

# Data Transmission

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August 19, 2020

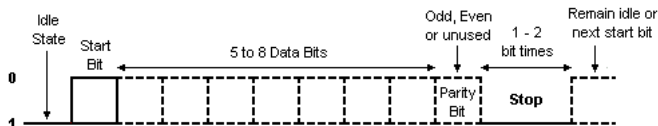
# Asynchronous and Synchronous Transmission

- To correctly interpret signals received from the sender, the receiver's bit intervals must correspond exactly to the sender's bit intervals.
- Receiver's clock is faster or slower, bit intervals are not matched
- Suppose the sender transmits 1 bit every  $\mu\text{s}$  according to its clock.
- The receiver samples every bit at its center, according to its own clock.
- Suppose the receiver clock is 0.01% faster. The first sampling will be 0.01 of a bit time away from the center of the bit
- After 50 samples, the center will be  $50 \times 0.01 = 0.5 \mu\text{s}$  away and it may be in error!
- Eventually, the receiver becomes completely out of step with the transmitter
- Synchronization is essential for correct reception

# Two approaches for synchronization

- Asynchronous transmission

- Avoid the timing problem by not sending long, uninterrupted streams of bits.
- data are transmitted one character at a time, where each character is 5 to 8 bits in length (with start and stop bits)
- Timing or synchronization must only be maintained within each character; the receiver has the opportunity to resynchronize at the beginning of each new character



Character format in asynchronous transmission

# Two approaches for synchronization

- Synchronous transmission: a block of bits is transmitted in a steady stream without start and stop codes. To prevent loss of synchronization:
  - provide a separate clock line between transmitter and receiver; send a pulse from one side to the other
    - works well over short distances, but over longer distances the clock pulses are subject to the same impairments as the data signal, and timing errors can occur
  - embed the clocking information in the data signal
- In both the cases, to allow the receiver to determine the beginning of a block of data, a preamble is sent and to determine the end of a block of data, a postamble is sent.



# Transmission Impairments

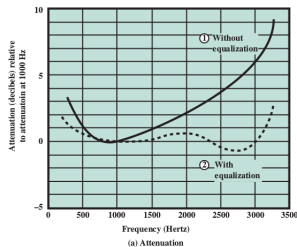
- Attenuation and attenuation distortion
- Delay distortion
- Noise

# Attenuation

- Guided media: attenuation is exponential - expressed as number of decibels per unit distance
- Wireless: a more complex function of distance and the makeup of the atmosphere
- Considerations:
  - 1 A received signal must have sufficient strength so that the electronic circuitry in the receiver can detect and interpret the signal.
  - 2 The signal must maintain a level sufficiently higher than noise to be received without error.
  - 3 Attenuation is greater at higher frequencies, and this causes distortion.

# Attenuation

- Points 1 and 2: Use of amplifiers and repeaters
- Attenuation distortion: attenuation is different for different frequencies
- Solution: Equalize attenuation over a band of frequencies - amplify high frequencies



$$N_f = -10 \log_{10} \frac{P_f}{P_{1000}}$$

Attenuation distortion can present less of a problem with digital signals. Why?

# Attenuation

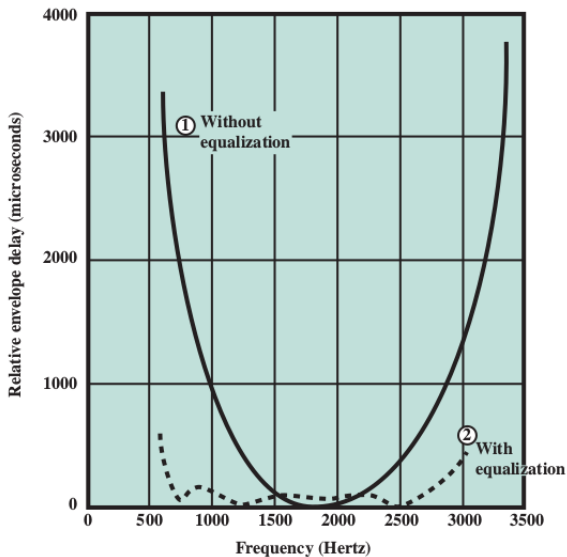
The loss in a cable is usually defined in decibels per kilometer (dB/km). If the signal at the beginning of a cable with  $0.3\text{dB}/\text{km}$  loss has a power of 2mW, what is the power of the signal at 5 km?



