

I HLD

Client Server Architecture

- Request Response model

- every computer / process in a network functions as a server or client.
 - client: any comp. that requests something from server
 - server: any comp. designed to serve the requests to client
- ~~if~~ → Client, web browser access website & sends request to server, web server → retrieves necessary webpage & resources & sends back to browser for display

- online gaming → game client running on player's device interacts with game server
- Can app server act as client?
 - Yes!
 - eg - DB connectivity → app server need to retrieve / update data from DB server
 - app server - client sending DB queries / requests to DB server
 - app servers might communicate with external services / APIs to retrieve / send data.

Communication / Networking in HLD

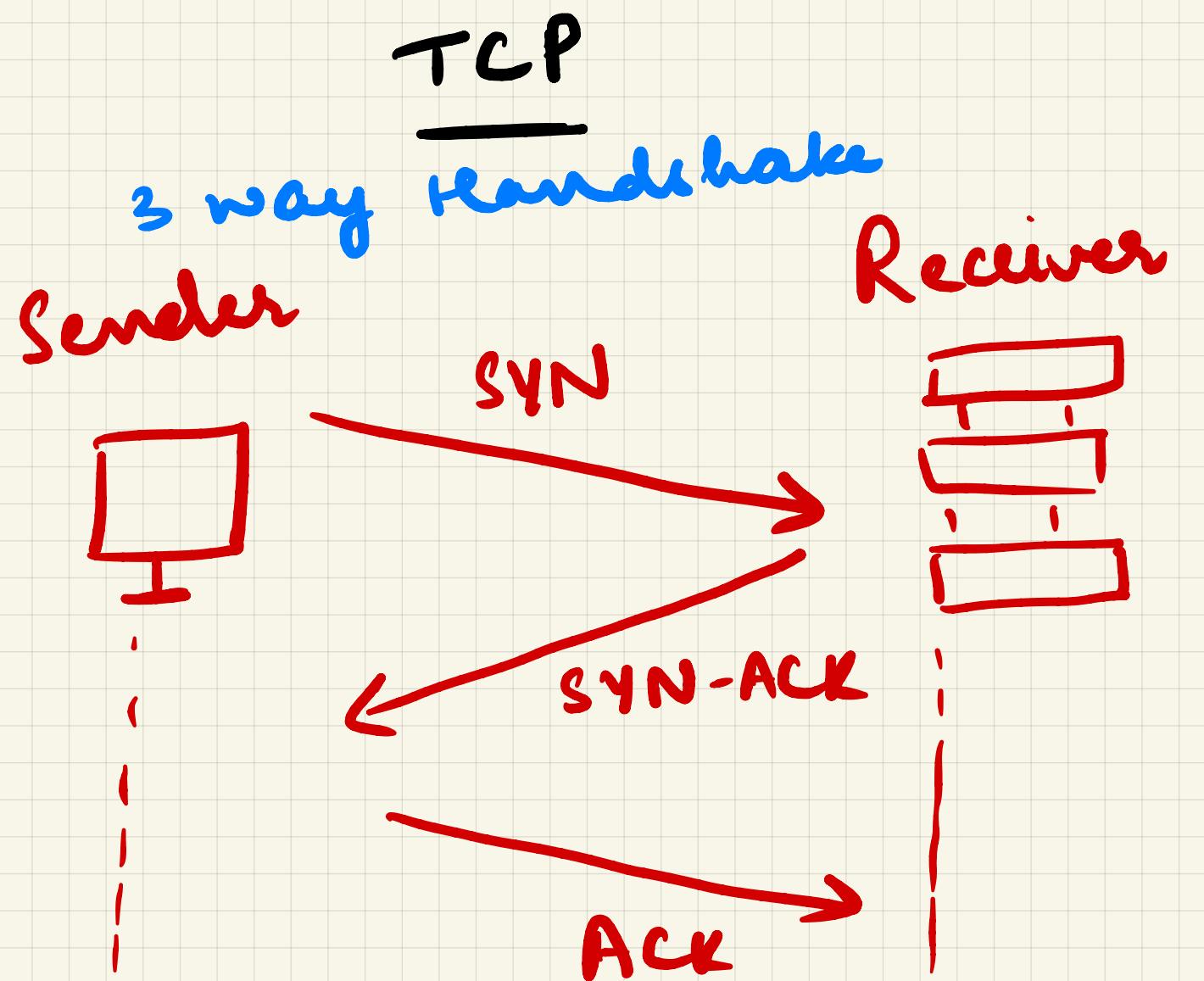
- Divide complex networking tasks into smaller, modular components, each responsible for specific set of functions.
- Network protocol stack - stack of layers
 - each layer provides services to the layer above it & uses services of layer below it → creating a hierarchical structure for n/w communication
 - Higher layers → app specific tasks
 - lower layers → lower level func closer to physical transmission of data
- OSI → open systems interconnect - 7 layers } conceptual framework
- TCP/IP → 4 layers - simplified, condensed model

