

# RECURSION CLASS-3

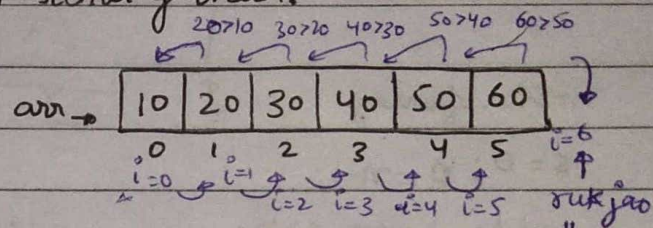
Date 13/10/23

Q check the given array is sorted or not, in ascending order.

$f(arr, size, index)$

(II) if  $arr[index] > arr[index-1]$

→ True  
→  $f(arr, size, index+1)$   
→ false  
→ return false



if  $(index \geq size)$   
→ return true

(I) main  $f(arr, size, index)$

arr size, (1) → for comparing it with previous ones

Code:-

```
bool isSorted (int arr[], int size, int index) {
```

// Base Case

```
if (index >= size) return true;
```

// Processing

```
if (arr[index] > arr[index-1]) {
```

// Aage check Krna padega ab

// Ab recursion sambhalega

```
bool aageKaAns = isSorted (arr, size, index+1);
```

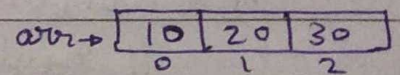
```
return aageKaAns;
```

```
} else {
```

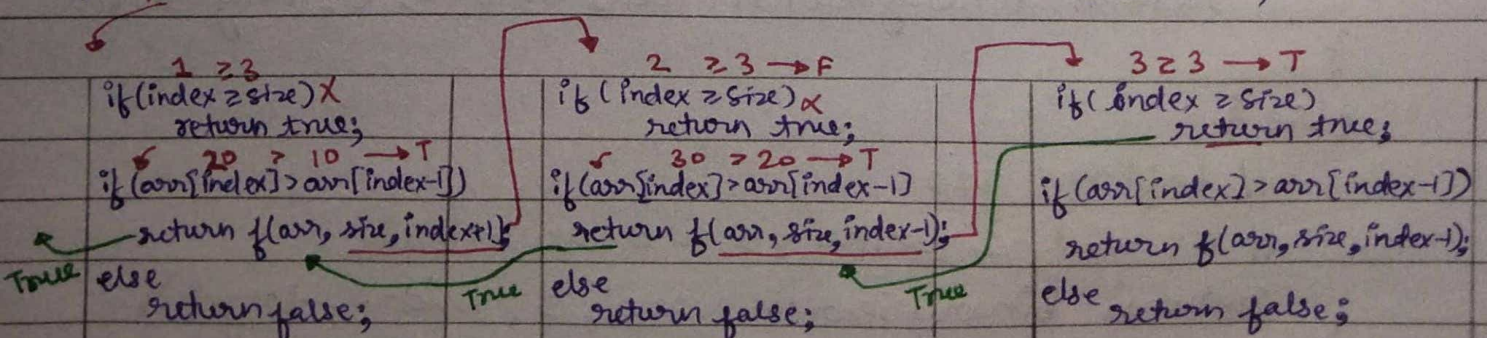
```
return false; // sorted nahi h
```

```
}
```

```
} main
```

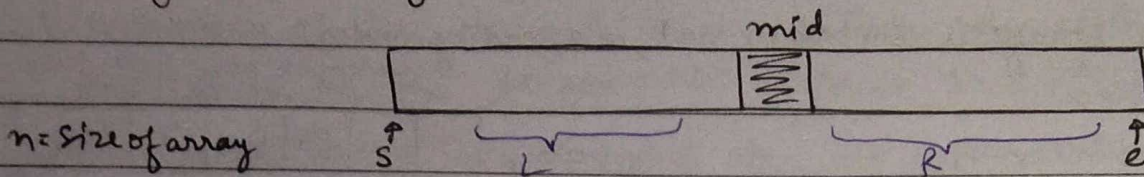


main →  $f(arr, size, index)$





## # Binary Search using Recursion:-



n = size of array

 $s = 0; e = n - 1;$  $mid = s + (e - s) / 2;$ while ( $s > e$ ) // Base caseif ( $arr[mid] == target$ ) // found wala case  
return mid;if ( $arr[mid] < target$ ) → Right me jao  
return f( $arr, \frac{mid+1}{s}, \frac{e}{e}$ )

else → Left me jao

return f( $arr, \frac{s}{s}, \frac{mid-1}{e}$ )