# Delay distortion

- Occurs only in transmission cables, not in air
- Velocity of propagation of a signal through a cable is different for different frequencies
- For a signal with a given bandwidth, the velocity tends to be highest near the center frequency of the signal and to fall off toward the two edges of the band.
- Effect: various components of a signal will arrive at the receiver at different times
- Critical for digital data: some of the signal components of 1 bit position will spill over into other bit positions, causing intersymbol interference
- Limits bit rate

# Delay distortion



(b) Delay distortion

- Transmitted signal distorted by transmission medium + noise
- noise: unwanted signals inserted between transmission and reception

- Thermal noise
- Intermodulation noise
- Crosstalk
- Impulse noise

# Thermal noise

- Due to thermal agitation of electrons
- A function of temperature
- Uniformly distributed across the bandwidths typically used in communications systems : white noise
- Cannot be eliminated and therefore places an upper bound on communications system performance

# Thermal noise

The amount of thermal noise to be found in a bandwidth of 1 Hz in any device or conductor is

$$N_0 = kT(W/\text{Hz})$$

$N_0$ : noise power density in watts per 1 Hz of bandwidth

$k$ : Boltzmann's constant =  $1.38 * 10^{-23}$  J/K

$T$ : temperature, in Kelvins (absolute temperature) where the symbol K is used to represent 1 Kelvin

- The dBW (decibel-watt) is used extensively in microwave applications.
- The value of 1 W is selected as a reference and defined to be 0 dBW. The absolute decibel level of power in dBW is defined as:

$$Power_{dBW} = 10 \log \frac{Power_w}{1\text{ W}}$$

- Question: Express a power of 1000 W in dBW
- Question: Express a power of 1 mW in dBW
- Another common unit is decibel-milliwatt.

$$Power_{dBm} = 10 \log \frac{Power_{mW}}{1\text{ mW}}$$

# Thermal noise

What is the thermal noise power density at room temperature ( $17^{\circ}\text{C}$ )?  
Express in  $\text{W/Hz}$  and in  $\text{dBW/Hz}$ .  $k = 1.38 * 10^{-23} \text{ J/K}$



# Thermal noise

- Thermal noise is assumed to be independent of frequency.
- The thermal noise in watts present in a bandwidth of B hertz is:  
$$N = kTB$$
- Given a receiver with an effective noise temperature of 294 K and a 10-MHz bandwidth, the thermal noise level at the receivers output is

# Intermodulation noise

- Due to signals at different frequencies sharing the same transmission medium
- Effect: to produce signals at a frequency that is the sum or difference of the two original frequencies or multiples of those frequencies
- Example: A signal at 4000HZ and one at 8000 Hz sharing the same transmission facility may produce energy at 12000 Hz.
- Caused because the transmitter, the medium or the receiver are not linear — their output is not the input multiplied by a constant.

# Crosstalk

- electrical coupling between nearby twisted pairs (example: a change in current through one wire induces a voltage across the ends of the other wire through electromagnetic induction)
- Microwave antennas picking up unwanted signals
- Usually less than or equal to thermal noise

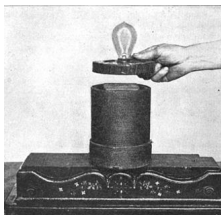


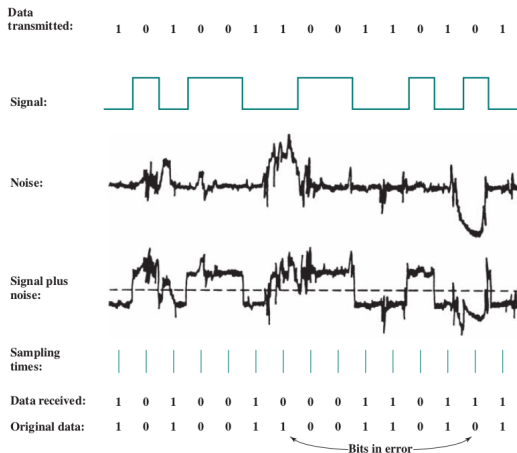
Figure: Example of inductive coupling, Wikipedia

# Impulse noise

The type of noises so far have relatively constant and predictable magnitudes

- Impulse noise is noncontinuous
- consists of irregular pulses or noise spikes of short duration and of relatively high amplitude
- A variety of causes: external electromagnetic disturbances (example?), faults in the system etc.
- Less impact to analog data
- The primary source of error in digital communication

# Effect of noise on a periodic signal



A sharp spike of energy 0.01s on a digital data stream at 56kbps will cause data of 560 bits to be removed.

Question: to what extent do these impairments limit the data rate that can be achieved?

