework versions, are summarized in Appendix D.

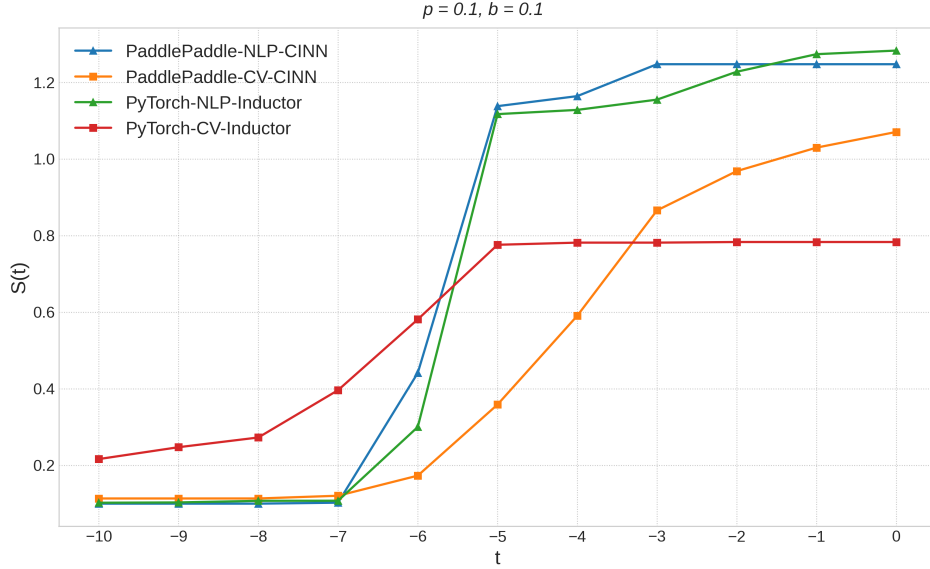


Figure 2: Speedup Score  $S_t$  on NVIDIA H20 for CV and NLP workloads.

As shown in Figure 2, the  $S_t$  trajectories exhibit distinct trends across workloads and compiler configurations. For both PaddlePaddle-NLP (blue) and PyTorch-NLP (green),  $S_t$  remains low under strict tolerances ( $t \leq -7$ ) but rises sharply within  $t \in [-6, -5]$ , indicating that most samples pass the correctness check once the tolerance is slightly relaxed. Both compilers converge to similar peak values around  $S_t \approx 1.2$  at  $t = 0$ , with PaddlePaddle achieving a marginally higher score.

In contrast, PaddlePaddle-CV (orange) exhibits a smoother performance increase from  $t = -7$  to  $t = 0$ , whereas PyTorch-CV (red) reaches a peak  $S_t \approx 0.8$  at  $t = -5$ . These results suggest that PyTorch-CV no longer experiences accuracy-related failures beyond  $t = -5$ ; but its performance remains constrained by an upper bound.

### 3.2.2 Metrics for Compiler Development

For compiler developers, error information from incorrectly executed samples is also crucial. To encode such information into the evaluation metric, we reinterpret the positive domain of the tolerance parameter  $t \in (0, +\infty)$ : each error type is assigned a discrete code  $c \in \{1, 2, 3\}$  for accuracy, runtime, and compilation errors, respectively, and  $t$  now denotes the tolerance level corresponding to these discrete error codes. Specifically,  $t \geq 1$  tolerates accuracy errors,  $t \geq 2$  tolerates execution crashes, and  $t \geq 3$  tolerates compilation failures.

Based on these discrete tolerance levels and error codes, we extend the fixed penalty factor  $b$  in  $S_t$  (Eq. 1) to a tolerance-dependent form  $\gamma_t$ :

$$\gamma_t = \prod_c (b^{\mathbb{1}(t < c)})^{\pi_c} = b^{\sum_c \pi_c \mathbb{1}(t < c)}, \quad (2)$$

where  $c \in \{1, 2, 3\}$  is the error code,  $\pi_c$  denotes the proportion of error code  $c$  among all erroneous samples, and  $\mathbb{1}(\cdot)$  is an indicator function. As  $t$  increases, more error types are tolerated and  $\gamma_t$  monotonically increases from  $b$  to 1. With  $\gamma_t$  defined, we obtain the **Error-aware Speedup Score**:

$$ES_t = \alpha^\lambda \cdot \beta^{\lambda \eta p} \cdot \gamma_t^{1-\lambda}. \quad (3)$$

For  $t \leq 0$ ,  $ES_t$  reduces to the original  $S_t$ . As  $t$  grows over the positive domain, more categories of errors are tolerated and  $ES_t$  increases monotonically. When all errors are tolerated ( $t \geq 3$ ),  $ES_t$  reaches its maximum, representing the theoretical upper bound of achievable compiler performance.

