

## Cider

# A Case for Block Level Variable Redundancy on a Distributed Flash Array

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\*This work does not reflect the policy of NSF in any way



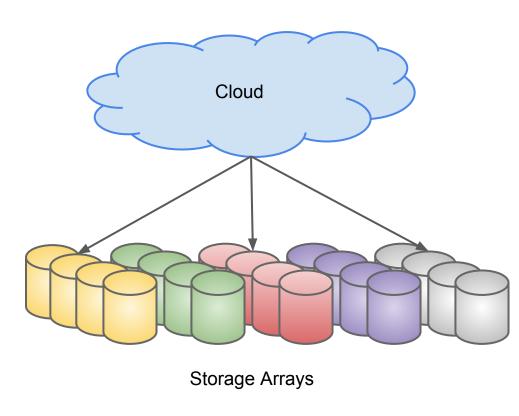
## Welcome to the Petabyte World

- Data is exploding and hoarded
- To large for traditional data centers
- Enterprise data migrating to cloud
- Cloud: Lower costs & higher reliability



## So, is cloud just a large data center?







#### **Cloud Data Centres**

- Different from traditional data centers
- Much larger capacity
- Infrastructure shared by many clients with different needs
  - o QoS, Security, etc.
- Need high flexibility



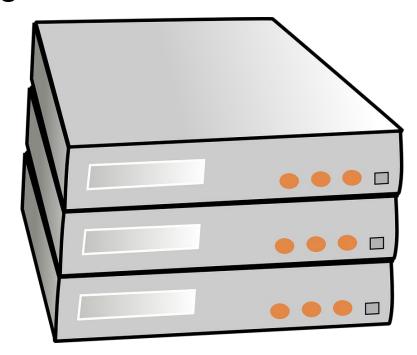






#### **Traditional Solutions: Features**

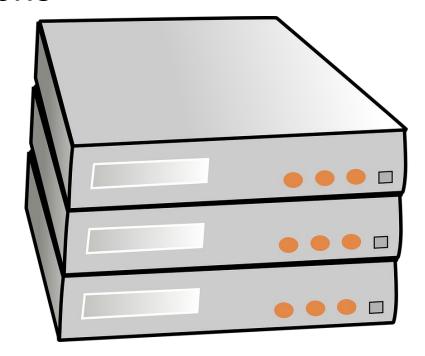
- RAID for redundancy
- SAN based block storage or NAS servers
- Interface used for traditional POSIX applications
- Replication across data centers or storage racks





#### **Traditional Solutions: Limitations**

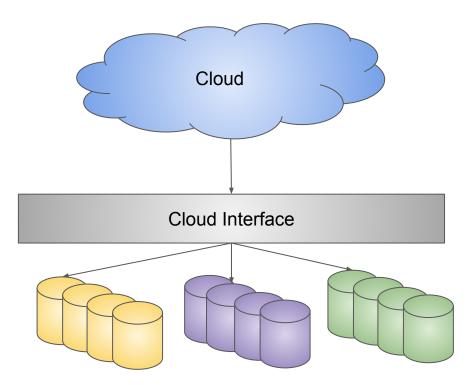
- Rigid: Cannot change the redundancy level once created
- Failure domain cannot span across multiple data centers
- High reconstruction time
- Unrecoverable Read Error





## Cloud Storage

- Redesign entire stack
- Often built on top of traditional data centers addressing some of the shortcomings
- Erasure coding and full replication
- Cost/Byte increases with reliability guarantee



**Traditional Storage Solutions** 



## User/Application's View of Data





## Example: Financial Data



- Maximum reliability for 7 years (Legal requirement)
- Medium reliability for 7-10 years
- Low reliability for 10+ years



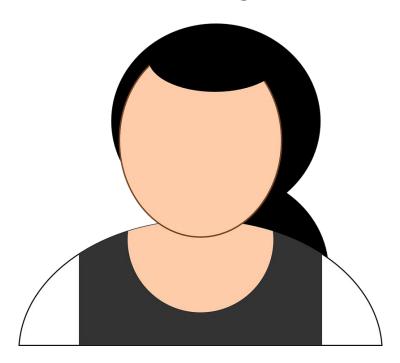
## Example: Data from Internet of Things