Base case

if ( $s > e$ )  
return -1;

Code:-

int binarySearch (int arr[], int s, int e, int target) {

// Base case

if ( $s > e$ ) return -1;

// Processing → 1 case khud solve kro

int mid =  $s + (e - s) / 2;$ if ( $arr[mid] == target$ )

return mid;

// Baaki recursion sambhal lega

if ( $arr[mid] < target$ ) {

// right me jao

//  $s = mid + 1$  karke jyanghe so we can write

return binarySearch (arr, mid+1, e, target);

}

else {

// Left me jao

//  $e = mid - 1$  karke so we can write

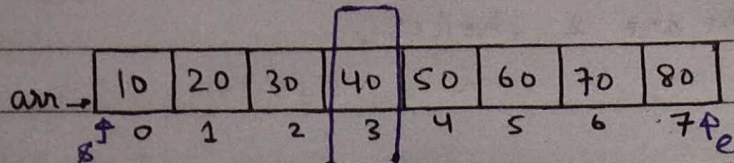
return binarySearch (arr, s, mid-1, target);

}

}



Dry Run:-



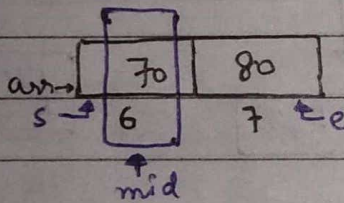
target = 80

call  
 $f(arr, 0, 7, 80)$   $\rightarrow$  6 returned  
 $f(arr, s, e, target)$   
 //B.C  $s > e \rightarrow F_x$

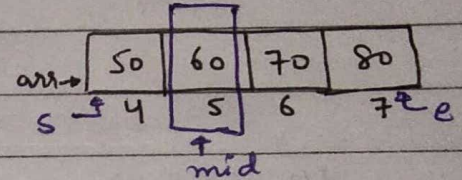
//Processing  $mid = \frac{0+7}{2} = 3$   
 $40 == 80 \rightarrow F$

//R.R if  $arr[mid] < target$   
 $40 < 80 \rightarrow True$

$\rightarrow$  return  $f(arr, mid+1, e, target)$   
 return 6



$f(arr, 4, 7, 80)$   
 //B.C  $s > e \rightarrow 4 > 7 \rightarrow F_x$   
 //Process  $mid = \frac{4+7}{2} = 5$   
 $60 == 70 \rightarrow F$



//R.R if  $arr[mid] < target$ .  
 $60 < 70 \rightarrow T$

//right me jao,  $s = mid+1$   
 $\rightarrow$  return  $f(arr, mid+1, e, target)$

return 6

$f(arr, 6, 7, 80)$

//B.C  $s > e \rightarrow 6 > 7 \rightarrow F_x$

//Processing  $mid = \frac{6+7}{2} = 6$

if  $arr[mid] == target$

$70 == 70 \rightarrow True$

return 6  $\rightarrow$  mid ki value

### # Subsequences of String :- [Include - Exclude Pattern]

From the given string, we can include or exclude any character in our output string. But ordering of characters must be same as it's in input string.

Ex:- i/p  $\rightarrow$  "abc"  $\Rightarrow$

- subsequences
- $\checkmark \times \times \rightarrow a$
  - $\times \checkmark \times \rightarrow b$
  - $\times \times \checkmark \rightarrow c$
  - $\checkmark \checkmark \times \rightarrow ab$
  - $\times \checkmark \checkmark \rightarrow bc$
  - $\checkmark \times \checkmark \rightarrow ac$
  - $\checkmark \checkmark \checkmark \rightarrow abc$
  - $\times \times \times \rightarrow ""$

i/p  $\rightarrow$  "xy"  $\Rightarrow$

- subsequences
- $\checkmark \times \rightarrow x$
  - $\times \checkmark \rightarrow y$
  - $\checkmark \checkmark \rightarrow xy$
  - $\times \times \rightarrow ""$

$\therefore n = \text{length of string} \rightarrow \text{Total Subsequences} = 2^n$



for every character, there are 2 choices.

