

ASK/Need/INFO

22 February 2018 22:19

C++, OOPS, IPC, Thread

Multi threaded programming, synchronization and IPCs
Boost, Advanced C++ Library STL
OO Concepts, Design Patterns & UML

C programming skills with working experience on Common Data structure and algorithms
Expertise in ACLs, TCAM management, 802.1x/MAC Authentication, Access Security features like DHCP snooping, IP source guard, DDos mitigation

Knowledge of Embedded C & Assembly programming
- Working experience in communication protocol CAN, LIN, Flexray, Ethernet
- Working experience diagnostic protocol (UDS)
- Working experience with CANoe, CANalyzer
- Working experience with AUTOSAR architecture, MCAL drivers
From <<https://www.naukri.com/mnjuser/inbox>>

Excellent knowledge of Openstack, Dockers, Neutron, Nova.
Experience in NVF, SDN Projects, SDN Applications, SDN Controllers (e.g. OpenDaylight), SDN Data Plane and Openflow or related development.
TCP/IP networking, Routing, Switching
Knowledge on RESTful services and Message Broker system
Ansible, Salt, Puppet, Chef atleast any one
Atleast on of the frameworks like Django, Flask etc.

Worked on L2/L3 protocols like VLAN, OSPF, IPSec, RSVP, Metro Ethernet, LLDP, LAG etc

Multi-threaded environment development experience

From <<https://www.naukri.com/mnjuser/inbox>>

Layer 2, Layer 3, TCP/IP technologies, VLAN, IP, QOS, DHCP, 802.1x, IP forwarding, Routing Protocols.

Hands on experience in networking Switching, DPDK, Network Switches, SDN
Knowledge on Open Network Linux (ONL), ONOS, Open switch Specification is highly desirable

Strong in Data structures, Algorithms

Exposure to Carrier Ethernet
Experience in Ethernet, Layer 2, Layer 3, TCP/IP technologies, VLAN, IP, QOS, DHCP, 802.1x, IP forwarding, Routing Protocols.
Develop and execute feature and/or system integration test cases for networking protocols in L2 and L3 layers like LLDP, ARP, RIP, OSPF etc., Knowledge of switch debugging, topology creation, innovative testing to reproduce customer defects.

From <<https://mail.yahoo.com/d/folders/1>>

Development with Real time OS, Embedded Linux, OSE

Worked in Agile development environment

C/C++; Java/Python

Scripting Skills in Ansible-playbook / Ruby / puppet/ REST/ Platform tools

Networking/SDN

1 year or more development experience in Networking (Switching, SDN)
Hands on experience in networking – Switching, DPDK, Network Switches, SDN
Python programming (basic/advance/REST-API/Jira/Flask is okay)

From <<https://mail.yahoo.com/d/folders/1>>

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HP: +91 9108506743 | Phone: +91 80 41247333.

Mav

GNU toolchain over Linux platform.

OOAD and problem solving skills.
Good understanding of multi-threading, IPC, memory management and other OS level concepts.
Knowledge of IPv4/IPv6 networking and transport layer protocols such as TCP/UDP/SCTP.
Knowledge of telecom protocols like S1AP, NAS, Diameter, Radius, GTPC, RANAP, SS7 etc.

should have good IE level understanding of MIB/SIB/RRC attach procedure/HO procedure/re-establishment/cell reselection etc
procedures like : UL/DL processing chain, CQI/RI/PMI reporting, DCI formats, MIMO, CA etc.
RLC/PDCP is also good to have skillset
Knowledge of L1 testing using test tools from Agilent, R&S (CMW500)
• Familiarity with JTAG, Lauterbach debugging or modem platform debugging environment.
Solid knowledge of LTE Access Stratum protocol layers and procedures required.
Hands on Testing Experience in any of recognized Router/Switching/FW platforms

Knowledge of TCP/IP, SNMP, IPSec, Radius
knowledge of test analysis tools like Wireshark , QXDM, VR5 , Spectrum Analyser etc

Desired Technology Knowledge on:

- Switches and IP Networking Fundamentals.
- Internet Protocols: TCP/UDP/IPv4/IPv6
- Expertise in one of the below technical areas
- **Layer 2:**
 - Data center & Switching technologies
 - STP,RSTP,MSTP
 - Access control, DHCP, Dot1X, POE
 - Virtual chassis, clustering
- **Layer 3:**

- Routing and routed protocols
 - MPLS, LDP, RSVP,L3VPN,L2VPN
 - OSPF,ISIS,BGP
- Platform specific features:
 - COS ,Firewall,ACL
 - Schedulers ,Rate limiters
 - Virtualization technologies
 - Multi chassis, clustering

From <<https://careers.juniper.net/careers/careers/jobdescription.html?iid=938337>>

From <<https://my.naukri.com/Inbox/viewConversation?id=&conversationId=781025094>>

Hands on experience in WLAN networks,RF,Microwave Technologies,L2 configurations,MPLS-VPN Networking,Transmission Technology,Configuration and troubleshooting of Wireless Modems,WAP,RF Radios,Switches and Routers.

>

Switch switches within the subnet, that is switching. In switching packets are transferred from source to destination using MAC address. Switching is done within the network.

Whereas Router routes between the network. Routing is a process which is done between two networks using IP addresses.

- 1 to 3 years of Test experience in LTE UE Protocols (RRC/PDCP/RLC/MAC/PHY) and various feature and respective call flows - Experience to test UE stability and performance and end-to-end integration with EPC.
- Knowledge of IMS, TCP, UDP, SCTP, IPSEC is highly desirable. - Knowledge on 5G is desirable.
- Experience in one of the scripting language e.g. shell, python, perl, etc., automation tools, test tools used in end-to-end LTE network testing - LTE/5G NR Protocol Testing, Functional Testing, Sanity test, Operator testing, Regression, Error verification etc.. - Perform the testing activities in network operators & vendor laboratories if needed. - Preliminary log analysis for failure test cases, Collect, analyze and compile FT results and log files and deliver these to project teams and Error Data Base. - Raise issues with accurate description, issues retest, ad-hoc issues retest in the correct environment and deliver accurate feedback to the project for the same.

Necessary Skills / Attributes

- Hands-on in generating test cases/test plans based out of requirements/3GPP Specifications. (In the above technical domain) - Hands on Planning/building test bed, test execution and debugging skills. - Ability to understand network scenario and come up with test cases beyond standard specification is
- Highly desirable. - Good Attitude towards learning new technologies and Aptitude skills - Mobile testing with telecom network in 2G/3G/4G and network lab equipment testing experience is highly desirable.

Mandatory skills - Python, any RDBMS, Data Structure and Algorithm
Optional skills - AWS, OOPS Concept, Redis, HTML, CSS, JavaScript, and JQuery

From <<https://www.naukri.com/mnjuser/inbox>>

Working experience in Embedded C and micro controller programming
• Working experience in projects based on AUTOSAR architecture.
• Working experience in communication protocol (CAN, LIN, SPI)
• Working experience with CANoe, CANalyzer and CAPL scripting

Broadcom experience in routing / switching architecture
Broadcom HAL experience in L3 routing and sync of routing tables
Preferable: Experience with HA (High availability). HA synchronisation with switch architecture

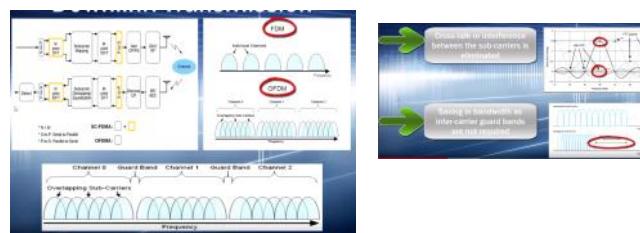
Lte physical layer

21 February 2018 07:10

Quantization:
is the process of mapping input values from a large set (often a continuous set) to output values in a (countable) smaller set. **Rounding** and **truncation** are typical examples of quantization processes.

Source encoder(transmitter):
Source coding is much about removing the redundant data or the data which is not important for the data to be conveyed to the receiver, I mean the extra data.
 Huffman Coding , Lempel-Ziv coding

Channel coding
Channel coding is more about adding some extra bits in the form of parity bits so that you can protect the data from becoming corrupt.
linear block codes, cyclic codes and convolutional codes



Block diagram of a communication system

DL	OFDMA
UL	SDFDMA

Inter symbol interference(ISI):
When multiple signal path combine at same time constructive interference and sometimes destructive

OFDMA:
Large of close spaced orthogonal subcarriers are used to carry data in several parallel data streams, offer works. OFDMA divide the incoming high - speed data stream into several low rate data streams and each data stream is mapped to one of the closest spaced orthogonal subcarrier and is used to modulate this subcarrier using varying level of QAM mode. (QPSK,QAM,16QAM) depend on signal quality

Le each carrier has its own bandwidth

each carrier in OFDMA is linear combination of the instantaneous symbol on each of sub carrier in the channel.

Subcarriers are chosen in such a way that they are harmonics (integer multiples of the lowest subcarrier => Fc, 2Fc, 3Fc, 4Fc, 5Fc)

Subcarrier spacing is constant.

Symbol duration = t = 1 / Fc

-cross talk is eliminated

-saving bandwidth in intercarrier guard band is not required

Each OFDMA symbol have cyclic prefix to prevent ISI caused by multiband delay spread Cyclic prefix represent the guard period at starting of

Duration of guard period > multipath delay spread then only ISI is prevented

Short and long both type of cyclic prefix is there

Limitation of OFDMA:

1. (Inter carrier interference (ICI)) At receiver during demodulation receive signal + local LOCO of receiver

Should both have identical frequency but not the case as internal LOCO drifts

Therefore base station send periodic synchronization signals to receiver.

2. High peak to average power ratio(PAPR):

The **peak-to-average power ratio (PAPR)** is the **peak amplitude squared** (giving the **peak power**) divided by the **RMS value squared** (giving the **average power**). It is the square of the crest factor... As it is a **power ratio**, it is normally expressed in decibel (db).

Frame structure:

FDD:

TDD:

d:downlink, U:uplink

5 frame: special

DwPTS: downlink pilot slot: syn+user data+dchannel for scheduling and control info

GP: Guard period

UpPTS: used to transmitting PRACH + sounding ref signal

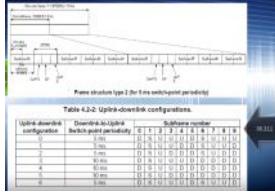
1ms=dwpts + gp + uppts

Physical resource block:

The smallest element of resource allocation that can be assigned by bTS scheduler to mobile is called PRB.

PRB:12 sub carrier duration of 0.5ms +

Normal=1 subcarrier = 7OFDMA symbols



FFT Size:
FFT size = $\lceil \text{IDFT}(\text{Tx})/\text{IDFT}(\text{Rx}) \rceil$

The number of spectral lines is always 1/2 of the selected FFT size.

Therefore :1024 point FFT produces 512 output spectral lines

The frequency resolution of each spectral line is equal to the **Sampling Rate**, divided by the FFT size.

For instance, if the FFT size is 1024 and the Sampling Rate is 8192, the resolution of each spectral line will be:

8192 / 1024 = 8 Hz.

Physical signals:

Downlink physical signal: do not convey information to or from higher layers

1. Reference signal: Used to determine channel impulse response (CIR), carrier offset estimate channel estimation, timing sync for coherent demodulation

CIR is calculated by reference signal if one sending channel response in MIMO then other is adhoc, when CIR is known data can be transferred by both simultaneously

Reference signal by every subcarrier

They are staggered in time and frequency domain

S10 unique ref signals

Specified identifier is assigned to each cell within a network and acts as cell specific identifier

when UE have to figure out downlink power it take power of this reference signal

2. Synchronization signals:

PSS(DwPTS synchronization signal);Zadoff chu sequence (0,1,2)

Zadoff-Chu sequence

Three sequences possible (0, 1, 2)

Mobile receives incoming signal for a time of at least 5 ms and compares with those three root sequences (N/2)=10

SSS(secondary synchronization signal)

Pseudo-random sequences known as Gold sequences

Different for the first and second transmission of the signal within the frame

Exact sequence indicates the cell identity group N(1)ID

Physical cell identify

$$N_{\text{ID}}^{\text{cell}} = 2N_{\text{ID}}^{(1)} + N_{\text{ID}}^{(2)}$$

A number between 0 and 503

Transmitted on synchronization signals

N(1)ID is the cell identity within the group, which runs from 0 to 2 and is signalled using the PSS

N(2)ID is the cell identity group, which runs from 0 to 167 and is signalled using the SSS

Uplink physical signals:

DRS: downlink reference signal

channel estimation

SRS: sounding reference signal

scheduling

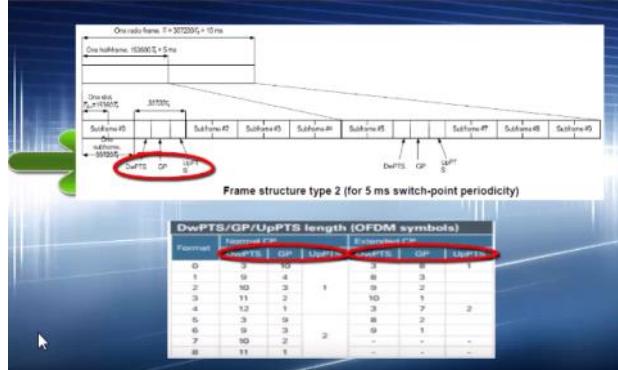
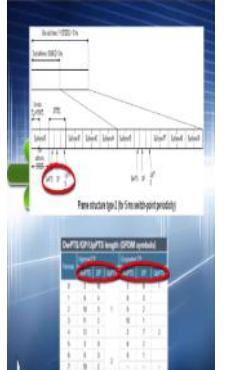
Uplink transmission:

Lte in UL uses SC-FDMA

->lower PAPR

Misleading term

A multicarrier scheme





Link Adaptation :
1.Power control without fixed rate(main target: voice/control , Goal: reliable communication)
2.Rate control without fixed power(main : data, Goal: Fast data rate)

Power and Rate control

Techniques:

Outer loop link adaptation

Channel quality indicator +HRQ+power control

Calculation of Slot Duration

Sampling frequency for 20 MHz channel bandwidth is 1500 Hz = 2048 (FFT size) = 30.72 MHz = F_s
Sampling time is: $T_s = 1/f_s = 1/(30.72 \text{ MHz}) = 33.33 \mu\text{s}$
20.72 MHz (Sampling Frequency in LTE = 20 MHz Bandwidth) = $8 \times 3.84 \text{ MHz}$ (Sampling frequency in UMTS)
Duration of a time slot should be duration of a 7 OFDM symbols + duration of 7 cyclic prefixes
OFDM symbol duration will be $= 1/1500 = 0.67 \mu\text{s}$
Duration of cyclic prefix = $5 \times 4.76 \mu\text{s}$ (sampling time)
For long cyclic prefix = $51.2 \mu\text{s}$ (sampling time)
Slot duration = duration of 7 OFDM symbols + duration of 7 cyclic prefixes
 $= 1/(1500 \times 2048) = 4.6 \mu\text{s}$
When antenna CP is used, size of first CP is: $160 \mu\text{s}$
Slot duration = $7 \times (1/1500) \text{ sec} + 160 \times (1/(1500 \times 2048)) \text{ sec} + 5 \times 4.64 \times (1/(1500 \times 2048)) \text{ sec} = 0.5 \text{ ms}$



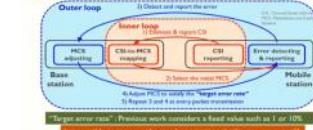
Table 6.6 - Physical signals

Signal	Repetition	Name	Use	Direction
DRS	BS	Dedication reference signal	Channel estimation	UL
SRS	BS	Scanning reference signal	Scheduling	UL
PSS	BS	Primary synchronization signal	Acquisition	DL
SSS	BS	Secondary synchronization signal	Acquisition	DL
BS	BS	Cell specific reference signal	Channel estimation and scheduling	DL
RS/R9	BS	UE specific reference signal	Channel estimation	DL
RS/R9	BS	MBMS reference signal	Channel estimation	DL
RS	BS	Positioning reference signal	Location services	DL
BS	BS	CM reference signal	Scheduling	DL

Table 6.5 - Physical control channels

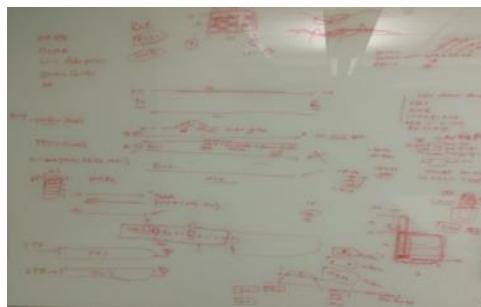
Channel	Release	Name	Information carried	Direction
PUCCH	R8	Physical uplink control channel	UCI	UL
PCFICH	R8	Physical control format indicator channel	CFI	DL
PICH	R8	Physical hybrid ARQ indicator channel	HR	DL
PDCCH	R8	Physical downlink control channel	DCI	DL
R-PDCCH	R10	Relay physical downlink control channel	DCI	DL
E-PDCCH	R11	Enhanced physical downlink control channel	DCI	DL

Overview of OLLA (Outer Loop Link Adaptation)



"Target error rate" - Previous work considers a fixed value such as 1 or 10%.

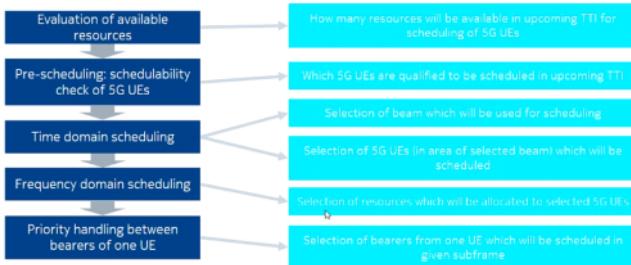
Is it possible to increase the rate by optimizing the "target error rate"?



Technical Details

Scheduling process overview

- All of presented steps are done every TTI and separately for downlink and uplink



challenges:
low cost components
synchronization of terminal
reduction in internal power consumption

LTE call flow
Handover,(intra rat, x2 ,x1)
Debugging,
Spec:3GPP Release 13
36.413->s1ap(enodeb and mme)
38->5g->



Mobile start listening

Send RACH: Random access procedure | hey I am here

Two part

RAPID=64 values	Random access preamble identity
RA-RNTI: decide time and frequency slot	RA radio network temp identity

RACH response	Rapid and time RA
Message 3	Very small
Message 5	Big message NAS req /PDN connection request

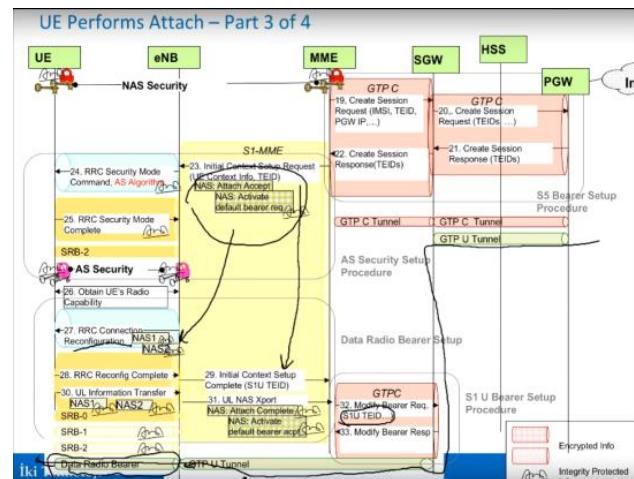
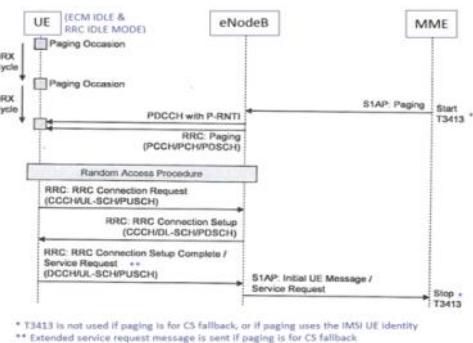
when RACH happens? prach(wcdma),channel request(gsm)

From <http://www.sharetechnote.com/html/RACH_LTE.html#Why_RACH>

initial access from rrc idle	
rrc contention reestablishment	
handover(contention based and non contention based)	

UE attach steps:

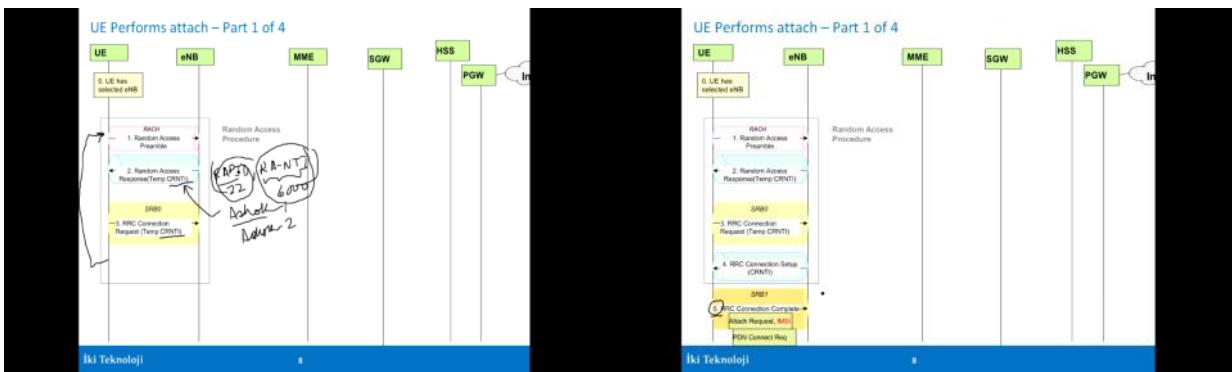
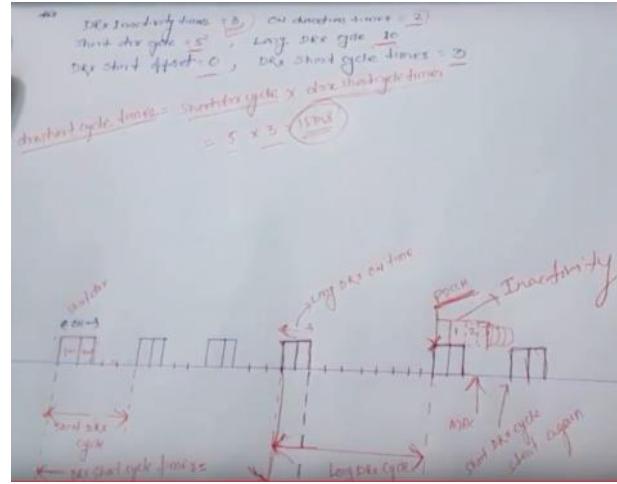
Power on UE	UE power on (MM/RR/PLMID) taken from USIM application)
Steps	
Random access procedure	1. UE sends Random access preamble: (64 preamble ,RA-RNTI if contention based, zadoff chu seq->FDD) 2. Random access response: Temp CRNTI(RAPID+RANTI) 3. RRC connection req (UE->enb): S-TMSI ,SRB-0 4. RRC connection setup(enb->ue):CRNTI 5. RRC connection setup Complete(UE->ENB):SRB-1,IMSI NAS PDU (EMM attach req: old GUTI or IMSI, ESM PDN connectivity req)
User authentication S6a(MME -> HSS)	1. s1ap initial attach(Enb->mme): 2. Diameter authentication info(MME->HSS):Req IMSI,PLMID 3. Diameter authentication answer(HSS->MME):RAND auth ,Sres,KASME 4. S1AP DL NAS Transport(MME->ENB):EMM attu req {RAND,AUTN} 5. RRC DL info transfer(ENB->UE):NAS EMM PDU 6. Ue ->UESIM:UESIM authenticate, to get SRES and C 7. RRC UL info transfer->NAS EMM Authentication uplink response 8. S1AP uplink transport(ENB->MME)->NAS EMM PDU MME compare SRES from UE and HSS
NAS security setup	1. EMM security mode command over RRC/s1ap transport(MME->UE) 2. EMM security mode complete over RRC/s1ap transport(UE->MME) ->algo for encryption and integrity protection
Authorization	Update location req (MME-HSS):IMSI, visited PLMID, RAT type Update location answer: subscription data
S5 bearer setup	MME does SGW and PGW selection 1. GTP-C Create session req(MME->SGW):(IMSI, APN, PDN SS/S5 address,F-Teid, PDN type, EPS Bearer ID, Default EPS QoS, Aggregate Maximum Bit Rate, PDN address (fwd?), PCO). 2. GTP-
AS security setup	1. S1AP-intial(S1 MME MME->ENB) context setup Request 2. RRC AS security mode command (ENB->UE):cyphering and integrity algo 3. RRC reconfiguration(ENB->UE):(EPS bearer id, DRBId, Radio configs) 4. RRC security mode complete 5. RRC connection reconfiguration complete 6. S1AP-intial context setup response complete(Enb->MME)
Data Radio bearer setup (Between UE and ENB) And S1U bearer setup	1. EMM attach complete over RRC/S1AP(UE->MME): (MME would trigger GTP-C procedure towards SGW to inform eNB related details to SGW.) 2. GTP-C Modify bearer req(MME->SGW): 3. Modify bearer res 4. User uplink data through DRB (UE->ENB): 5. User uplink data through GTP Tunnel (ENB->SGW):S1 6. User uplink data through GTP Tunnel (SGW->PGW):S5/s8
DHCP client ,server	This complete LTE initial attach procedure. UE is attached, default bearer active. RRC-Connected, EMM-Registered, ECM-Connected state.



