

STL: C++

1. Vector :-

1. Store element at contiguous position.
2. Dynamically changing its size in most efficient way (by reserving extra more space).
3. Uses Dynamically allocated memory to store elements.
4. Vector consume more memory than array.
5. Drawback :- Worse, when you try to delete or enter value at any other position than just 8 last pos.

Accessors :-

$v = \{5, 2, 9, 8, 1\}$

- $v[i] \Rightarrow$ set or get i th element $\Rightarrow O(1) \Rightarrow 2$ for $i=1$
- $v.at(i) \Rightarrow$ " with bound checking $\Rightarrow O(1) \Rightarrow 2$ for $i=1$
- $v.size() \Rightarrow$ set current no. of element $\Rightarrow O(1) \Rightarrow 5$
- $v.empty() \Rightarrow$ set True if vector is empty $\Rightarrow O(1) \Rightarrow \text{False}$
- $v.begin() \Rightarrow$ set random access iterator to start $\Rightarrow 1$ (element)
- $v.end() \Rightarrow$ set random access iterator to end $\Rightarrow 1$ (")
- $v.front() \Rightarrow$ set the first element $\Rightarrow 5$
- $v.back() \Rightarrow$ set the last " $\Rightarrow 1$
- $v.capacity() \Rightarrow$ set the max no. of element $\Rightarrow 8$

Modifiers :-

<code>v.push_back(val)</code>	Add val to end	$O(1)$
<code>v.insert(iterator, val)</code>	Insert val at the pos entered by iterator, val can be as many as we want	$O(n)$
<code>v.pop_back()</code>		
<code>v.pop_back()</code>	remove from end	$O(1)$
<code>v.erase(iterator)</code>	remove val index by iterator	$O(n)$
<code>v.erase(begin, end)</code>	remove element from beg to end	$O(n)$

```
int main() {
```

```
    vector<int> gl;
```

```
    vector<int>::iterator i;
```

```
    vector<int>::reverse_iterator ir;
```

```
    for (int i = 1; i <= 5; i++)
```

```
        gl.push_back(i);
```

```
    for (i = gl.begin(); i != gl.end(); ++i)
```

```
        cout << *i << " ";
```

```
    // 1 2 3 4 5
```

```
    for (ir = gl.rbegin(); ir != gl.rend(); ir++)
```

```
        cout << " " << *ir;
```

```
    // 5 4 3 2 1
```

```
}
```


Stack :-

It supports LIFO

Accessors :-

(5, 2, 5, 8, 1)		
s.top()	Ret top element	O(1)
s.size()	Ret current no. of element	"
s.empty()	Ret True if stack is empty	false

Modifiers :-

s.push(val)	Push val at top	O(1)
s.pop()	Pop val from top (does not ret anything)	O(1)

```
#include <iostream>
```

```
#include <stack>
```

```
using namespace std;
```

```
int main()
```

```
{  
    stack<int> s1;
```

```
    s1.push(10);
```

```
    , (20);
```

```
    " (20);
```

```
    " (5);
```

```
    " (1);
```

```
    stack<int> g = s1;
```

```
    while (!g.empty()) { → 1 5 20 30 10
```

```
        cout << g.top() << " ";
```

```
        g.pop();
```

```
    }
```

```
    cout << s1.size() << endl; → 5
```

```
    set<> a;
```

```
}
```

Queue :-

Date _____

is a type of container adaptor, which follows FIFO (insertion at last, removal at front)

5 standard container classes list below

Accessors :-

			(5, 2, 9, 8, 1)
q.front()	Get the front element	$O(1)$	5
q.back()	" " rear "	"	1
q.size()	" the current no. of "	"	5
q.empty()	True if empty.	"	False

Modifiers :-

q.push(val):	add val to end	$O(1)$	(5, 9, 2, 8) if val=8
q.pop():	remove from front	$O(1)$	(9, 2, 8)

Deque

Date _____

1. Provide functionality similar to vector but with efficient insertion & deletion at beg or end.
2. No guarantee for contiguous storage.
3. The element of deque can be scattered in diff chunks of storage.
4. Operation that require frequent removal & insertion at pos diff than beg & end. deque performs worse than list & forward list.
5. Deque can be expanded or contracted on both end (dynamic sizes).

Accessors:-

			(5, 2, 9, 8, 1)
<code>d[i]</code>	set or get i th element	$O(1)$	$\forall i = 2$
<code>d.at(i)</code>	" with bound checking	$O(1)$	$\forall i = 2$
<code>d.size()</code>	" current no. of ele	$O(1)$	5
<code>d.empty()</code>	True if empty	"	False
<code>d.begin()</code>	iterator to start	"	$\&d.begin()$
<code>d.end()</code>	iterator to end	"	$\&d.end() - 1$
<code>d.front()</code>	set first element	"	5
<code>d.back()</code>	set last "	"	1

Modifiers:-

<code>d.push_front(val)</code>	add val to front	$O(1)$
<code>d.push_back(val)</code>	" val to back	$O(1)$
<code>d.insert(iterator, val)</code>	same as vector	$O(n)$
<code>d.pop_front()</code>	remove val from front	$O(1)$
<code>d.pop_back()</code>	" val from end	$O(1)$
<code>d.erase(iterator)</code>	" val indexed by it	$O(n)$
<code>d.erase(begin, end)</code>	same as vector	$O(n)$

List :-

Date _____

1. are sequence container allow const. time insert & erase operation anywhere within a sequence & iterate in both direction.
2. implemented as double-link-list
3. they are like forward-list just iterate in both direction.
4. used in sorting algo becoz of const time insertion, deletion & moving element.
5. Drawback: No random access by pos. also they consume extra memory for linking.