Appendix C analyzes the sample-level meaning of  $ES_t$ .

**Experiment Result:** We obtain  $ES_t$  under the same conditions as in Section 3.2.1, as shown in Figure 3.

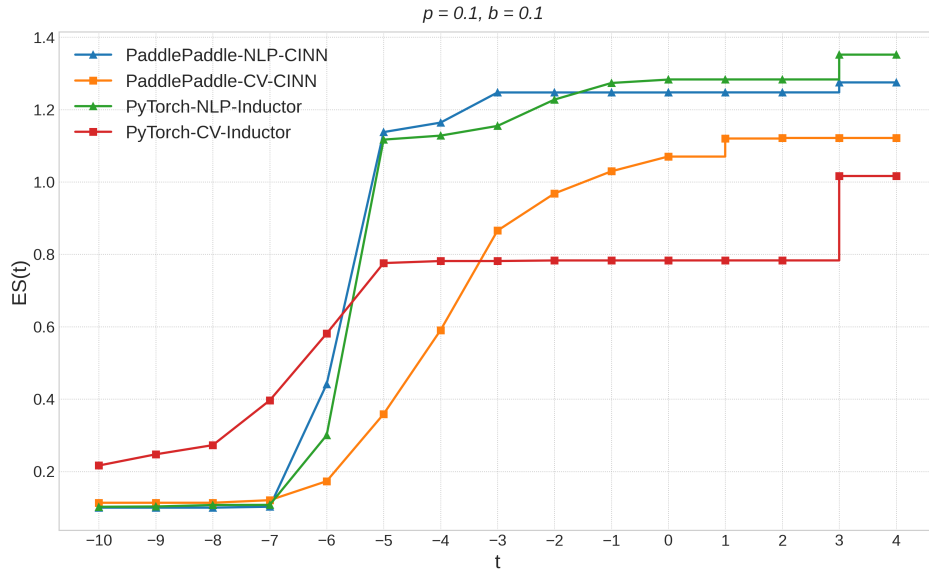


Figure 3: Error-aware Speedup Score  $ES_t$ .

In Figure 3, PaddlePaddle-CV exhibits a discrete jump from  $t = 0$  to  $t = 1$ , indicating that a portion of its failures are due to substantial accuracy violations. Once these tolerance constraints are fully relaxed, its  $ES_t$  score rises immediately. Jumps also occur at  $t = 3$  for both PyTorch-NLP and PaddlePaddle-NLP, revealing that a fraction of samples fail due to compilation errors. Notably, PyTorch-CV shows a sharp rise at  $t = 3$ , suggesting that a large number of its samples do not pass the compilation stage. At last, when  $t \geq 3$ ,  $ES_t$  no longer incorporates any penalties for accuracy violations, runtime crashes, or compilation errors, and thus reflects the raw speedup score—the theoretical upper bound of compiler performance.

While  $S_t$  and  $ES_t$  provide an aggregated view of compiler performance under varying tolerance levels, Figure 6 in Appendix D presents the raw distribution of per-sample speedups at  $t = 1$ , i.e., with accuracy checks fully relaxed. Furthermore, Tables 3–6 in the appendix provide detailed values of each component  $(\alpha, \beta, \lambda, \eta, \gamma)$ , offering a fine-grained view of how  $S_t$  and  $ES_t$  are constructed.

## 4 Construction of GraphNet

GraphNet provides a unified workflow for automated graph extraction, validation, and cross-backend performance evaluation. This section details the dataset requirements and construction methodology.

## 4.1 Dataset Constraints

We define five constraints applied to every computational graph in GraphNet to ensure overall dataset quality and cross-platform compatibility. These constraints act as validation requirements during dataset construction, ensuring consistency and usability at the source. Specifically, all computational graphs must satisfy the following constraints:

- **Runnable:** Each computational graph must successfully execute forward propagation under the designated framework without syntax errors, type mismatches, or runtime crashes.
- **Serializable:** Each sample and its associated metadata (e.g., input shapes, weight parameters) must be serializable into standard formats (e.g., JSON) and correctly de-serializable upon reloading.
- **Decomposable:** The entire computational graph must be decomposable into multiple non-overlapping subgraphs, where each subgraph represents an independent optimization unit. This supports compiler backends in performing fusion, scheduling, and other optimization tasks.
- **Static Analysis of all Operators:** The name, type, attributes, and dependency relations of all operators must be statically extractable (e.g., via `torch.fx`) without model execution.
- **Custom Operator Accessible:** If a sample includes user-defined custom operators, the corresponding source code for these operators must be traceable and accessible in a modular form.

