

RECURSION CLASS-3

Date 13.10.23...

Q check the given array is sorted or not, in ascending order.

f(arr, size, index)

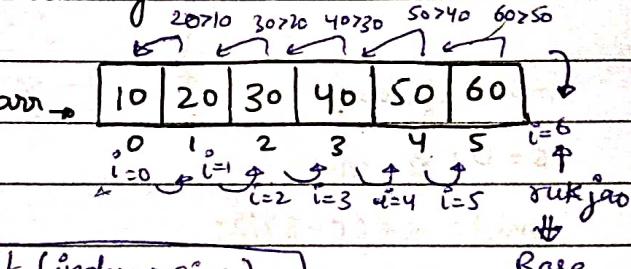
(i) if arr[index] > arr[index-1]

→ True

↳ f(arr, size, index+1)

→ false

↳ return false

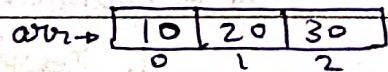


(E) main f(array, size, index)

arr[6], (1) → for comparing it with previous ones.

Code:-

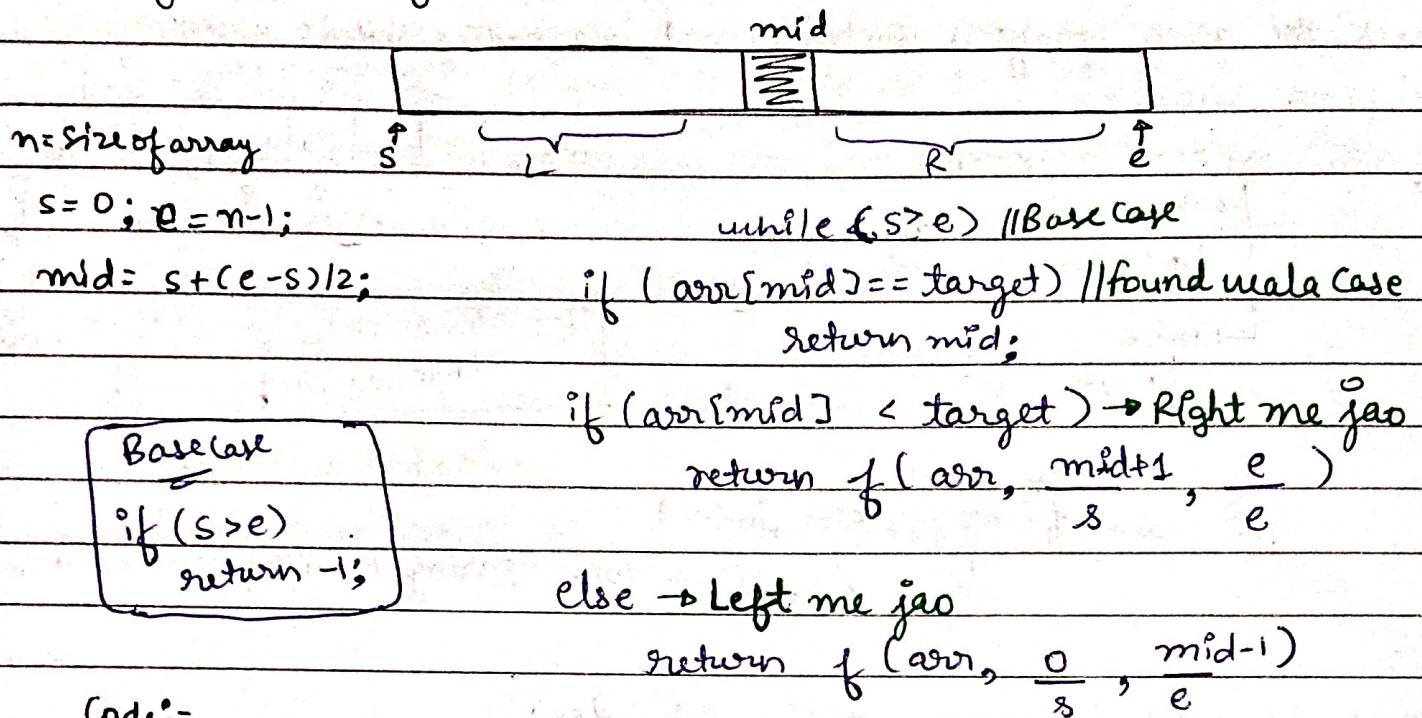
```
bool isSorted ( int arr[], int size, int index) {
    // Base Case
    if( index >= size) return true;
    // Processing
    if ( arr[index] > arr[index-1]) {
        // Aage check krna padega ab
        // Ab recursion sambhallega
        bool aageKaAns = isSorted (arr, size, index+1);
        return aageKaAns;
    } else {
        return false; // Sorted nahi h
    }
}
```