'k'

Include Exclude

Ex:- 

a	b	c
0	1	2

$f(\text{input}, \text{output}, \text{index})$

$f(\text{"abc"}, \text{" "}, 0)$

Include

Exclude

$f(\text{"abc"}, \text{"a"}, 1)$

Include

Exclude

$f(\text{"abc"}, \text{" "}, 1)$

Include

Exclude

$f(\text{"abc"}, \text{"ab"}, 2)$

Include

$f(\text{"abc"}, \text{"a"}, 2)$

Include

Exclude

$f(\text{"abc"}, \text{"b"}, 2)$

Include

Exclude

$f(\text{"abc"}, \text{" "}, 2)$

Include

Exclude

$f(\text{"abc"}, \text{"abc"}, 3)$

Ruk jao

$f(\text{"abc"}, \text{"ac"}, 3)$

Exclude

$f(\text{"abc"}, \text{"bc"}, 3)$

$f(\text{"abc"}, \text{"c"}, 3)$

$f(\text{"abc"}, \text{" "}, 3)$

$f(\text{"abc"}, \text{"ab"}, 3)$

$f(\text{"abc"}, \text{"a"}, 3)$

$f(\text{"abc"}, \text{"b"}, 3)$

Code:-

```
void findSubSequences (string input, string output, int index) {
```

// Base Case

```
if (index >= input.length()) {
```

// ans jo hai, wo output string me built ho chuka hai

// print krdo

```
cout << output << endl;
```

```
return;
```

```
}
```

// Processing

```
char ch = input[index];
```



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// include

output.push\_back(ch);

findSubSequences(input, output, index+1);

// exclude

output.pop\_back();

findSubSequences(input, output, index+1);

}

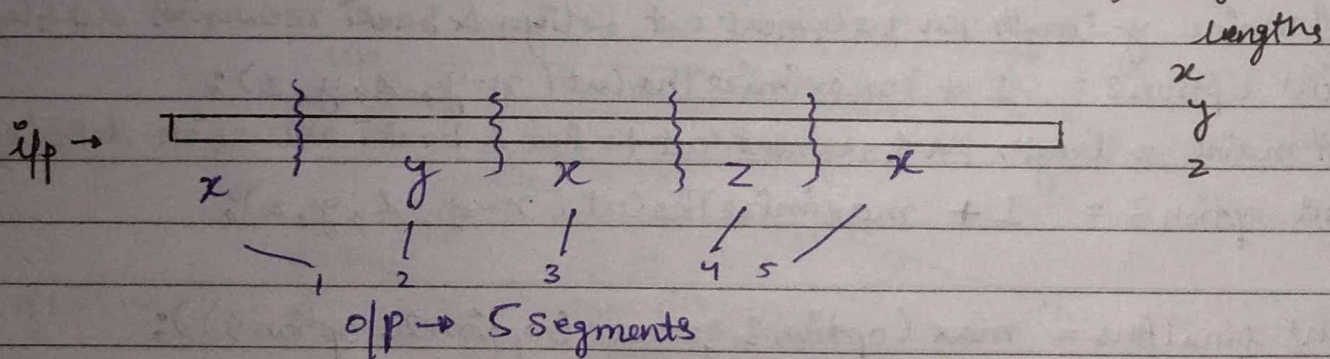
# Exploring all Possible ways Pattern

Q Cut into segments / maximize The Cut Segments (GFG)

i/p given a rod of  $n$  length. We can break it into ' $x$ ', ' $y$ ' & ' $z$ ' segments.

$x, y, z$  are integers.