## 4.2 Construction Methodology

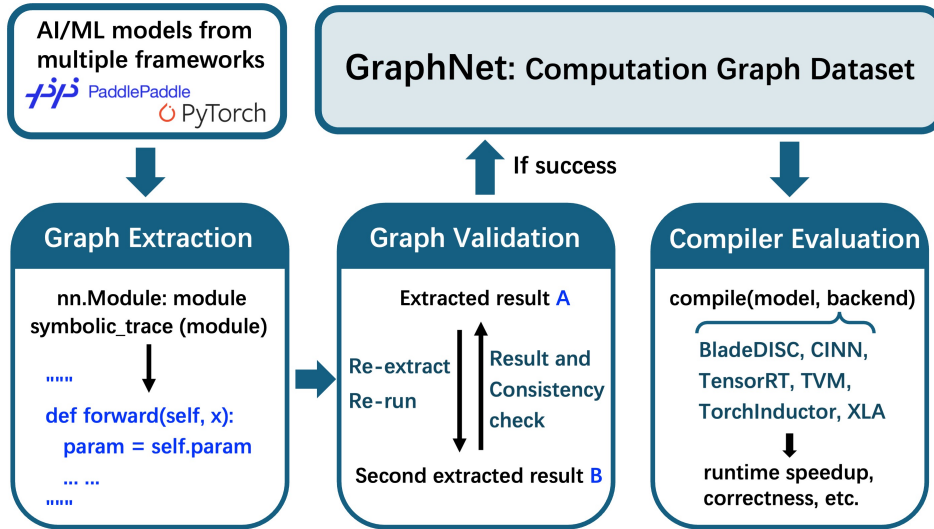


Figure 4: GraphNet Workflow Overview

GraphNet provides an automated workflow for collecting computational graphs across multiple platforms while enforcing the constraints introduced in Section 4.1. The dataset construction pipeline, illustrated in Figure 4, consists of two primary components: graph extraction and graph validation.

### 4.2.1 Graph Extraction

We design a lightweight extraction mechanism that captures dynamic computational graphs from real models and saves them as standardized GraphNet samples. First, we create Python decorator-based interfaces, including `graph_net.paddle.extract` and `graph_net.torch.extract`. Users can wrap the target model with the extractor, which automatically triggers the graph extraction process at runtime. During

execution, the extractor employs symbolic tracing and dynamic graph tracking mechanisms built into the framework to capture all operator invocations and tensor dependencies, thereby generating a complete dynamic computational graph. The captured graph is stored as a standardized set of files, as shown in Figure 5, encompassing the high-level IR of the computational graph in `model.py`, weights and input metadata, and optional custom operator implementations, thereby forming a complete GraphNet sample.

#### 4.2.2 Graph Validation

To ensure that the extracted data meets the requirements, we adopt a re-extraction and re-execution mechanism during validation stage. The workflow consists of four stages. First, starting from an extracted sample A (the original computational graph) generated by the extractor, the validator deserializes it into an executable Python function, reconstructing the model structure with its input and weight metadata. Second, the reconstructed model is executed to verify that the computational graph is runnable. Third, the validator performs the extraction process again on the reconstructed model, producing a second computational graph B. Finally, the validator repeats the reconstruction and execution on graph B, comparing its outputs with those from the second stage and checking structural and node-level consistency between graphs A and B.

This validation procedure inherently enforces all dataset constraints. Failures in serialization or deserialization terminate the reconstruction process; non-runnable models are detected during execution; and inaccessible custom operator code interrupts re-extraction. In addition, the consistency check guarantees that static analysis of all operators is performed completely and correctly.

GraphNet also incorporates a graph deduplication mechanism during data collection to eliminate redundancy. For each extracted computational graph, a unique *graph hash* value is generated from the models source code and graph topology. The validator then identifies and removes duplicate samples by comparing their hash values, ensuring that only distinct computational graphs are retained in the final dataset.

