University of Waterloo CS240 - Fall 2018 Assignment 1

Due Date: Wednesday September 19 at 5:00pm

Please read http://www.student.cs.uwaterloo.ca/~cs240/f18/guidelines.pdf for guidelines on submission. All problems 1 - 7 are written problems; submit your solutions electronically as a PDF with file name a01wp.pdf using MarkUs. We will also accept individual question files named a01q1w.pdf, a01q2w.pdf, ..., a01q7w.pdf if you wish to submit questions as you complete them.

There are 61 marks available; the assignment will be marked out of 60.

Problem 1 [3+3+3+3+3=15 marks]

Provide a complete proof of the following statements from first principles (i.e., using the original definitions of order notation). All logarithms are natural logarithms: $\log = \ln$.

a)
$$12n^3 + 11n^2 + 10 \in O(n^3)$$

d)
$$1000n \in o(n \log n)$$

b)
$$12n^3 + 11n^2 + 10 \in \Omega(n^3)$$

e)
$$n^n \in \omega(n^{20})$$

c)
$$12n^3 + 11n^2 + 10 \in \Theta(n^3)$$

Problem 2 [4+4=8 marks]

For each pair of the following functions, fill in the correct asymptotic notation among Θ , o, and ω in the statement $f(n) \in \sqcup(g(n))$. Provide a brief justification of your answers. In your justification you may use any relationship or technique that is described in class.

a)
$$f(n) = \sqrt{n}$$
 versus $g(n) = (\log n)^2$

b)
$$f(n) = n^3(5 + 2\cos 2n)$$
 versus $g(n) = 3n^2 + 4n^3 + 5n$

Problem 3 [6+6=12 marks]

Prove or disprove each of the following statements. To prove a statement, you should provide a formal proof that is based on the definitions of the order notations. To disprove a statement, you can either provide a counter example and explain it or provide a formal proof. All functions are positive functions.

a)
$$f(n) \not\in o(g(n))$$
 and $f(n) \not\in \omega(g(n)) \Rightarrow f(n) \in \Theta(g(n))$

b)
$$\min(f(n), g(n)) \in \Theta\left(\frac{f(n)g(n)}{f(n)+g(n)}\right)$$

Problem 4 [4 marks]

Derive a closed form for the following sum:

$$S(n) = \sum_{i=1}^{n} i/2^{i}.$$

Problem 5 [2+2+4+4=12 marks]

Consider the following procedure.

```
pre: n is a positive integer
pre: v[1..n] is a binary vector of length n,
     i.e., each entry is either 0 or 1
foo(v,n)
1.
     i := 1;
     while i \le n and v[i] = 0 do
3
         i := i+1
4
     od;
5.
     for j from 1 to i do
         print("Hello world!")
6.
7.
     od;
```

- a) How many inputs are there are of size n?
- b) What is the worst case number of calls to print? Give an exact formula in terms of n and justify your answer by giving an example of a worst case input of size n.
- c) For $i \in \{1, 2, ..., n\}$, let S_i denote the subset of inputs of size n for which the number of calls to print is i. Describe and enumerate S_i .
- d) What is the average case number of calls to print? Derive an exact closed form formula in terms of n.

Problem 6 [5 marks]

Prove that the following code fragment will always terminate.

```
s := 3*n // n is an integer
while (s>0)
  if (s is even)
    s := floor(s/4)
  else
    s := 2*s
```

Problem 7 [5 marks]

Analyze the following piece of pseudo-code and give a tight bound (i.e. Θ notation) on the running time as a function of n. Show your work. A formal proof is not required, but you should justify your answer.

```
1. mystery \leftarrow 0

2. \mathbf{for}\ i \leftarrow 1\ \mathbf{to}\ 3n\ \mathbf{do}

3. mystery \leftarrow mystery \times 4

4. \mathbf{for}\ j \leftarrow 1388\ \mathbf{to}\ 2010\ \mathbf{do}

5. \mathbf{for}\ k \leftarrow 4i\ \mathbf{to}\ 6i\ \mathbf{do}

6. mystery \leftarrow mystery + k
```