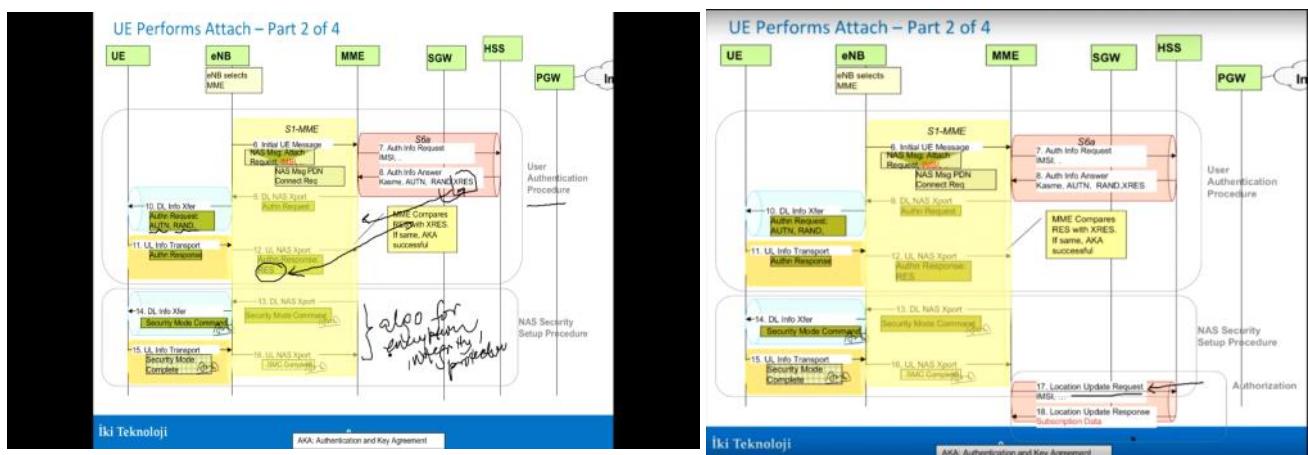
DRX: discontinuous reception
UE need to wake up and monitor PDCCH for every subframe. -> whether N/w is sending data for UE -> wastage of energy

How we know when N/W send data?
N/w decide when UE wake up or sleep using Transmitted by SIB2 through RRC setup/reconfiguration message

DRX cycle on duration timer	When UE wake up -> why (downlink data, change of s1, cell reselection, incoming call, receive sms)
DRX cycle off duration timer	
DRX inactivity:	If UE receive PDCCH until in on state then inactivity timer starts. If UE receive DRX mac control reception during inactivity timer: first stop (on state timer and drx inactivity timer) If ACK NAK -> starts DRX retransmission timer
DRX in idle mode: Paging cycle Frame CYCLE Paging Occasion	
DRX Connected and dedicated mode DRX on duration cycle and timer DRX inactivity cycle and timer Short DRX cycle timer Long DRX cycle timer DRX Retransmission timer	
Paging: PCCH->PCH->PDSCH RRC procedure waking up UE when it is in idle mode if any data for it	MME is responsible for LTE paging procedure by forwarding S1AP paging mess to more than one eNodeB Paging is define in term of radio frame(rf): 32rf,64rf,128rf Gen 128rf->1.28ms
idle mode	UE do not have S1AP connectivity with MME UE location of MME in idle mode is known by mme on aper tracking area basis
Which channel contain Paging info	PDCCH->indication of paging comes on PDCCH (it has P-RANTI to wake up look for paging) PDSCH -> actual content of data/info on PDSCH
MIB(24 bit):PBCH Only in DL dir every 40ms	System BW :3bit PHICH configuration:3bit Frame no sync:8bit No of antenna:10bit
SIB:PDCCH->PDSCH 80ms 13 sib	Cell reselection handover Initial attach and mobility related info
SIB-1	Cell selection, PLMN identity, TAC, Cell id
SIB2	RRC information common for all UE, RACH, paging info, initial power, SRS, PUCCH, PUSCH
sIB3	
Cell reselection: SIB 4: cell reselection same frequency SIB5: cell reselection/Handover different frequency SIB 6: from 4g lte to umts3g SIB 7:lte to 2g SIB 8:lte to CDMA200	

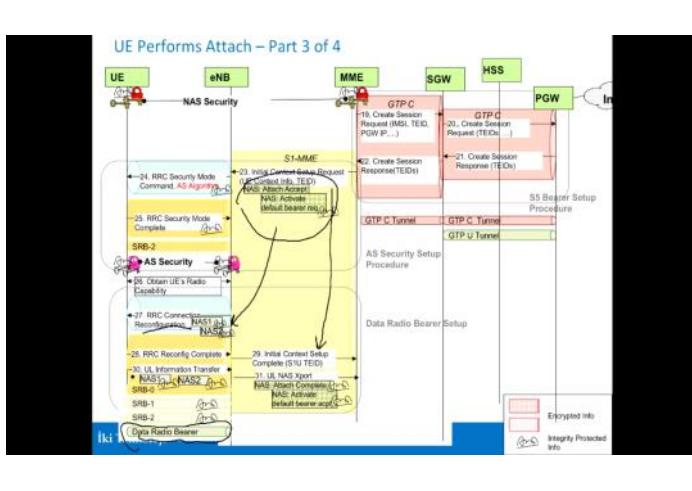


If XRES and Res is same then authentication completes
AFTER that integrity protection begins

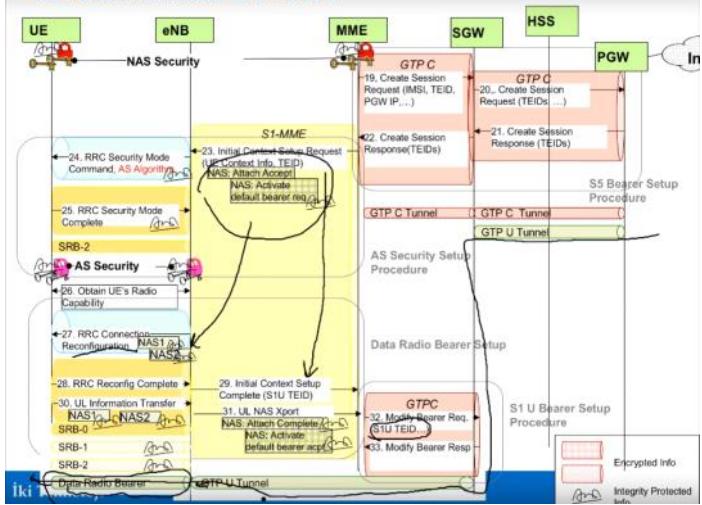


Subscription data	Data connectivity, sms or not, overall Gb
Location update	This mme is serving the mobile ,for ip not needed which mobile
Setting S5 and GTP C tunnels	

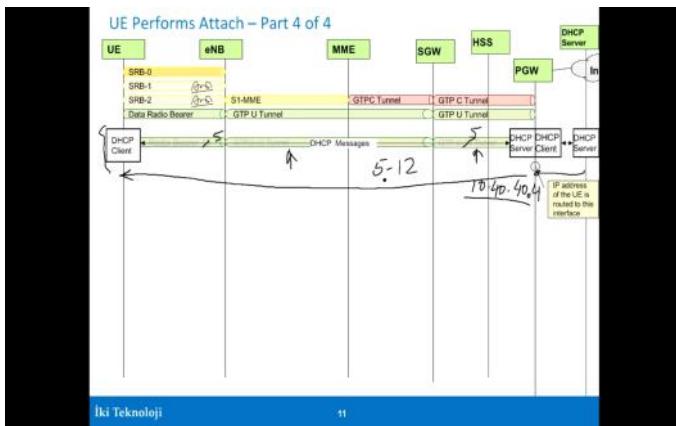
AKA: Authentication and Key Agreement	
Subscription data	Data connectivity, sms or not, overall Gb
Location update	This mme is serving the mobile ,for ip not needed which mobile
Setting S5 and GTP C tunnels	
GTP U	
After that setting data radio bearer	



UE Performs Attach – Part 3 of 4



After all this still we did not get the ip address of mobile,ip



RACH procedure	PRACH: no mac ,no rlc RACH res: no mac ,no rlc
RRC connection Request (UL SCH->ULCCH)	MAC ,RLC TM, no pdcp
RRC connection setup(DL CCCH->DL SCH)	MAC ,RLC TM, no pdcp
RRC connection setup complete(UL SCH->UL DCCH)	MAC ,RLC AM, PDCP
Authentication req	mac ,RLC am ,pdcp
Authentication res	mac ,RLC am ,pdcp

Mobility management in LTE:
EPS Bearer is a virtual connection between UE and PGW which identifies a data send and received between these two end points with specific QoS attributes. The procedure used to establish an **EPS Bearer** is called "**EPS Bearer Activation**".

Architecture Concept 2: Mobility Management (MM)

- What is **Mobility Management** in LTE?
 - It is concerned about the UE's **registration state** at the NAS layer.
 - There is an MM state-machine that runs in both the MME and UE. They key states are:
 - EMM-DEREGISTERED
 - EMM-REGISTERED
 - Examples of EMM Procedures
 - Attach
 - Authentication
 - Security Mode Command
 - GUTI reallocation
 - Tracking Area Update
 - Paging
 - Service Request
 - For each UE there is a **MM context**. The MM context consists of UE's security related parameters (keys, counters)
- Enhanced CS domain*
- EMM PS packet*
- Idle connected*

LTE always on ip connectivity until rel 12

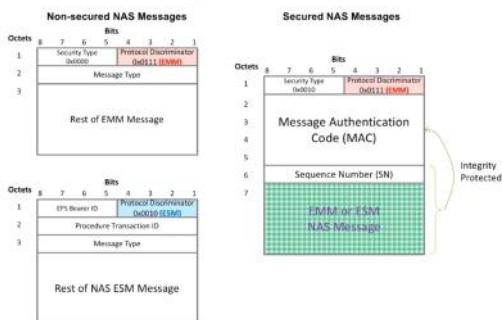
Architecture Concept 2: Mobility Management (MM) and Session Management (SM) Relation

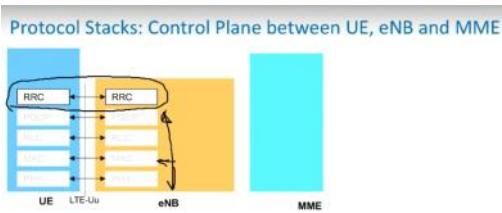
- There are separate NAS messages for EMM and ESM.
- EMM Context and ESM context are stored separately in the MME.
- LTE has the concept of "Always ON" IP connectivity.
 - The moment the UE attaches, a default bearer is setup for the UE to have connectivity. For UE to be in EMM-REGISTERED state, the UE MUST have at least one bearer context.
 - If the last default bearer is removed, UE is moved to detached state. If the last default bearer is removed, the UE automatically enters EMM-DEREGISTERED state.
- ESM procedures can be performed only if an EMM context has been established between the UE and the MME.
- Hence, there is coupling between EMM and ESM state in LTE.

Architecture Concept 2: Session Management (SM)

- What is **Session Management** in LTE?
 - It is related to UE's connectivity (**EPS bearers**) at the NAS layer.
- There is an ESM state-machine that runs in both the MME and UE for each EPS bearer. They key states are:
 - Bearer Context In-active
 - Bearer Context Active
- Examples of ESM procedures are:
 - PDN Connect Request/ Activate Default EPS bearer
 - PDN Disconnect Request/ Deactivate Default EPS bearer
 - Activate dedicated EPS bearer
 - Modify default/dedicated EPS bearer
 - ...
- For each UE there is a ESM context that includes context for all active bearers and includes parameters applicable to the bearers (Bearer ID, QoS parameters,...).

NAS EMM and ESM Message Formats





Radio Resource Control (RRC): The main control interface between eNB and UE

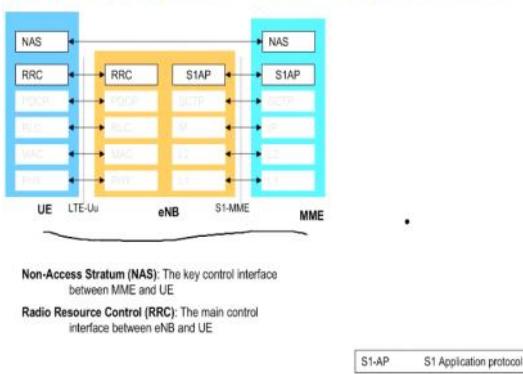
Packet Data Convergence Protocol (PDCP) Decapsulates/duplicates RLC PDUs

Radio Link Control (RLC): Segmentation/reassembly of logical channel data PDUs

Medium Access Control (MAC): Access control and arbitration

Physical Layer (PHY): Modulation, demodulation, MIMO

Protocol Stacks: Control Plane between UE, eNB and MME

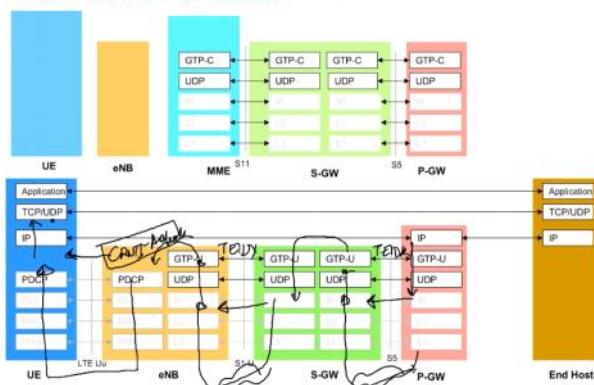


Non-Access Stratum (NAS): The key control interface between MME and UE

Radio Resource Control (RRC): The main control interface between eNB and UE

S1-AP S1 Application protocol

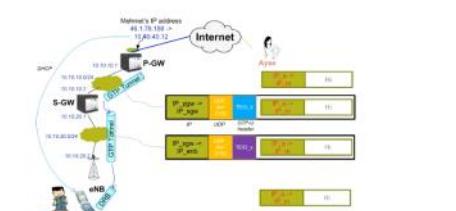
The GTP-C and GTP-U Protocols



GTP:	GPRS tunneling protocol ,
What is tunneling:	Putting one ip packet inside another ip packet
why	route to a path which is not topologically correct
What	

What is the purpose of GTP Tunneling?

- Route an IP packet along a path that is not topologically correct for the packet.



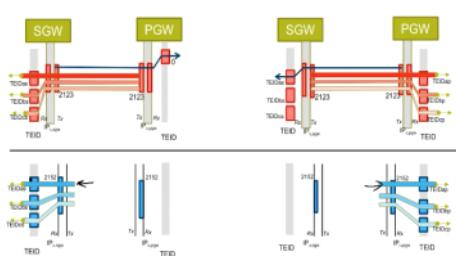
How does the GTP-U Tunnel get setup?

- GTP Protocol has two parts
 - Signaling part called GTP-C (GTP-Control)
 - User data part called GTP-U (GTP-User)
- GTP-C is used to setup GTP-U tunnel
- Both GTP-C and GTP-U run on top of UDP
- IP-in-UDP tunneling is only used for GTP-U
- GTP-C carries control/signaling messages

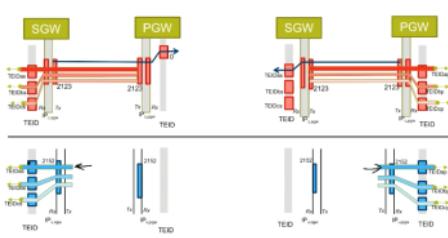


Full Teid (F-Teid) s11	ip address + TEID
GTPC ->	If ip address is zero it means that SGW and PGW is same

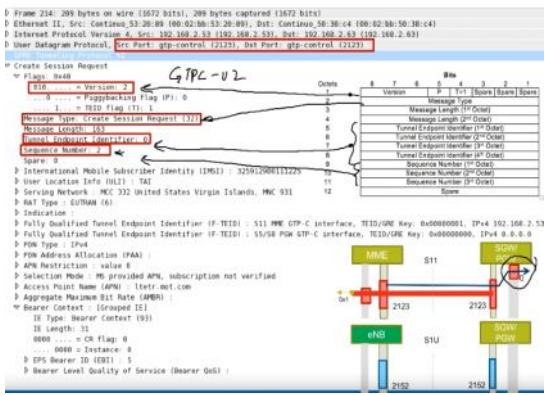
How does GTP-U Tunnel get setup: Example



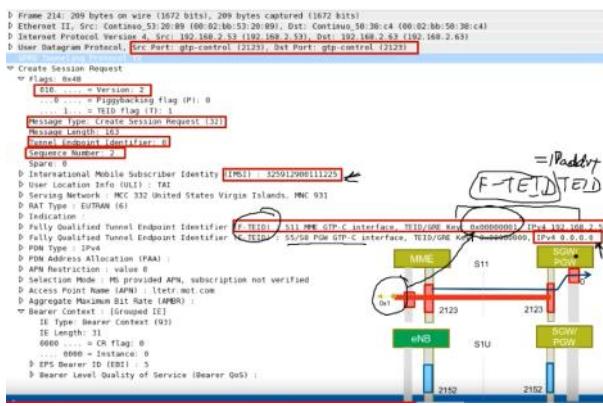
How does GTP-U Tunnel get setup: Example



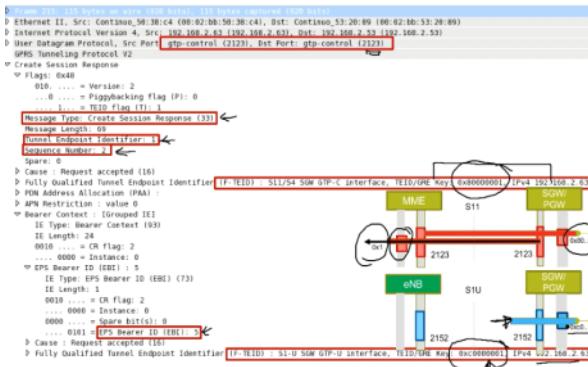
Packet Trace: Create Session Request (S11); MME->SGW



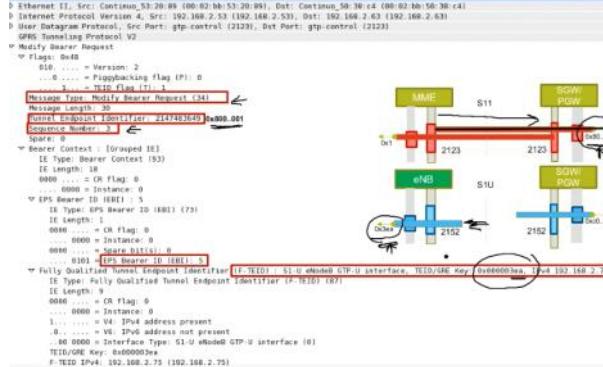
Packet Trace: Create Session Request (S11); MME->SGW



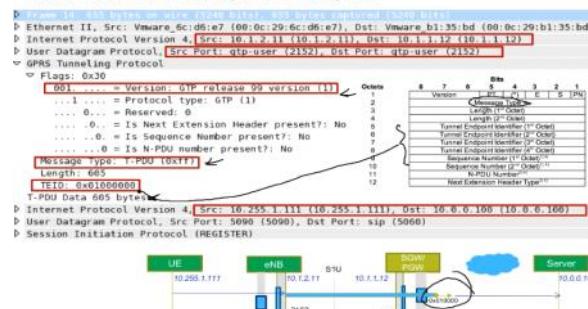
Packet Trace: Create Session Response (S11); SGW -> MME



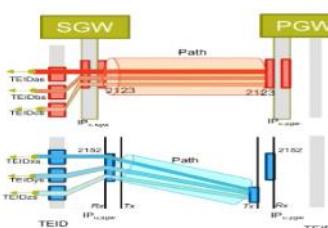
Packet Trace: Modify Bearer Request (S11); MME -> SGW



Packet Trace: GTP-U: S1-U; eNB -> SGW/PGW



GTP Path and Path Management



- Path is between two endpoints. Each end point is IP_address+UDP_Port#
- There can be several GTP tunnels on a path (each with different TEIDs).
- There is typically a “software process” that binds to each end-point.
- Path Management messages:** To ensure that a path is alive (both the physical link and process at the endpoint), periodic echo-request and echo-response are sent.
- Failure-detection and error-recovery mechanisms are defined.

Packet Traces: GTP-U Echo-request & Echo Response

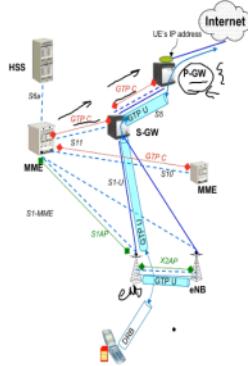
```

D Ethernet II, Src: VMware_6c:d6:e7 (00:0c:29:b1:35:bd), Dst: VMware_6c:d6:e7 (00:0c:29:b1:35:bd)
D Internet Protocol Version 4, Src: 10.1.2.11 (10.1.2.11), Dst: 10.1.2.12 (10.1.2.12)
D User Datagram Protocol, Src Port: gtp-user (2152), Dst Port: gtp-user (2152)
> GTP Tunneling Protocol
  > Flags: 0x32
    001... = Version: GTP release 99 version (1)
    ...1.... = Protocol type: GTP (1)
    ....0... = Reserved: 0
    ....0... = Is Next Extension Header present?: No
    ....1... = Is Sequence Number present?: Yes
    ....0... = Is N-PDU number present?: No
  Message Type: Echo request (0x01)
  Length: 6
  If ID: 0x00000000
  Sequence number: 0x0000
  Recovery: 0

D Ethernet II, Src: VMware_b1:35:bd (00:0c:29:b1:35:bd), Dst: VMware_6c:d6:e7 (00:0c:29:b1:35:bd)
D Internet Protocol Version 4, Src: 10.1.1.12 (10.1.1.12), Dst: 10.1.2.11 (10.1.2.11)
D User Datagram Protocol, Src Port: gtp-user (2152), Dst Port: gtp-user (2152)
> GTP Tunneling Protocol
  > Flags: 0x32
    001... = Version: GTP release 99 version (1)
    ...1.... = Protocol type: GTP (1)
    ....0... = Reserved: 0
    ....0... = Is Next Extension Header present?: No
    ....0... = Is Sequence Number present?: Yes
    ....0... = Is N-PDU number present?: No
  Message Type: Echo response (0x02)
  Length: 6
  If ID: 0x00000000
  Sequence number: 0x0000
  Recovery: 0

```

Interfaces on which GTP is used



- GTP-Cv2 is used not only for mobility management, but also for general signaling purposes, eg providing UE's cell information to the PGW and from there to operator's service network.