3 main

1 → 3	2 → 3 → F	3 → 3 → T
if(index >= size) X return true;	if(index >= size) X return true;	if(index >= size) X return true;
if(arr[index] > arr[index-1]) return f(arr, size, index);	if(arr[index] > arr[index-1]) return f(arr, size, index-1);	if(arr[index] > arr[index-1]) return f(arr, size, index-1);
True else return false;	True else return false;	True else return false;

Binary Search using Recursion:

Code:-

```

int binarySearch (int arr[], int s, int e, int target) {
    // Base case
    if (s > e) return -1;

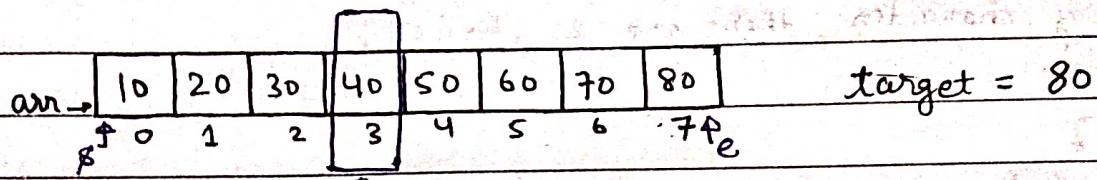
    // Processing → 1 case khud solve kro
    int mid = s + (e-s)/2;
    if (arr[mid] == target)
        return mid;

    // Baaki recursion sambhal lega
    if (arr[mid] < target) {
        // right me jao
        // s = mid + 1 karke jyange so we can write
        return binarySearch (arr, mid+1, e, target);
    }
    else {
        // Left me jao
        // e = mid - 1 karke so we can write
        return binarySearch (arr, s, mid-1, target);
    }
}

```

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Dry Run:-



(arr, 0, 7, 80) ↗ 6 returned mid
 $f(\text{arr}, s, e, \text{target})$
 $s > e \rightarrow$
 I.B.C 0 > 7 → F

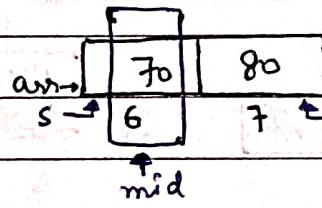
I.I Processing mid = $\frac{0+7}{2} = 3$
 $40 == 80 \rightarrow F$

I.I.R if($\text{arr}[mida] < \text{target}$)
 $40 < 80 \rightarrow \text{True}$

return 6 ↗ return $f(\text{arr}, mida+1, e, \text{target})$ I.I.R if($\text{arr}[mida] < \text{target}$).
 $60 < 70 \rightarrow T$

I.I.R right me jao, $s = mida + 1$

return $f(\text{arr}, mida+1, e, \text{target})$



f (arr, 6, 7, 80)

I.I.B.C → s > e → 6 > 7 → F

I.I Processing mid = $\frac{6+7}{2} = 6$

if ($\text{arr}[mid] == \text{target}$)

70 == 70 → True

return ⑥ → mid ki value.

Subsequences of String :- [Include - Exclude Pattern]

From the given string, we can include or exclude any character in our output string. But ordering of characters must be same as it's in input string.

Ex:- i/p → "abc" ⇒

✓ x x	→ a
x ✓ x	→ b
x x ✓	→ c
✓ ✓ x	→ ab
x ✓ ✓	→ bc
✓ x ✓	→ ac
✓ ✓ ✓	→ abc
✗ x x	→ "

i/p → "xy" ⇒

✓ x	→ x
x ✓	→ y
✓ ✓	→ xy
x x	→ "

Subsequences

∴ n = length of string → Total Subsequences = 2^n

for every character, there are 2 choices.

