

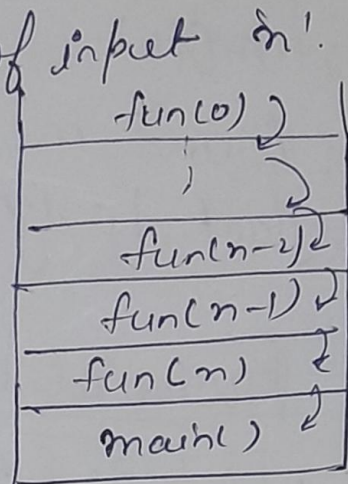
# \* Time & Space complexity of Recursive

## Solution :-

$T(n) \rightarrow$  Time Taken as a function of input  $n$ .

Recursive function maintain stack.

Base case reach hoga toh  
saare  $fun()$  ko unbind kar stack  
unbind karega.



```
void printArray(int a[], int n) {
    if (n == 0) return;
```

```
    cout << *a << " ";
    printArray(a+1, n-1);
}
```

T.C.  
Time Taken by any  
algorithm to with  
respect to its func<sup>n</sup>.

\* Formula Based Method

$$\begin{aligned} \textcircled{1} \quad F(n) &= k + F(n-1) \\ F(n-1) &= k + F(n-2) \\ F(n-2) &= k + F(n-3) \\ &\vdots \\ F(1) &= k + F(0) \\ + \quad F(0) &= k_1 \end{aligned}$$

$$\begin{aligned} T(n) &= O(nk + k_1) \\ &= O(n) \end{aligned}$$

\* Recursive Tree Method :-

$$\begin{aligned} PA(n) &\rightarrow k \\ &\downarrow \\ PA(n-1) &\rightarrow k \\ &\downarrow \\ n-2 &\rightarrow k \\ &\downarrow \\ n-3 &\rightarrow k \\ &\downarrow \\ n-2 &\rightarrow k \\ &\downarrow \\ 1 &\rightarrow k \\ &\downarrow \\ 0 &\rightarrow k_1 \end{aligned}$$

$$T(n) = nk + k$$

$$T_n = O(n)$$

$$F(n) = nk + k_1$$

$$T(n) = nk + k_1$$



Space Complexity :-

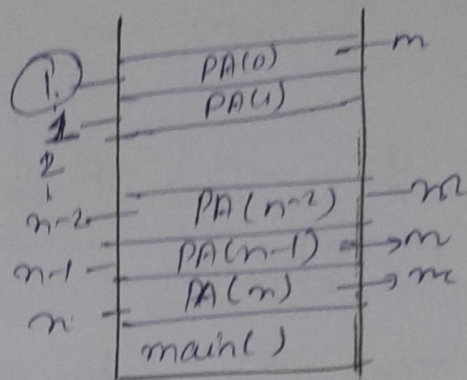
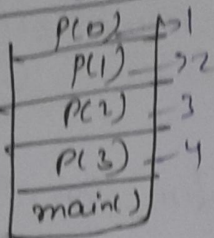
khud se kuch hum memory allocation nhi kar rhe hai DS stack me khud se manage kar rha hai

$$Sc \Rightarrow O(n+1) = O(n)$$

$n=3$

$$O(3+1) = O(4)$$

Upper bound me constant ko ignore kar dete hai



Factorial RE :-

```
int fact(int n){
    if(n==1)
        return 1;
    return n * fact(n-1);
}
```

$$T(n) = k + T(n-1)$$

$$T(n-1) = k + T(n-2)$$

$$T(n-2) = k + T(n-3)$$

$$T(1) = k$$

$$T(n) = nk$$

$$O(T_n) = O(nk)$$

$F(n) \rightarrow k$

$F(n-1)$

$F(n-2)$

$F(n-3)$

$N$

$$T(n) = nk$$

$$T = O(n)$$

Space Complexity

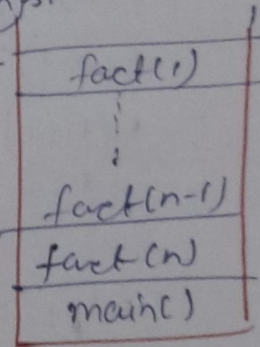
$$Sc \in O(nk)$$

constant

$$Sc = O(N)$$

memory hai har entry me space kharid rha hai

NAAM



# Binary Search RE

int BS(int a[], int k, int start, int end) {

if (start > end)  
return -1;

int mid = start + (end - start) / 2;

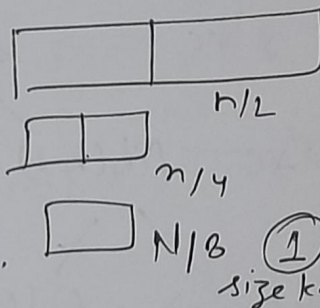
if (a[mid] == k) {  
return mid;

}  
else if (k > a[mid]) {  
return BS(a, k, mid + 1, end);

}  
else {  
return BS(a, k, start, mid - 1);

Iterative BS  
is no better  
quality than  
Recursive  
because it  
takes less  
memory.