To access a cell the UE needs to retrieve a Preamble for the system information. Which SIB contains the RA Preamble information?

- SIB8
- SIB4
- SIB2
- SIB1

Correct

That's right! You selected the correct response.

Next Question

The sequence for eNodeB acquisition is?

- Acquire Timing, Acquire Frequency, Physical Layer Cell Id, Read Sys Info
- Acquire Frequency, Acquire Timing, Physical Layer Cell Id, Read Sys Info
- Acquire Timing, Physical Layer Cell Id, Acquire Frequency, Read Sys Info
- Acquire Frequency, Physical Layer Cell Id, Acquire Timing, Read Sys Info

Correct

That's right! You selected the correct response.

Next Question

Which network entity allocates the UE the GUTI?

- P-GW
- PCRF
- MME
- S-GW

Correct

That's right! You selected the correct response.

Next Question

What is the main reason for S1 Release?

- Incoming data for UE
- Inactivity
- Failure in connections
- HSS failure

Correct

That's right! You selected the correct response.

Next Question

If the UE is in ECM-IDLE state, on Dedicated bearer creation, the MME will trigger?

- Network Triggered Service Request
- Location area update
- De-registration or
- Registration of th

Correct

That's right! You selected the correct response.

Next Question

In the Network Initiated Service request, Paging messages are sent to all eNodeBs within which area?

- Routing Area
- Location area
- Tracking Areas in Tracking Area list

Correct

That's right! You selected the correct response.

[Next Question](#)

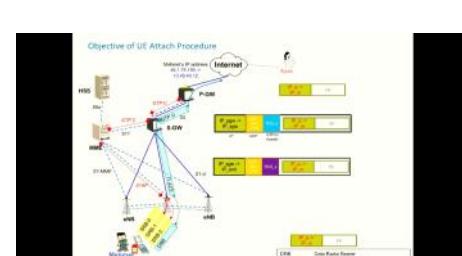
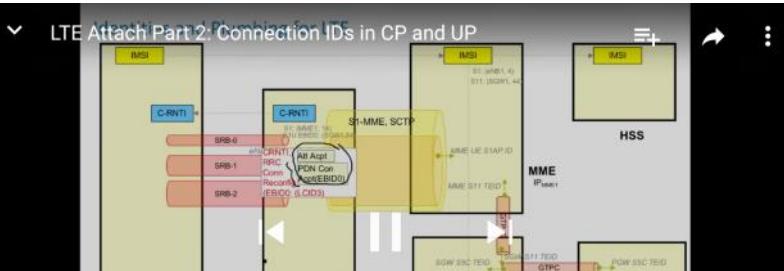
Data for the UE in a Network Initiated Service Request is buffered at which Network Element?

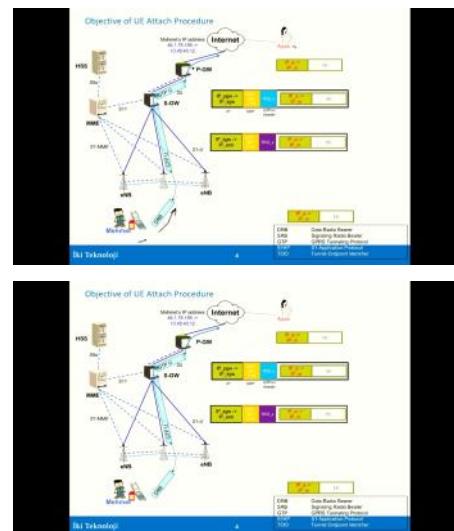
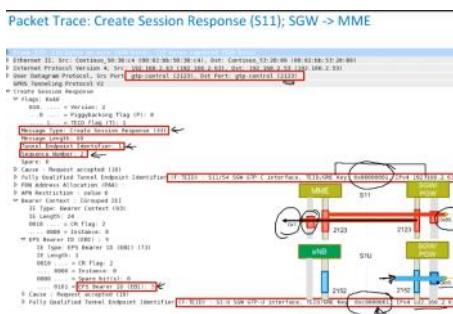
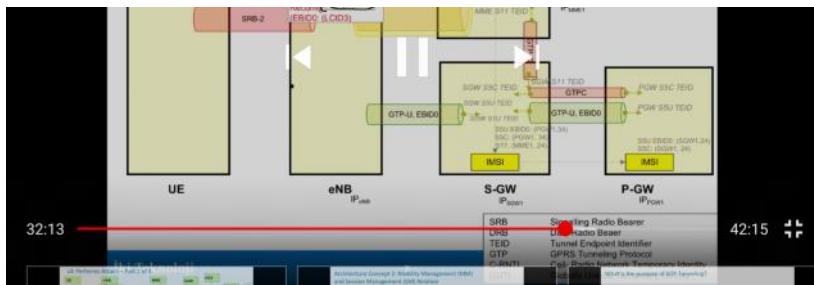
- S-GW
- P-GW
- eNodeB
- MME

Correct

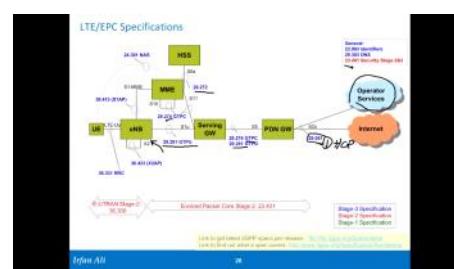
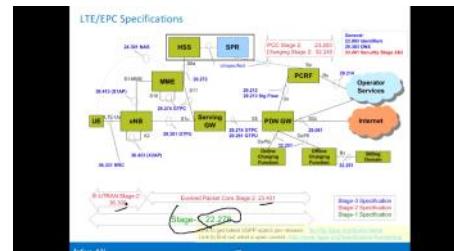
That's right! You selected the correct response.

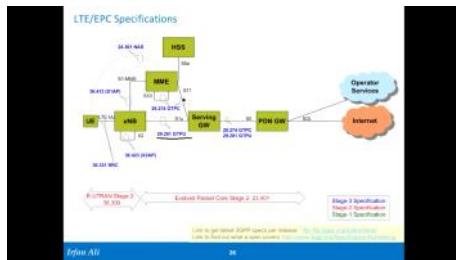
[Next Question](#)





Event	Existing and New Triggering Reasons	HO Usage
A1	Serving cell becomes better than absolute threshold	LTE
A2	Serving cell becomes worse than absolute threshold	LTE
A3	Neighbor cell becomes better than an offset relative to the serving cell	LTE
A4	Neighbor cell becomes better than absolute threshold	LTE
A5	Serving cell becomes worse than one absolute threshold and neighbor cell becomes better than another absolute threshold	LTE
A6	Serving cell becomes worse than one absolute threshold and neighbor cell becomes better than another absolute threshold due to in-device coexistence interference	LTE
A7	Neighbor cell becomes better than an offset relative to the serving cell due to in-device coexistence interference	LTE
A8	Serving cell becomes worse than absolute threshold due to in-device coexistence interference	LTE
B1	Neighbor cell becomes better than absolute threshold	Inter-RAT
B2	Serving cell becomes worse than one absolute threshold and neighbor cell becomes better than another absolute threshold	Inter-RAT
B3	Serving cell becomes worse than absolute threshold and neighbor cell becomes better than another absolute threshold due to in-device coexistence interference	Inter-RAT

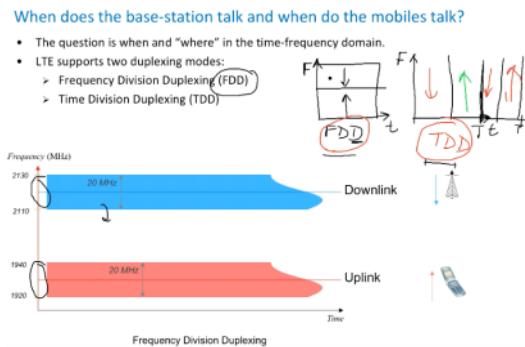




Lte frame

14 March 2018 18:32

When and where does the base station talks to mobile?



What if OFDM

- OFDM = Orthogonal Frequency Division Multiplexing
- What are orthogonal functions?
 - Two functions $\underline{h}_1(t)$ and $\underline{h}_2(t)$ are orthogonal over an interval $[0, T]$, if

$$\langle \underline{h}_1, \underline{h}_2 \rangle = \int_0^T h_1(t)h_2(t)dt = 0$$
 - Set of functions $\{h_1(t), h_2(t), \dots, h_n(t)\}$ are mutually orthogonal, if

$$\langle \underline{h}_i, \underline{h}_k \rangle = 0, \text{ if } i \neq k, \text{ and } \langle \underline{h}_i, \underline{h}_i \rangle = m_i$$

$$m_i = \text{energy} = 1$$

What if OFDM

- OFDM = Orthogonal Frequency Division Multiplexing
- What are orthogonal functions?
 - Two functions $\underline{h}_1(t)$ and $\underline{h}_2(t)$ are orthogonal over an interval $[0, T]$, if

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 - Set of functions $\{h_1(t), h_2(t), \dots, h_n(t)\}$ are mutually orthogonal, if

$$\langle \underline{h}_i, \underline{h}_k \rangle = 0, \text{ if } i \neq k, \text{ and } \langle \underline{h}_i, \underline{h}_i \rangle = m_i$$

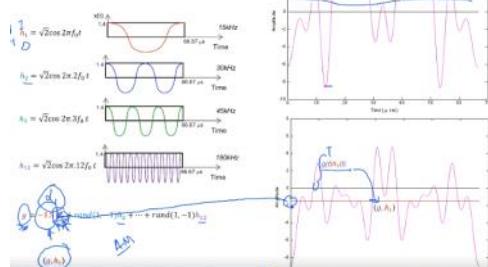
$$m_i = \text{energy} = 1$$
 - If $g(t) = \alpha_1 h_1(t) + \alpha_2 h_2(t) + \dots + \alpha_n h_n(t)$, for $t \in [0, T]$, then

$$\langle g, h_i \rangle = (\alpha_1 h_1, h_i) + (\alpha_2 h_2, h_i) + \dots + (\alpha_n h_n, h_i) \text{ for } i = 1..n$$

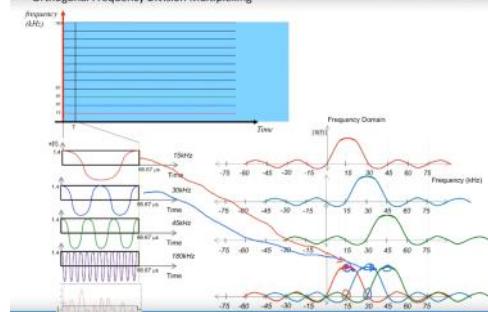
$$\sqrt{\alpha_i^2} = 1$$

Orthogonal cosine functions

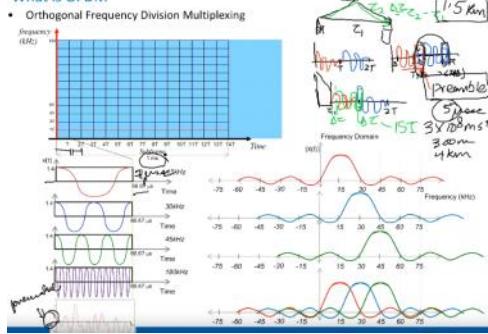
- Harmonics (multiples) of cosine functions of frequency f_0 are
- Let $f_0 = 15\text{ kHz}$, $T = 66.67\mu\text{s}$



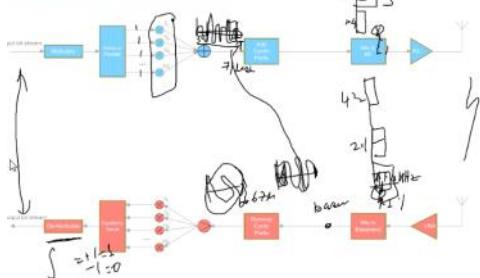
Orthogonal Frequency Division Multiplexing



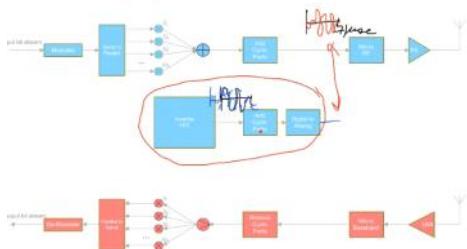
What is OFDM



How is OFDM signal generated?



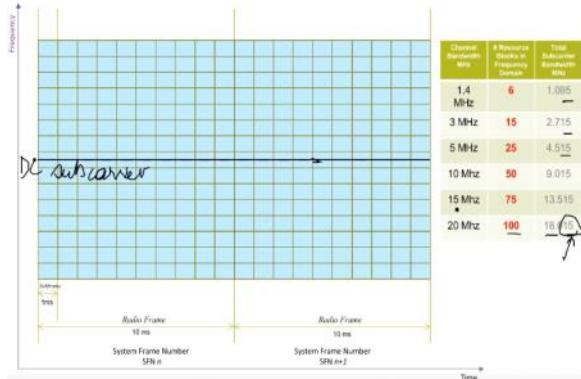
How is OFDM signal generated?



Why OFDM?

- The OFDM symbols duration is relatively long (66.67 μ s), which allows one to add time-gap (preamble) to handle relatively long **delay-spread** of the channel (5μ s \sim 1.5 km) without loosing much capacity.
 - Reduced inter-symbol interference
- Multiple sub-carriers (rather than a single carrier) over large bandwidths (20 MHz) enable to handle channel-fades over these large bandwidths.
- Increased processing capability.**

LTE Downlink Frame Structure



When mobile looks downlink radio frame it has very little idea of where exactly center frequency is ?

How many resource grid are there?

He is somewhere in between the frame?

\rightarrow To resolve/facilitate this PSS is provided.

It is in middle 62 subcarrier

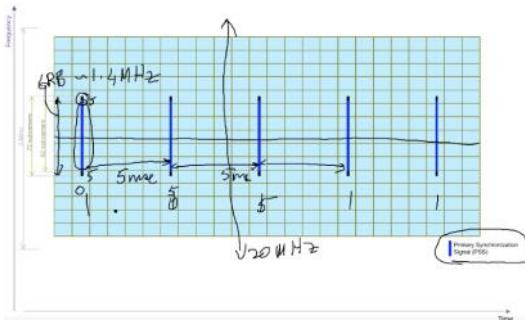
It is repeated every 5ms,

Sent at middle of dc frequency

5 resource at top and 5 at bottom with zero energy

72 carrier == 6RB i.e min transmitted in 1.4MHz

Synchronizing to DL Radio Frame



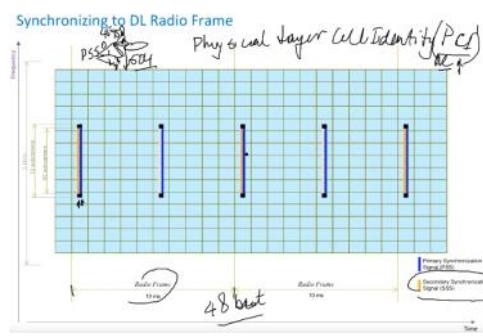
After PSS also UE would not know where is 10ms boundary
So SSS is one signal before PSS which will tell the cell identity.

PSS is transmitted on 6th symbol and SSS is transmitted on 5th symbol of slot zero and slot 10 of each radio frame on 72 subcarrier centered around DC.

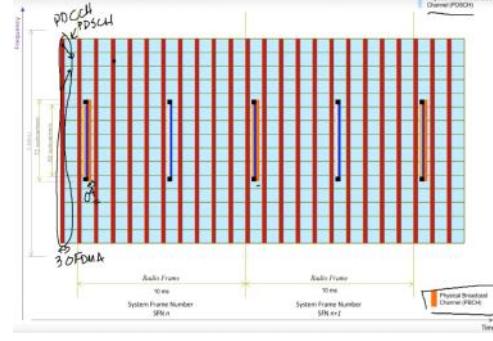
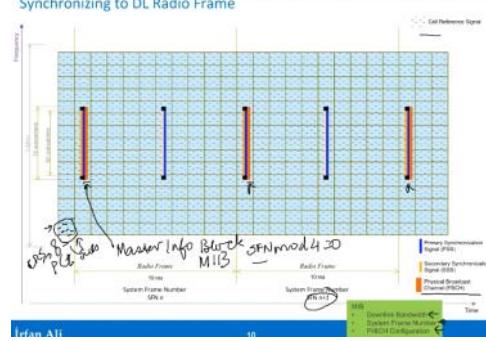
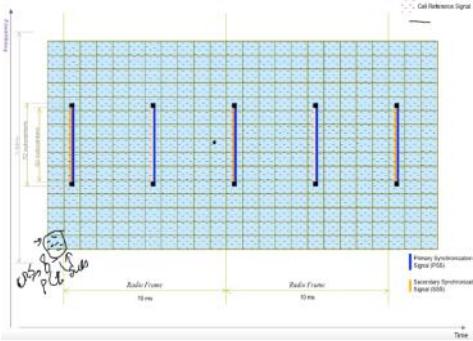
By PSS and SSS the UE will find cell identity $3(\text{PSS}) * 168(\text{SSS})$
Via cell identity UE will find cell reference signal

Once UE knows the PCI (physical layer cell identity) it can find out the Reference signal
PCI \rightarrow cell reference signal \rightarrow exact frequency and time structure \rightarrow there are 8 CRS in each subframe

Next imp thing is to find Physical broadcast channel
It is transmitted every 10ms



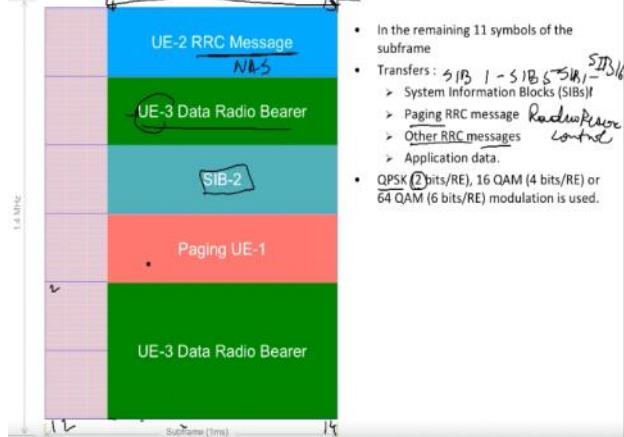
Synchronizing to DL Radio Frame



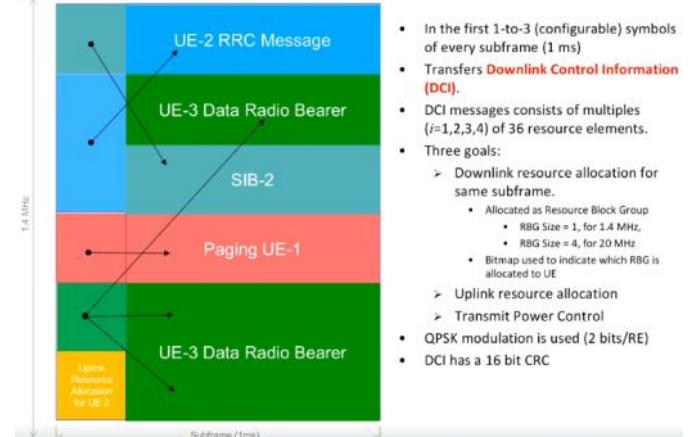
Now see PDSCH

1 frame(10ms) = 10 subframe(1ms) = 20 slot(0.5ms)
1slot=7symbol

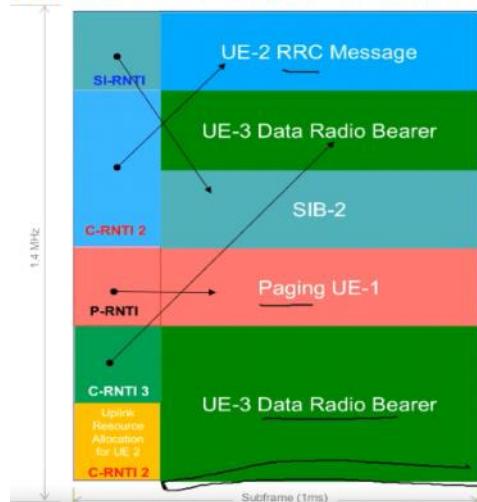
Physical Downlink Shared Channel (PDSCH)



Physical Downlink Control Channel (PDCCH)



How does a mobile know if there is a message for it in a subframe?



- There are four identities that a mobile searches for in the Downlink Control Information (DCI) in the PDCCH:
 - UE's unique cell-radio network temporary identity (C-RNTI)
 - Paging-RNTI, P-RNTI (0xFFFF) and System Information-RNTI, SI-RNTI (0xFFFF).
 - P-RNTI and SI-RNTI are the same for all mobiles.
 - The check for P-RNTI and SI-RNTI are not performed in every subframe, but on selected/ "paging-occasion" subframes, (once every DRX cycle).
 - During Random access
 - Random Access-RNTI (RA-RNTI): For Random access response message.
 - Temporary C-RNTI: For RRC Connection Setup message
- In the PDSCH, the MAC header tells the mobile, if the message is an RRC message or a data packet
 - Logical Channel ID = 0..2 \rightarrow SRB 0..2
 - Logical Channel ID = 3..10 \rightarrow DRBs

RACH

23 May 2018 07:12

- i) UE --> NW : RACH Preamble (RA-RNTI, indication for L2/L3 message size)
- ii) UE <-- NW : Random Access Response (Timing Advance, T_C-RNTI, UL grant for L2/L3 message)
- iii) UE --> NW : L2/L3 message
- iv) Message for early contention resolution

Typical 'Contention Free' RACH Procedure is as follows :

- i) UE <--NW : RACH Preamble (PRACH) Assignment
- ii) UE --> NW : RACH Preamble (RA-RNTI, indication for L2/L3 message size)
- iii) UE <--NW : Random Access Response (Timing Advance, C-RNTI, UL grant for L2/L3 message)

Exactly when and where Network transmit RACH Response

We all knows that Network should transmit RACH Response after it received RACH Preamble from UE, but do we know exactly when, in exactly which subframe, the network should transmit the RACH Response ? The following is what 3GPP 36.321 (section 5.1.4) describes.

Once the Random Access Preamble is transmitted and regardless of the possible occurrence of a measurement gap, the UE shall monitor the PDCCH for Random Access Response(s) identified by the RA-RNTI defined below, in the RA Response window which starts at the subframe that contains the end of the preamble transmission [7] plus three subframes and has length ra-ResponseWindowSize subframes.

It means the earliest time when the network can transmit the RACH response is 3 subframe later from the end of RACH Preamble. Then what is the latest time when the network can transmit it ? It is determined by ra-ResponseWindowSize. This window size can be the number between 0 and 10 in the unit of subframes. This means that the maximum time difference between the end of RACH preamble and RACH Response is only 12 subframes (12 ms) which is pretty tight timing requirement.