'K'

Include Exclude

Ex:-

a	b	c
0	1	2

$f(\text{input}, \text{output}, \text{index})$

$f(\text{"abc"}, "", 0)$

Include

Exclude

$f(\text{"abc"}, "a", 1)$

Include

Exclude

$f(\text{"abc"}, "ab", 2)$

$f(\text{"abc"}, "a", 2)$

$f(\text{"abc"}, "b", 2)$

$f(\text{"abc"}, "", 2)$

Include

Exclude

Ruk jao

$f(\text{"abc"}, "abc", 3)$

$f(\text{"abc"}, "bc", 3)$

$f(\text{"abc"}, "c", 3)$

$f(\text{"abc"}, "", 3)$

$f(\text{"abc"}, "ab", 3)$

$f(\text{"abc"}, "a", 3)$

$f(\text{"abc"}, "b", 3)$

Code :-

```
void findSubSequences (string input, string output, int index) {
```

// Base Case

```
if (index >= input.length()) {
```

// Ans jo hai, wo output string me built ho chuka hai

// print kro

```
cout << output << endl;
```

```
return;
```

}

// Processing

```
char ch = input[index];
```

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// include

output.push_back(ch);

findSubsequences(input, output, index+1);

// exclude

output.pop_back();

findSubsequences(input, output, index+1);

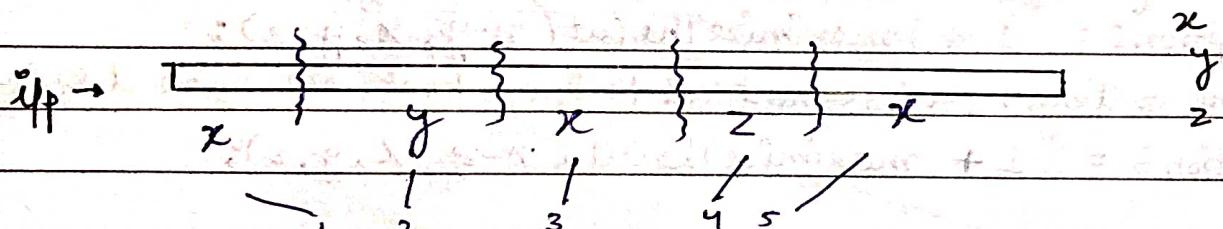
}

Exploring all Possible ways Pattern

Q Cut into segments / maximize The Cut Segments (GFG)

o/p given a rod of n length. We can break it into 'x', 'y' & 'z' segments.
x,y,z are integers.

o/p → return the maximum no. of possible segments of x, y & z lengths.



o/p → 5 segments

Keep in mind while applying this approach

→ Exploring All possible ways pattern (EK case tum solve karo baki recursion)
samjh胎 lega

options available

y₁ to x length

y₁ to y length

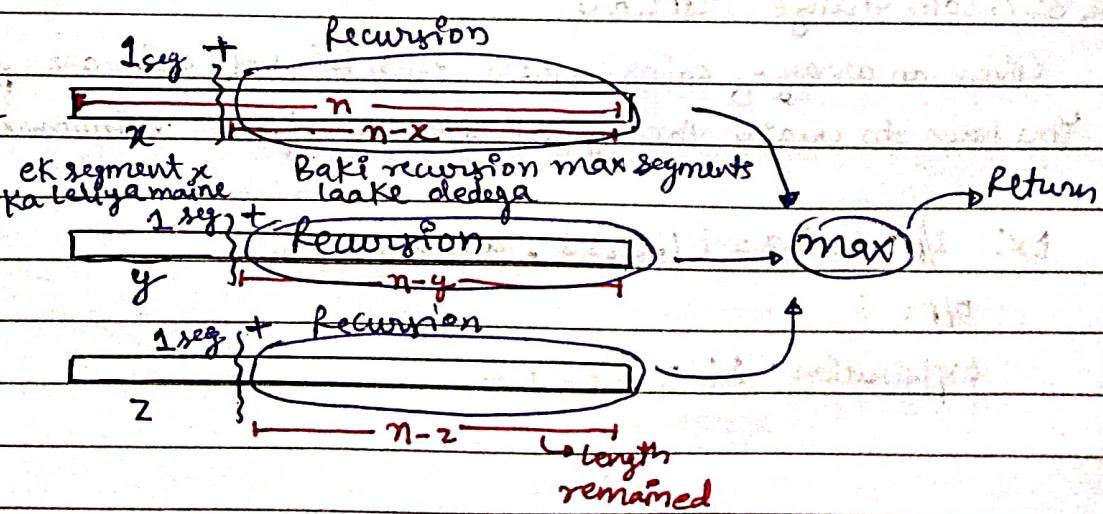
y₁ to z length

cut kr skte hain

option1

option2

option3



after making 1 segment
from given options

Special

Code:-

```

int maximizeTheCuts (int n, int x, int y, int z) {
    // Base Case
    if (n == 0) {
        // Jab length hi nahi hai rod ki to ans kya return krange go
        return 0; // Length of rod = 0 ke liye no. of segments = 0
    }
    if (n < 0) {
        // Jab < 0 wali condition aye to max me wo consider hina ho
        return INT_MIN;
    }
}

```

// maine x length ka 1 segment cut kriya and baaki recursion dekh lega

int option1 = 1 + maximizeTheCut(n-x, x, y, z);

// maine y length ka 1 segment cut kriya & baaki recursion dekh lega

int option2 = 1 + maximizeTheCut(n-y, x, y, z);

// maine z length ka 1 segment cut kriya & baaki recursion dekh lega

int option3 = 1 + maximizeTheCut(n-z, x, y, z);

int finalAns = max(option1, max(option2, option3));

return finalAns;

}

Q 322. Coin Change (Leetcode)

Given an array of coins where each kind of coin are present in infinite number.

