# Thursday June 16, 2016 Lecture 19 Number Theory

# Notables Homework #11 Page 329, Problem 4 Page 330, Problem 16 Page 330, Problem 18 Page 357, Problem 2 Page 358, Problem 8 Due Tuesday June 21 Read Chapters 5 and 8

Public Key for Confidentiality

Scenario: A wants to send a confidential message to B

A encrypts the message using B's Public Key
B decrypts message using B's Private Key

Note that
B does not need to know A's Public Key
A does not need to have it's own Keys
Only B can read the message
Message
Is confidential
May not be authentic (i.e., Not really from A)

Public Key For Authenticity

Scenario: A wants to send an authentic message to B
A encrypts the message using its own Private Key and sends it to B.
B decrypts the message with A's Public Key
Note that
B needs to know A's Public Key
B does not need to have it's own Keys
Anyone can read the message
Message
Is not confidential
Is authentic (i.e., Really from A)

Public Key For both Confidentiality and Authenticity

Scenario: A wants to send an authentic and confidential message to B

A encrypts the message using its own Private Key and then encrypts the result using B's public key, and sends it to B.

B decrypts the message with B's private Key, and then decrypts it further using A's public key.

Note that

B needs to know A's Public Key

A needs to know B's Public Key

Message

Is confidential

Is authentic (i.e., Really from A)

# RSA (special case) Relies on laboriousness of finding prime factorization. The Public Key is just a number k which is the product of two private prime numbers The Private Key is a number which is computed using factors of k Heavy math is involved here. What is presented here is a special case of RSA code. Think of the message as a "number"

## RSA Public key system

- The *Public Key* is public to everyone!
- Sender encrypts using the Public Key
- Only receiver knows how to decrypt

# RSA (special case)

- 1. Select two different prime numbers p and q such that  $p \mod 3 = 2$  and  $q \mod 3 = 2$
- 2. Compute  $s = \frac{2(p-1)(q-1)+1}{2}$
- 3. The Public Key is k, where k = pq
- 4. The Private Key is s.
- 5. Note that RSA can be broken if we know p and q.

## RSA-640

- P = 163473364580925384844313388386509085984178 367003309231218111085238933310010450815121 2118167511579
- Q = 190087128166482211312685157393541397547189 678996851549366663853908802710380210449895 7191261465571

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RSA Example

- p = 5
- q = 11
  - □ The corresponding public key is k = 55,
  - □ The corresponding private key s = 27

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# RSA Example (confidentiality)

- My public key is k = 55
- If you to send me a secret message, say 4, encrypt it using my public key as follows:

$$C = (4**3) \% 55 = 9$$

and then you send message 9 to me

 I will find your secret message by decrypting your message using my private key as follows:

$$T = (9**27) \% 55 = 4$$

- Note how big of a number this is:
  - 9\*\*27 = 58149737003040059690390169

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# RSA Example (authenticity)

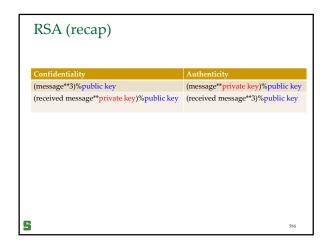
- My public key is k = 55
- If I want to send you a message, say 4, that you can be certain it came from me, I encrypt it using my private key as follows:

$$C = (4**27) \% 55 = 49$$

and then send you message 49

To make sure message 49 came from me, you (or any body) can decrypt the message using my public key as follows:

$$T = (49**3) \% 55 = 4$$





- 6nfX01TUfFaliu1wit5RJ5JQNFBzxWSePsviIml PKReIFSjpktWW6RbGk4pNj+fqh2DOWquaMz dXI27YFVuFJQ==
- This is a number in base 64, using the following symbols
  - □ 0-25 is 'A'-'Z'
  - □ 26-51 is 'a'-'z'
  - □ 52-61 is '0'-'9'
  - □ 62 is '+'
  - □ 63 is '/'
  - □ Pad is '='
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# Induction and Recursion Chapter 5

# Example

- Consider the predicate
  - $P(n) = 11 + 3 + \dots + (2n-1)$  is equal  $n^2$
- Let's verify the truth value of P(n) for some n:
  - $P(1) = 1 = 1 = 1^2$
- $P(2) = 1 + 3 = 4 = 2^2$
- $P(3) = 1 + 3 + 5 = 9 = 3^2$
- · ....
- $P(9) = 1 + 3 + 5 + 7 + 9 + 11 + 13 + 15 + 17 = 9^{2}$
- ο.
- How can we show that  $\forall nP(n)$  is true?

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### Example... (no calculators!)

- Suppose, without verifying, that
  - $P(21) = 1 + 3 + 5 + ... + 39 + 41 = 21^2$
- Given above, can we prove that
- $P(22) = 1 + 3 + 5 + \dots + 39 + 41 + 43 = 22^{2}$ ?
- $P(22) = 1 + 3 + 5 + \dots + 39 + 41 + 43 = 22^{2}?$

 $+43 = 22^2$ ?

- P(22) = P(21)
- $P(22) = 21^2 + 43 = 22^2 ?$   $P(22) = (22 1)^2 + (22 \times 2 1) = 22^2 ?$
- $P(22) = 22^2 22 \times 2 + 1 + 22 \times 2 1 = 22^2 \checkmark$

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### Example (no calculators!)

- How general is our method?
- Let's replace 21 with a place holder *k*.
- Suppose, without verifying, that  $P(k) = 1 + 3 + 5 + ... + (2k-1) = k^2$
- Given this, can we prove that

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P(k+1) = 1 + 3 + 5 + \dots + (2(k+1)-1) = (k+1)^2?
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- $P(k+1) = 1 + 3 + 5 + \dots + (2k-1) + (2(k+1)-1) = (k+1)^{2}?$
- $P(k+1) = P(k) + (2(k+1)-1) = (k+1)^{2}?$
- $P(k+1) = k^2 + (2(k+1)-1) = (k+1)^2?$
- $P(k+1) = ((k+1)-1)^2 + (2(k+1)-1) = (k+1)^2?$
- $P(k+1) = (k+1)^2 2(k+1) + 1 + 2(k+1) 1 = (k+1)^2 \checkmark$

# Example...

- What have we accomplished so far?
  - □ We have shown that  $P(k) \Rightarrow P(k+1)$ , for any k. For instance:
  - $P(1) \Rightarrow P(2),$
  - $P(2) \Rightarrow P(3),$
  - $P(6) \Rightarrow P(7),$
  - · ...
  - $P(19) \Rightarrow P(20),$
  - $P(2000) \Rightarrow P(2001), \dots$
  - □ Note that, we do NOT know if, for example, *P*(6) or *P*(19) ... is true. All we know is that IF, for example, *P*(6) is true, then *P*(7) is also true.
- What do we need to show that P(n) is true for all n?
- $\Box$  All we need is to PROVE that P(1) is indeed true.

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## Mathematical Induction

- Let P(n) be some propositional function involving integer n.
  - P(n) = n (n+3) is an even number "
  - $P(n) = 1 + 3 + \dots + (2n-1) = n^2$
  - · .....
- To prove that *P*(*n*) is true for all positive integers *n*, we can do the following:
  - □ Give a proof (usually a straight verification) that *P*(1) is true
  - Give a proof that for an <u>arbitrary</u> k, **IF** P(k) is true THEN P(k+1) is true. That is, we validate the logical implication:  $P(k) \Rightarrow P(k+1)$

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