# TCP/IP Overview

# Helpful References

- Forouzan, Data Communications and Networking, 5ed
- <a href="http://tcpipguide.com/free/t\_toc.htm">http://tcpipguide.com/free/t\_toc.htm</a>
- https://www.geeksforgeeks.org/services-and-segmentstructure-in-tcp/

#### **Network Review**

- When it comes to designing, implementing, debugging, and testing networking software, we must be aware of
  - Network layers
  - Network headers
  - Network protocol flow

## Network Layer Review

- Physical Layer:
  - Responsible for physical links between adjacent devices
  - Determines how to represent data in electromagnetic form
- Data Link Layer:
  - Responsible for getting data from one device to a neighbouring device in the current network
  - Breaks up data into individual frames
  - Handles error detection and correction as required
- Ethernet encompasses physical and data link layers

## Network Layer Review

- Network Layer:
  - Responsible for determining network routes
  - Gets a packet from a source network to a destination network
  - IP
- Transport Layer:
  - Deals with sender and receiver issues: (De)multiplexing, packet ordering, flow control, etc.
  - TCP (reliable but slower), UDP (unreliable but faster)

#### **IPv4** Header Review

Version	Header Length		ECN	Header + Data Length	ID	Flags	Offset	TTL	Proto	Header Check sum	Src IP	Dst IP	Opts
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- Differentiated Services Code Point (DSCP) can be used to describe the payload (e.g., Voice over IP), allowing the router to make smarter choices
- Explicit Congestion Notification (ECN) can be used to alert the destination of impending congestion, which can then alert the source

#### **IPv4** Header Review

Version	Header Length	DSCP	ECN	Header + Data Length	ID	Flags	Offset	TTL	Proto	Header Check sum	Src IP	Dst IP	Opts

- ID + Flags + Offset used to enable/identify fragments
- Time To Live (TTL); usu. a per-router hop count; at 0, packet is discarded and an ICMP Time Exceeded error is transmitted to the sender
- Proto identifies the payload type

#### TCP Header Review

Src Port	Dst Port	Sequence Number	Ack. Number	Header Length	0	Flags	Window Size	Chksum	Urgent Pointer	Opts
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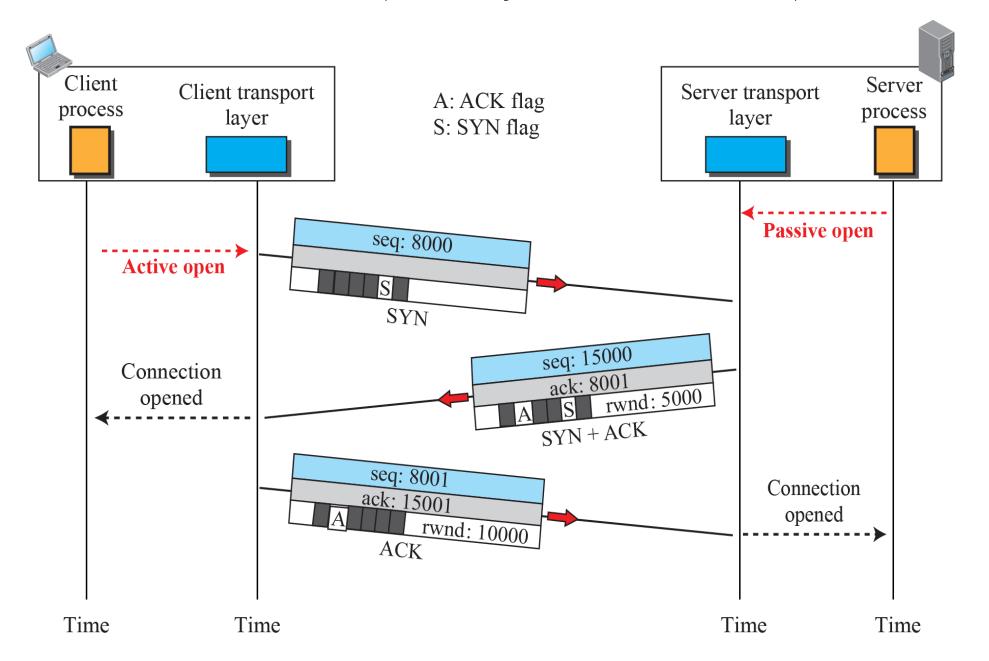
 Urgent Pointer points to the location of the urgent data in the payload section (if URG flag set)

