

Probability Analysis of Age of Information in Multi-hop Networks

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Introduction

- AGE OF information (Aol) is a metric that measures the information freshness from the perspective of the receiver monitoring a remote process. It is defined as the elapsed time since the generation of the freshest packet at the receiver.
- Aol increases until the arrival of a fresher status update at the receiver and drops upon its successful reception. Therefore, receiving status updates regularly plays a key role in sustaining information freshness.
- In addition, the staleness of a newly received update is characterized by the time it has spent in the network to reach the destination. As a result, a good Aol performance is achieved when status updates are delivered not only regularly but also timely.

Key words

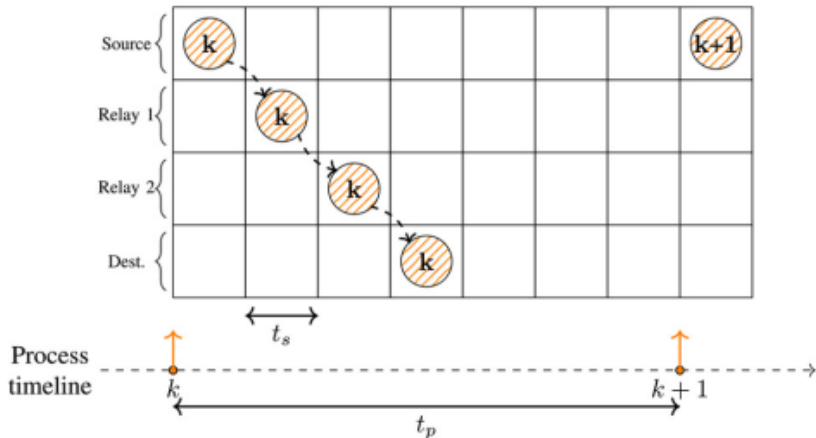
- Peak AoI
- preemptive Last Generated First Served (PLGFS)
- Sampling Event
- Sampling Period

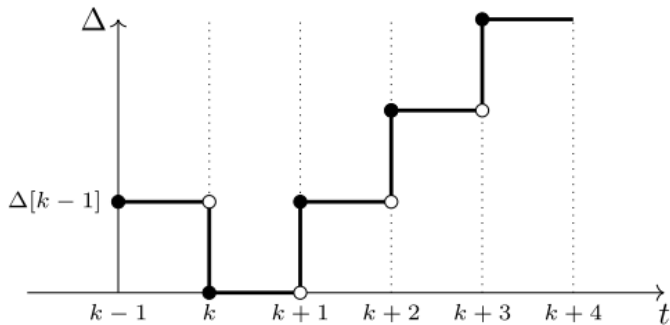
Our main motive for this work is:

- we consider a N-hop network with time-invariant packet loss probabilities on each link.
- We derive closed form equations for the probability mass function of Aol.

system model

- We consider a physical process located at the source node generating status updates periodically and a monitor located at the receiver node. The source and receiver nodes are N -hops away from each other.
- Each transmitter over the path, i.e., the source and the intermediate relay nodes, discards any older packet in the transmission queue upon the arrival of a new update.
- Each link n is prone to packet loss with time-invariant failure probability, i.e., $p_n(t) = p_n, \forall t$.
- Time is divided into slots of length t_s which is also the smallest time unit in our model. Each packet transmission starts at the beginning of a slot and completes within the same slot.





Analysis

Let $\gamma_n[k] \in \{0, 1\}$, $n \in \{1, 2, \dots, N\}$, indicate the outcome of the transmission on the n -th link in sampling period k with $Pr[\gamma_n[k] = 1] = 1 - p_n, \forall k$. Furthermore, let $\Delta_n[k]$, $n \in \{0, 1, 2, \dots, N\}$, be the AoI at hop n after the transmission slot that is allocated to the previous link.

$$\Delta_n[k] = \begin{cases} \Delta_{n-1}[k] & \gamma_n[k] = 1 \\ \Delta_n[k-1] + 1 & \gamma_n[k] = 0 \end{cases} \quad (1)$$

- we begin with a simple single hop network.