Transport layer - TCP, UDP

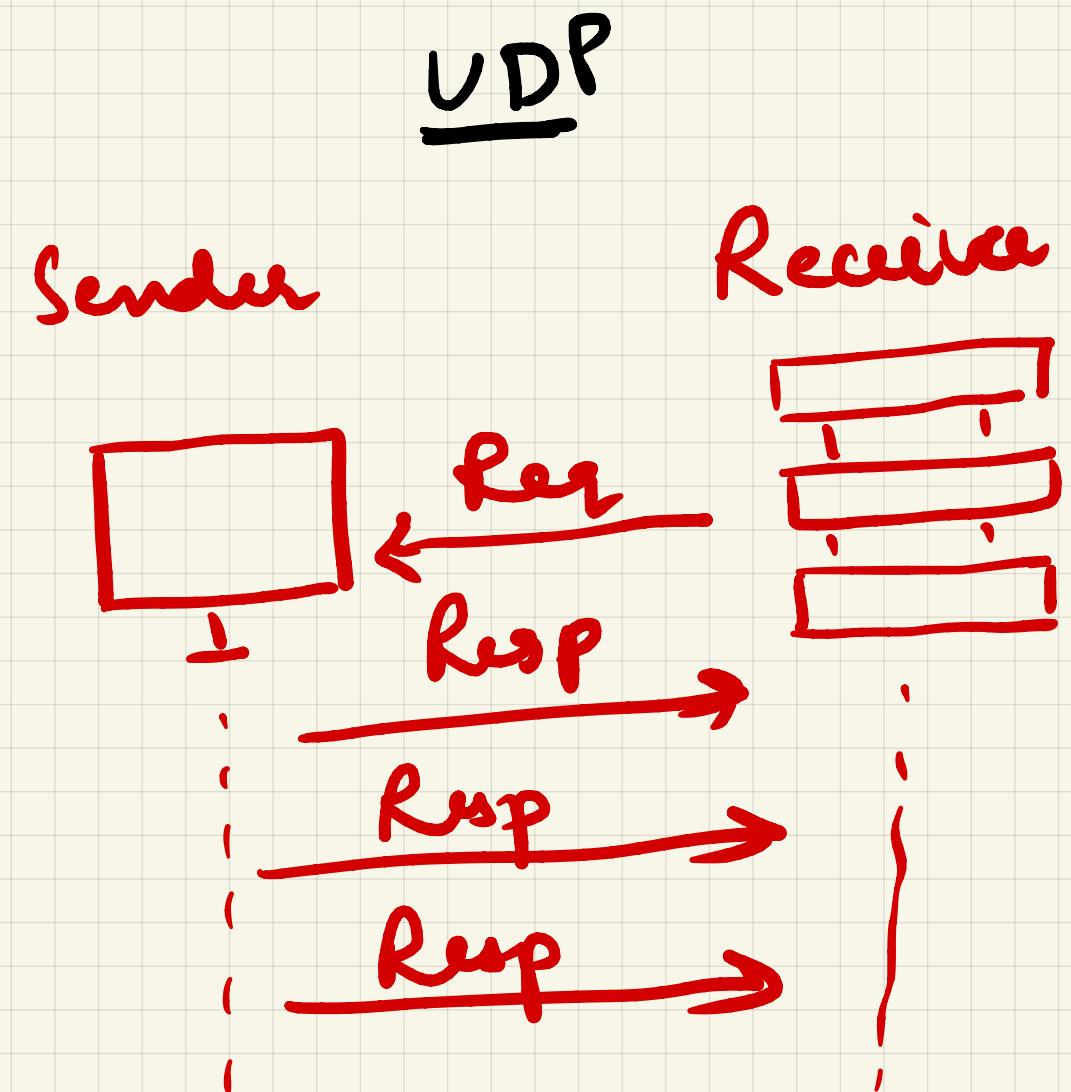
- provides reliable, end to end data delivery
- breaks down data into smaller segments, which can be transmitted over nw.