- High reliability to start with
- Reliability decreases temporally



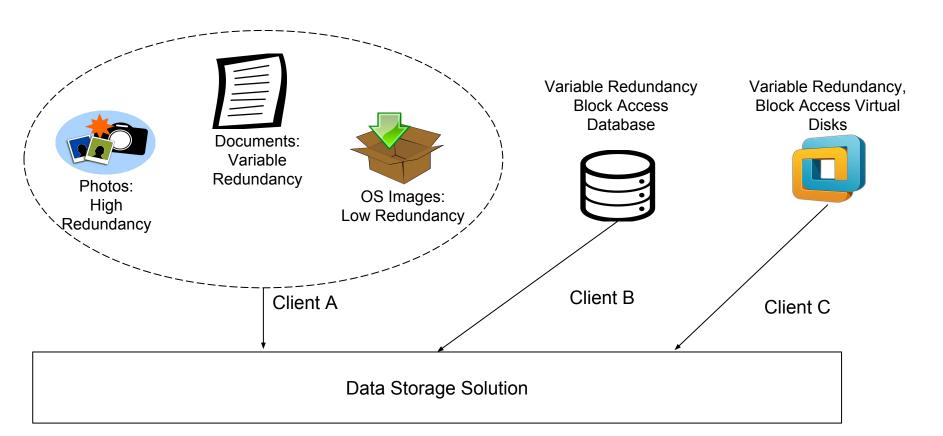
## Example: Regular Users



- Health records: Forever
- Personal photos: Highest reliability
- Other documents: Medium priority
- Downloaded files: Low priority



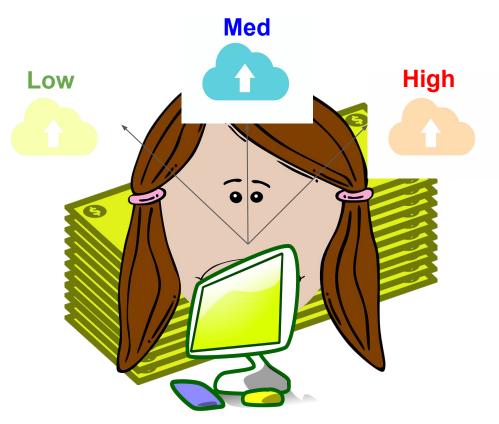
## Multiple Clients





#### User's Choice

- All data with the highest reliability!
- Different sets of data in different cloud stores
- Less than ideal





#### **Problem Statement**

Can we create a more flexible storage backend that the cloud can readily use?



#### Cider: Goals

Build a highly flexible storage backend suitable for the cloud.

- Similar semantics to existing backend (RAID, etc.)
- Variable redundancy
- Temporally variable redundancy
- Scalable
- Good performance

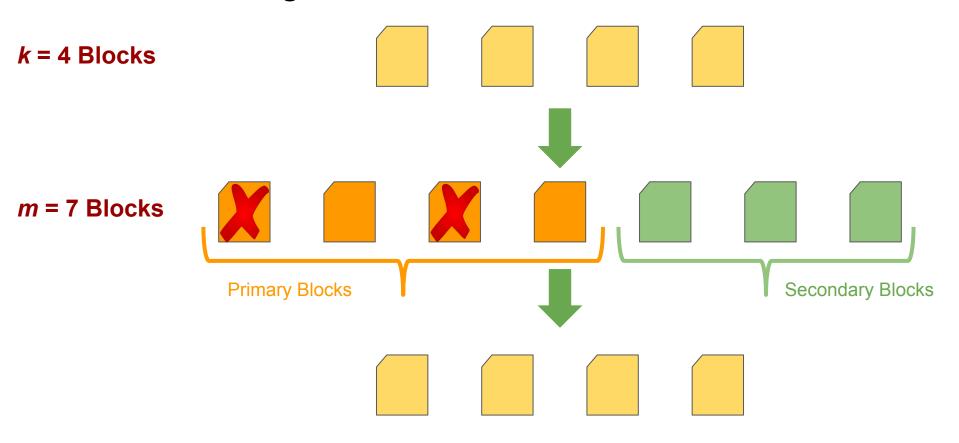


#### ...Rest of The Talk

- Recap on erasure coding in storage
- Variable redundancy blocks
- Adopting variable redundancy to flash array
- Prototype implementation
- Future work



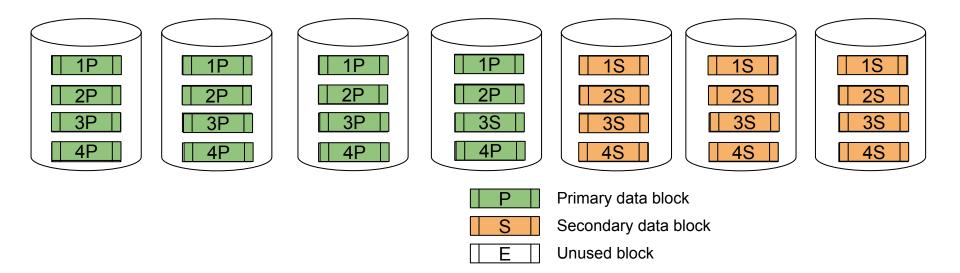
## **Erasure Coding**





## **Erasure Coding with Constant Scheme**

Scheme: 4 primary + 3 secondary





## **Constant Redundancy Distribution**

- Translation from Virtual Block numbers to Physical Block numbers is very straightforward
- Virtual to Physical address map can be computed statically
- Examples of systems using this technique at file level and less commonly at block level.



