

PSA Class - 5

Recursion

Need for recursion

Base Condition: knows how to solve the problem

Other condition: Invokes other conditions & base condition
i.e. Breaking down the problem

Ex: Factorial: (Iterative)

Factorial using iteration

```
int F(int n) {  
    int s = 1  
    for (int i = 2; i <= n; i++)  
        s *= i  
    }  
    return s;  
}
```

→ here it works for Python, but in C/C++ int has limited space so we need to have ^{large} infinite space to accommodate large n

n	s	i
5	1	1
5	2	2
5	6	3
5	24	4
5	120	5

Time complexity = $\Theta(n)$
Space complexity = $\Theta(1)$

Factorial : (Recursive)

Factorial using Recursion

```
int F(int n) {
    if (n < 2) {
        return 1;
    }
    return n * F(n-1);
}
```

Stack

return

n = 1	120 * 1
n = 2	120 * F(2-1)
n = 3	60 * F(3-1)
n = 4	24 * F(4-1)
n = 5	5 * F(5-1)

Time Complexity = $\Theta(n)$

Space Complexity = $\Theta(n)$

Stack overflow
can occur

$$T(n) = T(n-1) + C$$

n
↓

→ C

$$T(1) = 1$$

$$T(0) = 0$$

n=1

→ C

↓

n=2

→ C

↓

⋮

↓

n-(n-1) → C

$$\text{Time Complexity} = nC$$

$$= \Theta(n)$$

Print digits of a number in reverse order iteration

```
void P(int n) {
    if (n == 0) : print(n)
    while (n) {
        print(n % 10)
        n = n / 10
    }
}
```

n = 1986

print

6

n = 198

8

n = 19

9

n = 1

1

Time complexity = $\Theta(\log_{10} n)$
 Space complexity = $\Theta(1)$

Print digits of a number in reverse order recursion

```
void R(int n) {
    if (n < 10) { print(n); return; }
    print(n % 10)
    return R(n / 10)
}
```

n get out stack

n get out	stack
1	return
19	R(1)
198	R(19)
1986	R(198)

Time complexity = $\Theta(\log_{10} n)$
 Space complexity = $\Theta(\log_{10} n)$

$$T(n) = T(n/10) + C$$

$$T(0 \dots 9) = 1$$

$$\frac{n}{10^k} = 1$$

$$n = 10^k$$

$$\log_{10} n = k$$

n
 ↓
 n/10
 ↓
 n/100
 ↓
 n/10^k

Print digits of a number recursion

```
void A (int n) {
    if (n < 10) { print(n); return }
    n = A(n/10)
    print(n%10)
    return;
}
```

n	std out	Stack
1	1	9
19	9	A(1)
198	8	A(19)
1986	6	A(198)

Time complexity = $\Theta(\log_{10} n)$ Space complexity = $\Theta(\log_{10} n)$

Print digits of a number iteration

```
void A (int n) {
    a = []
    if n == 0 : print(n)
    while (n) {
        a.append(n%10)
        n = n/10
    }
    for ai in a[::-1] // a reversed
        print(ai)
    return;
}
```

n → 1986

a → 6

n → 198

a → 6 8

n → 19

a → 6 8 9

n → 1

a → 6 8 9 1

Time complexity = $\Theta(\log_{10} n)$ Space complexity = $\Theta(\log_{10} n)$

Print digits of a number iteration (constant space)

```

int R(int n) {
    s = 0
    while (n) {
        s = s * 10 + (n % 10)
        n = n / 10
    }
    return s;
}

```

memorized for entire procedure

$n \rightarrow 1986$
 $s \rightarrow 6$
 $n \rightarrow 198$
 $s \rightarrow 68$
 $n \rightarrow 19$
 $s \rightarrow 689$

Time complexity = $\Theta(\log_{10} n)$
 Space complexity = $\Theta(1)$

$n \rightarrow 1$
 $s \rightarrow 6891$

```

int R1(int n) {
    int s = 0
    int x = R(n, s)
}

```

Need for helper function

n	s
1986	0

helper function:

```

int R(int n, int s) {
    s = s * 10 + (n % 10)
    if (n < 10) { return s; }
    int x = R(n / 10, s)
    return x;
}

```

$n=1$
 $s=689$
 $s = 6891$
 $x=$
 $n=19$
 $s=68$
 $s = 689$
 $x = R(1, 689)$
 $n=198$
 $s=6$
 $s = 68$
 $x = R(19, 68)$
 $n=1986$
 $s=0$
 $s = 6$
 $x = R(198, 6)$
 6891

Time complexity = $\Theta(\log_{10} n)$
 Space complexity = $\Theta(\log_{10} n)$

HOP : Return no. of hops

	0	1	2	3	4	5
a	5	1	0	4	2	3

Find no. of hops to get to

Ex: 3

len of list is unknown

there is no repetition of numbers

$$a[3] = 4$$

$$a[4] = 2$$

$$a[2] = 0$$

$$a[0] = 5$$

$$a[5] = 3$$

No. of hops = 4

Hop.ipynb Iteration Solution

Stack Trace

```
def hop_easy(self, a: List[int], f: int) -> int:
```

```
    t = f
```

a -> 5 1 0 4 2 3

t -> 3

h -> 0

```
    h = 0 // number of hops
```

t -> 4

h -> 1

```
    while (True):
```

```
        if a[t] == f:
```

t -> 2

h -> 2

```
            return h
```

```
        else:
```

t -> 0

h -> 3

```
            t = a[t]
```

```
            h = h + 1
```

t -> 5

h -> 4

```
    return h
```

t -> 3

h -> 4

Time Complexity = $O(n)$ Space Complexity = $O(1)$

Hop.ipynb Recursion Solution

```
def hop(self, a: List[int], f: 'int') -> 'int':
    if (a[f] == f): return 0
    t = a[f]
    a[t], a[f] = a[f], a[t]
    result = 1 + self.hop(a, t)
    a[f], a[t] = a[t], a[f]
    return result
```

Stack Trace