You have to create the given amount from minimum no. of coins.

Ex:- I/p: coins = [1, 2, 5], amount = 11

O/p: 3

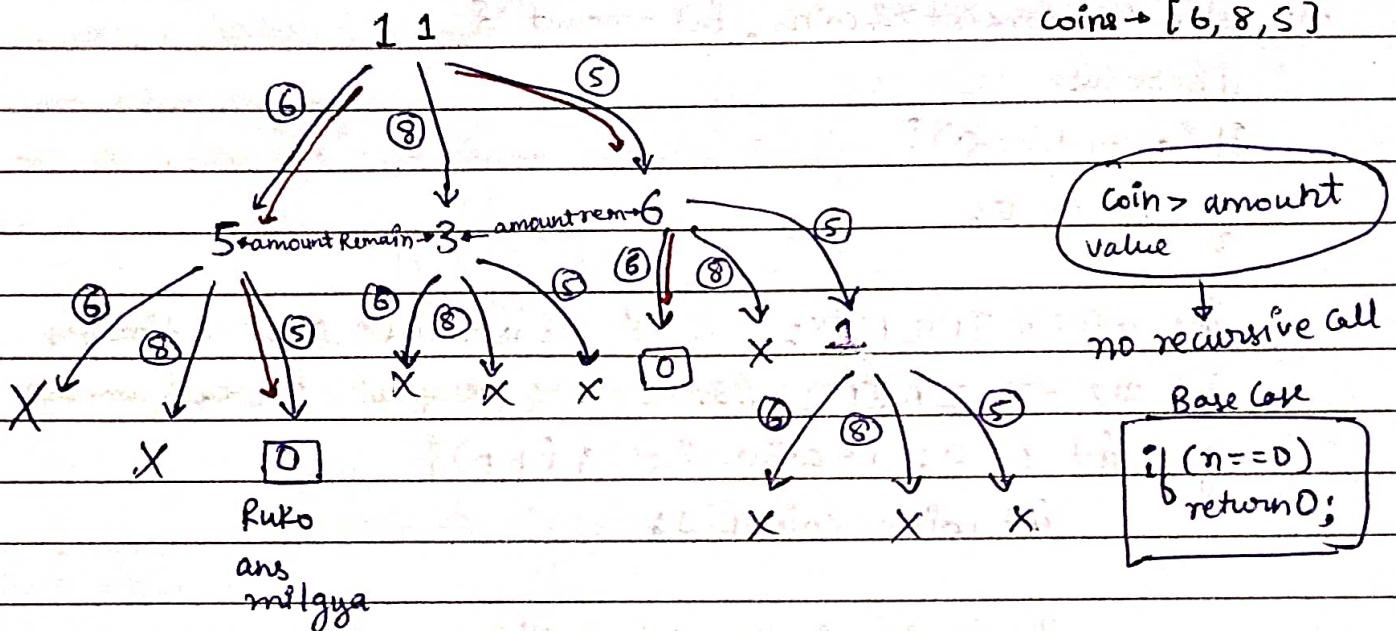
Explanation: $11 = 5 + 5 + 1$

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∴ Recursive Tree

Amount = 11

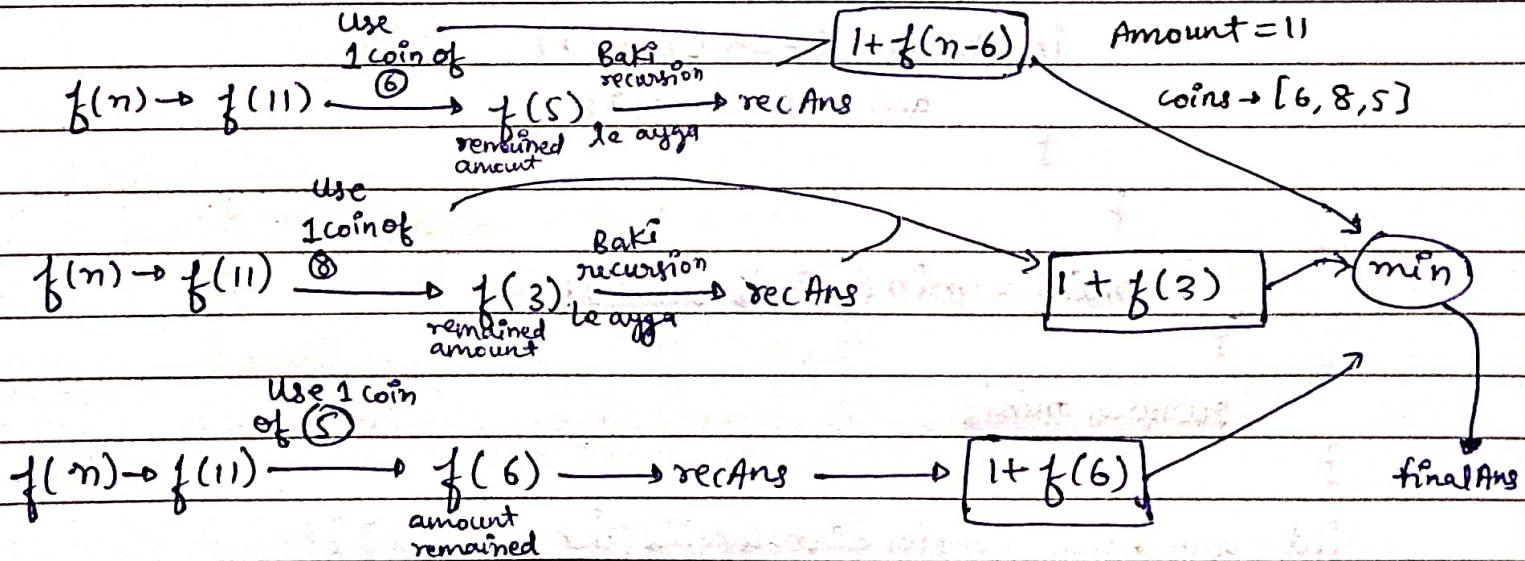
Coin → [6, 8, 5]



$$5 + 6 \rightarrow 11 \text{ 2 coins}$$

$$6 + 5 \rightarrow 11 \text{ 2 coins}$$

min → 2 coins



(Code:-)

```

int solve (vector<int>& coins, int amount) {
    // Base Case
    if (amount == 0) {
        return 0;
    }

    int mini = INT_MAX; // It's the variable to store finalAns
    int ans = INT_MAX; // Store ans of every coin to create amount
    for (int i = 0; i < coins.size(); i++) {
        int coin = coins[i];

        // current coin ko sirf tabhi use krange
        // Jab uski value <= amount hogi
        if (coin <= amount) {
            int recAns = solve (coins, amount - coin);
            if (recAns != INT_MAX) {
                ans = 1 + recAns;
            }
        }
        mini = min (mini, recAns);
    }

    return mini;
}