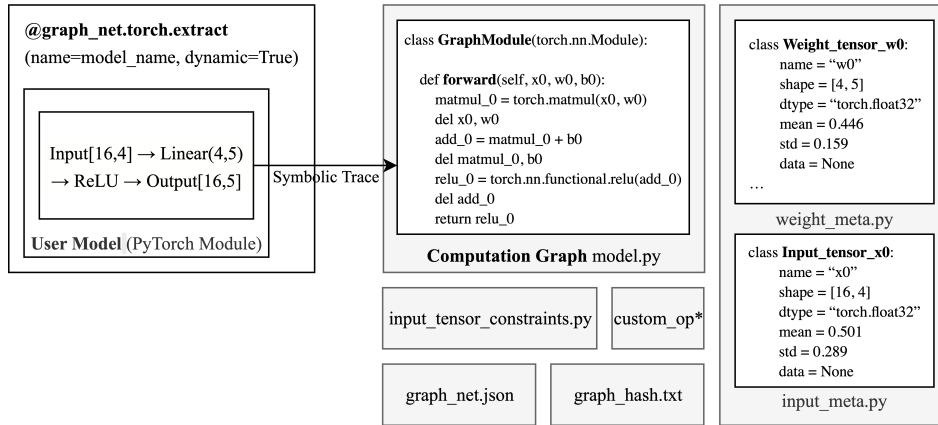


Figure 5: GraphNet Sample Composition

## 5 Background and Related Works

**Tensor Compiler:** Tensor compilers transform high-level computation graphs into optimized device-specific kernels by means of IR lowering and scheduling. Systems such as TVM [1] and Ansor [14] represent search-based compilers that rely on auto-tuning with cost models, while XLA [6] and Glow [11] follow a heuristic-based approach with graph-level optimizations. Recent works include MetaSchedule [12], Hidet [4], and BladeDISC [15], which improve scheduling abstraction, tuning efficiency, and dynamic-shape support. Industry systems include framework-specific compilers such as CINN (PaddlePaddle) and TorchInductor



(PyTorch), which are tightly integrated into their respective DL frameworks, as well as vendor-specific tools like TensorRT, which are specialized for the NVIDIA ecosystem.

**Performance Benchmarks:** DL model performance evaluation can be traced back to DeepBench [8], which measured the performance of fundamental operations such as matrix multiplication across different hardware platforms. This direction was later extended by the industry-wide MLPerf [7] suite, standardizing end-to-end evaluation of training and inference workloads. More recently, CompilerGym [3] introduced benchmark datasets and reinforcement learning environments with the goal of leveraging ML to improve compiler optimization, while KernelBench [10] proposed a benchmark focusing on the correctness and speedup of LLM-generated GPU kernels. In parallel, Furuta et al. [5] evaluated graph compilers across heterogeneous settings and introduced NGraphBench to address the gap between theoretical performance and actual deployment, emphasizing the importance of incorporating compiler effects into ML research.

## 6 Conclusion and Future Work

In this paper, we introduced GraphNet along with  $S_t$  and  $ES_t$  metrics to enable reproducible evaluation of tensor compilers across tasks and frameworks. GraphNet provides an open and well-structured resource for kernel research at the computational graph level, ensuring broad coverage of real-world workloads. Through experiments, we show that these tools allow researchers to gain an objective and comprehensive perspective on compiler optimization capability and to better identify potential performance bottlenecks.

To better align with the needs of compiler researchers and developers, we list our future roadmap here.

### 6.1 Completing GraphNet

**Framework Expansion** GraphNet currently supports only two deep learning frameworks, and its evaluation is limited to NVIDIA GPUs (e.g., H20, A100). We plan to expand the dataset by adding more samples from additional frameworks such as TensorFlow, JAX, and MindSpore, and extend evaluation to a wider range of hardware platforms, including TPUs and NPUs.

**Task Category Refinement** We will introduce more fine-grained categories within existing six major domains (CV, NLP, Audio, Multimodal, Scientific Computing, and Other). This will enable more targeted evaluation of compilers on specific application scenarios.

**Sample Feature Enhancement** We aim to enrich sample features by enabling graph decomposition, allowing full graphs to be split into disjoint subgraphs. We also seek to extend support for custom operators to preserve model-specific functionality.

**Distributed Scenario Support** It is also necessary to incorporate distributed computing scenarios, so that GraphNet can capture computation graphs with communication operators and support the evaluation of compiler optimizations in large-scale distributed systems.

### 6.2 Applying GraphNet

**Systematic Compiler Evaluation** Beside the limited experiment on CINN and TorchInductor, GraphNet can be extended to benchmark a broader range of tensor compilers under a unified metric. From the users perspective, this enables selecting compilers based on task categories and framework requirements. From the compiler developers perspective, it helps quickly identify worst-case scenarios and uncover optimization bottlenecks.

**High-level IR Translation** GraphNet includes computation graphs from multiple deep learning frameworks, providing a foundation for high-level IR translation. Such translation unifies graphs across ecosystems, ensuring that compiler evaluations are based on fully aligned datasets.



**AI for Compiler Research** GraphNet can serve as training and evaluation data for AI-generated compiler passes and kernels. These AI-generated optimizations can be directly applied to GraphNet samples and systematically compared with existing compiler backends under the  $ES_t$  metric. This enables a fair measurement of their speedup and correctness, and positions GraphNet as a benchmark platform for advancing AI-driven compiler research.