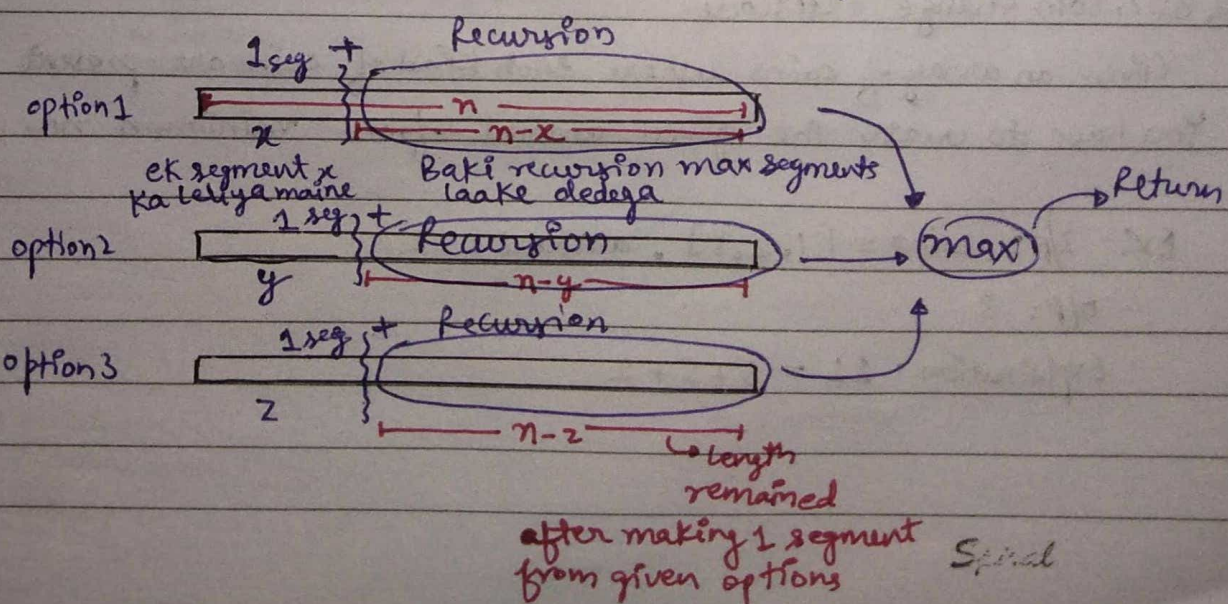
o/p  $\rightarrow$  return the maximum no. of possible segments of  $x, y$  &  $z$  lengths.



Keep in mind while applying this approach

$\rightarrow$  Exploring all possible ways pattern (EK case tum solve karo baki recursion) Samthal lega

options available  
ya to  $x$  length  
ya to  $y$  length  
ya to  $z$  length  
cut kr sktte hai





Code:-

```
int maximizeTheCuts (int n, int x, int y, int z) {
```

```
    // Base Case
```

```
    if (n == 0) {
```

```
        // Tab length hi nahi hai rod ki to ans kya return krnge so
```

```
        return 0;    // Length of rod = 0 ke liye no. of segments = 0
    }
```

```
    if (n < 0) {
```

```
        // Tab < 0 wali condition aye to max me wo consider hina ho
```

```
        return INT_MIN;
```

```
    }
```

```
    // maine x length ka 1 segment cut krliya and baaki recursion dekhlega
```

```
    int option1 = 1 + maximizeTheCuts(n-x, x, y, z);
```

```
    // maine y length ka 1 segment cut krliya & baaki recursion dekhlega
```

```
    int option2 = 1 + maximizeTheCuts(n-y, x, y, z);
```

```
    // maine z length ka 1 segment cut krliya & baaki recursion dekhlega
```

```
    int option3 = 1 + maximizeTheCuts(n-z, x, y, z);
```

```
    int finalAns = max(option1, max(option2, option3));
```

```
    return finalAns;
```

```
}
```

Q 322. Coin Change (Leetcode)

Given an array of coins where each kind of coin are present in infinite numbers. You have to create the given amount from minimum no. of coins.