How is the RACH Preamble Power determined ?

Regardless of the cases, the PRACH Power (P_{PRACH}) is determined by the following equation.
 $P_{\text{PRACH}} = \min\{P_{\text{CMAX}}, \text{PREAMBLE_RECEIVED_TARGET_POWER} + PL\}$

PL stands for Path Loss between eNB Tx antenna and UE Rx Antenna. PREAMBLE RECEIVED TARGET POWER is the PRACH power that eNB expect to receive.

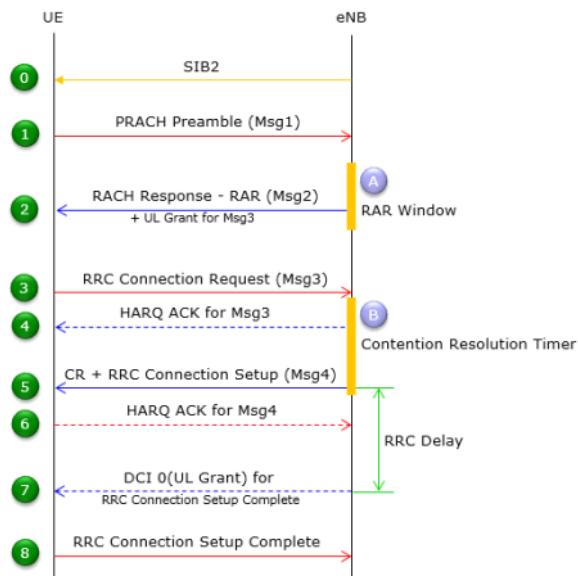
the MAC scheduler has control over the OFDM modulation in the sense that it decides, according to information received from other LTE network components, how much bandwidth each UE receives at any given moment.

Typically, a MAC Scheduler can be programmed to support one scheduling algorithm with many parameters.

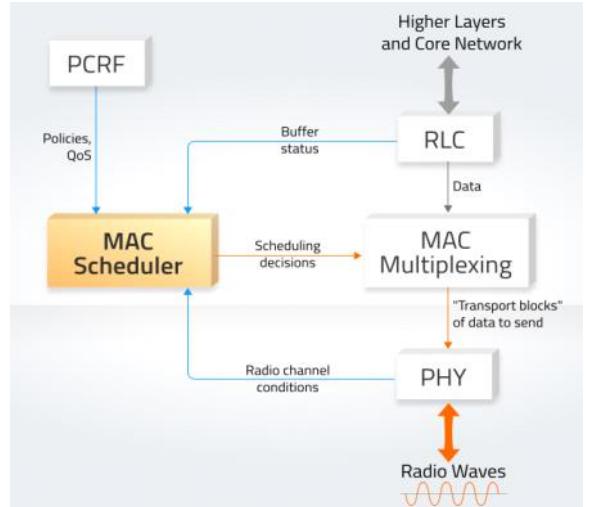
Here are some examples of scheduling algorithms:

- Round Robin – used for testing purposes and uses equal bandwidth for all UEs without accounting for channel conditions
 - Proportional Fairness – tries to balance between the QoS priorities and total throughput, usually preferred in commercial networks
 - Scheduling for Delay-Limited Capacity – guarantees that the MAC Scheduler will always prioritize applications with specific latency requirements
 - Maximum C/I – guarantees that the Mac Scheduler will always assign resource blocks to the UE with the best channel quality
- One of the key features of LTE is the ability to control and prioritize bandwidth across users. It is the MAC scheduler that gives LTE this capability.

< RACH Procedure during Initial Registration >



From http://www.sharetechnote.com/html/RACH_LTE.html#Why_RACH



RRC/data layer different types of messages

- 1.radio bearer: SRB0(RRC connection setup), SRB1, SRB2
- 2.control signaling: paging, MIB(PBCH), SIB(PDSCH)
- 3.multiple data radio berer:DRB1 (1p), DRB2(voice over ip)

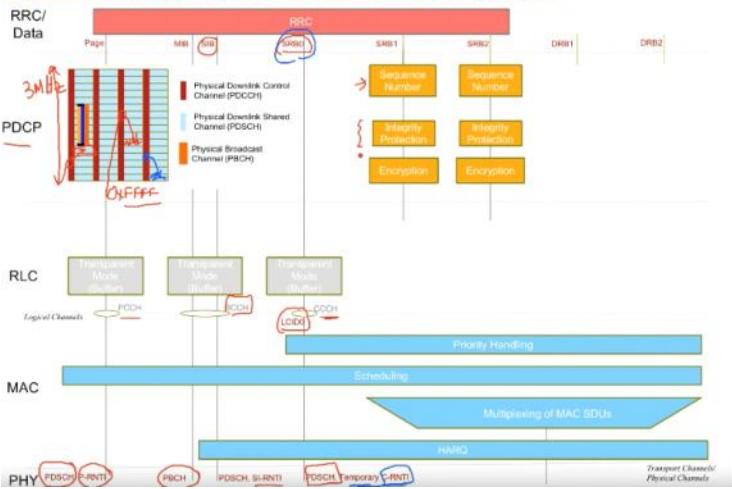
RRM

E1 interface	upcu, cpcu
Upf	Simmilar to sgw
Amf	mme

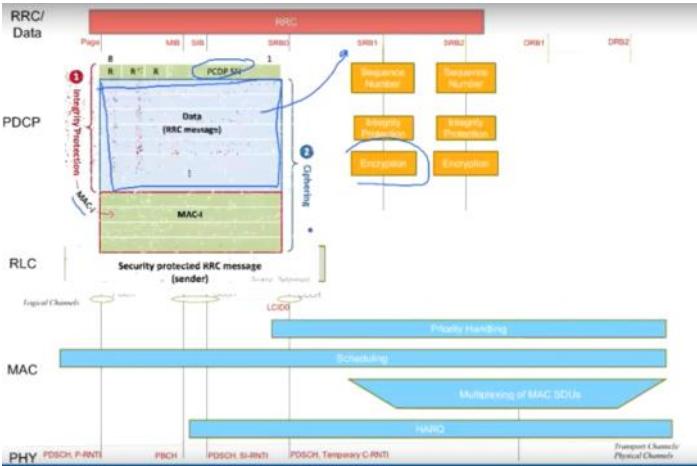
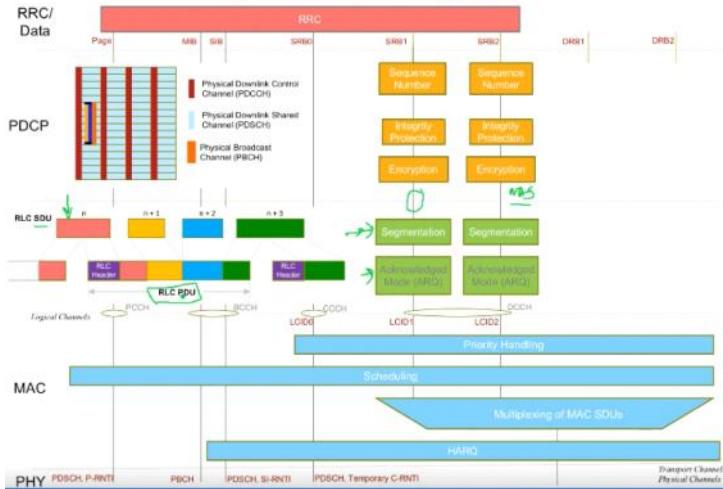
SRB1:Sequence no ,
At PDCP layer integrity protection before encryption

PDCP	->putting SN->take state from SRB1->integrity protection add MAC which is integrity check then take entire thing (DATA+MAC) and do cyphering DRB->SN->ROHC(robust header compression)->encapsulation Separate PDCP entity for each radio bearer and signaling bearer ->so sequence no are independent, and one per data bearer, so RLC don't concatenate packet from drb1 and drb2 But MAC will able to concatenate across this radio bearer and channels
RLC	Segmentation (gen srb2 NAS mess) , Functionality for acknowledge mode ip packet DL dir->unack mode Do error correction and duplicate detection only in AM mode through ARQ TM UM: real time app ,uni dir ,err insensative AM: bidirectional,delay tolterant ,err sen ,non real time, status pdu

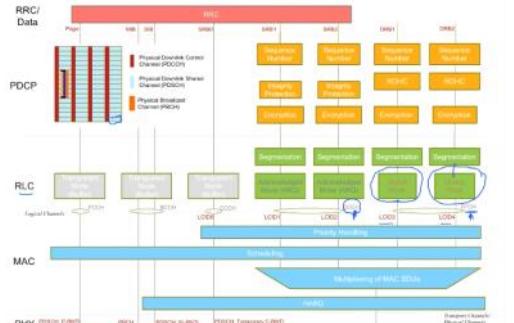
Downlink Protocol Layers and Channel Mapping in eNB



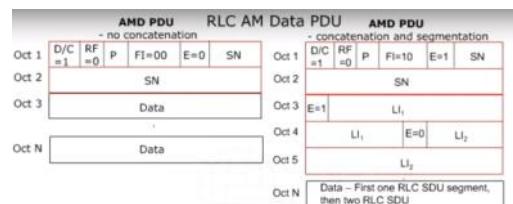
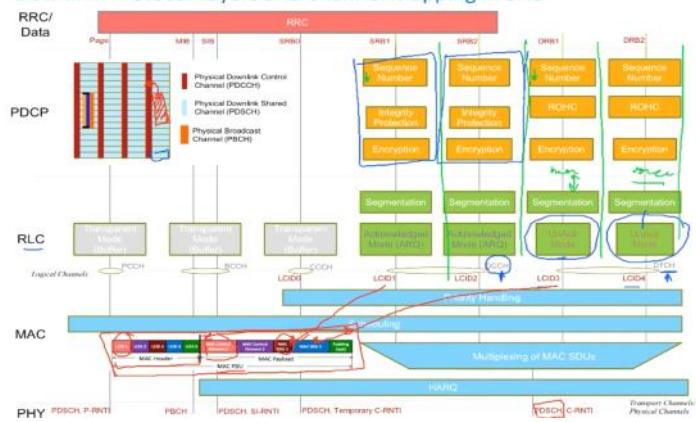
Downlink Protocol Layers and Channel Mapping in eNB



Downlink Protocol Layers and Channel Mapping in eNB

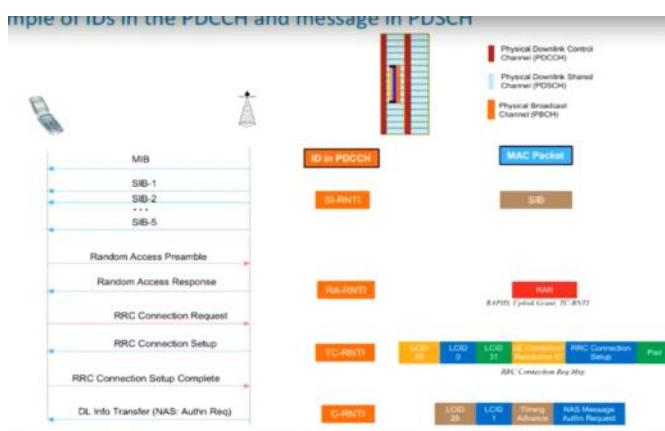
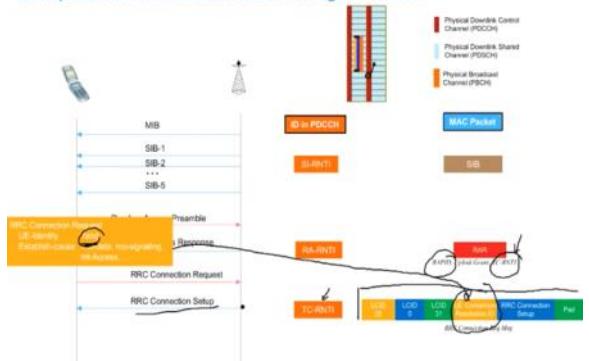
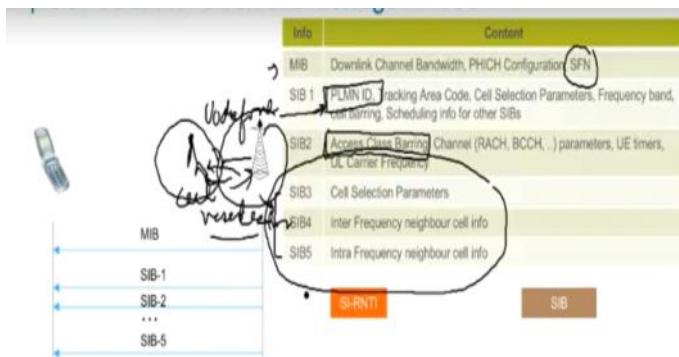


Downlink Protocol Layers and Channel Mapping in eNB



E	Extension bit
SN	Seq no
d/c	Data/ control
Rf	Re-segmentation flag ,0:complete pdu,1:pdu seg
P	Polling bit ,1=status report requested Polling procedure triggers STATUS reporting from receiving side AM RLC entity
Li	Length indicator
ISF	Last segment flag

SIB1	Time:Sent every 20 mili sec but changes every 80 ms Freq: SFN%8 ,sib1 changes to new sib1
PRACH	RAPID(0,63) + RARNTI(sub frame id)
RACH response	RAPID value ,Uplink grant(when u send next message,eg 4-5 sub frame down up,transmit in particular RB) ,TC-RNTI(radio label identity)
RRC connection req	UE identity,rand,Estb cause: mobile ori signalling



- Cell Reference Signals
 - Known reference signals are inserted at regular intervals within the OFDM time-frequency grid.
 - There are four resource elements per resource block that are dedicated to Reference Signal.

CRS	<p>Cell reference signal, inserted at regular interval in OFDMtime -freq grid</p> <p>There are 4 resource ele per resource block</p> <p>Cell ref sig transmitted much high power then other, and location depends on physical layer cell identity(PSS,SSS)</p> <p>The reference signals are derived from the product of a two-dimensional pseudo-</p>	
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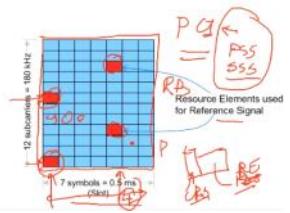
There are 4 resource ele per resource block
Cell ref sig transmitted much high power than other, and location depends on physical layer cell identity(PSS,SSS)

The reference signals are derived from the product of a two-dimensional pseudo-random sequence and a two-dimensional orthogonal sequence. There are 170 different pseudo-random sequences corresponding to 170 cell-identity groups, and three orthogonal sequences each corresponding to a specific cell identity within the cell identity group

Resource block 12 sub carrier, 7 OFDM sym=slot=5ms

Reference Signal.

- The location of Reference Signals depends on the Physical layer cell identity of the cell.
- The Primary and Secondary Synchronization Signals the Physical Layer Cell Identity

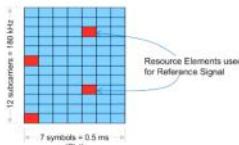


Reference Signal Received Power (RSRP)

- The RSRP is the average power (in watts) received from a single Reference Signal (RS) resource element
- RSRP measures only the RS power and excludes all noise and interference power.
- Knowledge of absolute RSRP enables mobile to calculate downlink path-loss.
- The maximum RSRP is based on maximum input power to UE of -25dBm (0.0032 mWatts). In 1.4 MHz BW with 6 RBs (72 Resource Elements), max RSRP is -44 dBm.
- The minimum value is -140 dBm (has 6 dBm of margin from minimum possible received power at UE).

$$RSRP = \frac{1}{K} \sum_{k=1}^K P_{rs,k}$$

where, $P_{rs,k}$ is the estimated received power (in Watts) of the k th Reference Signal Resource element.

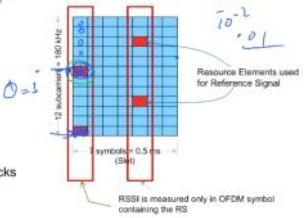


Measurement 2: Reference Signal Received Quality (RSRQ)

- RSRQ does not give an indication of signal quality, i.e. the strength of the reference signal compared to overall energy in the channel (aka received signal strength indicator (RSSI))
- The RSSI parameter represents the entire received power including the wanted power from the serving cell as well as all co-channel power and other sources of noise.
- Measuring RSRQ becomes particularly important near the cell edge when decisions need to be made, regardless of absolute RSRP, to perform a handover to the next cell.
- The maximum value of RSRQ is -3 dB (One reference signal has 50% energy in the RB)
- The minimum value of reported RSRQ is -19.5 dB. (One reference signal RE has only 1% of energy in RB)

$$RSRQ = \frac{RSRP}{(RSSI / N_{RB})}$$

where N_{RB} is the number of Resource blocks ($N_{RB} = 6$ for 1.4MHz Bandwidth)



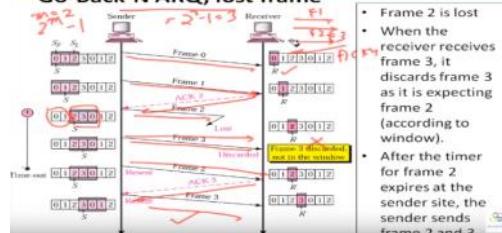
References

- Specifications:
 - TS 36.300: RAN Architecture
 - TS 36.331: RRC
 - TS 36.323: PDCP
 - TS 36.322: RLC
 - TS 36.321: MAC
- Other References:
 - LTE in Bullets ← color
 - www.sharetechnote.com ↗ informative communication
 - www.youtube.com/lte4g ↗ informative communication

ARQ



Go-Back-N ARQ, lost frame



- Frame 2 is lost
- When the receiver receives frame 3, it discards frame 3 as it is expecting frame 2 (according to window).
- After the timer for frame 2 expires at the sender site, the sender sends frame 2 and 3.

S1ap

16 March 2018 07:05

The S1 Application Protocol (S1AP) provides the control plane signalling between E-UTRAN and evolved packet core (EPC). The used interface is S1-MME which is located between eNB and MME. The services provided by the S1AP are divided to UE associated and non UE associated. The UE associated services are related to one UE and the non UE associated services are related to the whole S1-MME interface. The S1AP is transferred over SCTP.

From <<https://www.synopsys.com/software-integrity/security-testing/fuzz-testing/defensics/protocols/s1ap-server.html>>

S1AP messages:

- Handover Required
- Handover Request Acknowledge
- Handover Failure
- Handover Cancel
- E-RAB Setup Response
- E-RAB Modify Response
- E-RAB Release Response
- Initial Context Setup Response
- Initial Context Setup Failure
- Reset
- S1 Setup Request
- UE Context Modification Response
- UE Context Modification Failure
- eNB Configuration Update
- Handover Notification
- E-RAB Release Indication
- Initial UE Message
- Uplink NAS Transport
- NAS non delivery indication
- Error indication
- UE Context Release Request
- Uplink S1 CDMA2000 Tunneling
- UE Capability Info Indication
- eNB Status Transfer
- eNB Direct Information Transfer
- eNB Configuration Transfer
- Uplink UE Associated LPPa Transport
- Uplink Non UE Associated LPPa Transport

Tested NAS messages

- Attach complete
- Attach request
- Authentication failure
- Authentication response
- Detach request
- EMM status
- Extended service request

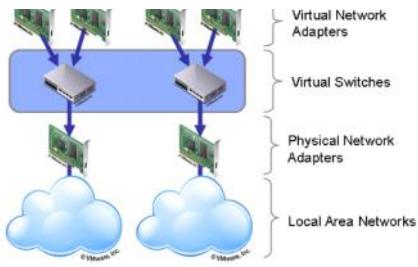
Security mode complete
Security mode reject
Security protected NAS message
Service request
Tracking area update complete
Tracking area update request
ctivate dedicated EPS bearer context accept
ctivate dedicated EPS bearer context reject
ctivate default EPS bearer context accept
earer resource allocation request
Bearer resource modification request
Deactivate EPS bearer context accept
ESM information response
Modify EPS bearer context accept
Modify EPS bearer context reject
PDN connectivity request

Handovers

23 February 2018 12:10

X2: between enbs

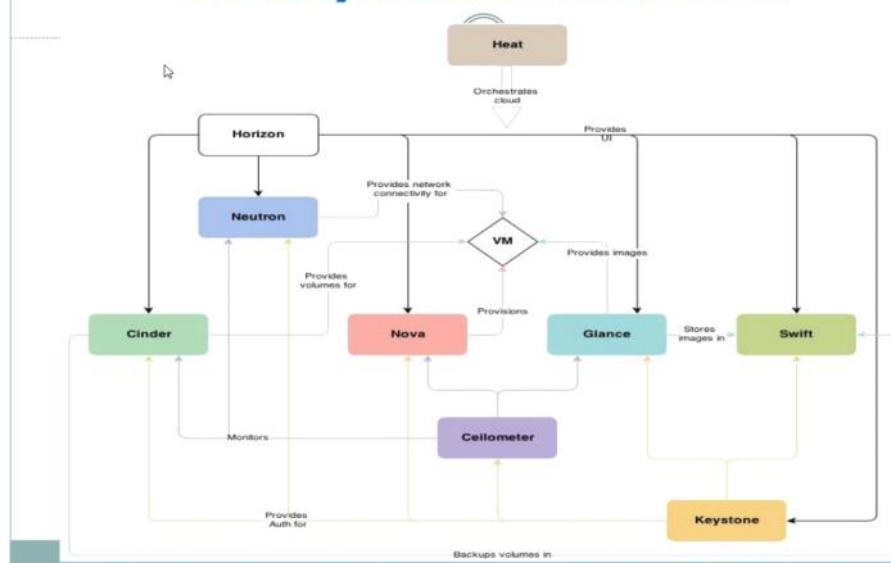
S1: between enb and SgNB



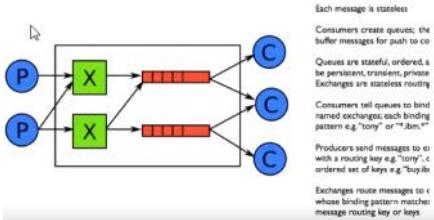
REST Methods

HTTP Method	Action	Examples
GET	Obtain information about a resource	http://example.com/api/orders (retrieve order list)
GET	Obtain information about a resource	http://example.com/api/orders/123 (retrieve order #123)
POST	Create a new resource	http://example.com/api/orders (create a new order, from data provided with the request)
PUT	Update a resource	http://example.com/api/orders/123 (update order #123, from data provided with the request)
DELETE	Delete a resource	http://example.com/api/orders/123 (delete order #123)

Conceptual Architecture



RabbitMQ/QPID - Everything is a message



Best Practice – Log Management

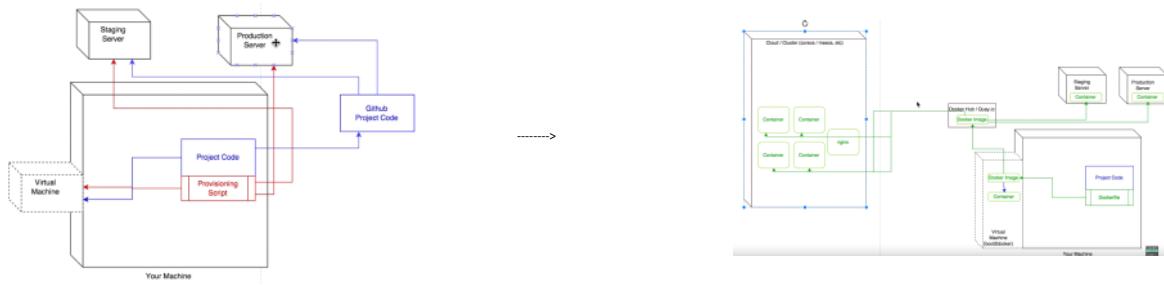
- Use log management and indexing
 - Logstash with Kibana
 - ElasticSearch
 - rsyslogd
- Allows to easily follow complete message flows that go across multiple servers and services



How To Watch The Movement

Role type	Service	Log location
Cloud controller	nova-*	/var/log/nova
Cloud controller	glance-*	/var/log/glance
Cloud controller	cinder-*	/var/log/cinder
Cloud controller	keystone-*	/var/log/keystone
Cloud controller	neutron-*	/var/log/neutron
Cloud controller	horizon	/var/log/apache2/
All nodes	misc (swift, dnsmasq)	/var/log/syslog
Compute nodes	libvirt	/var/log/libvirt/libvirtd.log
Compute nodes	Console (boot up messages) for VM instances	/var/lib/nova/instances/<instance_id>/console.log
Block Storage nodes	cinder-volume	/var/log/cinder/cinder-volume.log

- Logs, logs and more logs
- Each Openstack Service has its own logs
- To follow a flow, or to find problems you often need to look across several logs



RSTP

17 February 2018 01:33

RSTP evolved from the original STP IEEE 802.1D protocol to provide faster spanning-tree reconvergence after a switch port, switch, or LAN failure. Where STP took up to 50 seconds to respond to topology changes, RSTP responds to changes within the timeframe of three hello BPDUs (bridge protocol data units), or 6 seconds

why RSTP:

layer 2 topological loop

1. mac address table corruption
2. broadcast storm unknown unicast and broad cast (switch process)

HELLO time:

Specifies, in seconds, how often the system broadcasts HELLO frames to other members of the spanning tree.

