

## Tuesday June 21, 2016 Lecture 21

Solving recurrence relations



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### Notables

- Homework #12
  - Page 525, Problem 12
  - Page 525, Problem 14
  - Find the solution to  $a_n = 2a_{n-1} + 2n^2$  with  $a_1 = 4$ .
  - Page 396, Problems 2, 4, 6, 12, 36
  - Page 413, Problem 6
  - Page 414, Problem 22
  - Due Thursday June 23
- Read Chapter 6

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## Solving Recurrence Relations

Section 8.2



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### Example

- The *Fibonacci* numbers,  $f_0, f_1, f_2, \dots$ , are defined recursively by:
  - $f_0 = 0$ ,
  - $f_1 = 1$ ,
  - $f_n = f_{n-1} + f_{n-2} \quad n \geq 2$ .
- Suppose we want to find  $f_6$ 
  - $f_2 = f_1 + f_0 = 1 + 0 = 1$
  - $f_3 = f_2 + f_1 = 1 + 1 = 2$
  - $f_4 = f_3 + f_2 = 2 + 1 = 3$
  - $f_5 = f_4 + f_3 = 3 + 2 = 5$
  - $f_6 = f_5 + f_4 = 5 + 3 = 8$
- Is there a better way?
  - Can we come up with a closed-form formula for  $f_n$ ?

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### Recurrence Relations

Consider the following equation involving a recursive function  $f_n$

$$c_0 f_n + c_1 f_{n-1} + c_2 f_{n-2} + \dots + c_r f_{n-r} = g(n)$$

where each  $c_i$  is a constant. Clearly if the value of  $r$  consecutive  $f$ 's in the sequence  $f_{k-r}, f_{k-r+1}, \dots, f_{k-1}$  are known for some  $k$ , then the value of  $f_k$  can be computed. The goal is to find a closed form solution for  $f_n$ . We will study how to solve such recurrence relations, in particular, when  $g(n) = 0$ . Example:

$$\begin{cases} f_n - f_{n-1} - f_{n-2} = 0 & \text{for } n \geq 2 \\ f_0 = 0, f_1 = 1 \end{cases}$$

Note that in the above example,  $c_0 = 1, c_1 = -1$ , and  $c_2 = -1$ .

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### Solving Homogeneous Recurrence Relations

Consider the following homogeneous recurrence relation

$$c_0 f_n + c_1 f_{n-1} + c_2 f_{n-2} + \dots + c_r f_{n-r} = 0$$

It is known that the solution to the above equation has the generic form  $\alpha^n$ .

We proceed as follows:

- (1) Construct the **characteristic equation**  $c_0 \alpha^r + c_1 \alpha^{r-1} + c_2 \alpha^{r-2} + \dots + c_{r-1} \alpha + c_r = 0$ .
- (2) Find all the roots of the characteristic equation, call them  $\alpha_1, \alpha_2, \dots, \alpha_r$ .
- (3) If all the roots are distinct, then  $f_n = A_1 \alpha_1^n + A_2 \alpha_2^n + \dots + A_r \alpha_r^n$ . We will then use the initial conditions to find  $A_1, A_2, \dots, A_r$ .
- (4) If all the roots are not distinct we do the following. Suppose a root, say  $\alpha_1$ , is a  $k$ -multiple root. Then the corresponding term for  $\alpha_1$  becomes:

$$(A_1 n^{k-1} + A_2 n^{k-2} + \dots + A_{k-2} n^2 + A_{k-1} n + A_k) \alpha_1^n$$

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### Example

Solve the following recurrence relation:

$$\begin{cases} f_n - f_{n-1} - f_{n-2} = 0 & \text{for } n \geq 2 \\ f_0 = 0, f_1 = 1 \end{cases}$$

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Solve the following recurrence relation:

$$\begin{cases} f_n - f_{n-1} - f_{n-2} = 0 & \text{for } n \geq 2 \\ f_0 = 0, f_1 = 1 \end{cases}$$

Solution: The characteristic equation is  $\alpha^2 - \alpha - 1 = 0$  which has the distinct

$$\text{roots } \alpha_1 = \frac{1+\sqrt{5}}{2}, \text{ and } \alpha_2 = \frac{1-\sqrt{5}}{2}. \text{ So, } f_n = A_1 \left( \frac{1+\sqrt{5}}{2} \right)^n + A_2 \left( \frac{1-\sqrt{5}}{2} \right)^n.$$

We now use the initial conditions to find A's

$$\begin{cases} f_0 = A_1 \left( \frac{1+\sqrt{5}}{2} \right)^0 + A_2 \left( \frac{1-\sqrt{5}}{2} \right)^0 = A_1 + A_2 = 0. \\ f_1 = A_1 \left( \frac{1+\sqrt{5}}{2} \right)^1 + A_2 \left( \frac{1-\sqrt{5}}{2} \right)^1 = 1 \end{cases}$$

These equations would give  $A_1 = \frac{1}{\sqrt{5}}$ , and  $A_2 = -\frac{1}{\sqrt{5}}$ . Thus,

$$f_n = \frac{1}{\sqrt{5}} \left( \frac{1+\sqrt{5}}{2} \right)^n - \frac{1}{\sqrt{5}} \left( \frac{1-\sqrt{5}}{2} \right)^n$$

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Solve the following recurrence relation.

$$\begin{cases} f_n + 6f_{n-1} + 12f_{n-2} + 8f_{n-3} = 0 & \text{for } n \geq 3 \\ f_0 = 1 \\ f_1 = -2 \\ f_2 = 8 \end{cases}$$

Solution: The characteristic equation is:

$$\alpha^3 + 6\alpha^2 + 12\alpha + 8 = (\alpha + 2)^3 = 0$$

Its triple root is  $-2$ .