## 7 Acknowledgment

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# Appendix

## A Configuration of atol(t) and rtol(t)

Data Type	atol(t)	atol(-5)	atol(0)
float16	$10^t$	1e-5	1
bfloat16	$10^t$	1e-5	1
float32	$10^t$	1e-5	1
float64	$10^{t \cdot 7/5}$	1e-7	1
complex32	$10^t$	1e-5	1
complex64	$10^t$	1e-5	1
complex128	$10^{t \cdot 7/5}$	1e-7	1
quint8	$10^t$	1e-5	1
quint2x4	$10^t$	1e-5	1
quint4x2	$10^t$	1e-5	1
qint8	$10^t$	1e-5	1
qint32	$10^t$	1e-5	1
others	0.0	0.0	0.0

Table 1: atol configuration

Data Type	rtol(t)	rtol(-5)	rtol(0)
float16	$10^{t \cdot 3/5}$	1e-3	1
bfloat16	$10^{t \cdot 1.796/5}$	1.6e-2	1
float32	$10^{t \cdot 5.886/5}$	1.3e-6	1
float64	$10^{t \cdot 7/5}$	1e-7	1
complex32	$10^{t \cdot 3/5}$	1e-3	1
complex64	$10^{t \cdot 5.886/5}$	1.3e-6	1
complex128	$10^{t \cdot 7/5}$	1e-7	1
quint8	$10^{t \cdot 5.886/5}$	1.3e-6	1
quint2x4	$10^{t \cdot 5.886/5}$	1.3e-6	1
quint4x2	$10^{t \cdot 5.886/5}$	1.3e-6	1
qint8	$10^{t \cdot 5.886/5}$	1.3e-6	1
qint32	$10^{t \cdot 5.886/5}$	1.3e-6	1
others	0.0	0.0	0.0

Table 2: rtol configuration

Building on PyTorch’s default testing settings, we develop a log-linear interpolation scheme to construct a (atol, rtol) configuration table. Specifically, we perform log-linear interpolation between two reference points, for instance,  $\text{atol}_{\text{fp32}}(-5) = 10^{-5}$  and  $\text{atol}_{\text{fp32}}(0) = 1$ . In essence, this means that  $\lg(\text{atol}(t))$  and  $\lg(\text{rtol}(t))$  vary linearly with  $t$ , leading to the unified representation  $\text{atol}(t), \text{rtol}(t) = 10^{kt}$ . This formulation provides a coherent framework that encompasses floating-point (float16, bfloat16, float32, float64), complex, and quantized integer types. As a result, the tolerance  $t \in (-\infty, 0]$  is mapped smoothly onto tolerance bounds across diverse precisions.

## B Sample-level Interpretation of $S_t$

In this section, we show that the macro-level Speedup Score  $S_t$  (Eq. 1) can be equivalently expressed as the geometric mean of per-sample rectified speedups.

**Definition B.1** (Rectified Speedup). For each test sample  $i$ , the *rectified speedup* under tolerance  $t$  is defined as:

$$\tilde{s}_{t,i} = \begin{cases} s_i, & \text{if } \text{correct}_{t,i} \wedge s_i \geq 1, \\ s_i^{p+1}, & \text{if } \text{correct}_{t,i} \wedge s_i < 1, \\ b, & \text{if } \neg \text{correct}_{t,i}, \end{cases} \quad (4)$$

where  $s_i$  denotes the raw speedup ratio,  $\text{correct}_{t,i}$  indicates whether the execution satisfies the tolerance criterion  $t$ ,  $p \in (0, 1)$  is the degradation penalty coefficient, and  $b \in (0, 1)$  is the failure penalty. This formulation ensures that: (i) correct executions with speedup retain their measured gain; (ii) correct executions with slowdown are exponentially penalized; and (iii) failed executions incur a fixed penalty.

We further define the *Geometric Mean Rectified Speedup* (GMRS) as:

$$\text{GMRS}_t = \left( \prod_{i=1}^N \tilde{s}_{t,i} \right)^{1/N} \quad (5)$$

where  $N$  is the total number of evaluated test samples.