## **UDP Header Review**

Source Port	Destination Port	Header + Data Length	Checksum (Optional if used with IPv4)
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# Forouzan, Data Communications and Networking, 5ed

# TCP Connection (3-way Handshake) Review



#### Connection Establishment Notes

- ACK = Acknowledgement
- rwnd = Receive window (used for flow control)
- SYN = Synchronize sequence numbers
- SYN and SYN/ACK segments do not carry data
- An ACK without data does not consume a sequence number
- Full duplex requires two receive windows

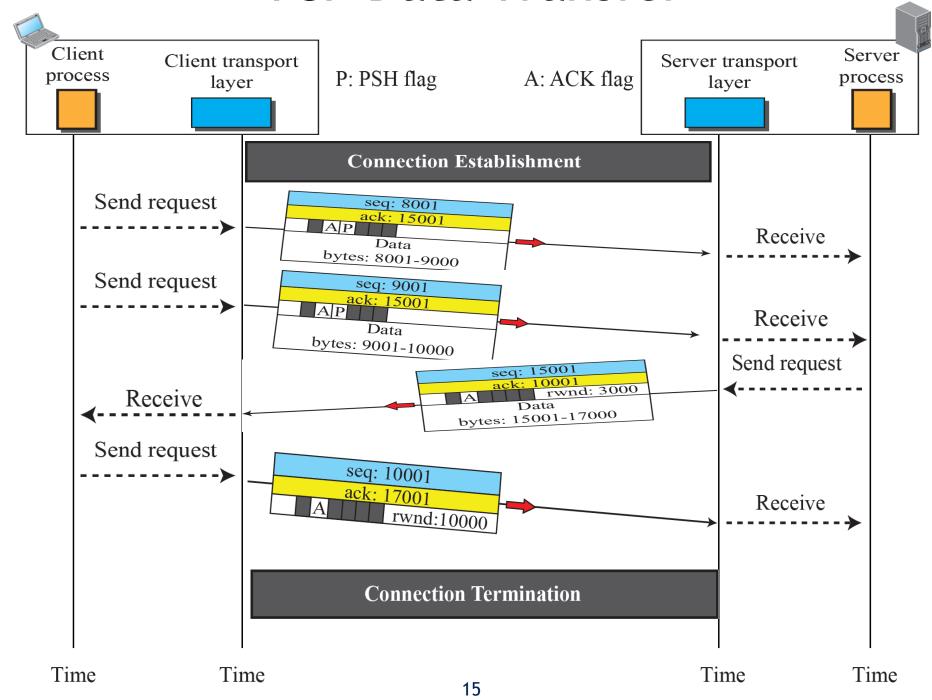
# Sequence Numbers

- Initial number is random (Wireshark obscures this by default to make analysis easier)
- Subsequent numbers depend on number of bytes transmitted
  - e.g., assume first number is 10001
  - If 1000 bytes are transmitted, the sequence number range will be 10001 to 11000
  - The next segment would start at 11001

A TCP connection is established using a

During TCP connection establishment, if the server receives a packet with the SYN flag set and a sequence number set to 18573, it must reply with the \_\_\_\_\_\_ flag(s) set and the sequence number set to \_\_\_\_\_ and the acknowledgement number set to \_\_\_\_\_

