

# Encouraging Peer Grading in MOOCs

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# About Us



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    - Distributed Computing
    - Molecular Computing
    - Computational aspect of Games and Mechanisms

# Algorithmic Game Theory and Mechanism Design

- Understanding and developing **efficient** algorithms for:
  - analyzing games: **Algorithmic Game Theory**
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- Algorithmic Game Theory vs. Game Theory
  - Inefficiency of equilibria
  - **Computational complexity** of finding equilibria
- Algorithmic Mechanism Design vs. Mechanism Design
  - Care about **universal and worst-case** results
  - Care about **implementability**, e.g., polynomial-time computable mechanisms

# Outline

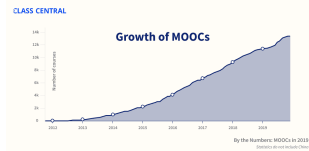
## Motivation

## Model

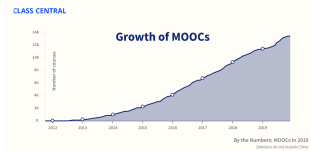
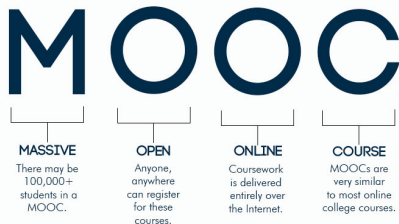
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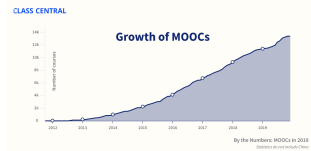
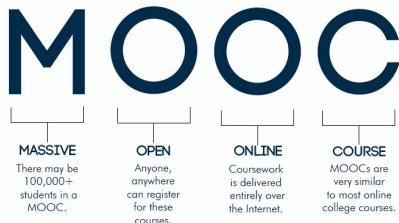


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- Number of engaged learners reached 110M in 2019.<sup>1</sup>
- 13.5K courses by 900+ universities by 2019.

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# Peer Grading/Assessment in MOOCs



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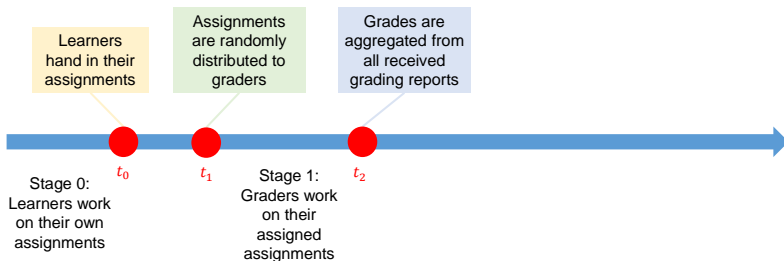
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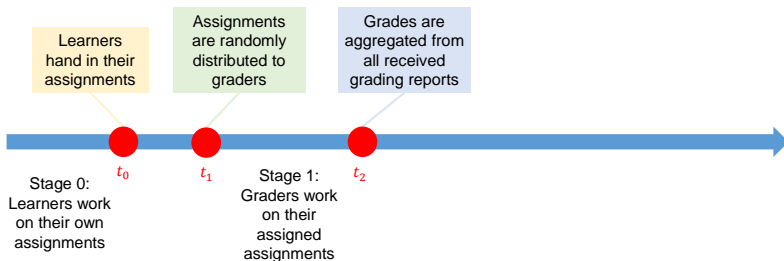


- Extremely high amount of learners overload course staff
- Deployed and analyzed in Coursera.org in large-scale as early as [Piech et al., 2013]
- Still **the only practical solution** to grading **high-level assignments** in MOOCs.

# Basic Peer Grading Mechanism

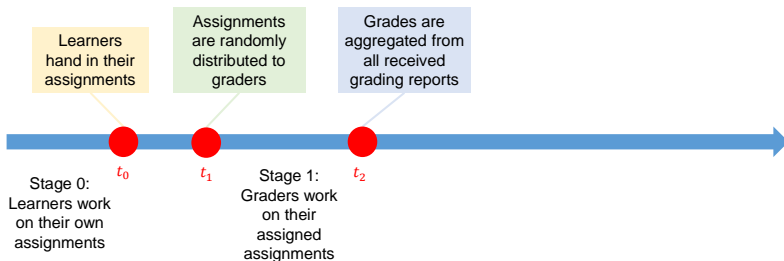


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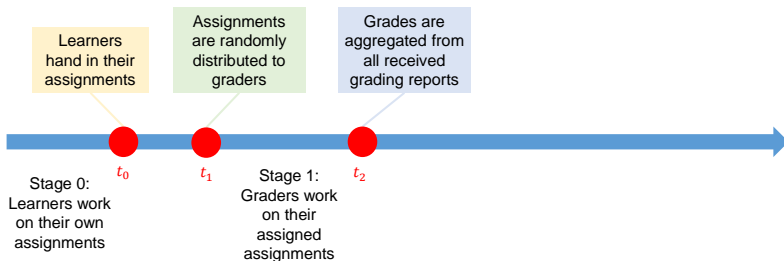
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  - Incentives