The failure probability on the first link is denoted with $p_1 \in [0, 1]$. Thus, the probability of the Aol at the relay being $\delta_1 \in \mathbb{Z} \geq 0$ can be written as:

$$Pr[\Delta_1[k] = \delta_1] = (1 - p_1)p_1^{\delta_1} \forall k \quad (2)$$

$$E[\Delta_1] = \sum_{\delta_1=0}^{\infty} Pr[\Delta_1[k] = \delta_1] \cdot \delta_1 \quad (3)$$

$$= \frac{p_1}{1 - p_1} \quad (4)$$

- Now we consider a 2 hop network
- The loss probabilities on two links are independent, we can treat each link independently.
- The contribution of the source-to-relay link to Δ_1 , is analogous to the single-hop case. However, in contrast to the source node, the information at the relay is now δ_1 periods old.

$$\Delta_{n-1}[k_n] = \delta_{n-1} \quad (5)$$

$$Pr[\Delta_2[k] = \delta_2 | \Delta_1[k_2] = \delta_1] = \begin{cases} 0 & \delta_2 < \delta_1 \\ (1 - p_2)p_2^{\delta_2 - \delta_1} & \delta_2 \geq \delta_1 \end{cases} \quad (6)$$

$$Pr[\Delta_2[k] = \delta_2] = \sum_{\delta_1=0}^{\delta_2} Pr[\Delta_2[k] = \delta_2 | \Delta_1[k_2] = \delta_1] \times Pr[\Delta_1[k_2] = \delta_1] \quad (7)$$

$$Pr[\Delta_2[k] = \delta_2] = (1 - p_1)(1 - p_2) \cdot \frac{p_2^{\delta_2+1} - p_1^{\delta_2+1}}{p_2 - p_1} \quad (8)$$

$$E[\Delta_2] = \frac{p_1}{1 - p_1} + \frac{p_2}{1 - p_2} \quad (9)$$

Age at nth link

- we extend our results to n-hop

$$\begin{aligned} &Pr[\Delta_n[k] = \delta_n] \\ &= \sum_{\delta_{n-1}=0}^{\delta_n} Pr[\Delta_n[k] = \delta_n | \Delta_{n-1}[k_n] = \delta_{n-1}] \times Pr[\Delta_{n-1}[k_n] = \delta_{n-1}], \forall k \end{aligned} \quad (10)$$

- Closed form for higher number of hops can also be obtained using a similar analysis. For instance for 3-hops we obtain:

$$Pr[\Delta_3[k] = \delta_3] = \frac{\prod_{i=1}^3 (1 - p_i)}{p_2 - p_1} \sum_{j=1}^2 (-1)^j \cdot p_j \cdot \frac{p_3^{\delta_3+1} - p_j^{\delta_3+1}}{p_3 - p_j} \quad (11)$$

Algorithm

Algorithm 1 Recursive Age Function: $f(\delta_n, n, \mathbf{p}^n) = o$

Input: δ_n age, n number of hops, \mathbf{p}^n vector of loss probabilities for n hops

Output: o the probability of age δ_n with n hops

Initialize: $o \leftarrow 0$

if $n = 1$ **then**

return $(1 - p_n) \cdot p_n^{\delta_n}$

else

for $\delta_{n-1} \in [0, \delta_n]$ **do**

$o \leftarrow o + \left((1 - p_n) \cdot p_n^{\delta_n - \delta_{n-1}} \right) \cdot f(\delta_{n-1}, n - 1, \mathbf{p}^{n-1})$

end for

return o

end if

Simulation

- Each scenario is simulated for $T_{sim} = 100000$ sampling periods and repeated 100 times
- We measure the expected Aol as:

$$E[\Delta_3] = \frac{1}{T_{sim}} \sum_{k=1}^{T_{sim}} \Delta_3[k] \quad (12)$$

- we can also use

$$E[\Delta_3] = \sum_{i=1}^3 \frac{p_i}{1 - p_i} \quad (13)$$