Ex:- I/p: coins = [1, 2, 5], amount = 11

O/p: 3

Explanation:  $11 = 5 + 5 + 1$

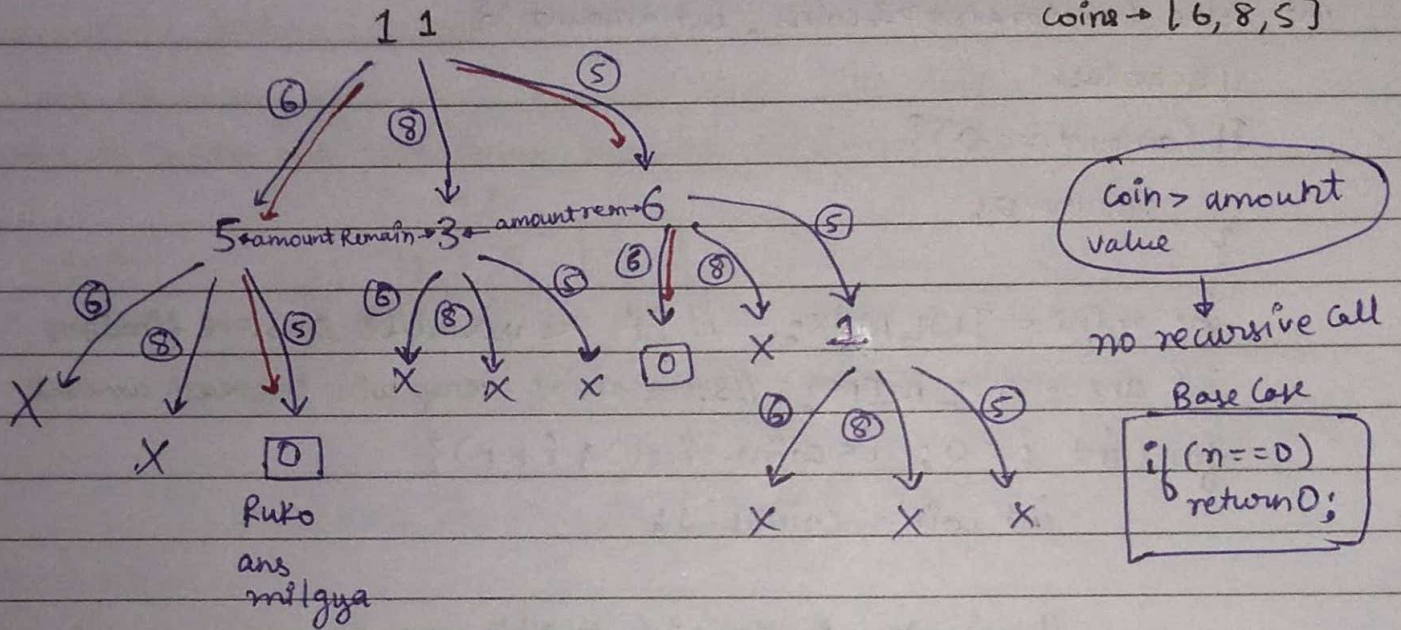


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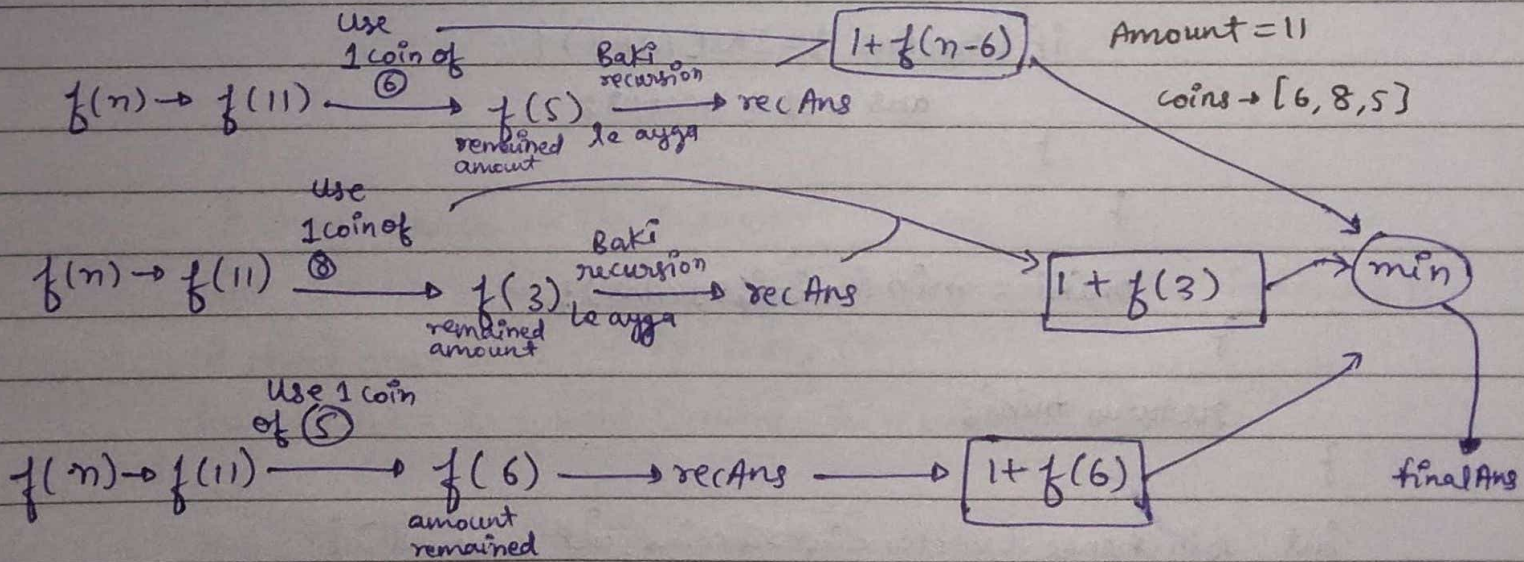
### ∴ Recursive Tree