Accessors :-

(5, 2, 5, 0, 1)

<code>l.size()</code>	ret current no. of element	$O(1)$	5
<code>l.empty()</code>	true is empty	"	false
<code>l.end()</code>	ret bidirectional iterator to last	"	5
<code>l.begin()</code>	" " " to end	"	1
<code>l.front()</code>	ret first element	$O(1)$	5
<code>l.back()</code>	ret back element	$O(1)$	1

Modifiers :-

<code>l.push_front(val)</code>	} Same as previous
<code>l.push_back(val)</code>	
<code>l.insert(iterator, value)</code>	
<code>l.pop_front()</code>	
<code>l.pop_back()</code>	
<code>l.erase(it)</code>	
<code>l.erase(beg, end)</code>	
<code>l.remove(value)</code>	remove all occurrence of val, $O(n)$
<code>l.remove_if(test)</code>	remove all element that satisfy test (class or function & return 0 or 1)

Date _____

P.reverse() Reverse the list $O(n)$
P.sort() sort the list $O(n \log n)$
P.sort(comparison) sort with comparison function, take 2 element $O(n \log n)$
as i/p if set 0, means swap the items
bool comparison (int a, int b) {
 return (a > b);
} // this give reverse sorting.
P.merge() merge sorted list

bool test (const int & val) {
 return (value > 4);
} // remove val > 4

P.unique() remove duplicates.

Priority Queue :-

Date _____

1. It is a container adapter.
2. First element is greatest of all element.
3. It is based on heap.
4. container can be vector, deque
5. uses make_heap, push_heap & pop_heap;

Accessors :-

<code>q.top()</code>	returns biggest element	$O(1)$
<code>q.size()</code>	same	"
<code>q.empty()</code>	same.	"

Modifiers :-

<code>q.push(val)</code>	max-heap insertion	$O(\log n)$
<code>q.pop()</code>	remove biggest value	$O(\log n)$

Set :-

Date _____

1. It is a container that store unique element following a specific order.
2. The val in set can't be modified but they can be inserted or remove from the container.
3. element in a set are always sorted.
4. access individual element by their key.
5. They are implemented as binary search tree.

Accessors :-

s.find(key) Ret pointer pointing to s or $O(\log n)$
s.end() if not found

s.lower_bound(key) (3, 5, 8, 9) ret 8 for key = 8 $O(\log n)$
set 5 for key = 4 ✓

s.upper_bound(key) ret 8 for key = 6
set 3 for key = 8 ✓

both return a iterator

s.equal_range(key) Return pair (lower_bound(key), upper_bound(key)) $O(\log n)$
pair (int, int) (5, 8) for key = 7

s.count(key) ret 1 or 0 if $O(\log n)$
key exists or not

s.size() size of tree $O(1)$

s.empty() true or false "

s.begin() iter for first element. "

s.end() ret iter one past last element. "

element inserted in order: 5, 8, 9, 8, 3, 5
set = {3, 5, 8, 9} multiset {3, 5, 5, 8, 8, 9}



1. Modifiers:

→ does not matter what it points.

s.insert (iter, key) key is inserted in sorted order & iter to that key is returned. $O(\log n)$

s.insert (key) return a pair, an iter of key & a bool val if key is inserted in set. $O(\log n)$

s.erase (key) if present remove. $O(\log n)$

Map :-

Date _____

1. are associative container that store elements formed by a combination of a key value & a mapped value.
2. key are sorted & unique.
3. Type of key & value may differ.
4. order of insertion is maintained & slower than unordered map.
5. implemented as binary search tree.

Accessors:-

m[key] Ret val for key or add default val for key if no key is present in map (i.e 0 or null) $O(\log n)$

m.find(key) Ret a iter point to pair of (k,v) or .end() if no key found.
 $p = *(m.find(key))$
p.first or p.second $O(\log n)$

m.lower_bound(key) for map $O(\log n)$
(2,1), (3,5), (5,9), (8,4)
(5,9) for iter pointing if k=5
(2,1) not iter " k=0
(8,4) " " k=8
m.end() " " k=9

m.upper_bound(key) $O(\log n)$
(8,4) not iter if key = 5
(2,1) " if key = 0
m.end() " if key = 8

m.equal_range(key) Ret a pair containing upper & lower bound for key.

S (8, 4) for item's first } k=2
 (2, 1) for item's second
 m.end()

S. m.size()

same

O(1)

m.empty()

"

"

m.begin()

"

"

m.end()

"

S Modifiers:-

m[key] = val

O(log n)

m.insert(pair)

to insert (k, v) pair

"

m.erase(key)

if k is present remove
 key, value pair

"