Transmit Hold Count

Specifies the maximum number of spanning tree frames the system can transmit on a port within the Hello Time interval.

Forward Delay

Specifies, in seconds, the length of time for which an interface is blocked from forwarding network traffic

after the spanning tree topology has been modified. This property is more useful for STP than RSTP or MSTP.

Maximum Age

Specifies, in seconds, the length of time for which spanning tree information from other bridges is considered valid.

When you change the value of the Maximum Age option, you change the amount of time, in seconds, that spanning tree information received from other bridges is considered valid. The default value is 20, and the valid range is 6 to 40.

Note that when running RSTP, you must maintain the following relationships between the Maximum Age and the Hello Time and Forward Delay options:

Maximum Age $\geq 2 * (\text{Hello Time} + 1)$

Maximum Age $\leq 2 * (\text{Forward Delay} - 1)$

Destination mac: 01:80:c2:00:00:00

Protocol version identifier: 0->stp; 2->rstp

Default bridge priority: 32,768

Bridge id: bridge id(16 bit)+mac(48 bit)

Root selection:

lowest bridge ID becomes the root bridge

Port Roles Determine Participation in the Spanning Tree

Root port:

The port closest to the root bridge (has the lowest path cost from a bridge). This is the only port that receives frames from and forwards frames to the root bridge.

Designated port:

The port that forwards traffic away from the root bridge toward a leaf. A designated bridge has one designated port for every link connection it serves. A root bridge forwards frames from all of its ports, which serve as designated ports

Alternate port

A port that provides an alternate path toward the root bridge if the root port fails and is placed in the discarding state. This port is not part of the active spanning tree, but if the root port fails, the alternate port immediately takes over.

Backup port

A port that provides a backup path toward the leaves of the spanning tree if a designated port fails and is placed in the discarding state. A backup port can exist only where two or more bridge ports connect to the same LAN for which the bridge serves as the designated bridge. A backup port for a designated port immediately takes over if the port fails.

Disabled port

Port States Determine How a Port Processes a Frame

Discarding—The port discards all BPDUs. A port in this state discards all frames it receives and does not learn MAC addresses.

Learning—The port prepares to forward traffic by examining received frames for location information in order to build its MAC address table.

Forwarding—The port filters and forwards frames. A port in the forwarding state is part of the active spanning tree.

In stp :

Blocking - A port that would cause a switching loop if it were active. No user data is sent or received over a blocking port, but it may go into forwarding mode if the other links in use fail and the spanning tree algorithm determines the port may transition to the forwarding state. BPDU data is still received in blocking state. Prevents the use of looped paths.

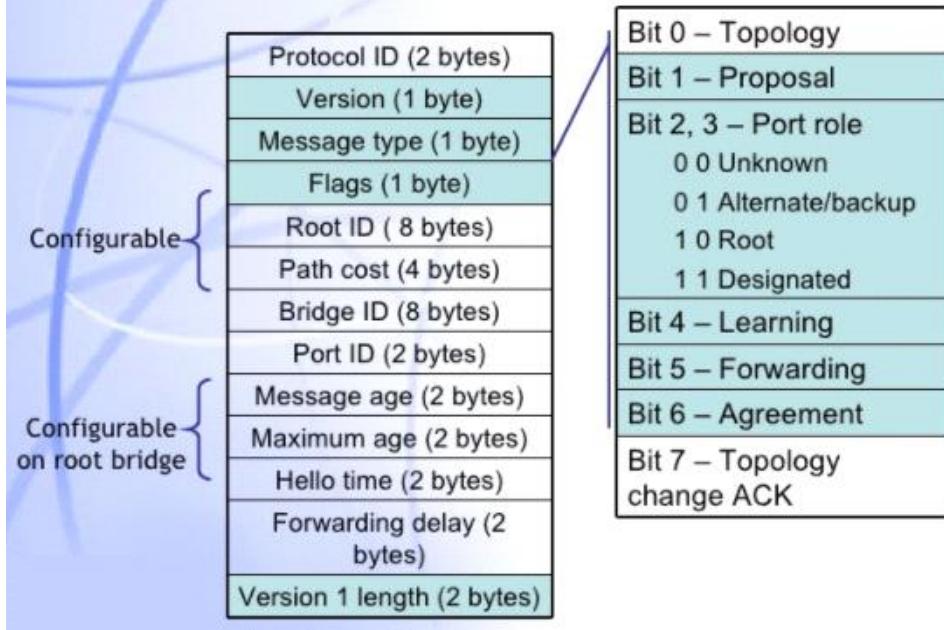
Listening - The switch processes BPDUs and awaits possible new information that would cause it to return to the blocking state. It does not populate the MAC address table and it does not forward frames.

Learning - While the port does not yet forward frames it does learn source addresses from frames received and adds them to the filtering database (switching database). It populates the MAC address table, but does not forward frames.

Forwarding - A port receiving and sending data in Ethernet frames, normal operation. The Forwarding port monitors incoming BPDUs that would indicate it should return to the blocking state to prevent a loop.

Disabled - A network administrator has manually disabled a switch port

RSTP BPDU Format



```

# Frame 54: 52 bytes on wire (416 bits), 52 bytes captured (416 bits)
# IEEE 802.3 Ethernet
#   Destination: Spanning-tree-(for-bridges)_00 (01:80:c2:00:00:00)
#     Address: Spanning-tree-(for-bridges)_00 (01:80:c2:00:00:00)
#       .... .1 .... .... .... = IG bit: Group address (multicast/broadcast)
#       .... .0. .... .... .... = LG bit: Globally unique address (factory default)
#   Source: Nokiasie_45:62:7c (00:40:43:45:62:7c)
#   Length: 38
# Logical-Link control
#   DSAP: Spanning Tree BPDU (0x42)
#   IG Bit: Individual
#   SSAP: Spanning Tree BPDU (0x42)
#   CR Bit: Command
#   Control field: u, func=UI (0x03)
#     000.00.. = Command: Unnumbered Information (0x00)
#     .... .11 = Frame type: Unnumbered frame (0x03)
# Spanning Tree Protocol
#   Protocol Identifier: Spanning Tree Protocol (0x0000)
#   Protocol Version Identifier: Spanning Tree (0)
#   BPDU Type: Configuration (0x00)
#   BPDU flags: 0x00
#     0.... .... = Topology Change Acknowledgment: No
#     .... .0 = Topology Change: No
#   Root Identifier: 32768 / 0 / 00:40:43:45:62:7c
#     Root Bridge Priority: 32768
#     Root Bridge System ID Extension: 0
#     Root Bridge System ID: 00:40:43:45:62:7c
#     Root Path Cost: 0
#   Bridge Identifier: 32768 / 0 / 00:40:43:45:62:7c
#     Bridge Priority: 32768
#     Bridge System ID Extension: 0
#     Bridge System ID: 00:40:43:45:62:7c
#     Port identifier: 0x8006
#     Message Age: 0
#     Max Age: 20
#     Hello Time: 2
#     Forward Delay: 15

```

Rapid spanning tree in use

Bridge configuration

ID priority: [0...15]

Bridge address:

Forward delay: s [4...30]

Max age: s [6...40]

Transmit hold count: [1...10]

Force protocol version: RSTP

Hello time: ms

Port configuration

EIF	<input checked="" type="checkbox"/> RSTP in use	Port number [1...4095]	Priority [0, 16, ...240]	Path cost [1...200000000]
FSM 1	<input checked="" type="checkbox"/>	<input type="text" value="1"/>	<input type="text" value="128"/>	<input type="text" value="20000"/>
FSM 2	<input checked="" type="checkbox"/>	<input type="text" value="2"/>	<input type="text" value="128"/>	<input type="text" value="20000"/>
FTIF 1	<input checked="" type="checkbox"/>	<input type="text" value="3"/>	<input type="text" value="128"/>	<input type="text" value="20000"/>
FTIF 2	<input checked="" type="checkbox"/>	<input type="text" value="4"/>	<input type="text" value="128"/>	<input type="text" value="20000"/>
FTIF 3	<input checked="" type="checkbox"/>	<input type="text" value="5"/>	<input type="text" value="128"/>	<input type="text" value="20000"/>
FTIF 4	<input checked="" type="checkbox"/>	<input type="text" value="6"/>	<input type="text" value="128"/>	<input type="text" value="20000"/>

Port info

EIF	Port ID	Role	State
FSM 1		N/A	N/A
FSM 2		N/A	N/A
FTIF 1		N/A	N/A
FTIF 2		N/A	N/A
FTIF 3		N/A	N/A

OSPF

17 February 2018 03:01

7 stages of OSPF (Darling It Test Some Extremely Large Farts) layer3

link state protocol
hop count and BW

DR
all flood LSA to DR
and DR send/forward LSA to all other routers
all router from nebi data base with BDR also (flood LSA to BDR)

1. down stage	network cmd hello message A sends router id to 244.0.0.5 (multicast router id add for ospf every other is listening to the add).
2. init	B reply with hello unicast message to A
3. Two way stage	establish neighbor rel and fill neb table (DR,BDR elected)
4. EXSTART State	who will start transmission i will start my router id is ---- no i will start my router id is ---- (master/slave rel setup)
5. Exchange start	start transmission here is summary of LSDB(link state database) by DBD(Database decision packet) (exchange DBD)
6. Loading state	router A checks its own database and neighbor summary database ensuring that both are same by LSACK packet if not same or entry is missing send LSR(link state res) to send complete entry for route B LSU link state update () LSAck
7. FULL State	same sync database calculate OSPF cal BW: $10^8 / \text{BW}$

====

LSA types(link state advertisement)

type 1->backbone area ->router area
type 2->standard area -> network lsa (DR generated)
Type 3->Stub area ->Summary LSA (Abi summary route) ->block type 5
Type 4->Totally stub area->Summary loc (ASBR location)->block type 3,4,5
Type 5->not so stubby area(used type 7 router ripped it to type 5)->external LSA (ASBR summary route)

LSA type 1:Router LSA

when plug in new router tell new router is added.
using LSA type 1 every router advertise connected link info to other router belong to same area
based on router id

LSA 2:originated by DR

tells who is DR
it contains list of attached routers on transmit link and include subnet mask

LSA 1 and LSA 2 will not cross there area boundary within the area between the router

LSA type 3:summary lsa

Summary LSA to area outside originating area
initialized by the ABR of originating area

LSA type4:

summary LSA
location of ASBR
Lsa is used to advertise an ASBR to all other areas in the AS
generated by the ABR of originating area and regenerated by all subsequent ABRs.

LSA type 5:

external LSA
advertise by ASBR used to advertise networks from other autonomous systems.
advertise routerid ASBR unchanged throughout as
Type 4 LSA is need to find ASBR

To control the advertisement of external routes into an area, OSPF uses stub areas. By designating an area border router (ABR) interface to the area as a stub interface, you suppress external route advertisements through the ABR. Instead, the ABR advertises a default route (through itself) in place of the external routes and generates network summary (Type 3) link-state advertisements (LSAs).

A stub area that only allows routes internal to the area and restricts Type 3 LSAs from entering the stub area is often called a *totally stubby area*.

The screenshot shows a configuration window for OSPF. At the top, there's a checkbox for "OSPF in use" which is checked. Below it are fields for "Area ID" (empty), "Stub flag" (unchecked), "Router ID" (set to 20.20.20.20), and several configuration parameters for OSPF cost values and hold times. On the right side, there's a section for "Traffic shaping type" set to "WFQ". It includes fields for "Total shaper information rate" (0.5 Mbit/s), "Total shaper burst size" (4000 octets), and "Shaper information rate" (1000.0 Mbit/s). There are also fields for "Committed information rate" (0.0 Mbit/s) and checkboxes for "Enable QoS" and "Include Ethernet overhead". Under "OSPFv2 parameters", the "Interface area" is set to "OSPF disabled". Other parameters like "Configured interface cost", "Hello interval", "LSA retransmission interval", "Router dead interval", and "Packet transmission delay" are listed with their respective values. At the bottom, there are checkboxes for "MTU mismatch detection in use" and "OSPF with BFD", and a field for "BFD destination addresses".

OSPF IPv6 LSA Types :

Routing Networks Is Indeed A Great New Learning Initiative

Router LSA – Network LSA – InterArea Prefix LSA – InterArea Router LSA – AS External LSA – Group Membership LSA – NSSA LSA (Not so stubby Area) – Link LSA – IntraArea Prefix LSA

From <<http://zahid-stanikzai.com/mnemonics-for-networking/>>

Selection of DR/BDR

router id :name of router is highest ip address of active physical interface of the router
if logical interface(loopback interface never go down) is configured then highest ip add of logical interface

preference

- 1.manually router id command
- 2.highest loopback (eg:192.168.1.1) (as loopback will not crash unless you off the router)
- 3.highest IP

share LSA directly if point to point
 All will send LSA to designated router on 224.0.0.6(all dr address) if it is broadcast (not directly connected)
 and then designated router send info on 224.0.0.5 (all spf address) to all routers
 network also select a backup designated router.

DR and BDR is elected per multicast network not area.

BY default priority of router is one.
 if priority is zero it will not take part in DR and BDR election

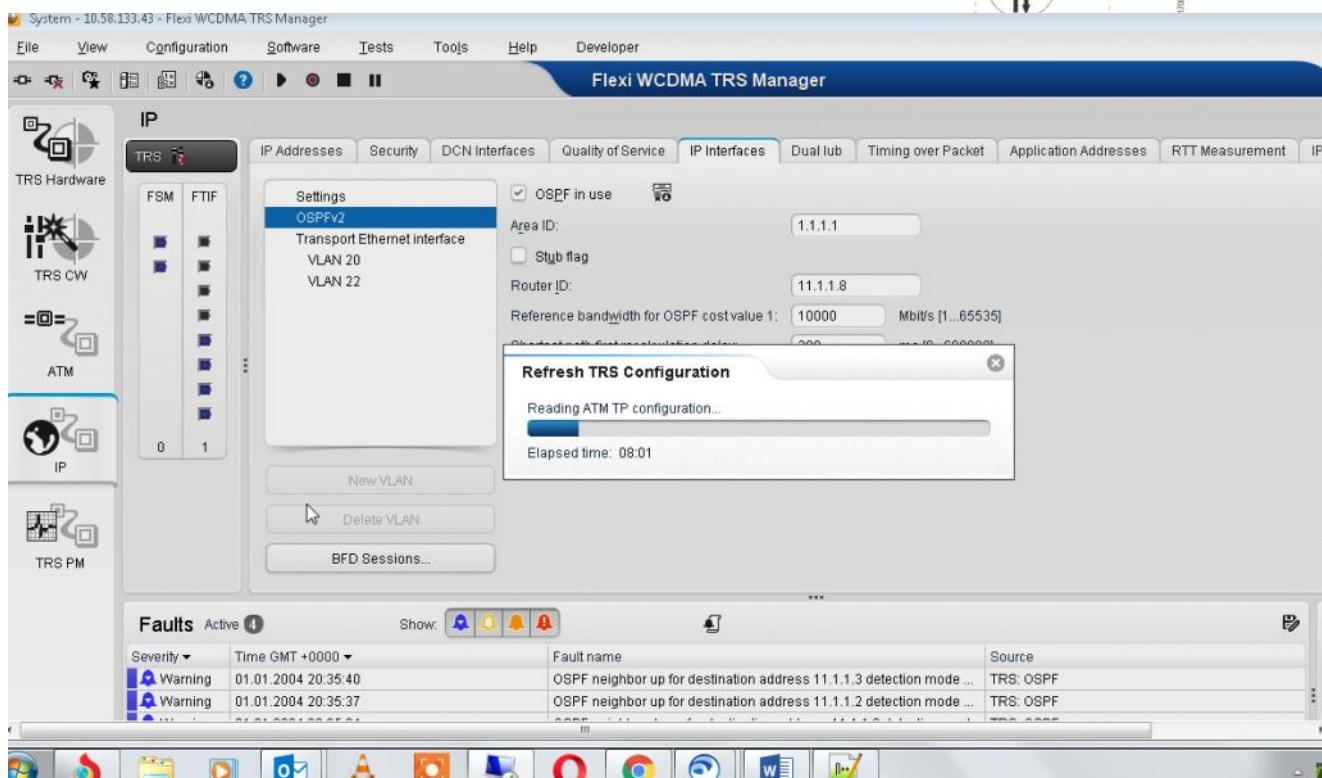
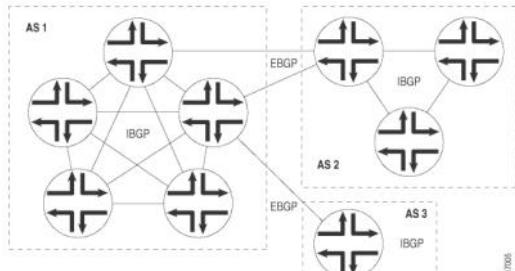
DIJKSTRA's ALGO:

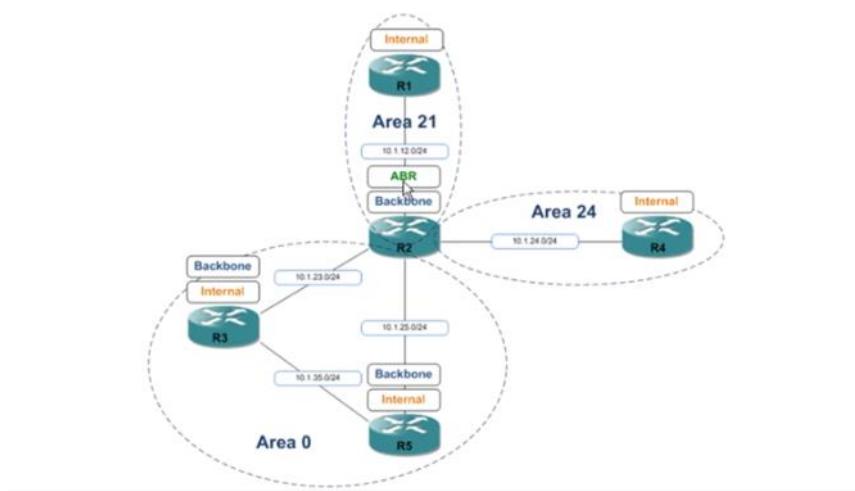
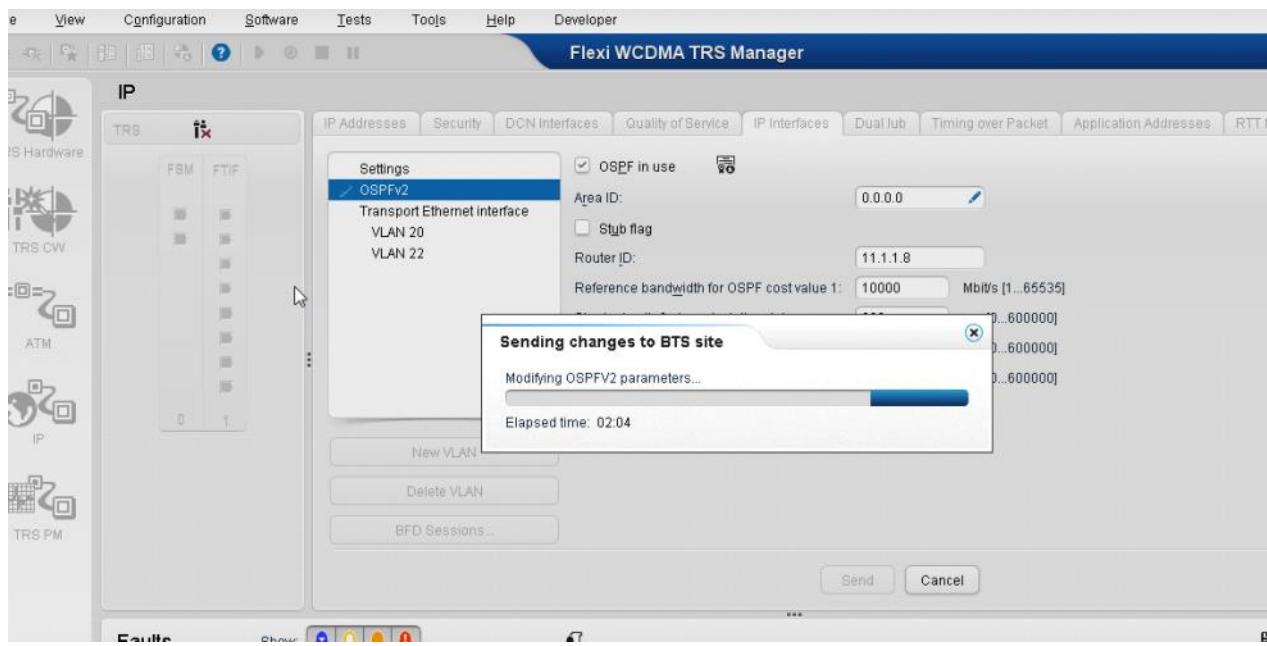
First select one as source (0) all other path(infinity)-->find path to all neigbour vertexes
 -->go to second vertex(with shortest path value) self path+ other path from nei node
 if path to other node is short the prev update it

select the shortest city first

[d(0,2)]