# Channel capacity

The **maximum** rate at which data can be transmitted over a given communication path, or channel, under given conditions, is referred to as the **channel capacity**.

We are trying to relate the four concepts below:

- Data rate: The rate, in bits per second (bps), at which data can be communicated
- Bandwidth: The bandwidth of the *transmitted signal* as constrained by the transmitter and the nature of the transmission medium, expressed in cycles per second, or hertz
- Noise: The average level of noise over the communications path
- Error rate: The rate at which errors occur, where an error is the reception of a 1 when a 0 was transmitted or the reception of a 0 when a 1 was transmitted



# Channel capacity

- The greater the bandwidth, the greater the cost
- All transmission channels have limited bandwidth
- Therefore, we need to make efficient use of bandwidth

For digital data: as *high a data rate* as possible at a particular *limit of error rate* for a given *bandwidth*

# Nyquist Bandwidth

*Consider a noise-free channel*

- If the signals to be transmitted are binary (have only 2 voltage levels) the highest data rate  $C_r$  that can be supported if the bandwidth is  $B$  is  $2B$  bps.
- In general,  $C_r = 2B \log_2 M$ , where  $M$  is the number of discrete voltage levels.
- For example, if 4 voltage levels are used ( $M=4$ ), each level can represent two bits (00, 01, 10, 11).
- For a given bandwidth, the data rate can be increased by increasing the number of different signal elements
  - The receiver has to distinguish more levels
  - Noise and other impairments will limit the value of  $M$

Example: A voice channel of bandwidth 3100Hz is used to transmit data. What is the maximum data rate that can be supported if there are 2 levels? If there are 8 levels?

- Does the Nyquist bit rate agree with the intuitive discussion we had earlier?
- We need to distinguish between the bandwidth of a signal and the bandwidth of the channel carrying that signal
- Case I : We take  $f$  and the first two harmonics,  $3f$  and  $5f$   
Bandwidth of the signal =  $5f - f = 4f = 4 \text{ MHz}$ , assuming  $f = 1 \text{ MHz}$ .  
Period  $T = 1/f = 1 \mu\text{s}$ . Assuming we need to transmit 1010.., 1 bit occurs every  $0.5 \mu\text{s}$ . Data rate =  $1 \text{ bit} / 0.5 \mu\text{s} = 2 \text{ Mbps}$ .
- Let us use only  $f$  and no harmonic frequencies (the worst case, that is, the worst signal). On a *channel* with  $B = 1 \text{ MHz}$ , we can transmit only  $f = 1 \text{ MHz}$  and therefore, the maximum data rate that can be achieved =  $2 \text{ Mbps}$

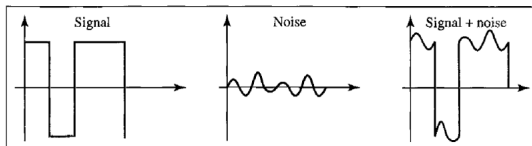
- On a *channel* with  $B=2\text{MHz}$ , we can transmit  $f=2\text{MHz}$  and therefore, the maximum data rate that can be achieved = 4 Mbps. We can also transmit a signal with a data rate of 2Mbps on the same channel, with higher harmonics (better quality).
- Conversely, the *minimum* bandwidth of a channel required to carry a 2 Mbps data rate is 1MHz.
- **The discussion on this slide and the previous one assumes that  $M=2$  and that the channel is noise-free.**

# Noisy Channels

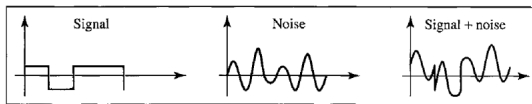
Consider a channel that has noise.

- If data rate is increased, then more bits will occur during the interval of a noise spike, and hence more errors will occur.
- For a given level of noise, we would expect that a greater signal strength would improve the ability to receive data correctly in the presence of noise.
- signal-to-noise ratio (SNR, or S/N) is:  
$$SNR_{dB} = 10 \log_{10} \frac{\text{averagesignalpower}}{\text{averagenoise power}}$$
- SNR is measured at the receiver
- This expresses the amount, in decibels, that the intended signal exceeds the noise level

# Question



A



B

The above picture shows

- (A) A-High SNR, B-High SNR
- (B) A-High SNR, B-Low SNR
- (C) Hard to say
- (D) A-Low SNR, B-High SNR

# Question

The values of SNR and  $SNR_{dB}$  for a noiseless channel are

- (A) infinity, infinity
- (B) 0, infinity
- (C) infinity, 0
- (D) 0,0