Application layer

- topmost layer in networking stack
- provides services, protocols & interfaces for end-users & apps to interact over the network.
- HTTP, HTTPS, FTP, SFTP, RTMP, SMTP, POP3, IMAP, DNS, DHCP, SNMP, SSH



- connect^{or} oriented
- protocol connect^{or}
- devices should establish connect before transmitting data & close after that



- connectionless protocol
- no overhead for opening, maintaining / terminating connect

- reliable, guarantees delivery of data
 - extensive error-checking mechanism
 - sequencing of data i.e. packets arrive in order at the receiver
 - Slower
 - Broadcasting not supported
 - retransmission of lost packets ✓
 - safe & trustworthy communication procedure
 - email, web surfing, military services
- delivery not guaranteed, unreliable
- basic error-checking using checksum
- no sequencing of data, if required - manage at app layer
- Faster, simpler
- Broadcasting supported (multicasting)
- no retransmission
- quick communication is necessary, not dependability
- game streaming, video & music streaming

3-Way TCP Handshake - Connection Establishment

- ① SYN (Synchronize) → client sends SYN packet to server
→ indicating intent to establish connection
→ packet contains seq. no. to identify the data
- ② SYN-ACK (Synchronize - Acknowledge)
→ Ack. client's req & indicating willingness to establish connect.
- ③ ACK (Acknowledgement)
→ confirm receipt of SYN-ACK
→ connect established

Reliable Data Transfer

- ① Segmentation → divide data to TCP segments
 - each segment has seq. no, control flags, acknowledgements etc.
- ② Transmission → send segments
- ③ Acknowledgment → receiver sends ACK packets to sender
 - if sender does not receive ACK within a certain time, it retransmits the segments
- ④ Flow Control → "sliding window" mechanism to manage flow control
 - sender & receiver negotiate a window size which determines no. of segments that can be sent before sending ACK, prevents sender from overwhelming the receiver with data.

Connect Termination - Graceful Connection Release

- either client or server can initiate the connection termination process by sending a FIN packet to other party, indicating intention to close the connection
- ① FIN (Finish) → Receiving party sends ACK packet to ack
 - ② ACK → Receiving party sends its own FIN packet
 - ③ FIN-ACK → Receiving party sends its own FIN packet to terminate the connection
initiating party responds with ACK
 - ④ FINAL ACK → Initiating party acknowledges the termination
 - ⑤ TIME-WAIT → Both sides enter a time-wait state to ensure that any delayed packets related to closed connection are handled.