**Proposition B.1.** The macro-level Speedup Score  $S_t$  defined in Eq. 1 is equivalent to the geometric mean of per-sample rectified speedups:

$$S_t = \alpha^\lambda \cdot \beta^{\lambda\eta p} \cdot b^{1-\lambda} = \text{GMRS}_t$$

*Proof.* Consider a benchmark containing  $N$  test samples in total. Among them, let  $M$  be the number of correctly executed samples, and  $K$  be the subset of those whose raw speedup is less than one (i.e., slowdowns). Formally, we define:

$$\begin{aligned} M &= \#\{i \mid \text{correct}_{t,i}\}, & \lambda &= \frac{M}{N}, \\ K &= \#\{i \mid \text{correct}_{t,i}, s_i < 1\}, & \eta &= \frac{K}{M}. \end{aligned}$$

Thus, the sample space can be partitioned as:

- $(M - K)$  correctly executed and accelerated samples ( $s_i \geq 1$ );
- $K$  correct executions with slowdown ( $s_i < 1$ );
- $(N - M)$  failed or incorrect samples.

Expanding Eq. (5) under this partition gives:

$$\text{GMRS}_t^N = \prod_{i:\text{correct}, s_i \geq 1} s_i \cdot \prod_{i:\text{correct}, s_i < 1} s_i^{p+1} \cdot \prod_{i:\neg \text{correct}} b \quad (6)$$

Define the geometric means:

$$\alpha = \left( \prod_{i:\text{correct}} s_i \right)^{1/M}, \quad \beta = \left( \prod_{i:\text{correct}, s_i < 1} s_i \right)^{1/K}$$

Hence,

$$\prod_{i:\text{correct}} s_i = \alpha^M, \quad \prod_{i:\text{correct}, s_i < 1} s_i = \beta^K,$$

and consequently:

$$\prod_{i:\text{correct}, s_i \geq 1} s_i = \frac{\alpha^M}{\beta^K}$$

Substituting into Eq. (6):

$$\begin{aligned} GMR S_t^N &= \frac{\alpha^M}{\beta^K} \cdot (\beta^K)^{p+1} \cdot b^{N-M} \\ &= \alpha^M \cdot \beta^{Kp} \cdot b^{N-M} \end{aligned}$$

Taking the  $N$ -th root gives:

$$GMR S_t = \alpha^{M/N} \cdot \beta^{Kp/N} \cdot b^{(N-M)/N}$$

Substituting  $\lambda = M/N$  and  $\eta = K/M$  yields:

$$GMR S_t = \alpha^\lambda \cdot \beta^{\lambda\eta p} \cdot b^{1-\lambda} \tag{7}$$

This matches exactly the macro-level formulation in Eq. 1, completing the proof.  $\square$

## C Sample-level Interpretation of $ES_t$

This section provides the sample-level interpretation of the Error-aware Speedup Score  $ES_t$ . We show that the aggregated error penalty  $\gamma_t$  in Eq. (3) is equivalent to the geometric mean of per-sample penalty factors, and that  $ES_t$  can be viewed as the geometric mean of per-sample *error-aware rectified speedup*.

**Definition C.1.** For each erroneous sample  $i$ , let  $c_i \in \{1, 2, 3\}$  denote its error code (1 for accuracy errors, 2 for execution crashes, 3 for compilation failures). Given tolerance level  $t$ , we define the sample-level *penalty factor* as

$$r_{t,i} = \begin{cases} b, & t < c_i, \\ 1, & \text{otherwise,} \end{cases} \quad (8)$$

where  $b \in (0, 1)$  is the base penalty introduced in Section 3.2.1. This factor equals  $b$  if the current tolerance level  $t$  does not forgive the error type  $c_i$ , and 1 otherwise.

**Proposition C.1.** The aggregated penalty  $\gamma_t$  can be written as the geometric mean of  $\{r_{t,i}\}$  over all erroneous samples:

$$\gamma_t = \left( \prod_{i=1}^E r_{t,i} \right)^{1/E} \quad (9)$$

where  $E$  is the total number of erroneous samples.

*Proof.* From the definition of  $r_{t,i}$ , each term in the product  $\prod_i r_{t,i}$  contributes  $b$  iff  $t < c_i$ :

$$\prod_{i=1}^E r_{t,i} = \prod_{i=1}^E b^{\mathbb{1}(t < c_i)} = b^{\sum_{i=1}^E \mathbb{1}(t < c_i)}$$

Grouping by error code  $c$  and letting  $E_c$  denote the number of samples with code  $c$ , we have

$$\sum_{i=1}^E \mathbb{1}(t < c_i) = \sum_{c=1}^3 E_c \mathbb{1}(t < c)$$

thus

$$\prod_{i=1}^E r_{t,i} = b^{\sum_{c=1}^3 E_c \mathbb{1}(t < c)}$$

Taking the  $E$ -th root gives

$$\left( \prod_{i=1}^E r_{t,i} \right)^{1/E} = b^{\frac{1}{E} \sum_{c=1}^3 E_c \mathbb{1}(t < c)} = b^{\sum_{c=1}^3 \pi_c \mathbb{1}(t < c)}$$

where  $\pi_c = E_c/E$  is the fraction of error code  $c$  among all erroneous samples. The right-hand side is exactly the definition of  $\gamma_t$  in Eq. (2).  $\square$