## TCP Data Transfer

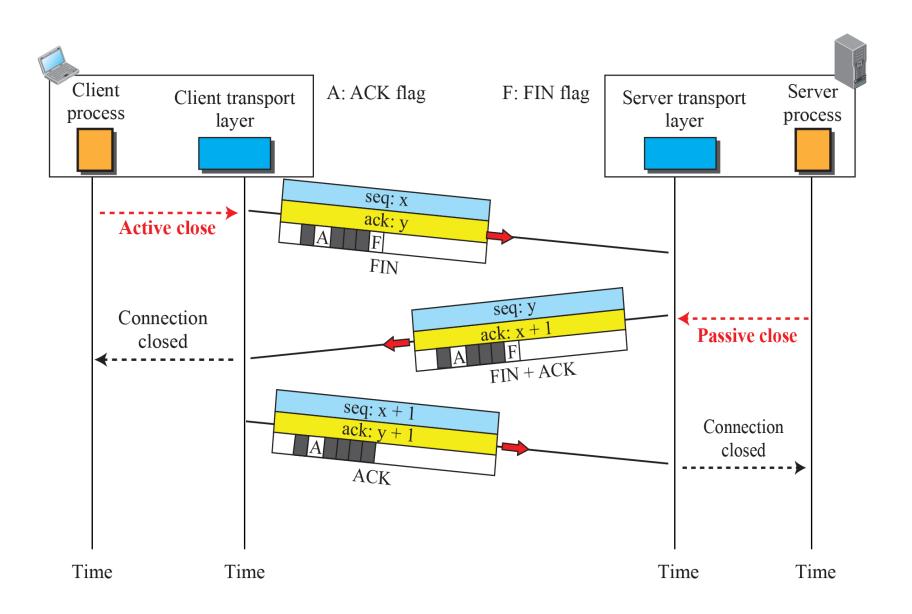


#### **Data Transfer Notes**

- The sequence number is referring to the first byte in the segment
- Note the window size update and the ACK number
- The push flag indicates to the destination that buffered data should be made available to the application right away, rather than delay for more data to accumulate
- Can emphasize urgency even more by setting the URG flag; points to urgent data in the data stream and causes a signal to be sent to the listening application

• If the TCP segment sequence number is 58293 and 100 bytes are sent, the acknowledgement number must be \_\_\_\_\_

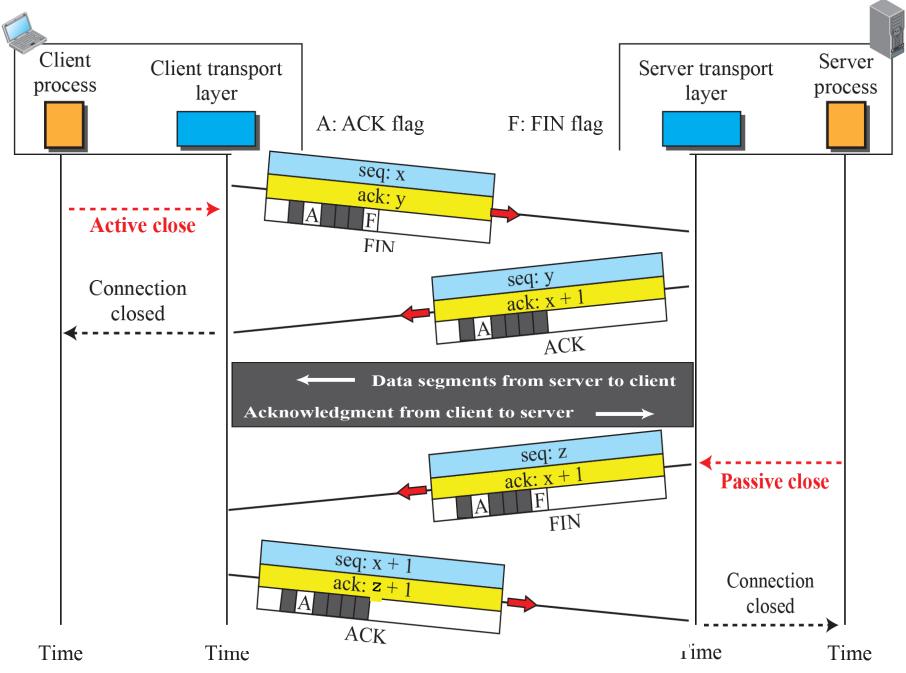
## **TCP Connection Termination**



#### **Connection Termination Notes**

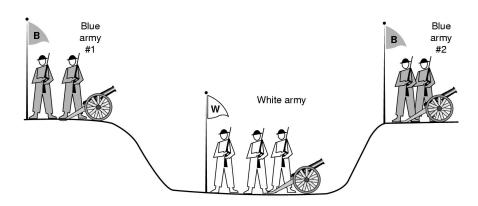
- FIN/FIN+ACK consumes 1 sequence number even if it contains no data
- Last ACK contains no data
- But what if more data must be sent? (e.g., a database query takes time to retrieve?)