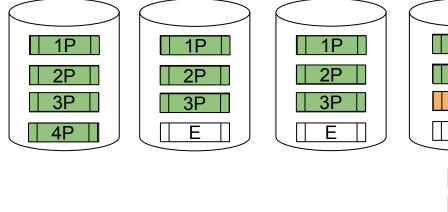
#### Variable Redundancy Distribution

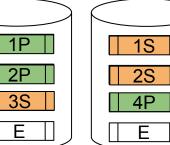
1<sup>st</sup> Block: 4 + 3

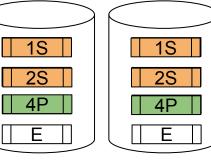
2<sup>nd</sup> Block: 4 + 2

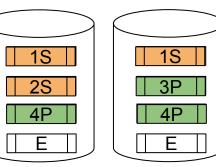
3<sup>rd</sup> Block: 4 + 1

4<sup>th</sup> Block: 4 + 0









P Primary data block

Secondary data block

E Unused block



#### Variable Redundancy Distribution

#### Features

- Allows heterogenous data to be stored with different redundancy
- Maximum freedom to configure granularity of redundancy
- Erasure coding allows failure domains can span across multiple servers
- Flexibility with reconstruction times

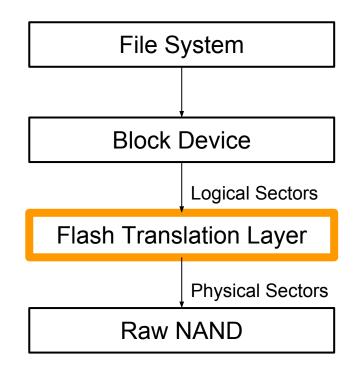
#### Challenges

- Flexibility comes at an extra cost of a lookup or address translation
- Additional storage overhead for metadata associated with mapping
- Block allocations have to carefully managed
- How do we mitigate the extra overhead?



## Rethinking Variable Redundancy for Flash

- Flash Translation Layer
  - Address translation
    - Block allocation
  - Garbage collection
  - Wear-Leveling
- Blocks are converted from Virtual block to physical blocks



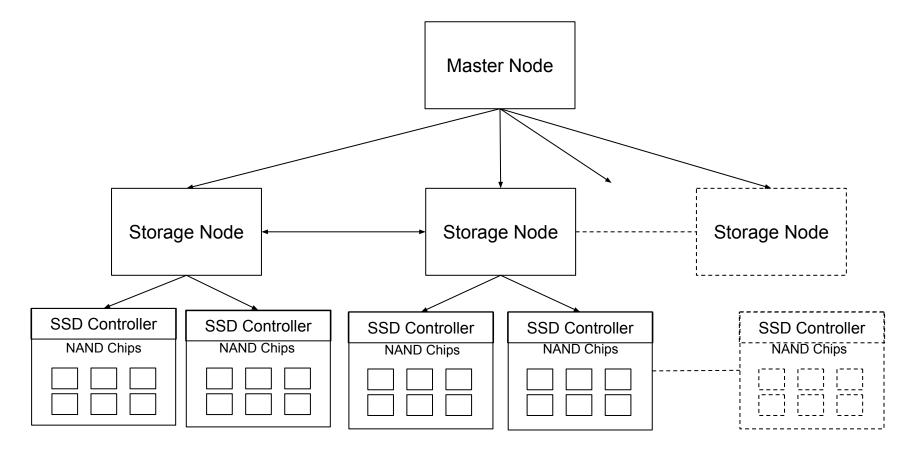


## Case Study: Application Managed Flash

- SSD based arrays v/s HDD based arrays
- FTL: Highly restrictive, unpredictable
- Application managed FTL
  - Combining with deduplication, etc.
  - Higher control
- Challenges: Extracting internal parallelism
  - NAND chips, dies, channels, planes
- Cider with FTL



#### Case Study: Globally Managed Flash Array





## Use Case Study: SSD Simulator

- Modified FlashSim to enable a globally managed FTL
- 5 SSDs in a RAID Class
- Local translation, garbage collection and block allocation is disabled
- Global FTL as pluggable module