Amount = 11

coins  $\rightarrow [6, 8, 5]$


$$5 + 6 \rightarrow 11 \quad 2 \text{ wings}$$
$$6 + 5 \rightarrow 11 \quad 2 \text{ coins}$$

min  $\rightarrow$  2 coins





Code:-

```
int solve (vector<int>&coins, int amount){
```

// Base Case

```
if (amount == 0){
```

```
    return 0;
```

```
}
```

```
int mini = INT_MAX; // It's the variable to store finalAns
```

```
int ans = INT_MAX; // Store ans of every coin to create amount
```

```
for (int i = 0; i < coins.size(); i++){
```

```
    int coin = coins[i];
```

// current coin ko sirf tabhi use krnge

// Jab uski value <= amount hogi

```
if (coin <= amount) {
```

```
    int recAns = solve (coins, amount - coin);
```

```
    if (recAns != INT_MAX) {
```

```
        ans = 1 + recAns;
```

```
    }
```

```
}
```

```
    mini = min (mini, recAns);
```

```
}
```

```
return mini;
```

```
}
```

```
int coinChange (vector<int>&coins, int amount) {
```

```
    int ans = solve (coins, amount);
```

```
    if (ans == INT_MAX)
```

```
        return -1;
```

```
    else
```

```
        return ans;
```

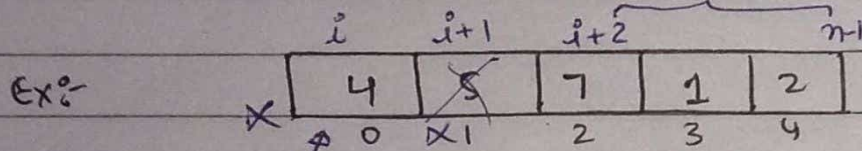
```
}
```



# Max sum of non-adjacent elements

Q 198. House Robber (leetcode)

Given an array of houses. We can only rob adjacent houses. We have to return the maximum amount which can be rob.



choori karta hu  $\rightarrow 4 + f(i+2, n-1)$  final  
max  $\rightarrow$  Ans

choori nahi karta  $\rightarrow 0 + f(i+1, n-1)$

// B.C

$i \geq \text{size} \rightarrow \text{ruk jao}$

Code:- `int solve (vector<int> &nums, int size, int index) {`

`// Base Case`

`if (index  $\geq$  size) {`

`return 0;`

`}`

`// chori karo  $\rightarrow$  ith index pr`

`int option1 = nums[index] + solve(nums, size, index+2);`

`// chori mat karo  $\rightarrow$  ith index pr`

`int option2 = 0 + solve(nums, size, index+1);`

`int finalAns = max(option1, option2);`

`return finalAns;`

`}`

`int rob (vector<int> &nums) {`

`int size = nums.size(); int index = 0;`

`int ans = solve(nums, size, index);`

`return ans;`

`}`