Figure 1: ASes, EBGP, and IBGP

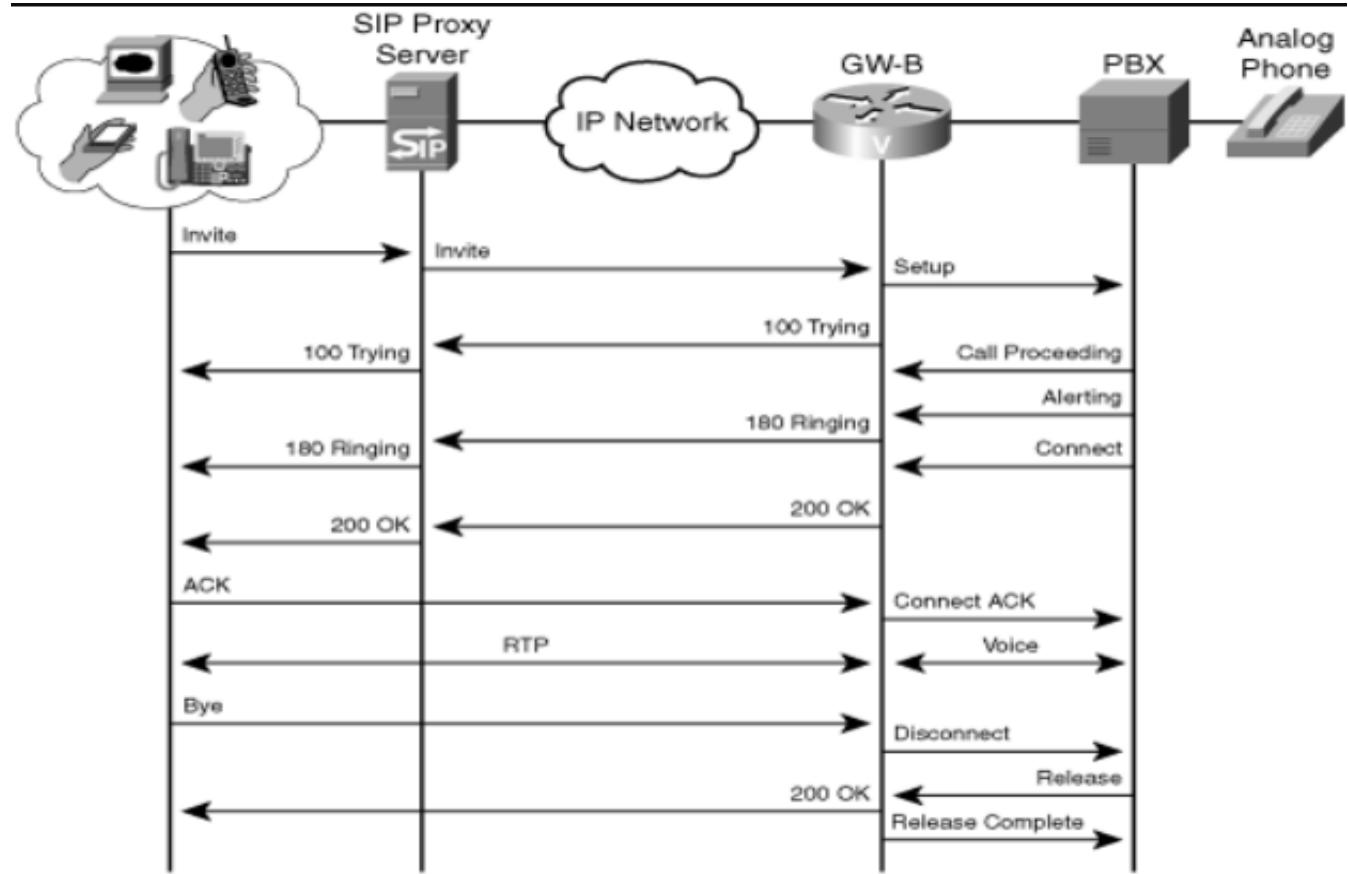
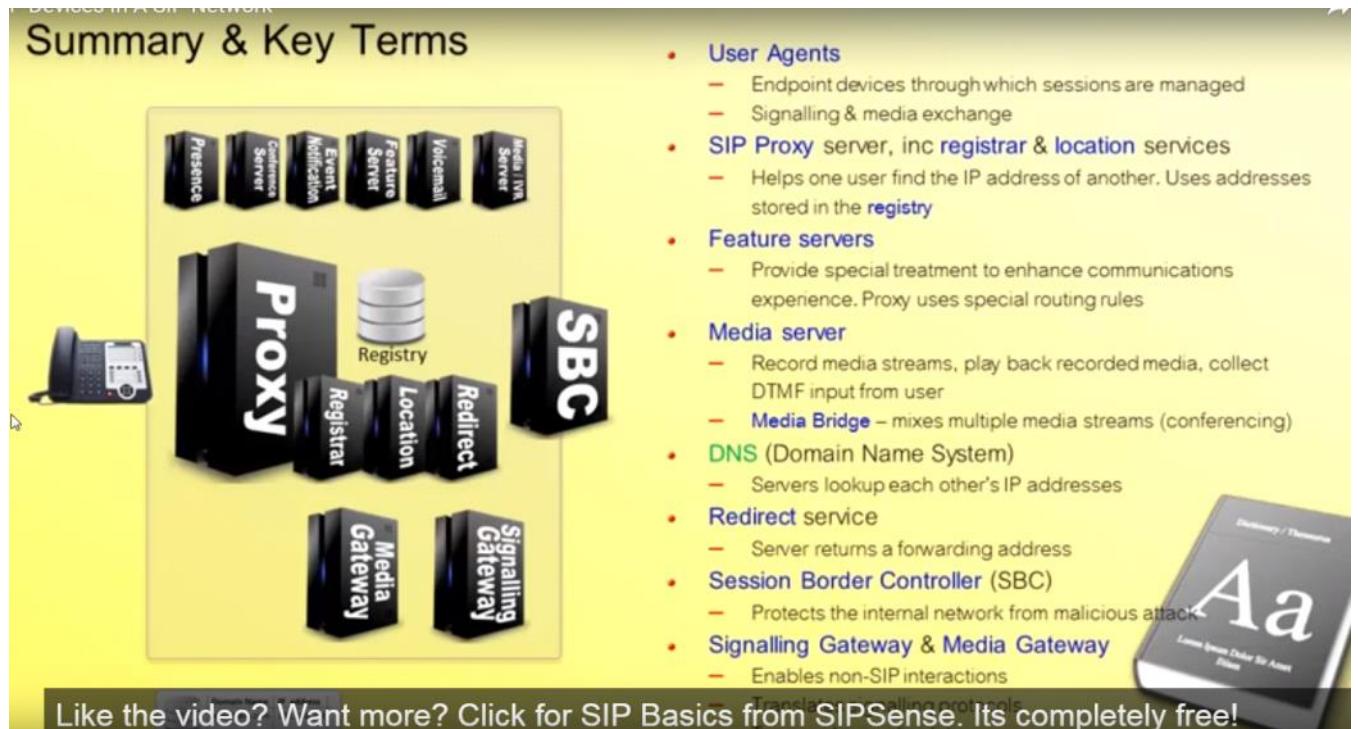




SIP

16 February 2018 22:50

Sip is a signalling protocol to setup ip communication





Invite: All calls start with an invite message. The direction of the arrow will show who originated the call. In this case it was originated by the PBX. Also note the letters SDP. This packet included codec options for setting up the voice path. The PBX informs the carrier that it can accept g729, g711A, and g711U.

100 Trying: This is a status message that states that something is happening to begin processing this call. It is not a confirmation message, that the invite has been received by the remote party.

180 Ringing: This is a confirmation that the invite message was received and that the user is being notified of the call.

183 Session Progress: In this packet the carrier has selected the g729 codec and is informing the PBX of the decision. This is codec choice is usually based on a preprogrammed priority table. This message could indicate to the PBX to begin the ring back tone.

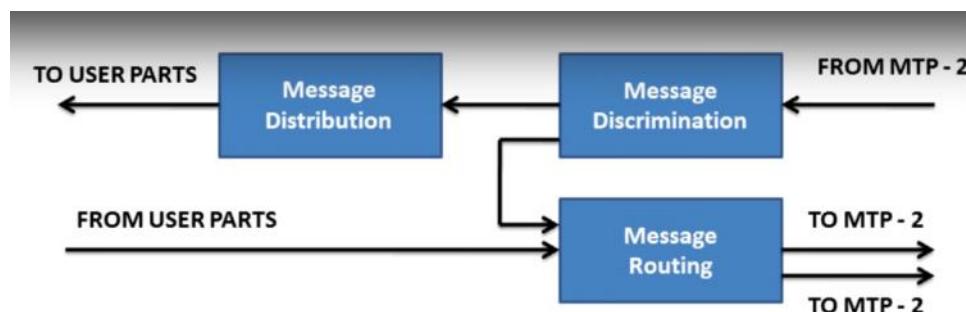
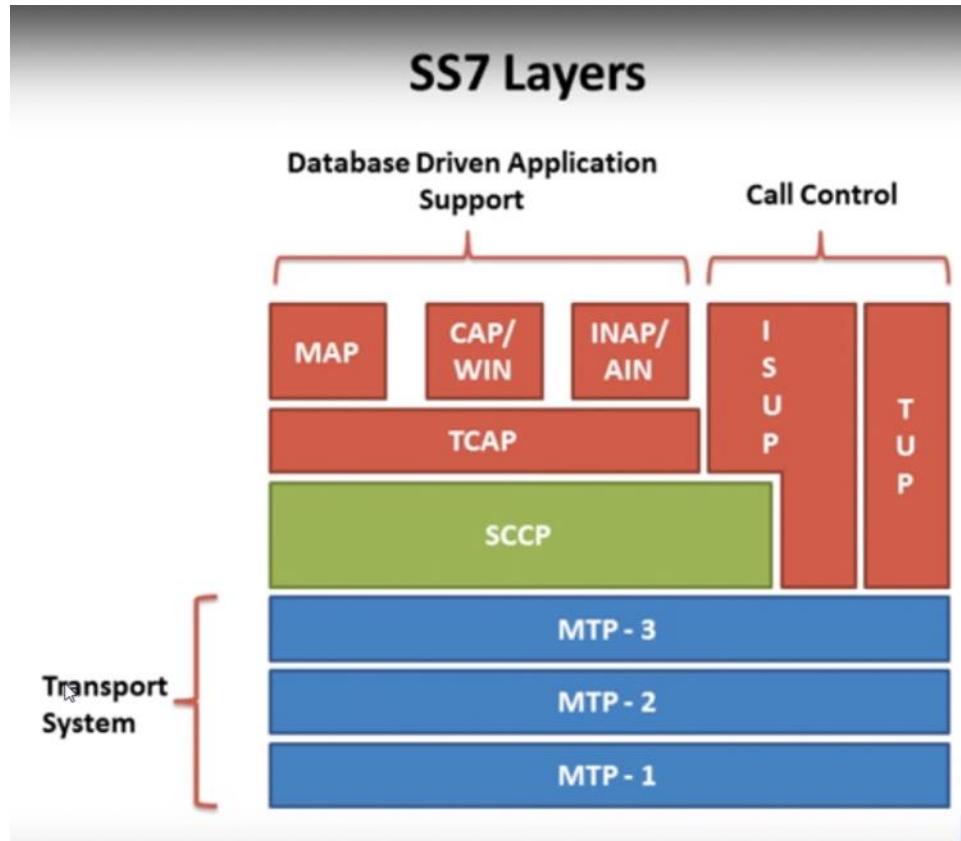
200 OK SDP: This message provides acknowledgement of the invite and provides the media properties..

ACK: This message acknowledges that the 200 OK message was received by the PBX.

RTP: This is the audio stream between the carrier and your media gateway.

Bye: This is a message to signal the carrier that I have terminated the call from the PBX side. The direction of the arrow can be used to determine who originated the termination of the call.

200 OK: This is the acknowledgement from the carrier, that the Bye message was received.



Message Handling is Based Upon :

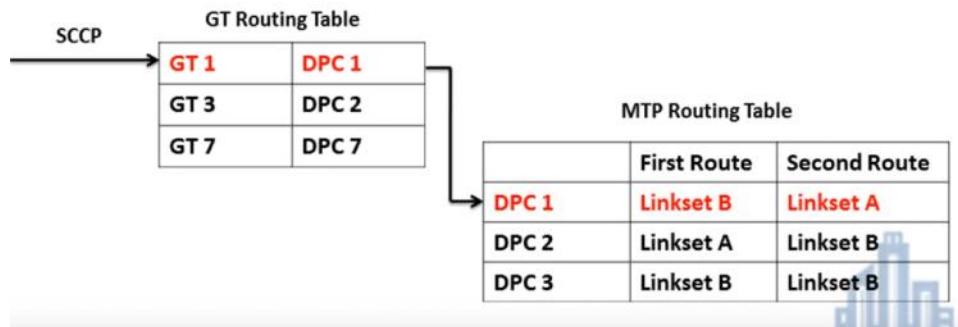
- SIO – Service Information Octet – NI, SI.
- Routing Label – OPC, DPC, Signalling link Selection.
- Some Important Service Indicators (SI)
 - SCCP – 0011
 - TUP – 0100
 - ISUP – 0101



Additional Features of SCCP

- More Flexible Routing based on GT (Global Title)
- Can be used to run OSI Layers
- Support for TCAP,MAP Layers

Global Title Translation



3. User Parts

- ISUP – ISDN User Part
- TUP – Telephone User Part
- TCAP - Transaction Capabilities Application Part
 - AIN
 - INAP
 - CAP
 - MAP

Ethernet header

01 June 2018 16:03

Destination
source
vlan tag:
type:

Preamble:
SOF

Des mac:48 byte --> manufacture ID(24bits) and serial no() L2 add/hardware address
source mac
type: 2 byte: message to network layer (0x800:IPv4, 0x86DD:IPv6)
data : varies
FCS:32 bits
uses CRC for Hash ,used for error checking

value match -> has integrity

VLAN:
virtual local area network

problem : host in nw(switches) 1 cannot--> connect to network 2
Vlan: connect both in one switch, logical segmentation of networks

possibility to communicate between vlan --<>-- needed router

VTP:
switch->
VTP server :
make a change revision no will be change or inc

VTP client : synchronize with server

VTP transparent: will not sync

Domain name (all switch belong to same switch)
problem
If VTP client deattach (and all vlan is deleted) and attach in network then may advertise revision higher than server and then server sync with it and delete all vlan

IPv4

23 February 2018 04:54

IPv4 is a connectionless protocol for use on packet-switched networks.

It operates on a best effort delivery model

the most significant bits are the network prefix, which identifies a whole network or subnet, and the least significant set forms the host identifier, which specifies a particular interface of a host on that network

Version

IP Header Length : generally 20bytes to 60 (variation is possible) ipv6:40bytes(fixed)

Differentiated Services Code Point (DSCP)(6 bit field)

total length: ip and its payload (0-65535 depends on MTU 1500)

(it is different from total frame size why?

618-604=18 bytes) it will not include ethernet header

DS mac:6bytes, SC mac:6bytes, Type:2bytes(0800 -->IP)

identification: unique no to each packet

flags:

reserved :01

x:0x80 (evil bit)

DF : 0->fragmentation, 1->don't fragment

MF: 0->last segment, 1->more fragment is required

fragment offset

router broke Ip packet and reassemble at far end.

each fragment have same identification no but different offset no(no of bytes)

Time to live: (default windows 128,unix 255)

if time to live value is zero the router will delete the packet and send ICMP message to a originator

ttl value exceeded

protocol:

UDP(17)

ICMP(1,2)

TCP(6)

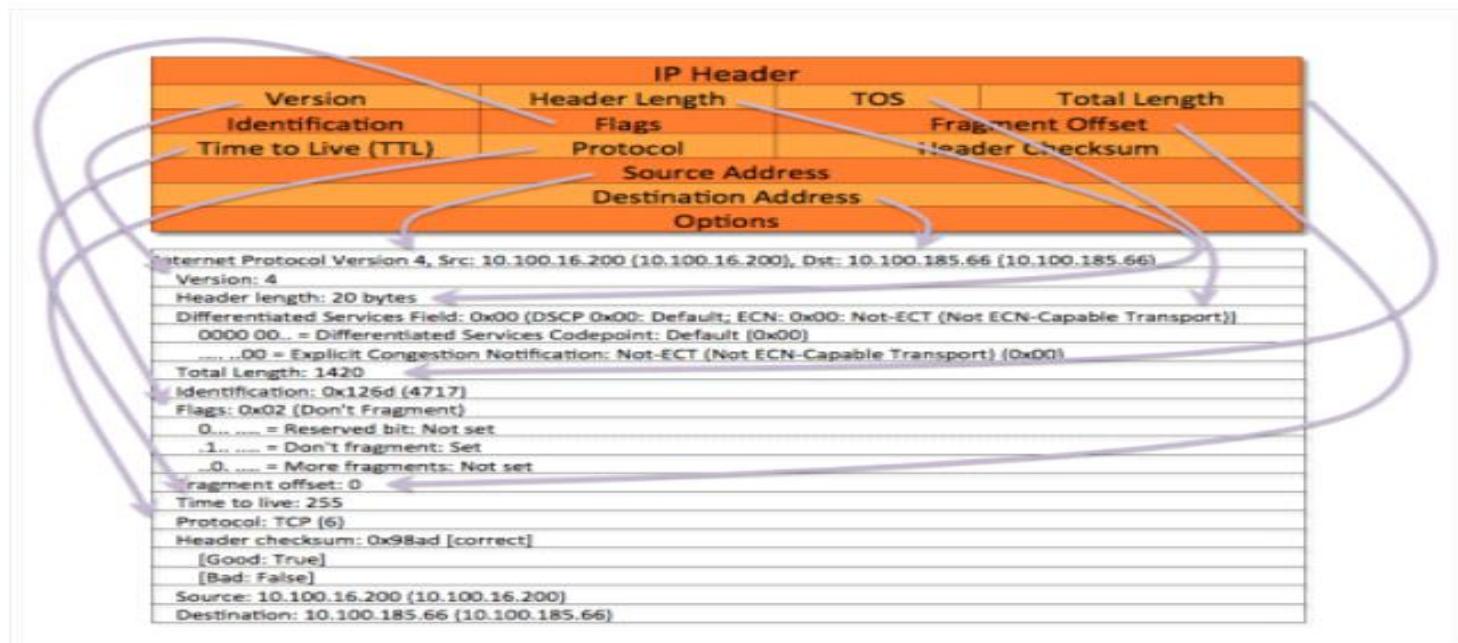
IGRP()

OSPF(89)

ESP(15)

AH(11)

header checksum: --->removed from ipv6(flag,offset,identification)



Ipv6

23 February 2018 04:54

TWAMP

23 February 2018 05:00

TWAMP | Two way active measurement protocol

Two-way measurements are helpful because round-trip delays do not require host clock synchronization and remote support might be a simple echo function.

1. Enable
2. configure terminal

ip sla server twamp	Device(config)# ip sla server twamp Configures the device as a TWAMP server and enters TWAMP server configuration mode.
port port-number	Device(config-twamp-srvr)# port 9000 Configures the port to be used by the TWAMP server to listen for connection and control requests.
timer inactivity seconds	Device(config-twamp-srvr)# timer inactivity 300 Configures the inactivity timer for a TWAMP control session

For reflector: no port no is define

1. enable
2. configure terminal
3. ip sla responder twamp
4. timeout seconds
5. end

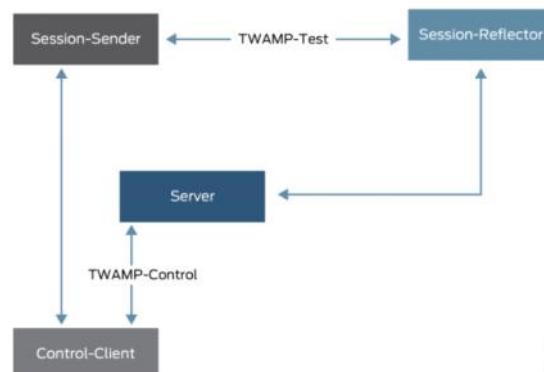
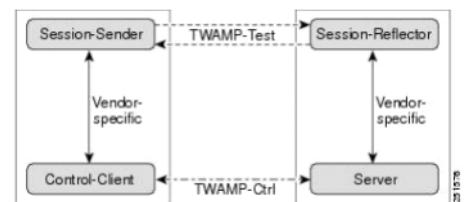


Figure 1: TWAMP Architecture



The screenshot shows the BTS Site Manager interface. The top menu includes File, View, Configuration, Software, Tests, Tools, Antenna, and Help. The left sidebar has icons for BTS Hardware, TRS Hardware, Commissioning, BTS PM, IP, and TRS PM. The main area has tabs for IP, TRS, IP Addresses, Security, Quality of Service, IP Interfaces, CES over PSN, Iub IP, Timing over Packet, Application Addresses, and RTT Meas. In the IP tab, there's a Sessions section with Session 1 selected. Fields include Source IP address (55.1.1.2 - VLAN 55), Destination IP address (55.1.1.1), Destination port (5018), DSCP (34), Message size (100 octets [69...1400]), 15 min packet loss ratio fault threshold (100.00 % [0.00, 0.01,...100.00]), and 1 min round trip time fault threshold (145666 µs [0, 100,...1000000]). Below this are Lost TWAMP messages, Transmitted TWAMP messages, and Statistics. The Faults tab shows 5 active faults: Minor (BTS time not corrected (0026)) and Major (BTS PNC/Flav Direct interface signalling link failure). The Details section shows State: Active, Started: 13.01.2017 10:12:12 GMT +0530, and Cleared:.

rtamp7_dp_ul.pcap [Wireshark 1.7.0-RL50_120212 (SVN Rev 37656 from /trunk)]

File Edit View Go Capture Analyze Statistics Telephony Tools Internals Help

Filter: Expression... Clear Apply

No.	Time	Source	Destination	Protocol	Length	Sequence number @ Sender	Sequence number @ Reflector	Info
4	0.100394	55.1.1.1	55.1.1.2	UDP/TWAMP	118	893	139742	Session-Refl
5	0.200334	55.1.1.2	55.1.1.1	UDP/TWAMP	118	894	0	Session-Refl
6	0.200542	55.1.1.1	55.1.1.2	UDP/TWAMP	118	894	139743	Session-Refl
7	0.300453	55.1.1.2	55.1.1.1	UDP/TWAMP	118	895	0	Session-Refl

Transmit Timestamp @ Reflector: Jan 1, 1970 00:00:00.000000000 UTC
 [Integer part of seconds: 0]
 [Fractional part of seconds: 0]
 Error Estimate: 0x0005
 [0... = Error Estimate S bit: no notion of external synchronization (0)]
 [.0... = Error Estimate Z bit: MBZ bit (0)]
 [.00 0000 = Error Estimate Scale: 0]
 [. 0000 0101 = Error Estimate Multiplier: 5]
 MBZ: 0x0000
 Receive Timestamp @ Reflector: Not representable
 [Integer part of seconds: 14735]
 [Fractional part of seconds: 610558555]
 Sequence number @ Sender: 894
 Transmit Timestamp @ Sender: Not representable
 [Integer part of seconds: 1484658265]
 [Fractional part of seconds: 2461763584]
 Sender Error Estimate: 0x0005
 [0... = Error Estimate S bit: no notion of external synchronization (0)]
 [.0... = Error Estimate Z bit: MBZ bit (0)]

0000	00 0f bb a9 e6 29 60 a8 fe 4d 5e a3 81 00 20 37) .MA... 7
0010	08 00 45 88 00 64 00 00 40 00 ff 11 0a fc 37 01	.E..d. @....7.
0020	01 01 37 01 01 02 13 9a 13 89 00 50 f9 d9 00 02	.7.....P....
0030	21 df 00 00 00 00 00 00 00 00 05 00 00 00 00 00	I.....