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- Efforts put in grading others' work were **not reflected in** concrete **reward**
- Only absent graders, instead of ineffective ones, were punished
- An improved mechanism that **incentivizes effort** in peer grading is necessary.



## Proposed approaches

- Through **punishment**: a Stackelberg game that limited TA resources are efficiently allocated to double check the grading [Carbonara et al., 2015]

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- Through **reward**: Crowdgrader<sup>2</sup>: incentivizing **accurate** graders directly [de Alfaro and Shavlovsky, 2014]

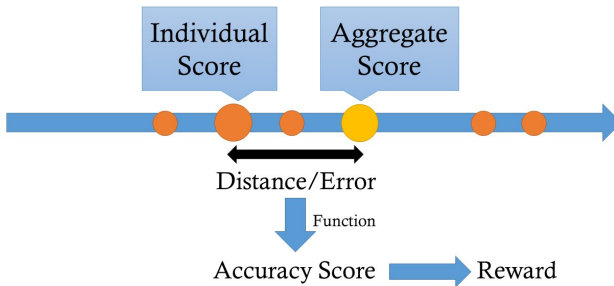
crowdGrader

CrowdGrader lets students submit and collaboratively grade their solutions to homework assignments.

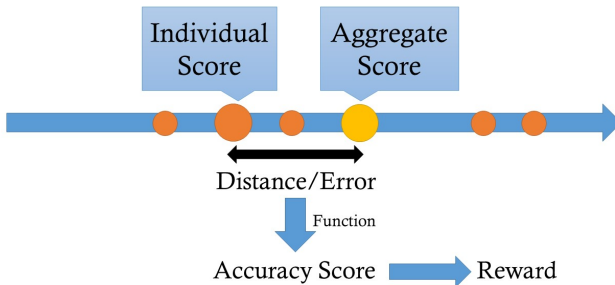
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<sup>2</sup><https://www.crowdgrader.org/>

# Crowdgrader in a Nutshell



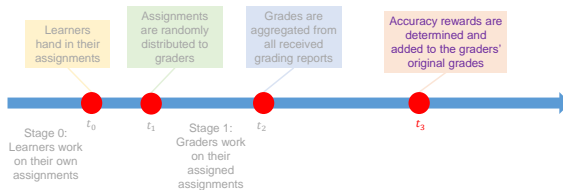
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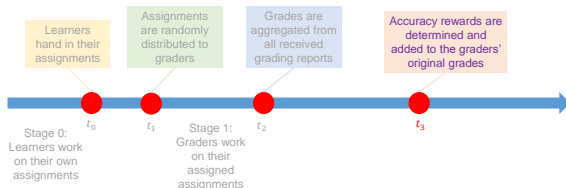
- Crowd-grading is found **as effective and satisfactory** (to the learners) **as TA-grading** [de Alfaro and Shavlovsky, 2014]



# Revamped Peer Grading Mechanism

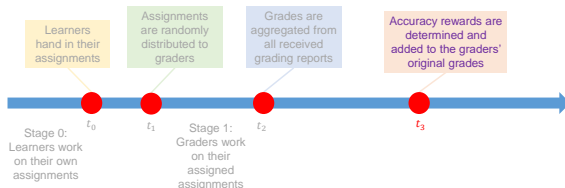


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  - While the settings in Crowdgrader work empirically, are they **theoretically robust** as well?

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  - Finding **real settings** that satisfy the necessary conditions

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- Each submission is graded by exactly  $k$  learners



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- Efforts are **unobservable**

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- We assume **truthfulness**, i.e.,  $a_j$  directly reports  $S_j^i$  to the mechanism without manipulation

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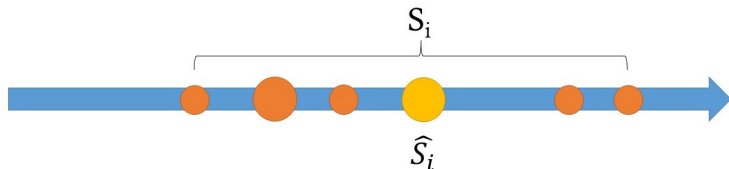
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  - Examples: Average, Median, Olympian Average, etc.