3 kuch k processing left hai  
and fir dubara  
se call kar  
det hai n/2  
ke liye.



$$F(n) = k + F(n/2)$$

$$T(n/2) = k + T(n/4)$$

$$T(n/4) = k + T(n/8)$$

$$T(n/8) = k + T(n/16)$$

$$T(n/16) = k + T(n/32)$$

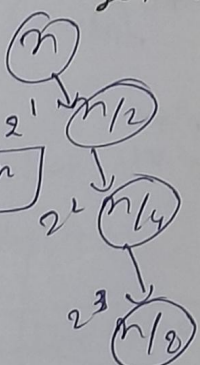
$$T(n/32) = k$$

$$\frac{n}{2^a} = 1$$

$$n = 2^a$$

$$a = \log n$$

$$T(n) = \log n$$



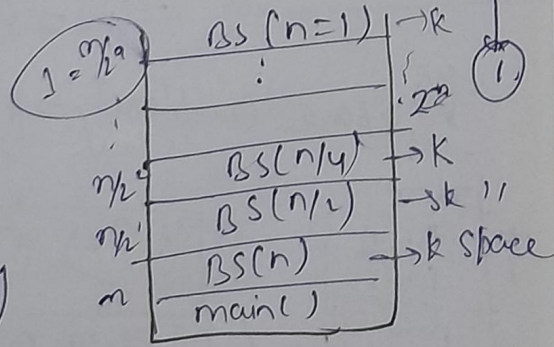
$$T(n) = a + k$$

Space-Complexity

$$a = \log n$$

$$O(Tn) = O(k \log n)$$

$$S.C \Rightarrow O(\log n)$$



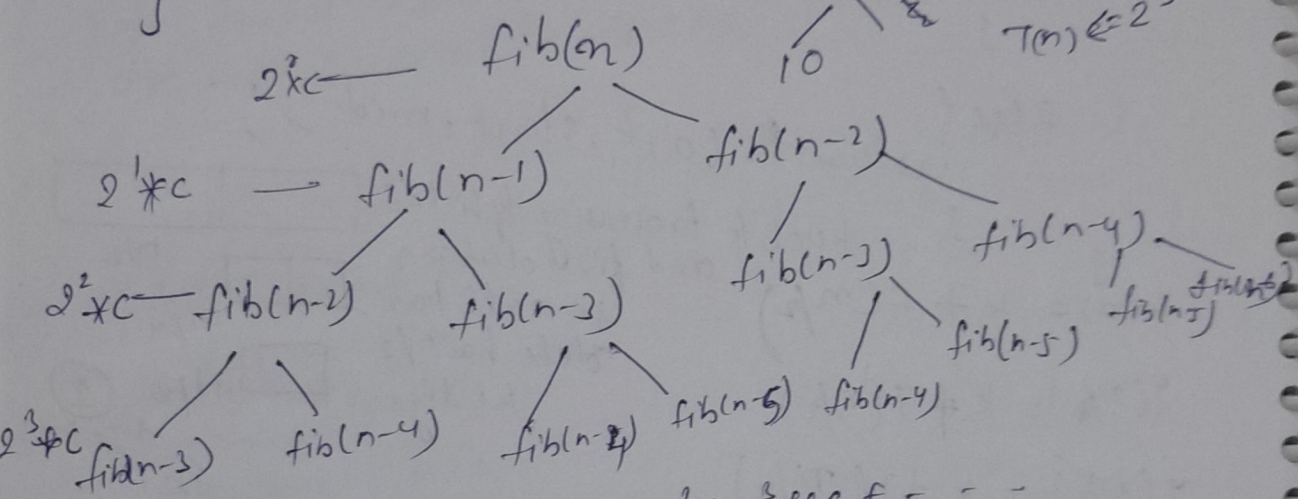
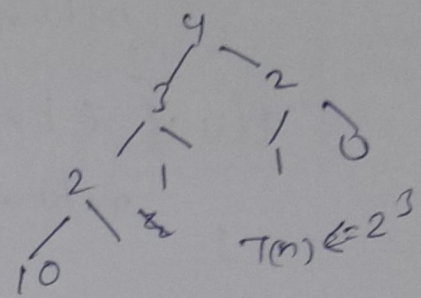


Recursive Algor ka drawback hote hai vo stack pe space le ke aa jate hai. Two method To find complexity of RE

# \* fib onacci Series RE

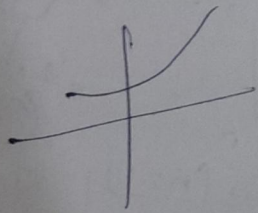
- ① Ek Visual Tree Mathematical Tarika hai
- ② Bikan Bekar Tarika hai ye

```
int fib(int n) {
    if (n == 0 || n == 1)
        return n;
    return fib(n-1) + fib(n-2);
}
```



$T(n) \leq 2^0 + 2^1 + 2^2 + 2^3 + \dots + 2^{n-1}$

$T(n) \leq [2^0 + 2^1 + 2^2 + \dots + 2^{n-1}]$



$T(n) \leq 2^n - 1$   
 $T(n) = O(2^n)$

Crazy high T.C.

\* Space Complexity  $\rightarrow$  kitna depth me jaa rha hai

fib(4)
fib(5)
main(5)

BC  $\Rightarrow O(n)$

(n == 1 or n == 0)  
 return n..