#### TCP Half Close



# Two Army Problem

- Assume the final ACK gets lost
- How do we know that the left hand side won't wait forever for a FIN?
- Problem: Do we ACK the ACK?

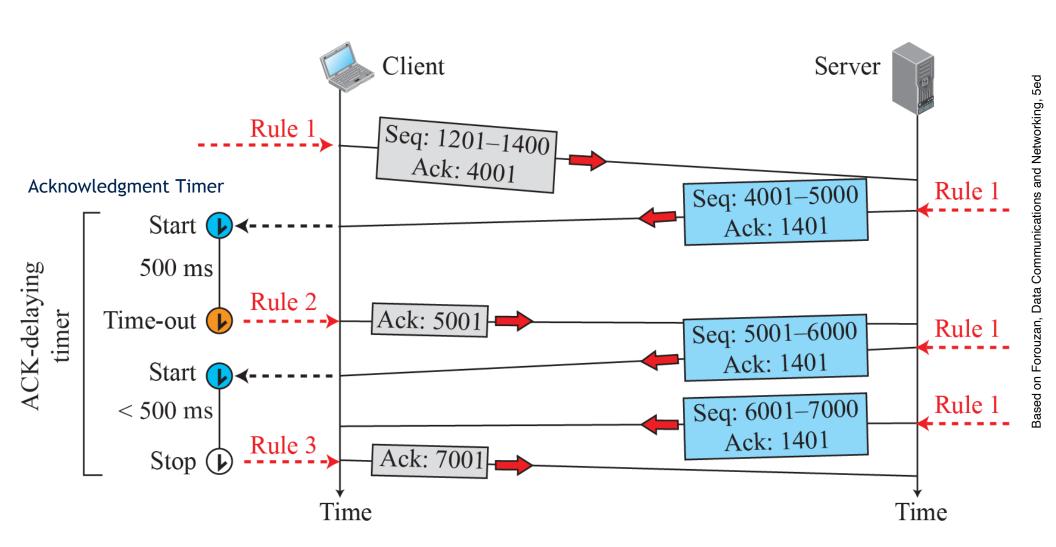


Source: Andy Tanenbaum, Computer Networks, 4th ed., Prentice-Hall

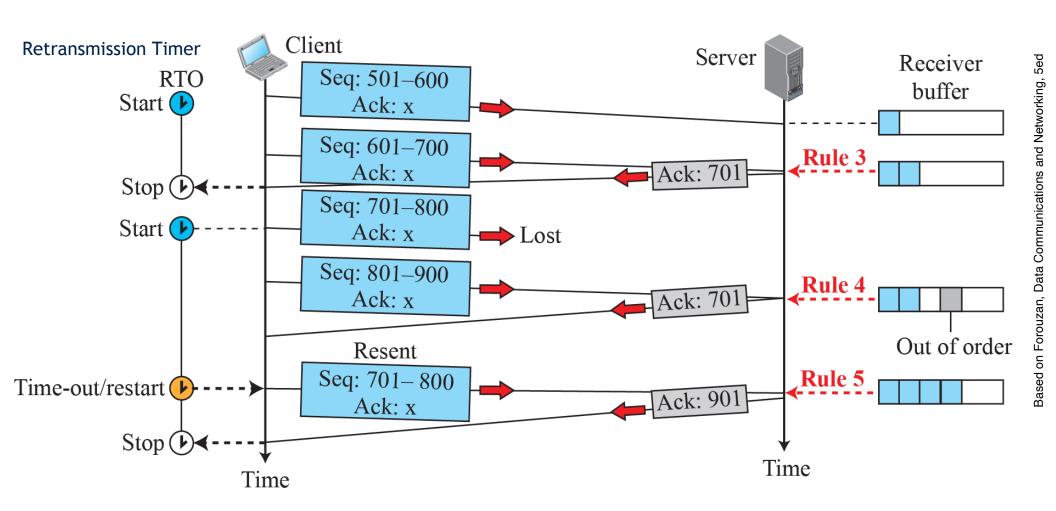
# Two Army Problem

- TCP "solution" to this problem uses timeouts
- Usually works in practice

- To handle lost or damaged segments, timers are set by both sides to determine whether or not to retransmit data or ACKs
- Rule 1: When there is data to be sent, the sender must include an ACK, with the acknowledgement number set to the next sequence number it expects to receive from the receiver