10.58.167.33 - Remote Desktop Connection

rtamp7_dp_ul.pcap [Wireshark 1.7.0-RL50_120212 (SVN Rev 37656 from /trunk)]

File Edit View Go Capture Analyze Statistics Telephony Tools Internals Help

Filter: Expression... Clear Apply

No.	Time	Source	Destination	Protocol	Length	Sequence number @ Sender	Sequence number @ Reflector	Info
16	0.701311	55.1.1.1	55.1.1.2	UDP/TWAMP	118	899	139748	Session-Refl
17	0.801290	55.1.1.2	55.1.1.1	UDP/TWAMP	118	900	0	Session-Refl
18	0.801440	55.1.1.1	55.1.1.2	UDP/TWAMP	118	900	139749	Session-Refl
19	0.901452	55.1.1.2	55.1.1.1	UDP/TWAMP	118	901	0	Session-Refl
20	0.901621	55.1.1.1	55.1.1.2	UDP/TWAMP	118	901	139750	Session-Refl
21	1.001642	55.1.1.2	55.1.1.1	UDP/TWAMP	118	902	0	Session-Refl
22	1.001861	55.1.1.1	55.1.1.2	UDP/TWAMP	118	902	139751	Session-Refl
23	1.101824	55.1.1.2	55.1.1.1	UDP/TWAMP	118	903	0	Session-Refl
24	1.102039	55.1.1.1	55.1.1.2	UDP/TWAMP	118	903	139752	Session-Refl

User Datagram Protocol, Src Port: commplex-link (5001), Dst Port: 5018 (5018)
 Two Way Active Measurement Protocol
 Sequence number @ Reflector: 0
 Transmit Timestamp @ Reflector: Jan 1, 1970 00:00:00.000000000 UTC
 [Integer part of seconds: 0]
 [Fractional part of seconds: 0]
 Error Estimate: 0x0000
 [0... = Error Estimate S bit: no notion of external synchronization (0)]
 [.0... = Error Estimate Z bit: MBZ bit (0)]
 [.00 0000 = Error Estimate Scale: 0]
 [. 0000 0000 = Error Estimate Multiplier: 0]
 MBZ: 0x0000
 Receive Timestamp @ Reflector: Not representable
 [Integer part of seconds: 14735]

0000	60 a8 fe 4d 5e a3 00 0f bb a9 e6 29 81 00 a0 37	'..MA...)....7
0010	08 00 45 88 00 64 00 00 40 00 ff 11 0a fc 37 01	.E..d. @....7.
0020	01 02 37 01 01 01 13 89 13 9a 00 50 00 00 00 00	.7.....P....
0030	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00

o 0.132724	22.1.1.2	22.1.1.1	UDP	118		Source port:
9 0.211555	22.1.1.2	22.1.1.1	UDP	118		Source port:
10 0.211705	22.1.1.1	22.1.1.2	UDP	118		Source port:
11 0.232991	55.1.1.2	55.1.1.1	UDP/TWAMP	118	5545	0 Session-Ref:
12 0.233148	55.1.1.1	55.1.1.2	UDP/TWAMP	118	5545	5545 Session-Ref:
13 0.311196	22.1.1.2	22.1.1.1	UDP	118		Source port:
14 0.311348	22.1.1.1	22.1.1.2	UDP	118		Source port:
15 0.333125	55.1.1.2	55.1.1.1	UDP	118		Source port:
16 0.333272	55.1.1.1	55.1.1.2	UDP	118		Source port:

[.... = Error Estimate Scale: 0]
[.... 0000 0101 = Error Estimate Multiplier: 5]
MBZ: 0x0000

Receive Timestamp @ Reflector: Not representable

[Integer part of seconds: 233729]

[Fractional part of seconds: 3741567299]

Sequence number @ Sender: 5545

Transmit Timestamp @ Sender: Not representable

[Integer part of seconds: 1484042192]

[Fractional part of seconds: 2384527188]

Sender Error Estimate: 0x0005

[0.... = Error Estimate S bit: no notion of external synchronization (0)]

IPSEC

23 February 2018 12:27

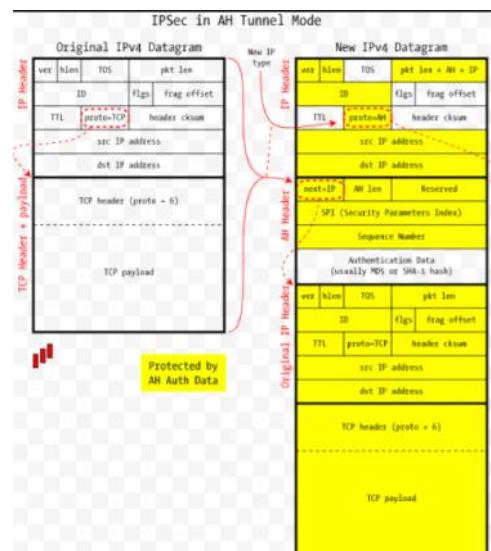
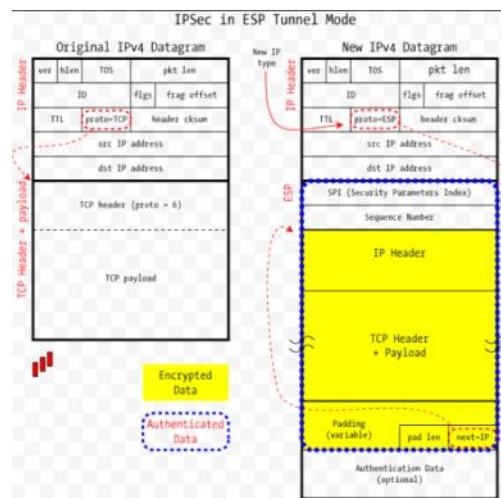
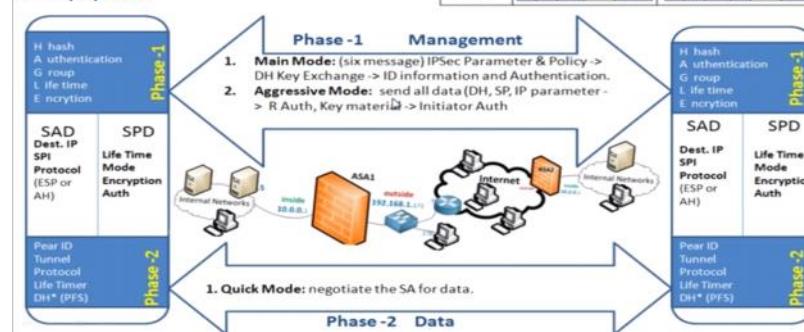
IKE	Internet key exchange,udp500,NAT-T udp 4500,	Key exchange and security parameter negotiation	Phase1:Management Main mode (6 message): ipsec parameter and policy-> DH key exchange(defey helman)-> id info and authentication Aggressive mode: Send all data in 1st message(DH,SP,IP parameter->R auth, key material->initiator auth) Association made here is called IKE association (IZM)	Phase2:DATA Negotiate the SA for data Association made here is called IPSEC association Two unidirectional association are made. One for incoming. Second for outgoing.
ESP	Encapsulating Security payload, PID50	Authentication, encryption, integrity anti reply	Encryption:DES,3DES,RCS	
AH	Authentication header,pid51	Authentication, integrity, and anti-reply	Authentication:Preshared ,digital signature,public key	
SAD	Security association database		Intregity:HMAC-SHA1	
SPD	Security policy database			

IKE (Internet Key Exchange) Key Exchange and security parameter negotiation [ISAKMP & OAKLEY], UDP 500, (NAT-T UDP 4500)

IPSec

ESP (Encapsulating Security Payload) authentication, encryption, integrity, anti reply. PID 50

AH (Authentication Header) authentication, Integrity, anti replay. PID 51



GPON

23 February 2018 18:18

GPON	Gigabit passive optical network
------	---------------------------------

Gigabit: data transmission speed

Passive :no optical power is required in GPON network only end point requires power

Optical n/w: optical/light signal used to carry data

Gpon support

Triple -play service(tv,voip and internet service)

High bandwidth

Long reach (up to 20km)

OLT	Optical line terminal
ODN	Optical distribution network
ONU/ONT	Optical network unit/ termination

Gpon uses optical wavelength division multiplexing so single fiber can be used as uplink and downlink.

Laser of wavelength (λ) uplink=1490nm

Downlink=1310nm

Downstream=transmitted in broadcast manner(AES encryption is used)

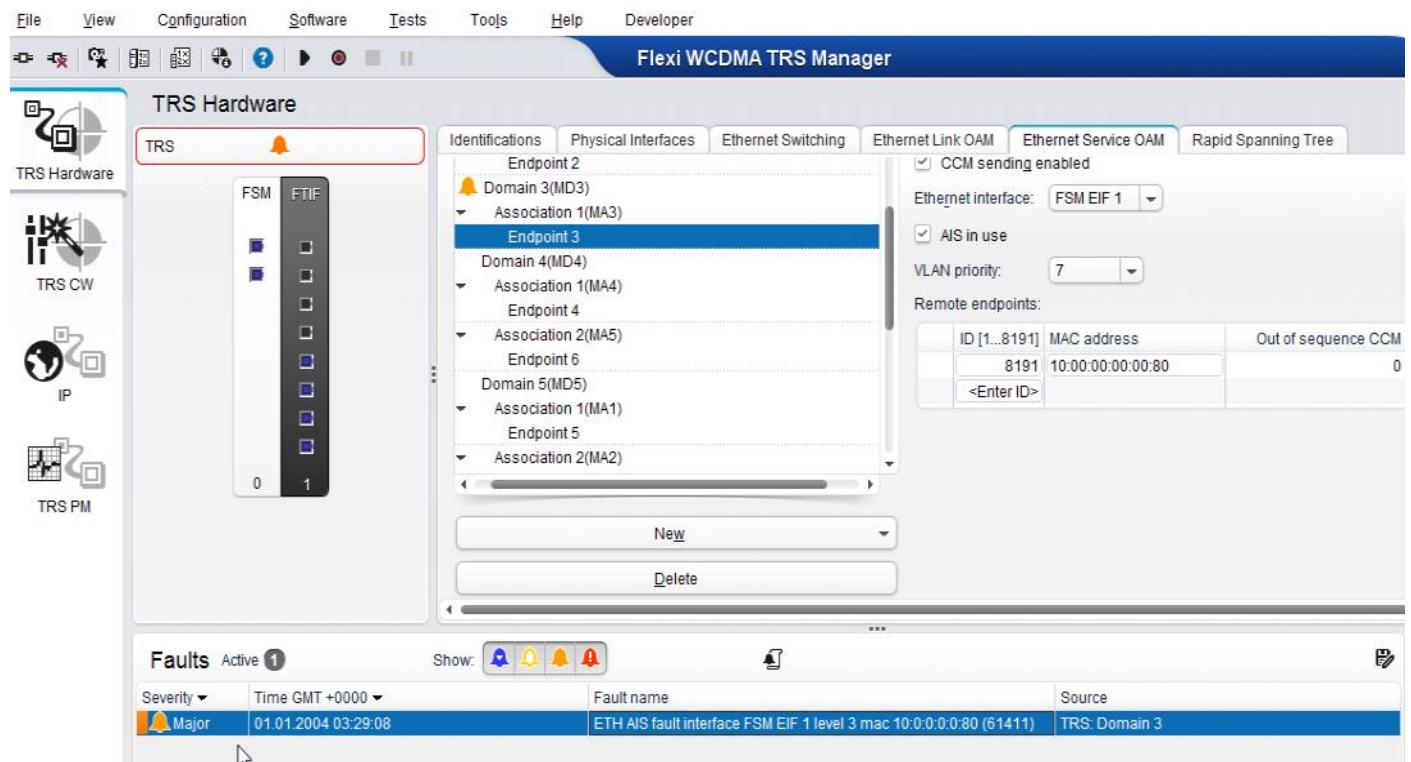
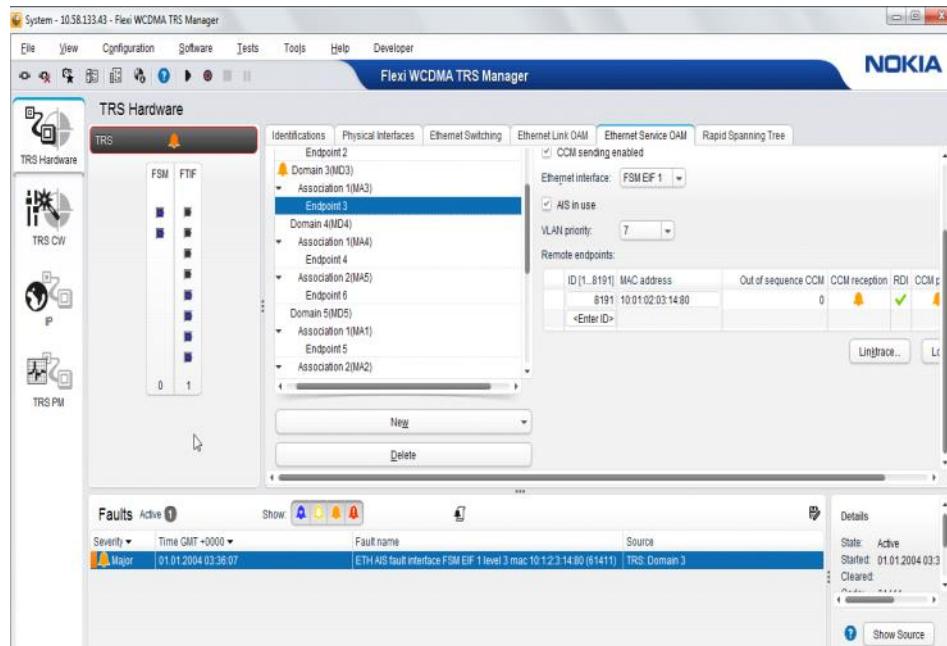
Upstream=packet transmitted in TDMA manner

X GPON: dl->10Gbits/sec

ul:->2.4Gbits/s

Soam

25 February 2018 09:14



TRS

TRS Hardware

FSM FTIF

TRSCW

IP

TRSPM

Ethernet Service OAM in use

Ethernet Service OAM performance monitoring in use

Settings

Domain 1(MD1)

Association 1(MA1)

Endpoint 2

Association 2(MA2)

Endpoint 2

Domain 2(MD2)

Association 1(MA1)

Endpoint 2

Domain 3(MD3)

Association 1(MA3)

Endpoint 3

Domain 4(MD4)

Association 1(MA4)

ID: 2

CCM sending enabled

Ethernet interface: FSM EIF 1

AIS in use

VLAN priority: 7

Remote endpoints:

ID [1...8191]	MAC address	Out of sequence CCM	CCM rec
1	00:00:00:00:00:02	0	0
3	00:00:00:00:00:03	0	0
<Enter ID>			

Faults Active 3

Show:

Severity	Time GMT +0000	Fault name	Source
Major	01.01.2004 17:59:59	CCM fault level - prio 5 cause MDL_MISMATCH (61410)	TRS: NE
Major	01.01.2004 17:59:58	ETH AIS fault interface FSM EIF 1 level 0 mac 00:00:00:00:00:02 (61...)	TRS: Domain 1
Major	01.01.2004 17:59:15	CCM fault level 0 prio 3 cause REMOTE_CCM_DEFECT (61410)	TRS: Domain 1 / Association 1

ICMP

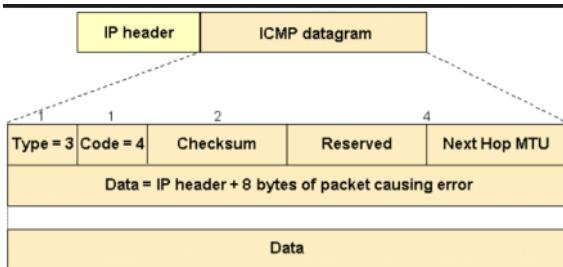
13 March 2018 10:42

ICMP is designed to overcome the following two problems in IP

1. no error reporting
2. lacks a mechanism for queries

ICMP messages are not directly passed to the data link layer. Instead, the message is first encapsulated inside IP datagrams before going to the lower layer.

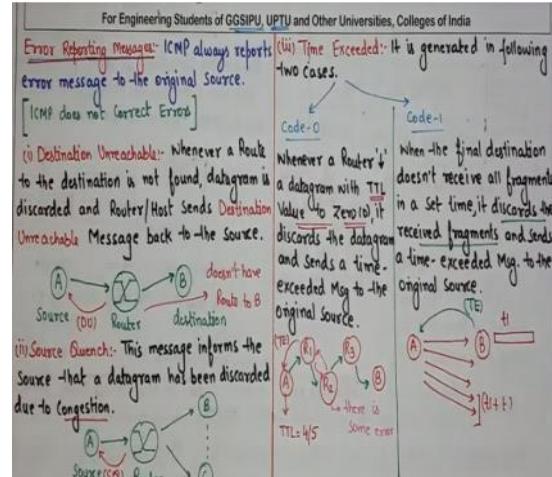
		icmp mes
	ip header	ip data
frame header	frame data	trailer



type	type of message, error or query
code	reason for message type
checksum	error checking
data section	error: finds the original packet with error query: carry extra info

ICMP always reports errors to the original source.

destination unreachable	
Source Quench:	informs source that a datagram has been discarded due to congestion
Time exceeded	code0: TTL value to zero code1: all fragments did not reach in time
parameter problem	code0: error/ambiguity in one of header fields code1: required part of option is missing
Redirection	It is sent from a router to a host on the same local network
ICMP query message	
echo request and reply	these messages are used to determine whether two systems can talk with each other
Timestamp request and reply	sending time = receive - original receiving time = returned - transmit round trip = sending + receiving



Machine learning

22 March 2018 10:32

Machine learning	Using data to solve problem Prediction Iterative correct		
Why	Ability to predict		
where	Search engine, youtube recomm		
Types	Supervised() Output is known Regression Classification	Unsupervised Output is unknown Clustering	Reinforcement
Data as input	Feature engineering	Algo	Output /model
FEATURE engg			
Dimension reduction			

38.300 full 5g spec

Higher fre->sub carrier spacing increase

Within 2.5 full sweep happen

Ss block

Pss max power sss hlf(3db down)

MIB:->sub carrier spacing data

N fft point

Fr1 range: 450mz->6GHz microcell

Fr2 range: 24,5GZ-52.6GHZ 120,240 Analog beam forming only

Fs=120

->fs 60

10000000

-4.84 guard

--

95160 khz

95160/12=793 sub carrier

793-1 =792 scince one dc sub

792/12

=66 max sub carrier

Cell search :section 4.1

Higher freq->less power ->less region->narrow beam

0-3GZ: max beam 4 (ss block)

3-6 : beam 8 {4,8,16,24}+28*n

6:- beam 64

125/14 *4

Prach:

->how far the ue is

Cp prach seq cp

Max cell radius:

Cp /30.72 =18.75 ->cyclic pre

18.75 -

Assume path profile=0

2.34*300/2=350meters

Ss/PBCH block

Tdd:4 ofdm

Fdd:576 sub carreier

576/4->

GNBlogical architecture in5G18A :

•Distributed Unit (gNB-DU) –contains L1, L2-RT, L2-NRT (RLC) baseband processing functions, also Radio

Unit (RU) functions are included in scope of gNB-DU

•Central Unit (gNB-CU) –contains L2-NRT (PDCP) baseband processing functions, OAM, C-Plane

FFT
Sin wave -> 1 fre
square ware->many fre

Sampling rate= subcarrier spacing* fft size :how many sample available in sec

Eailer ->A to D
IQ sample(time domain digital)->Analog cobham

3GPP
The Mobile
Broadband Standard
LTE



The screenshot shows the official 3GPP website. The header features the 3GPP logo and the text "The Mobile Broadband Standard LTE". Below the header is a navigation menu with links to "About 3GPP", "Specifications Groups", "Specifications", "3GPP Calendar", "Technologies", "News & Events", "Home", "Sitemap", and "Contact". A "Cover Story" section is visible on the left, and a "Search" bar is on the right. The main content area displays a large green sidebar with various 3GPP working group names like TSG RAN, TSG CT, TSG SA, etc., each with a brief description. A central column contains news items and a footer with links to recent news stories and awards.

3GPP management standards from wireless telecom R&D are an important another main

Links to some of the news stories that were recently on the home page and news pages.

3GPP Awards for 2017 and Lifetime Achievement

Logical unit	Physical unit	Logical functions	HW virtualization	HW implementation	Location
gNB-DU	RU	• RF	• Non-Virtualized HW	AirScale (MAA)	Radio site, typically with integrated antennas
		• L1 • L2-RT (RLC+MAC)	• Non-Virtualized HW	AirScale (ABIL)	Centralized/remote, has strict real-time requirements (~50us to antenna connector)
	RAU	• L2-NRT (RLC) • CP-RT • OAM	• Non-Virtualized HW	AirScale (ASIK)	
gNB-CU-U	RAC	• L2-NRT (PDCP)	• Virtualized HW	AirFrame	Centralized, no strict real time requirements
gNB-CU-C		• C-Plane	• Virtualized HW	AirFrame	
gNB-CU-M		• OAM	• Virtualized HW	AirFrame	

computer networks

12 May 2018 21:25

class A	1 ->126
class B	128->192
class C	192->126
class D	127 ->120
limited broadcast	255.255.255.255
directed broadcast	1XX.255.255.255

1. A host 200.100.1.1 want to send packet to all host in the same network. what is Source IP and DIP?
200.100.1.1 255.255.255.255
2. for loopback?
DIP: 127.0.0.X
3. How many bits are allocated to NID and HID in 23.192.157.234 address?
23 ->class A
class full and classless->depends on /no
here class full
NID=8 bits, HID =24 bits
In subnetting: big n/w to small n/w
applied for single n/w
improve security. (can mention access control list for a subnet)
bits are borrowed from host portion not n/w portion. NID|HID
In supernetting : small n/w to big n/w
security decrease.

IMS

18 May 2018 06:27

EUTRAN<-- IP--->EPC-> (packet based)does not have circuit switching
|- IMS(ip multimedia subsystem)->service network

IMS is SIP based.