int coinChange (vector<int>& coins, int amount) {
    int ans = solve (coins, amount);

    if (ans == INT_MAX)
        return -1;
    else
        return ans;
}

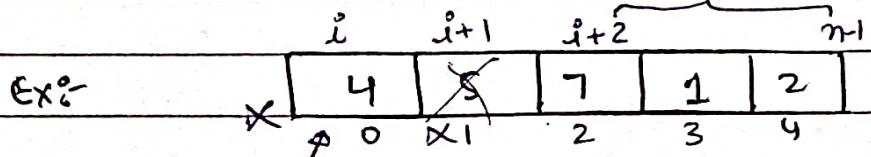
```

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Max sum of non-adjacent elements

Q. 198. House Robber (Leetcode)

Given an array of houses. We can only rob adjacent houses.
We have to return the maximum amount which can be rob.



chori Karta hu $\rightarrow 4 + f(i+2, n-1)$ final
 $\max \rightarrow \text{Ans}$

1/B.C chori nahi kri $\rightarrow 0 + f(i+1, n-1)$

$i \geq \text{size} \rightarrow \text{rukjao}$

Code:- int solve (vector<int>& nums, int size, int index) {

// Base Case

if (index \geq size) {

return 0;

}

// chori Karo \rightarrow ith index pr

int option1 = nums[index] + solve(nums, size, index+2);

// chori mat karo \rightarrow ith index pr

int option2 = 0 + solve(nums, size, index+1);

int finalAns = max(option1, option2);

return finalAns;

}

int rob (vector<int>& nums) {

int size = nums.size(); int index = 0;

int ans = solve (nums, size, index);

return ans;

}