- Rule 2: When the recipient has received a valid segment from the sender, and the recipient has no data to send in return, it must set a timeout after which an ACK must be returned to the sender, with the acknowledgement number set to the next expected sequence number
- Rule 3: When a second unacknowledged segment arrives, the recipient must return an ACK to the sender, with the acknowledgement number set to the next expected sequence number

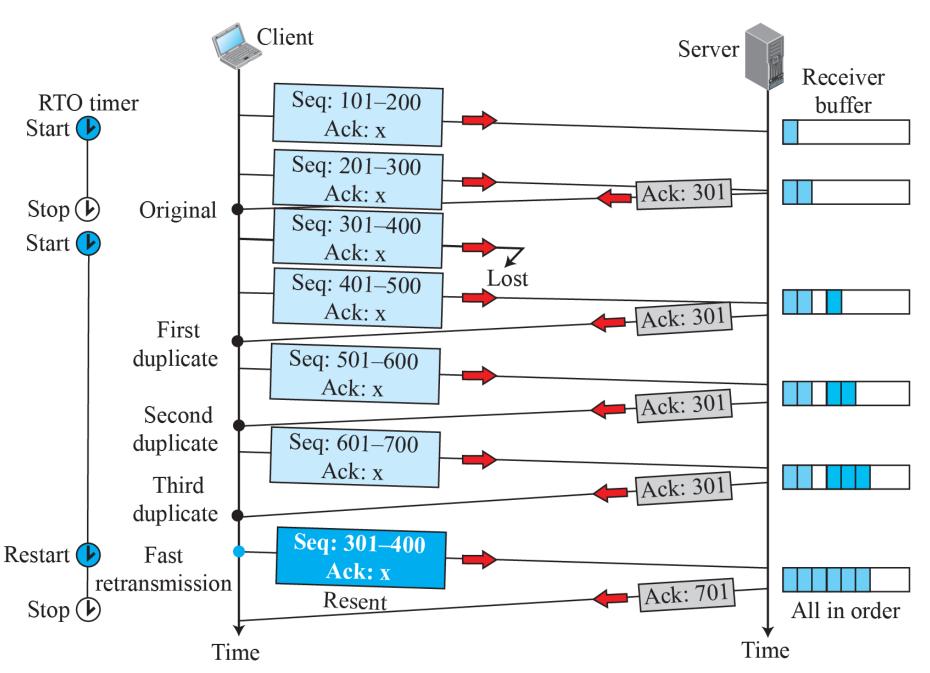


- Rule 4: When an out-of-order sequence number arrives, the segment is buffered and an ACK is returned to the sender, with the acknowledgement number set to the next expected sequence number
- Rule 5: When a missing segment arrives, the recipient returns an ACK to the sender, with the acknowledgement number set to the next expected sequence number