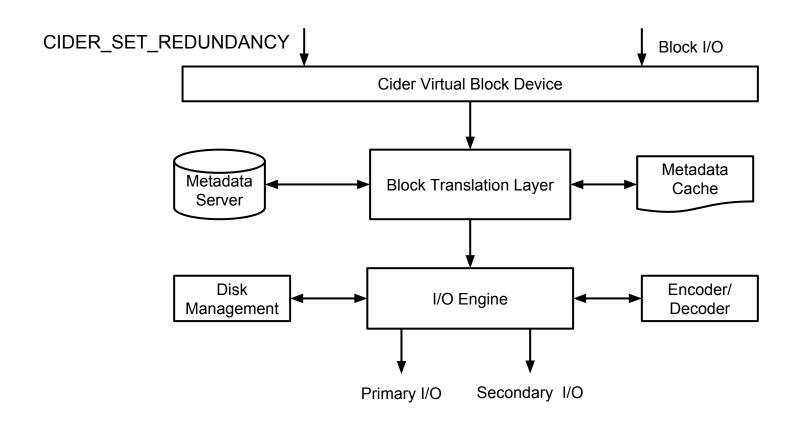


#### Cider: Architecture

- Virtual Block Device: Backed up off shelf physical storage devices
- ioctl(CIDER\_SET\_REDUNDANCY, redundancy)
- Writes
  - Uses the previously set redundancy level
  - ioctl/write should be automatically called to ensure correct redundancy
  - Changing redundancy through write with no data
- Reads
  - Uses the previously saved metadata
  - Parallel reads from backing disks

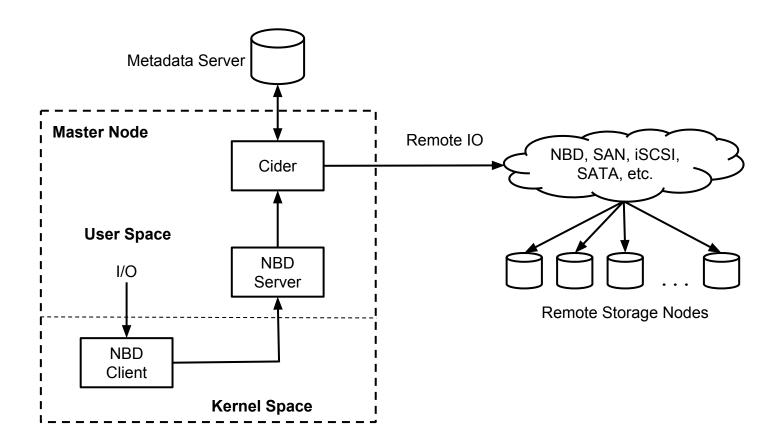


#### Cider Architecture





## Prototype Implementation



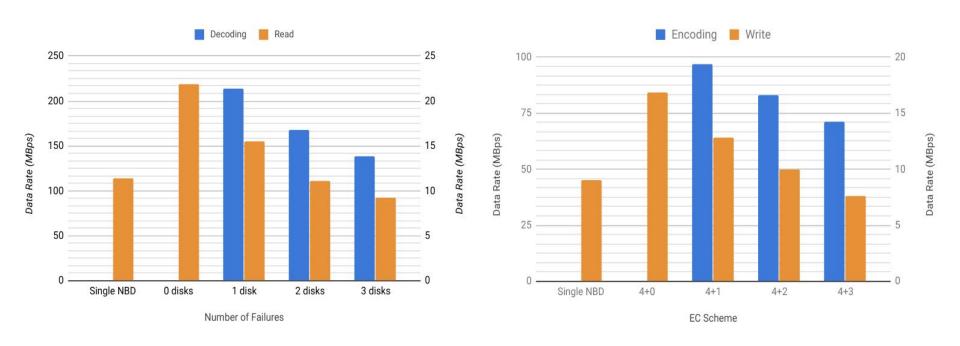


#### **Evaluation: Preliminary Results**

- Microbenchmarks on off-the-shelf hardware:
- 7 hard disks connected through NBD
- Writes throughputs with different EC schemes
- Reads throughputs with different failures
- Raw NBD device as a base performance



#### Microbenchmarks: Preliminary Results





#### Discussions and Future Work

- Disk recovery and handling failures
- Disk defragmentation
- Extracting full parallelism



#### Related Work

- Distributed and Block storage
  - o Petal, Blizzard, Network RAID
- Erasure Coding
  - Placement of chunks
  - Faster computations
- Flash and FTL
  - MinFlash, NoFTL



#### Conclusion

- Classify and store data based on their criticality
- Limitations of existing solutions
- Cider as an alternative for a block level



#### Thank You!

# A Case for Block Level Variable Redundancy on a Distributed Flash Array

(Patent Pending PCT/US2015/026267)

Presented at The 26th International Conference on Computer Communications and Networks (ICCCN 2017)

July 31 -August 3, 2017, Vancouver, Canada

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