So,  $f_n = (A_1 n^2 + A_2 n + A_3)(-2)^n$ . Using the initial conditions

we will find  $A_1 = \frac{1}{2}$ ,  $A_2 = -\frac{1}{2}$ , and  $A_3 = 1$ .

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### Example

Recall that

$$S_n = 1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$$

Let's formulate the above as a recurrence relation:

$$\begin{cases} S_n = S_{n-1} + n & \text{for } n \geq 1 \\ S_0 = 0 \end{cases}$$

Can we solve for  $S_n$  using the method we just learned?

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Solve the following recurrence relation:

$$\begin{cases} S_n = S_{n-1} + n & \text{for } n \geq 1 \\ S_0 = 0 \end{cases}$$

**Solution:**

$$S_n - S_{n-1} = n$$

$$\frac{S_{n+1} - S_n}{S_n - S_{n-1}} = \frac{n+1}{n}$$

$$\frac{S_{n+1} - 2S_n + S_{n-1}}{S_{n+1} - 2S_n + S_{n-1}} = 1$$

$$\frac{S_{n+2} - 2S_{n+1} + S_n}{S_{n+2} - 2S_{n+1} + S_n} = 1$$

$$S_{n+2} - 3S_{n+1} + 3S_n - S_{n-1} = 0$$

Characteristic equation is:  $\alpha^3 - 3\alpha^2 + 3\alpha - 1 = 0 \Rightarrow (\alpha - 1)^3 = 0$ .

So, we have a triple root,  $\alpha = 1$ . Thus,  $S_n = (A_2 n^2 + A_1 n + A_0)(1)^n$

Using the initial conditions we will find:  $A_2 = A_1 = \frac{1}{2}$ , and  $A_0 = 0$ .

$$S_n = \frac{1}{2} n^2 + \frac{1}{2} n = \frac{n(n+1)}{2}$$

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### Thinking Recursively; Applications of recurrence relations

- Tower of Hanoi
- Binary Search
- Gossiping Problem
- Quick Sort

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## Tower of Hanoi



- Let  $T_n$  be the minimum number of moves that will transfer  $n$  disks from one peg to another, according to the game rules. We have:
  - $T_1 = 1$
  - $T_2 = 3$
  - ...
  - $T_n = 2 T_{n-1} + 1$

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## Binary Search

- Given an ordered list of items, the objective is to determine if the list contains a given item:
  - List: 3, 7, 8, 10, 14, 18, 22, 34
  - Given item: 25
- Algorithms:
  - Linear Search
  - Binary Search
- Let  $C(n)$  be the number of two-item comparisons required to determine if the list contains the given item. If we do it via binary search, we have (worst case scenario) :
  - $C(1) = 1$
  - $C(n) = 1 + C(n/2)$

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## Gossiping Problem

- Consider a set of  $n$  people, where each person has a gossip to share with all others as follows. This is done by placing person-to-person "calls" whereby the two parties share all their gossips. What is the least number of calls, denoted  $G(n)$ , required so that everyone knows all the gossips?
  - $G(1) = 0$
  - $G(2) = 1$
  - $G(3) = 3$
  - ...
  - $G(n) \leq 1 + G(n-1) + 1$

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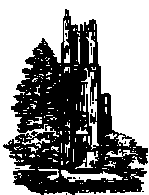
## Merge sort analysis

- Let  $M(n)$  be the cost of doing a merge sort on  $n$  numbers.
- $M(n) = 2M\left(\frac{n}{2}\right) + n, \quad M(1) = 0$
- Verify that  $f(n) = n \log n$  is indeed a solution

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## The Basics of Counting



Read Chapter 6  
Sections 6.1 – 6.4

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## Sample Counting Problems

- How many different passwords are there of length 5?
- How many different IP addresses are there for use on the Internet?
- How many one-to-one functions are there from a set of  $m$  elements to one with  $n \geq m$  elements?
- In how many ways can we select a person from a group of 35 students and 24 faculty?
- How many non-negative integer solutions are there to the equation:  $a + b + c + d = 17$  ?
- What size of a group does it take to have a better than 50% chance of two people having the same birthday?

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## Counting Principles

- The **Product Rule**:
  - If a procedure can be carried out by performing tasks  $T_1, T_2, \dots, T_n$  in sequence, where task  $T_i$  can be done in  $k_i$  ways after tasks  $T_1, T_2, \dots, T_{i-1}$  are done, then there are  $k_1 \times k_2 \times \dots \times k_n$  ways to carry out the procedure.
- The **Sum Rule**:
  - If there are  $n$  independent events  $E_1, E_2, \dots, E_n$  which can occur in  $k_1, k_2, \dots, k_n$  ways, respectively, then there are  $k_1 + k_2 + \dots + k_n$  ways in which *exactly one* of these events can occur.



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## Selection/Arrangement of Objects

- Many counting problems involve *selection* and/or *arrangement* of objects. Issues in solving such problems include:
  - Are objects **unique**?
  - Does the order **matter**?
  - Is **replacement** allowed?



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## Combinations & Permutations

- **Definition**: An  $r$ -combination of  $n$  objects is an *unordered* selection of  $r$  of these objects where **replacement** is **not** allowed. If the objects are **unique**, then the  $r$ -combination is just an  $r$ -element subset of the set of objects.
  - Example: From the set  $S = \{1, 2, 3, 4\}$  we can choose two elements in six different ways, namely,  $\{1, 2\}, \{1, 3\}, \{1, 4\}, \{2, 3\}, \{2, 4\}, \{3, 4\}$
- The notation  $C(n, r)$  denotes the number of  $r$ -combinations of  $n$  *distinct* objects.
  - Alternative notation is  $\binom{n}{r}$



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