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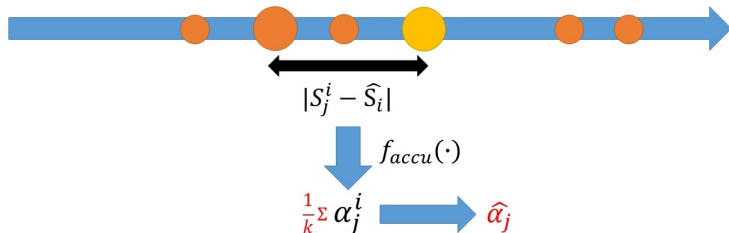
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- $\hat{\alpha}_j = \frac{1}{k} \sum_i (\alpha_j^i)$ :  $a_j$ 's average accuracy

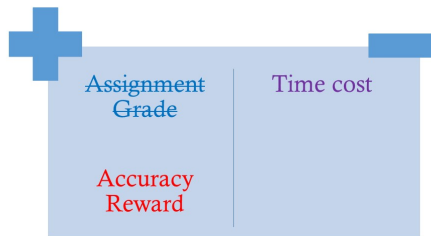
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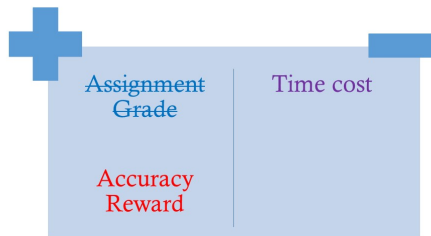


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- $\pi_i = \cancel{(1 - \lambda)M\hat{S}_i} + \lambda M\hat{\alpha}_i - r_i \sum_j t_i^j$ : final utility of learner  $i$

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## Definition: EC-1

A setting satisfies **EC-1** if:

Given the others' strategies  $T_{-j} = [t_1, t_2, \dots, t_{j-1}, t_{j+1}, \dots, t_k]$ ,  $\mathbf{E}[\alpha_j](T_{-j}, t_j)$  is **non-decreasing** and **concave** on  $t_j \in [0, U]$ .

More effort means more accuracy, with diminishing marginal increment.

## Definition: EC-2

A setting satisfies **EC-2** if:

For any pair of two strategy profiles  $T_{-j}$  and  $T'_{-j}$  s.t.:

- $t_p < t'_p$  for some  $p$ ,
- $t_q = t'_q \forall q \neq p$ ,

it holds that  $\frac{\partial}{\partial t_j} \mathbf{E}[\alpha_j](T_{-j}, t_j) < \frac{\partial}{\partial t_j} \mathbf{E}[\alpha_j](T'_{-j}, t_j), \forall t_j$ .

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- A positive reinforcement.

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## Proposition

Define  $f_{\text{avg}}(\mathbf{S}) = \frac{1}{k} \sum_j (S_j)$ . Suppose that  $f_g \sim N(v, g^2)$  where  $g = g(t_j)$  is a non-increasing convex function. If  $f_{\text{agg}}(\cdot) = f_{\text{avg}}(\cdot)$ ,  $f_{\text{accu}}(\cdot)$  is non-increasing piecewise continuous, then for any values of  $(M, k, \mathbf{r}, U, \lambda)$ ,  $(M, k, \mathbf{r}, U, \lambda, f_g(\cdot), f_{\text{agg}}(\cdot), f_{\text{accu}}(\cdot))$  satisfies both EC-1 and EC-2.

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- Closely related to Nash equilibria

# Encouraging Grading

## Theorem

- Assume  $k$  is fixed, both EC satisfied in  $G_1$  with  $\bar{r} = \bar{r}_1$
- $G_3$  differs with  $G_1$  only in  $\bar{r}_3 < \bar{r}_1$
- $T_1 = [t_{1i}]$  is an equilibrium in  $G_1$

$\implies$  There exists an equilibrium  $T_3$  in  $G_3$  where  $t_{3i} \geq t_{1i} \forall i$ .  
If the  $i$ -th equality holds, then  $t_{1i} = U$ .

- Decreasing  $\bar{r}$  **distorts the entire equilibria upwards!**



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  - If the learners are homogeneous (think of they are **all of the same type**), then every grader **gives identical level of effort** in pure NE.
  - If some graders have fixed strategies (think of TA paddings), then the equilibria holds as long as irrational players' strategies are public information.

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- Special case: homogeneous grading
- Extension: irrational graders, TAs, biased grading, etc

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- **Truthfulness is an assumption**
- Application is limited, while other similar scenarios exist: peer grading is actually very similar to other peer assessment scenarios, e.g., **peer reviewing** in academics, or **crowdsourcing**.





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