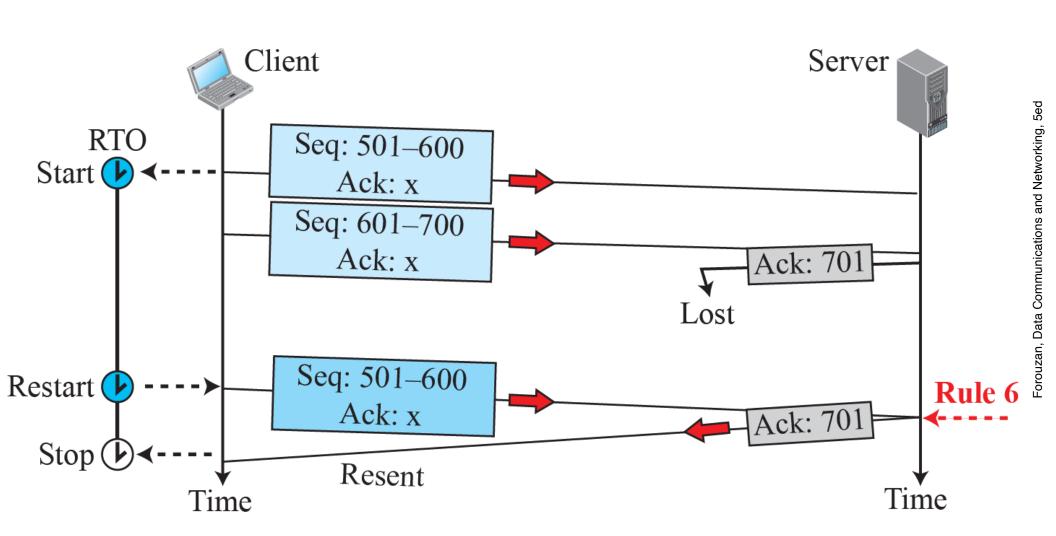
• Fast Retransmission Rule: If a 3rd duplicate ACK arrives, the subsequent segment is deemed lost, its retransmit timer is restarted, and the segment is retransmitted

#### Fast Retransmission



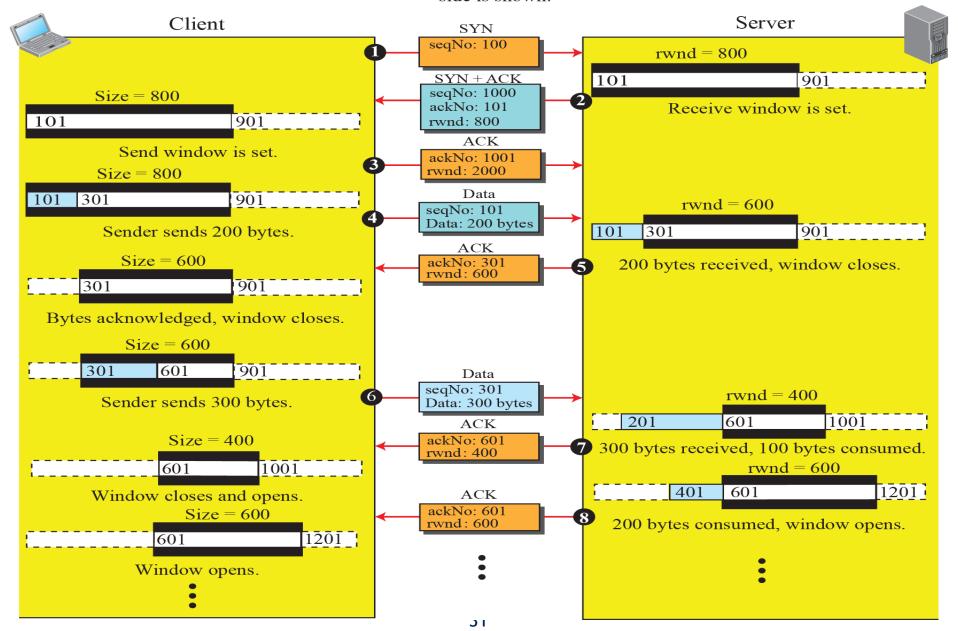
 Rule 6: When a duplicate segment is received, it is discarded, but the recipient returns an ACK to the sender, with the acknowledgement number set to the next expected sequence number

#### Lost ACK



#### TCP Flow Control

**Note:** We assume only unidirectional communication from client to server. Therefore, only one window at each side is shown.



 If a valid TCP segment is received from a sender, the recipient must \_\_\_\_\_\_, assuming the recipient has no valid data to send

 If a duplicate TCP segment is received from the sender, the recipient must \_\_\_\_\_\_, \_\_\_\_ and set the acknowledgement number to \_\_\_\_\_\_

Assuming the receive buffer already contains bytes 1000 - 1200 and a TCP segment arrives with sequence number 1400 and length 100, the recipient must \_\_\_\_\_ and set the acknowledgement number to \_\_\_\_ and save bytes \_\_\_\_\_

# Key Skills

• Explain TCP connection establishment, data transmission, error control, flow control, and connection termination