**Definition C.2.** (Error-aware rectified speedup). Extending the rectified speedup defined in Appendix B, we introduce the per-sample *error-aware rectified speedup*:

$$\text{error\_aware\_rectified\_speedup}_{t,i} = \begin{cases} \text{speedup}_i, & \text{correct}_{t,i} \wedge \text{speedup}_i \geq 1, \\ \text{speedup}_i^{p+1}, & \text{correct}_{t,i} \wedge \text{speedup}_i < 1, \\ r_{t,i}, & \text{otherwise.} \end{cases} \quad (10)$$

**Proposition C.2.** Since  $\gamma_t$  acts as the geometric mean of  $\{r_{t,i}\}$  for a given  $t$ , the macro-level metric  $ES_t$  in Eq. (3) can be trivially proved as the geometric mean of these error-aware rectified speedup values:

$$ES_t = \left( \prod_{i=1}^N \text{error\_aware\_rectified\_speedup}_{t,i} \right)^{1/N} \quad (11)$$

## D Details of Benchmark Experiments

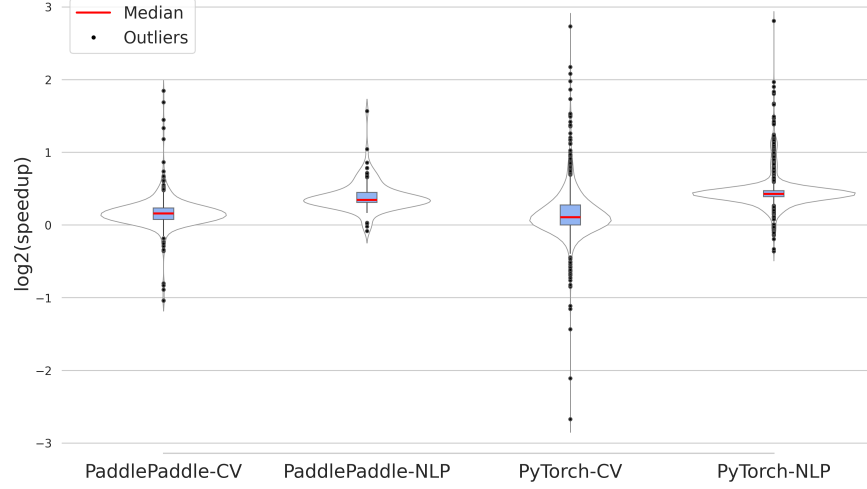


Figure 6: Violin Plots of Per-sample Speedups ( $\log_2(\text{speedup})$ )

Experiments were conducted on a server with dual Intel Xeon Platinum 8563C CPUs, 2 TB DRAM, and eight NVIDIA H20-3e GPUs (144 GB each). Only one GPU was used to ensure a single-device setting. The software environment included Python 3.11, PyTorch 2.8, and PaddlePaddle 3.2.

PaddlePaddle and PyTorch show broadly consistent dataset distributions. For graphs with fewer than 256 operators, the proportions are 6% vs. 8.5% in NLP (PaddlePaddle vs. PyTorch) and 13% vs. 5.6% in CV. For graphs with more than 256 operators, the proportions are 94% vs. 91.5% in NLP and 87% vs. 94.4% in CV. These residual differences will be resolved through a high-level IR transformation tool.

The four tables below (Tables 3-6) report detailed scores for the NLP / CV models tested on PaddlePaddle and PyTorch, corresponding to the statistical results behind each point of the  $S_t$  plots in Section 3.2.1 and the  $ES_t$  plots in Section 3.2.2. Since  $S_t$  does not consider the region  $t > 0$ , the corresponding entries are marked “-” to indicate “empty”.  $\gamma$  is used only when computing  $ES_t$ . For points that contain no correct samples,  $\eta$  is initialized to 0.  $\alpha$  and  $\beta$  are initialized to 1 if no correct samples exist.