Initiator

Established
connection

active close

FIN-WAIT-1

FIN-WAIT-2

TIME WAIT

closed

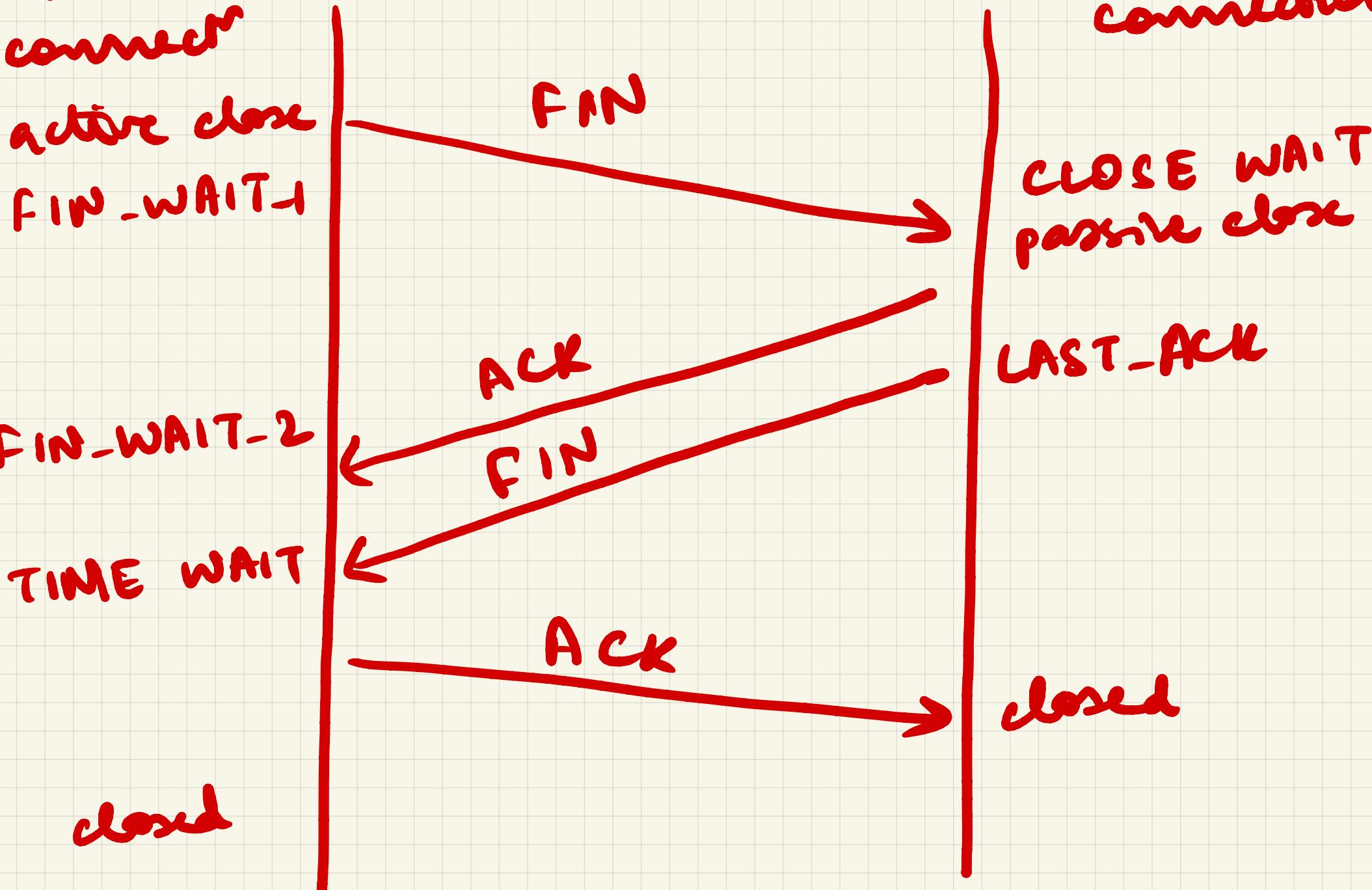
Receiver

Established
connection

CLOSE_WAIT
passive close

LAST-ACK

closed



- Why does DNS use UDP and not TCP? DNS is supposed to be reliable
- DNS msgs are v. small & fit within UDP segment -
- DNS is a widely used protocol, there is significant traffic that needs to be handled by DNS servers
- In large DNS responses i.e. when response exceeds max size of UDP packet, server may switch to TCP to transmit complete response