## **Team Cactus Project 2 Report**

Finally, produce a brief written project report *in PDF format*. Your report should be submitted as a checked-in file in GitHub. Your report should include the following:

1. Analyze your greedy algorithm code mathematically to determine its big-O efficiency class, probably  $O(n^2)$  or  $O(n \log n)$ .

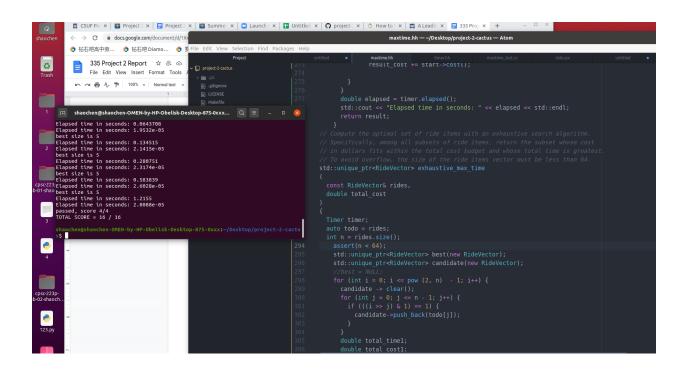
```
GREEDY
  Std: unique_ptrc Ride Vector) greedy_max_time (const Ride Vectors rides, double total-cost)
   - auto todo=rides; -
    Std: unique_ptrcPidevector> result (new Ridevector); --- 0
     double result_cost = 0;
     while (: todo. empty ()) {2
                                                         ? ((og(n)+1)
          duto start = todo(0)
          duto iter = todo, begin (); -
          auto current = iteri-
          for (auto c: todo) &
           if (c > rideTine()/c>cost() > stert > rideTine()/start > cost()) { 3

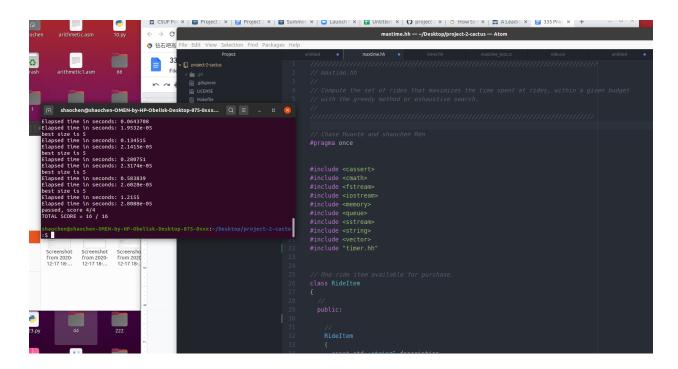
Start = c; 31
4
             iter++
          todo. erasa (carrent);
           if (result_cost + start > cost() (= total_cost) {
               result > push-back (stort);
                result_cost +- start > cost (); _
1= 2+max(2,0) = 2+2=4
2= 3+max (2,0) = 3+2=5
 3= (n-64+1) x (6+1) = 63n x 6= 378n
 4= (logn) x (63nx6) = 63nx6 logn
 5 = 2 + 63 nx6 log(n) = 0 (n log(n))
```

2. Analyze your exhaustive optimization algorithm code mathematically to determine its big-O efficiency class, probably  $O(2^n \cdot n)$ .

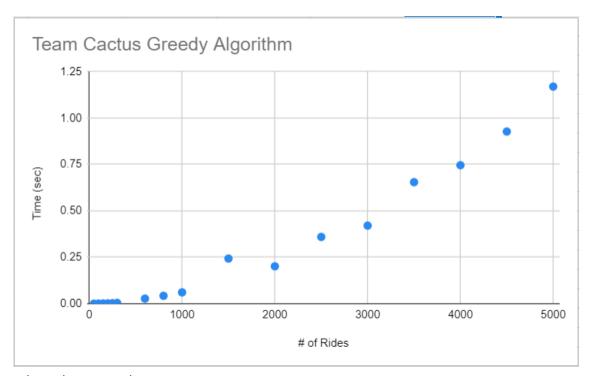
```
auto todo = rides; SC 1
      int n= rides.size(); SC 1
      if(n >= 64){SC 1
            return nullptr;SC 0
      else
            assert(n < 64);SC 1
            std::unique_ptr<RideVector> best(new RideVector); SC 0
            for 0 to 2^n doSC 2^n +1
                   candidate = empty(); SC 1
                   for 0 to n-1 do{SC n
                         if(((i >> j) \& 1) == 1)SC 3
                               candidate.push_back(todo[j])SC 1
                   if (total_cost(candidate) <= total_cost)SC 1
                         if(best.empty() | | total_time(candidate) > total_time(best))SC 2
                               *best = *candidate; SC 1
Step count: first for loop
(2^n + 1) * 1 * n * (3+1) * (1+2+1)
= (2^n+1)*16n
Proof:
(2^n+1)*16n
16n*2^n+16n belongs to O(2^n*n)
Find c > 0 and n0 > n st 16n*2^n+16n
Choose c = 16+16 = 32
16n*2^n+16n \le 32*2^n*n
16*2^n*n-16n*2^n+16*2^n*n-16n \ge 0 V n \ge 0 n0=1
Bu definition that 16n*2^n+16n belongs to O(2^n*n)
```

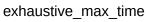
1. Your names, CSUF-supplied email address(es), and an indication that the submission is for project 2.

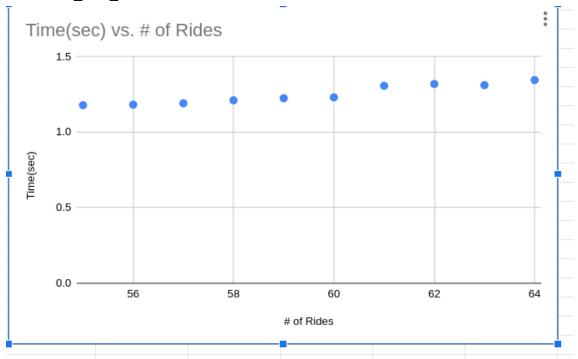




## 2. Two scatter plots meeting the requirements stated above. Greedy







3. Answers to the following questions, using complete sentences.

a. Is there a noticeable difference in the performance of the two algorithms? Which is faster, and by how much? Does this surprise you?

When comparing the 2 functions, greedy is considerably faster than exhaustive search. Reason being is because of exhaustive's exponential power to n which is why we had to bring the list down to 64. Anything longer and the search could be too demanding. Although this would be a surprising moment, we've been anticipating these results the whole time so no we aren't surprised by the time discrepancy.

b. Are your empirical analyses consistent with your mathematical analyses? Justify your answer.

When comparing the empirical analysis to the mathematical analysis they both correlate to a similar median with a few outside outliers. An example is how our greedy function is an O(nlogn) search which can relate to a graph when substituting enough n values for o(nlog(n)) with a slope of zero.

c. Is this evidence consistent or inconsistent with hypothesis 1? Justify your answer.

Consistent to some degree but ideally inconsistent. The reason being is because exhaustive search produces a more correct answer than greedy but with real world implications we cannot always use exhaustive search whether it be time or resource constraints.

d. Is this evidence consistent or inconsistent with hypothesis 2? Justify your answer.

Hypothesis 2 is correct. Exhaustive for example. Once we entered an exponential time complexity when looping all subsets of an n-element sequence, the time considerably jumped per ride. To the knowledge given in the class, o(c^n) is the highest time complexity.