

VIRTUAL MEMORY SIMULATOR

Generator.cpp

This cpp file generates the page reference string keeping in mind the locality of reference. We have taken inputs asking user to feed the probability of locality of reference and maximum pages and read probability. First of all we have randomly initialised a working set from the total number of pages and then we generate the next page reference number by checking if the random number given by "rand()%100" is less than the % locality of reference. If it is so then we output a page number from the working set otherwise output a page number which is not there in the working set. Similarly we have generated whether the page is to be read or written.

Paging.cpp

This cpp file outputs the actions to be taken against each page reference. We have two data structures one is page table which is stored as an array of integers and secondly a frame structure which stores the page number that the corresponding physical frame stores and the last time it was accessed.

Sequence of the whole process :

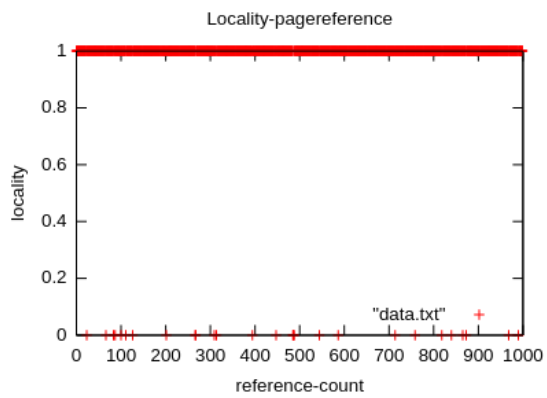
1. Given a page number to reference first of all we check whether that page number is there in the page table or not and if it is there, is its valid bit 1 or not which indicates whether that page has been allotted a physical frame or not.
2. If it is the valid bit one then we just do the operation (read or write) on that page.
3. If the page has not been allotted a physical frame then we search for an unallocated physical frame. If there is one we just allocate that frame and update our current frame counter.
4. If there is no unallocated physical frame available (this is done by checking if current frame count < total number of frames). If so then we apply the replacement_policy as given by the user to select the victim page.
5. In case of FIFO replacement_policy we have maintained a pointer to the current victim page (the page which will be sacrificed if there is no physical frame left). We unmap this page set the valid bit to 0 for its corresponding entry in the page table and increment the victim pointer by one.
6. If the replacement_policy is random then we random select a page and then unmap it from physical frame.
7. If the replacement_policy is LRU then we select a page that has the lowest value of counter and unmap that page.
8. If the replacement_policy is NRU then we keep on searching for a page in the physical frame which has its reference bit 0. If we are unable to find one then we sacrifice (last_use_ptr + 1)th page. All the reference bits are set to 0 after every 1000 MEMORY accesses or 10th page fault.
9. If the replacement_policy is second_chance then we keep on searching for a page in the physical frame which has its reference bit 0 and meanwhile we change the reference bit of pages which have their reference bit set to 1 during the scan.

Image file :

This is the graph corresponding to reference representing whether the locality of reference is preserved or not.

0 : Not preserved

1 : Preserved



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