t	$\alpha$	$\beta$	$\lambda$	$\eta$	S(t)	$\gamma$	ES(t)
-10	1.000	1.000	0.000	0.000	0.100	0.100	0.100
-9	1.000	1.000	0.000	0.000	0.100	0.100	0.100
-8	1.000	1.000	0.000	0.000	0.100	0.100	0.100
-7	1.441	1.000	0.009	0.000	0.103	0.100	0.103
-6	1.321	0.858	0.575	0.049	0.442	0.100	0.442
-5	1.284	0.878	0.953	0.050	1.139	0.100	1.139
-4	1.282	0.878	0.962	0.049	1.165	0.100	1.165
-3	1.278	0.878	0.991	0.048	1.248	0.100	1.248
-2	1.278	0.878	0.991	0.048	1.248	0.100	1.248
-1	1.278	0.878	0.991	0.048	1.248	0.100	1.248
0	1.278	0.878	0.991	0.048	1.248	0.100	1.248
1	1.278	0.878	0.991	0.048	-	0.100	1.248
2	1.278	0.878	0.991	0.048	-	0.100	1.248
3	1.278	0.878	0.991	0.048	-	1.000	1.276
4	1.278	0.878	0.991	0.048	-	1.000	1.276

Table 3: Values for PaddlePaddle-NLP

t	$\alpha$	$\beta$	$\lambda$	$\eta$	S(t)	$\gamma$	ES(t)
-10	0.981	0.927	0.010	0.800	0.102	0.100	0.102
-9	1.038	0.927	0.013	0.615	0.103	0.100	0.103
-8	1.170	0.840	0.029	0.321	0.107	0.100	0.107
-7	1.175	0.840	0.030	0.310	0.108	0.100	0.108
-6	1.344	0.822	0.423	0.064	0.300	0.100	0.300
-5	1.342	0.817	0.930	0.046	1.118	0.100	1.118
-4	1.345	0.817	0.933	0.046	1.129	0.100	1.129
-3	1.348	0.817	0.941	0.045	1.156	0.100	1.156
-2	1.354	0.818	0.963	0.045	1.229	0.100	1.229
-1	1.360	0.818	0.975	0.045	1.275	0.100	1.275
0	1.363	0.818	0.977	0.044	1.284	0.100	1.284
1	1.363	0.818	0.977	0.044	-	0.100	1.284
2	1.363	0.818	0.977	0.044	-	0.100	1.284
3	1.363	0.818	0.977	0.044	-	1.000	1.353
4	1.363	0.818	0.977	0.044	-	1.000	1.353

Table 4: Values for PyTorch-NLP



$t$	$\alpha$	$\beta$	$\lambda$	$\eta$	$S(t)$	$\gamma$	$ES(t)$
-10	1.056	0.917	0.054	0.739	0.114	0.100	0.114
-9	1.056	0.917	0.054	0.739	0.114	0.100	0.114
-8	1.056	0.917	0.054	0.739	0.114	0.100	0.114
-7	1.056	0.920	0.080	0.529	0.121	0.100	0.121
-6	1.117	0.906	0.228	0.216	0.173	0.100	0.173
-5	1.137	0.906	0.526	0.129	0.359	0.100	0.359
-4	1.131	0.915	0.732	0.109	0.591	0.100	0.591
-3	1.125	0.912	0.892	0.105	0.866	0.100	0.866
-2	1.123	0.911	0.939	0.105	0.969	0.100	0.969
-1	1.122	0.911	0.965	0.105	1.031	0.100	1.031
0	1.121	0.910	0.981	0.105	1.072	0.100	1.072
1	1.122	0.910	0.993	0.104	-	1.000	1.121
2	1.122	0.910	0.993	0.104	-	1.000	1.123
3	1.122	0.910	0.993	0.104	-	1.000	1.123
4	1.122	0.910	0.993	0.104	-	1.000	1.123

Table 5: Values for PaddlePaddle-CV

$t$	$\alpha$	$\beta$	$\lambda$	$\eta$	$S(t)$	$\gamma$	$ES(t)$
-10	0.949	0.825	0.345	0.618	0.217	0.100	0.217
-9	0.954	0.829	0.404	0.623	0.249	0.100	0.249
-8	0.957	0.821	0.447	0.609	0.274	0.100	0.274
-7	0.984	0.819	0.605	0.552	0.399	0.100	0.399
-6	1.006	0.814	0.766	0.505	0.586	0.100	0.586
-5	1.030	0.815	0.882	0.472	0.783	0.100	0.783
-4	1.028	0.815	0.886	0.474	0.788	0.100	0.788
-3	1.028	0.815	0.886	0.474	0.788	0.100	0.788
-2	1.028	0.815	0.887	0.473	0.790	0.100	0.790
-1	1.028	0.815	0.887	0.473	0.790	0.100	0.790
0	1.028	0.815	0.887	0.473	0.790	0.100	0.790
1	1.028	0.815	0.887	0.473	-	0.100	0.790
2	1.028	0.815	0.887	0.473	-	0.100	0.790
3	1.028	0.815	0.887	0.473	-	1.000	1.025
4	1.028	0.815	0.887	0.473	-	1.000	1.025

Table 6: Values for PyTorch-CV