

UVA 558

UVA 558 – Wormholes

This problem involves determining whether a scientist can travel indefinitely **back in time** using a network of wormholes. Wormholes are modeled as **directed edges** between star systems with **time shifts** as weights.

The core challenge is to detect if there exists a **negative weight cycle reachable from the starting star system (node 0)**. Such a cycle would allow the scientist to loop repeatedly, going further back in time with each traversal.

The solution models the star systems as **nodes in a directed graph** and each wormhole from system **x** to **y** with time change **t** as a **directed edge with weight t**. Using the **Bellman-Ford algorithm**, negative cycles are detected by relaxing all edges for **n iterations** (where n = number of nodes). If any distance can still be reduced in the nth iteration, it confirms a negative cycle reachable from the start, resulting in the output "**possible**". Otherwise, output "**not possible**".

Steps

1. Understand the Problem Requirements

- Determine if there exists a cycle of wormholes with **negative total time change**.
- The cycle must be **reachable from start node 0**.
- A negative cycle allows traveling **back in time indefinitely**.

2. Model the Scenario as a Graph

- Represent each star system as a **node** in a directed graph.
- Represent each wormhole as a **directed edge** with weight = time change **t**.
- Positive weights → forward in time, negative weights → backward in time.

3. Apply Bellman-Ford Concept

- Initialize distances: **dist[0] = 0**, all others = ∞ .

- Relax all edges **n-1 times**, updating the earliest reachable times for each node.

4. Detect Negative Cycles

- Perform an additional (nth) iteration over all edges.
- If any distance can still be reduced, it indicates a **negative cycle reachable from node 0**.

5. Interpret the Result

- If a negative cycle exists → print "**possible**".
- If no negative cycle → print "**not possible**".

Key Points

- Bellman-Ford works even with **negative weights**.
- Cycle must be **reachable from start node 0**, so distances initialized carefully.
- The **nth iteration** is critical to detect negative cycles.
- Complexity: $O(n * m)$, where $n = \text{nodes}$, $m = \text{edges}$.

Input:

4 4

0 1 10

1 2 20

2 3 30

3 0 -60

Execution:

Edges stored as (u, v, weight):

- Edge 0: (0, 1, 10)
- Edge 1: (1, 2, 20)
- Edge 2: (2, 3, 30)
- Edge 3: (3, 0, -60)

Step 1: Initialization

- $\text{dist}[0] = 0$
- $\text{dist}[1] = \text{INF}$
- $\text{dist}[2] = \text{INF}$
- $\text{dist}[3] = \text{INF}$

Step 2: First Iteration (i=0) - Relax all edges

- Edge 0: $\text{dist}[0] + 10 = 10 < \text{INF} \rightarrow \text{dist}[1] = 10$
- Edge 1: $\text{dist}[1] + 20 = 10 + 20 = 30 < \text{INF} \rightarrow \text{dist}[2] = 30$
- Edge 2: $\text{dist}[2] + 30 = 30 + 30 = 60 < \text{INF} \rightarrow \text{dist}[3] = 60$
- Edge 3: $\text{dist}[3] + (-60) = 60 + (-60) = 0 \geq \text{dist}[0] \rightarrow \text{no change}$

After iteration 1: $\text{dist} = [0, 10, 30, 60]$

Step 3: Second Iteration (i=1) - Relax all edges

- Edge 0: $\text{dist}[0] + 10 = 10 \geq \text{dist}[1] \rightarrow \text{no change}$
- Edge 1: $\text{dist}[1] + 20 = 10 + 20 = 30 \geq \text{dist}[2] \rightarrow \text{no change}$
- Edge 2: $\text{dist}[2] + 30 = 30 + 30 = 60 \geq \text{dist}[3] \rightarrow \text{no change}$
- Edge 3: $\text{dist}[3] + (-60) = 60 + (-60) = 0 \geq \text{dist}[0] \rightarrow \text{no change}$

After iteration 2: $\text{dist} = [0, 10, 30, 60]$ (no changes)

Step 4: Third Iteration (i=2) - Relax all edges

- All edges: no changes (same calculations as iteration 2)

After iteration 3: $\text{dist} = [0, 10, 30, 60]$ (no changes)

Step 5: Fourth Iteration (i=3) - Check for negative cycles

- Edge 0: $0 + 10 = 10 \geq 10 \rightarrow$ no change
- Edge 1: $10 + 20 = 30 \geq 30 \rightarrow$ no change
- Edge 2: $30 + 30 = 60 \geq 60 \rightarrow$ no change
- Edge 3: $60 + (-60) = 0 \geq 0 \rightarrow$ no change

Step 6: Analysis

- No distance updates occurred in the 4th iteration
- Therefore, no negative cycle reachable from node 0 exists

Output:

not possible

Pseudocode:

```
READ c
FOR each case:
    READ n, m
    edges = []
    FOR m times:
        READ u, v, w
        store edge (u, v, w) in edges

    dist = [INF] * n
    dist[0] = 0

    REPEAT n times:
        FOR each edge (u, v, w) in edges:
            IF dist[u] + w < dist[v]:
                dist[v] = dist[u] + w

    negative_cycle = false
    FOR each edge (u, v, w) in edges:
```

```
IF dist[u] + w < dist[v]:  
    negative_cycle = true
```

OUTPUT "possible" if negative_cycle else "not possible"

Code:

https://github.com/shanto470/algorithm_plm/blob/main/BellmanFord/UVA%20558/uva558.cpp