IMS (5 component):

CSCF setup processing

Volte:(1-2 sec)
HD voice call
video calling

Volte handset

SIM(GSM)--USIM(LTE)---ISIM(VoLTE)

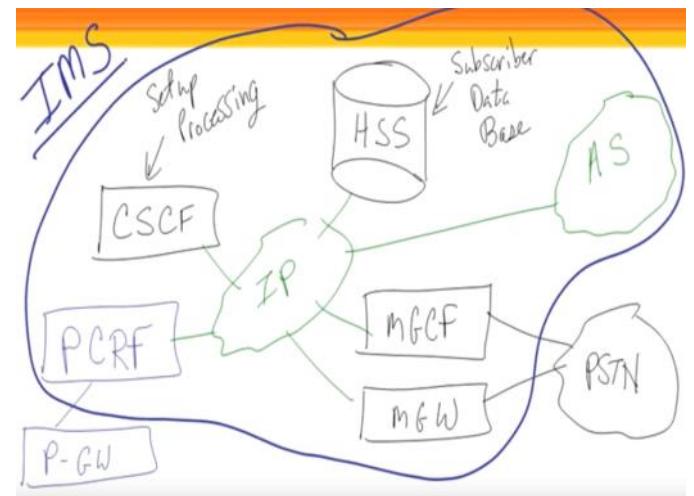
IMPI	private entity global identity allocated by home nw
IMPU	public identity like telephone no(SIP URI,)

Handset:
SIP User Agent

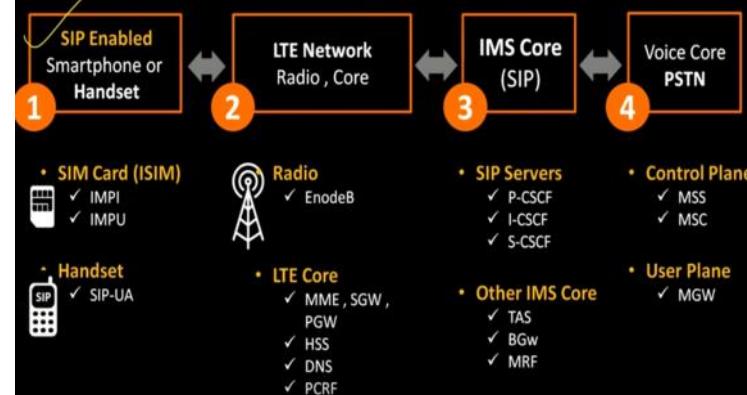
(SIP-UA) UAC(user agent client): send SIP request
UAS(server):receive requests and send response

ims core -> directly pgw with Sgi
lte is used as carrier
wifi fixed broadband
DSL

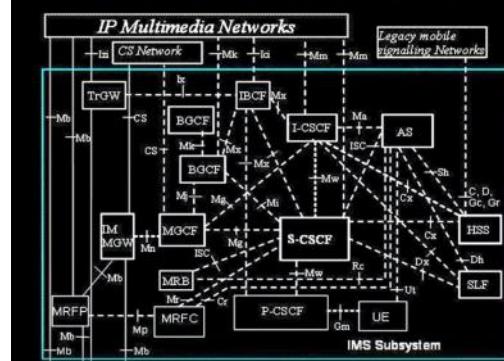
SIP server
media Gateway



VoLTE Network - Building blocks



IMS VoLTE Architecture & Specs



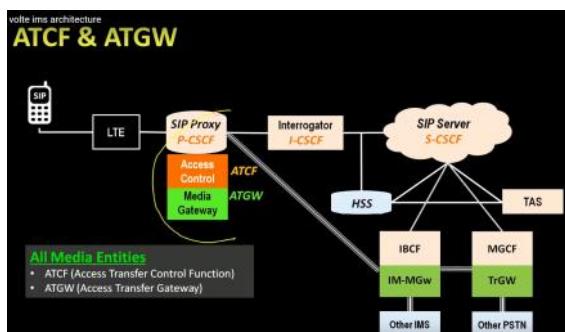
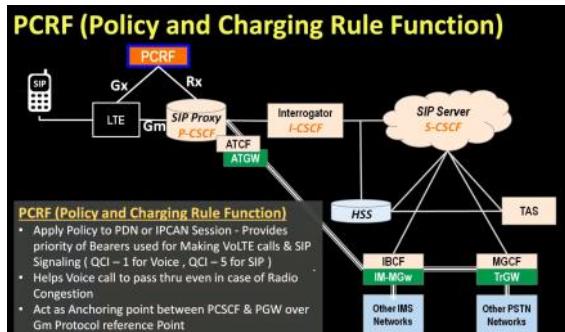
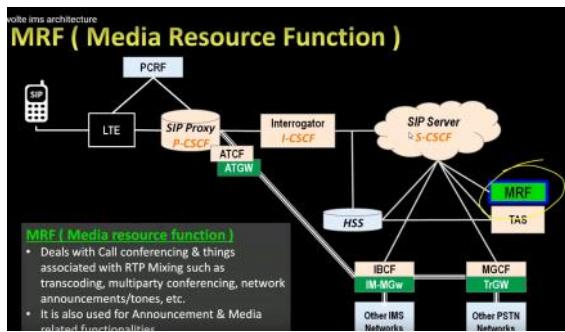
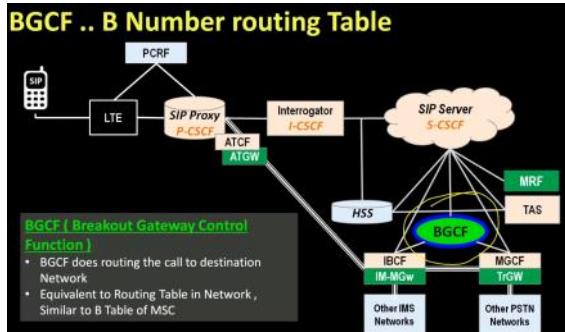
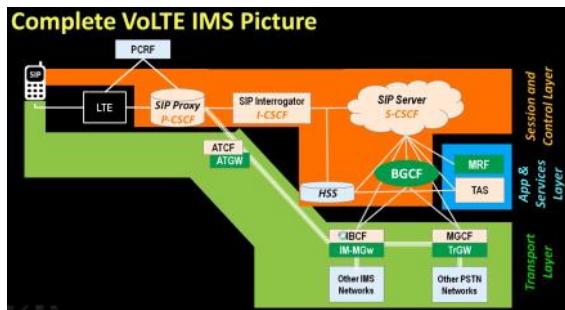
3

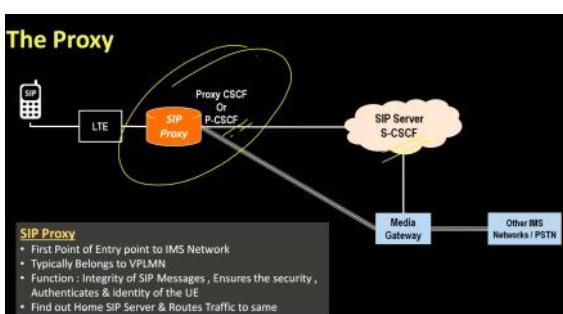
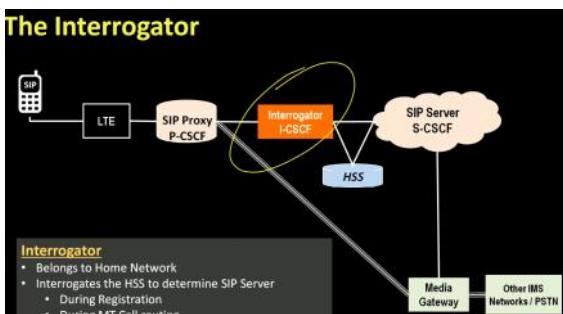
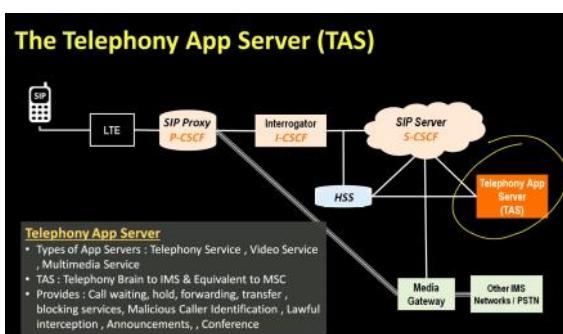
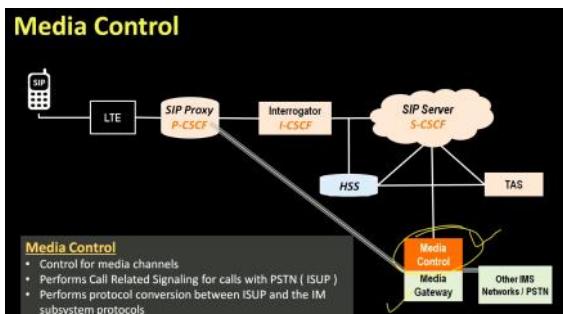
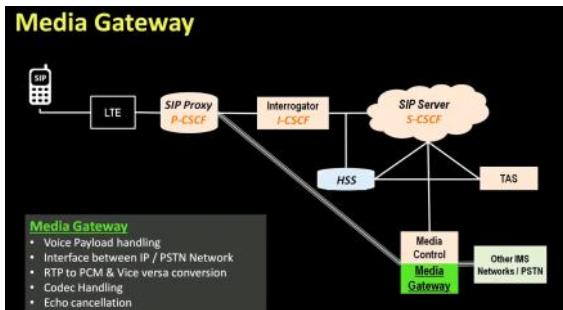
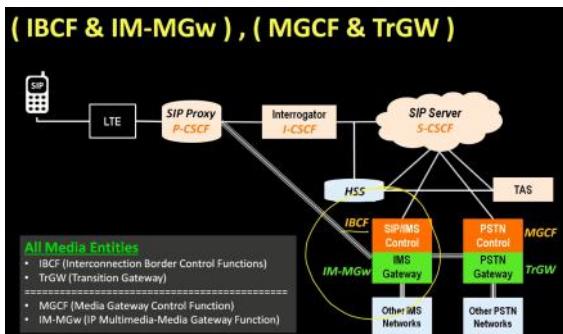
3GPP
IMS SPECS
3GPP 23.228

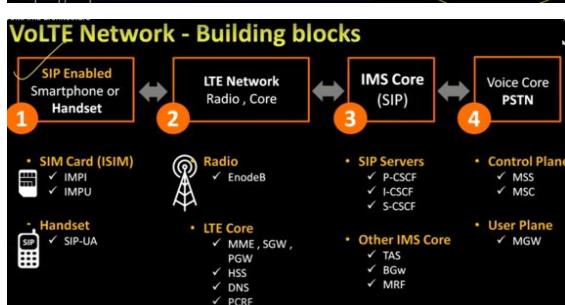
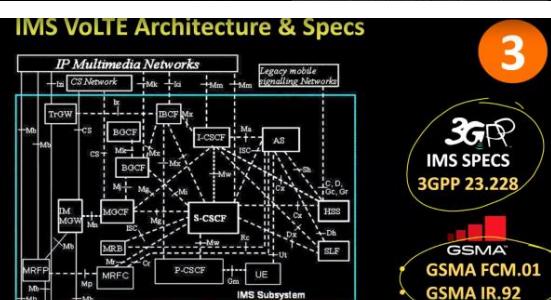
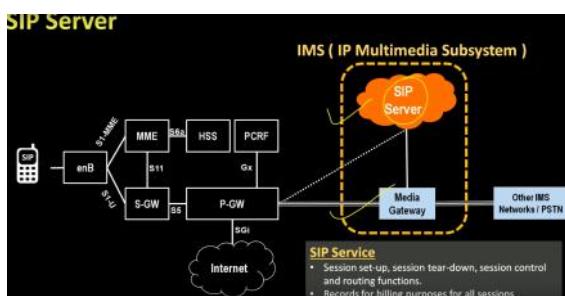
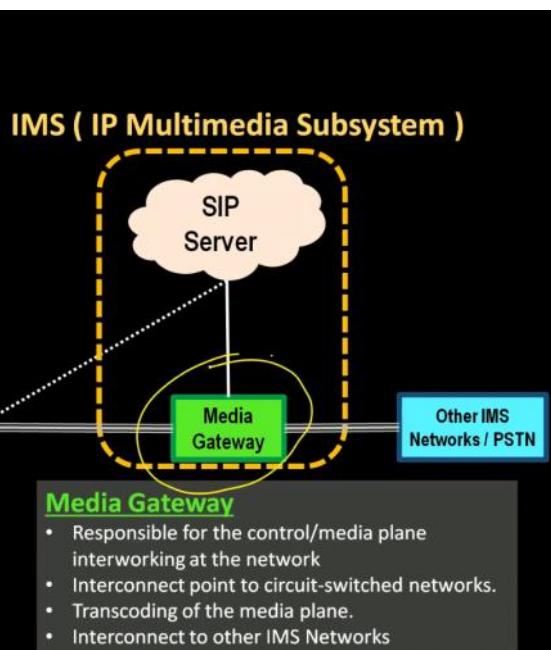
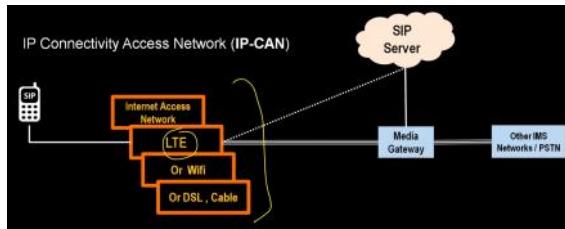
GSMA
GSMA FCM.01
GSMA IR.92

Changes in Core Network to support VoLTE

Packet Core SGW / PGW	Packet Core MME	Voice Core MSS / MSC
<ul style="list-style-type: none"> Default & dedicated Bearer Enable IMS APN IP Pool PGW - SBC / PCSCF Reachability Disable Gy for IMS 	<ul style="list-style-type: none"> Sv Link – SRVCC IMS APN SGW / PGW Selection Paging Policy - VoLTE Calls QCI – 1,5 Testing & Validation 	<ul style="list-style-type: none"> Sv Link – SRVCC IMS & MGW – IP Reachability Break in & Break out Calls Config Common Codec between MSS & IMS







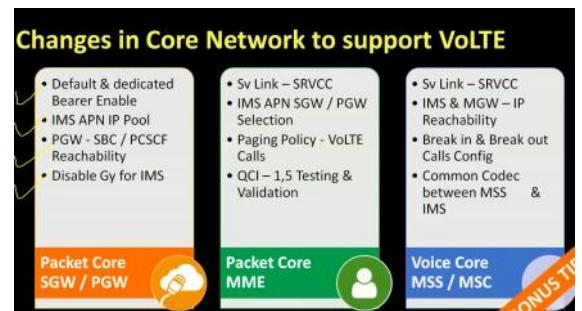
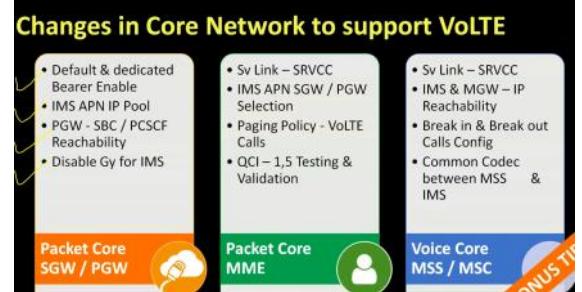
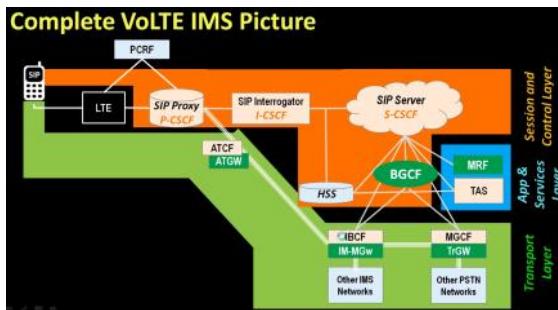
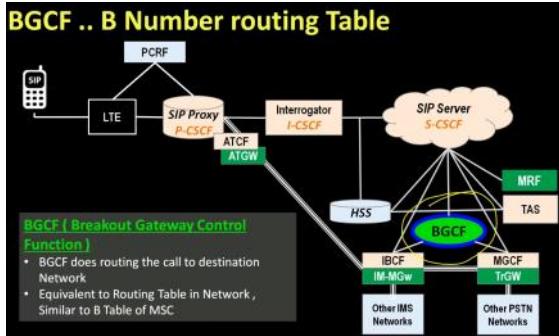
I don't have I-SIM in my Network ?

Do existing 3G or LTE customer using USIM need to go for SIM Swap for availing VoLTE Services ?

USIM itself can be used to Avail VoLTE Services , Deriving the IMPU and IMPI for authentication can be done by VoLTE stack/client residing on the UE (ISIM is not Mandatory)

Derived
IMPI :- <IMSI>@ims.mnc<3-digit MNC>.mcc<3-digit MCC>.3gppnetwork.org

Derived
IMPU :- <MSISDN>@ims.mnc<3-digit MNC>.mcc<3-digit MCC>.3gppnetwork.org



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Derived
IMPU :- <MSISDN>@ims.mnc<3-digit MNC>.mcc<3-digit MCC>.3gppnetwork.org

SCTP

18 May 2018 07:28

SCTP: stream control transmission Protocol

RFC :2960

connection oriented transport protocol

build to overcome specific problems in TCP

Uses sequences no and acknowledgements to provide delivery guarantees.

Uses a "receiver window" between the peers to indicate the amount of data can be receive in buffer

layer4 layer

source port

destination port

verification tag: zero for first packet

initial TSN: random number

forward TSN:

init(chunk type:1)

init_ack

COOKIE_ECHO

COOKIE_ACK

Data

sack

data and sack have same tsn

ASCONF(chunk type:193) :change or tear down association n middle

shutdown

shutdown ack

shutdown complete

error:

init --> abort : other side ip address not configured

wire shark filters:

sctp_sack_rtt=

sctp_chunk_type=(9==all error chunks,4==heart beat,5==heartbeat sack,1==init

Word Offset	Byte 0	Byte 1	Byte 2	Byte 3
0x0000	Hardware Type (0x01)	Protocol Type (0x80)		
0x0010	HLEN (0x06)	PLEN (0x04)	Operation	
0x0020		Sender Hardware Address		
0x0030			Sender Protocol Address	
0x0040				
0x0050		Target Hardware Address		
0x0060			Target Protocol Address	
0x0070				

Wireshark Filters for SCTP

Error Types:

- | | |
|----|--|
| 1 | Invalid Stream Identifier |
| 2 | Missing Mandatory Parameter |
| 3 | Stale Cookie Error |
| 4 | Out of Resource |
| 5 | Unresolvable Address |
| 6 | Unrecognized Chunk Type |
| 7 | Invalid Mandatory Parameter |
| 8 | Unrecognized Parameters |
| 9 | No User Data |
| 10 | Cookie Received While Shutting Down |
| 11 | Restart of an Association with New Addresses |
| 12 | User Initiated Abort |
| 13 | Protocol Violation |

Session Setup: The Four-Way Handshake

1. INIT - Sent by client. Includes its "Verification Tag" value.
2. INIT ACK - Response to INIT, includes its own Verification Tag as well as a "State Cookie".
3. COOKIE ECHO - Client sends the State Cookie back to server.
4. COOKIE ACK - Final acknowledgment.

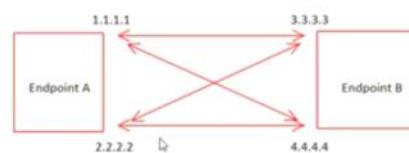
- Multiple data chunks in a single IP packet are very common. This means that TSN numbers and acknowledgements can be harder to follow per-packet than TCP streams. It is even common for multiple applications to be carried over one association, making it harder to trace at a packet level (especially true when SCTP is carrying an SS7 protocol stack).

SCTP Protocol Terminology

- Chunk: An individual SCTP message sent within a packet. A "data chunk" is equivalent to a TCP data segment.
- Path: A connection between two IP endpoints.
- Association: Overall connection between two SCTP endpoints. Equivalent to a TCP session.
- TSN: Transmission Sequence Number, roughly equivalent to TCP sequence number. An increment of data chunks, not bytes.
- Heartbeat: Like a TCP keepalive. Sent on a per-path basis.
- Shutdown: Equivalent to FIN flag in TCP
- Abort: Equivalent to RST flag in TCP

The Association

- An association is formed between two endpoints.
- Paths are established between IP addresses owned by those endpoints.



Word Offset	Byte 0	Byte 1	Byte 2	Byte 3
0x0000	Hardware Type (0x01)	Protocol Type (0x80)		
0x0010	HLEN (0x06)	PLEN (0x04)	Operation	
0x0020		Sender Hardware Address		
0x0030			Sender Protocol Address	
0x0040				
0x0050		Target Hardware Address		
0x0060			Target Protocol Address	
0x0070				

Wireshark Filters for SCTP

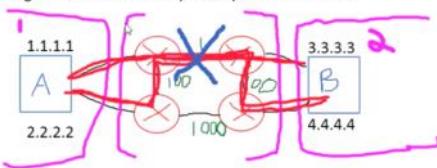
Error Types:

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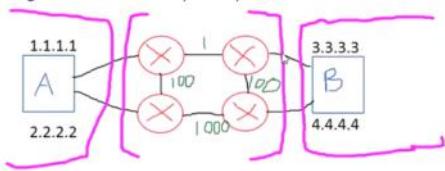
Problems in SCTP:

- Multiple data chunks in a single IP packet are very common. This means that TSN numbers and acknowledgements can be harder to follow per-packet than TCP streams. It is even common for multiple applications to be carried over one association, making it harder to trace at a packet level (especially true when SCTP is carrying an SS7 protocol stack).

- Layer 4 deciding where data is sent and received reduces path predictability for the network administrators working in Layers 1-3. Packets might not be where everyone expects them to be.



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Possible Problems

- NAT can create conflict between IP information in L3 and L4.

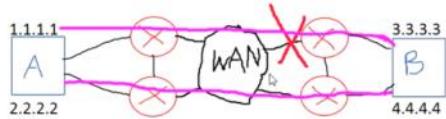
```

Internet Protocol Version 4, Src: 192.168.0.101 (192.168.0.101),
Stream Control Transmission Protocol, Src Port: 6666 (6666), Dst
Source port: 6666
Destination port: 9999
Verification tag: 0x48e63127
Checksum: 0x500d88fe (not verified)
ASCONF chunk
  Chunk type: ASCONF (193)
  Chunk flags: 0x00
  Chunk length: 32
  Sequence number: 0xa1104d8a
  ■ IPv4 address parameter (Address: 192.168.0.101)
  
```

- Not all firewalls that support rules based on Layer 4 port numbers may support this particular Layer 4 protocol.

Advantages over TCP

- Four-way handshake protects from SYN floods.
- Supports multiple IP addresses on both sides of a session. The protocol is not bound to one stream or even IP version.
- Supports explicit error codes and message types.
- Selective ACKs are supported in all SCTP connections.
- Window size is 32-bit value with no scaling window.
- Supports "Partially Reliable" delivery options (Forward TSN).
- Multiple IPs from multiple subnets gives potential for L4 convergence to happen before L3.



Retransmissions

- If a receiver notices packet loss, it sends a selective ACK indicating the missing TSNs. Sends three times to trigger retransmission, referred to as a "fast retransmission".
- If RTO expires without acknowledgements, retransmission is triggered. **RTO is per destination IP address.**
- The "Forward TSN" chunk can be used to stop retransmission logic for a TSN and essentially give up on it.

OSI

20 May 2018 17:06

Big	bits layer1	mac/physical
Fat	frames layer2	data link
People	packets layer3	network
Snort	segments layer4	transport
Donuts	datagrams layer5 through 7	session, presentation, application

All People Say That Nerds Dress Poorly
All People Seem To Need Data Processors
please do not touch susu arts

bridge selection process in Switches.

The normal order is :

- Step 1 : Lowest Root Bridge ID (BID)
 - Step 2 : Lowest Path cost to Root Bridge
 - Step 3 : Lowest Sender BID
 - Step 4 : Lowest Port ID
- Ruined Puzzles Sacrifice Programs

From <<http://zahid-stanikzai.com/mnemonics-for-networking/>>

From <<http://zahid-stanikzai.com/mnemonics-for-networking/>>

Multicast

20 May 2018 17:32

Protocol Independent Multicast PIM – SM v2 Message types:

High Roads Reaching Jungle Blocked All Cars

1. Hello
2. Register
3. Register-Stop
4. Join/Prune
5. Bootstrap
6. Assert
7. Candidate RP Advertisement

Remembering the list of Cisco IOS log severity levels.

Every Alley Cat Eats Watery Noodles In Doors

- Emergency (0)
- Alert (1)
- Critical (2)
- Error (3)
- Warning (4)
- Notifications (5)
- Information (6)
- Debug (7)

To remember BGP Best Path Algorithm

We Love Oranges AS Oranges Mean Pure Refreshment.

- Weight (Highest is Better)
- LOCAL_PREF (Highest is Better)
- Originated Locally
- AS_PATH (Shortest Wins)
- ORIGIN Type (IGP is Lower than EGP and EGP is Lower than Incomplete)
- MED (Lowest is Better)
- Paths (External > Internal)
- RID (Router ID – Lowest is Better)

From <<http://zahid-stanikzai.com/mnemonics-for-networking/>>

DHCP

20 May 2018 17:33

DHCP process. DORA, Discover DHCP server by broadcasting on the LAN segment, the DHCP server with Offer up an IP address, we will Request the IP address, and finally Acknowledge the IP address request.

From <<http://zahid-stanikzai.com/mnemonics-for-networking/>>

dynamic host configuration protocol

assign ip address to hosts
server address , default gateway, DNS
comes as client and server

ip address should be unique in network
each computer run dhcp client,

DHCP discover:	if computer attach to a network first it try find DHCP server via broadcast message all other device then DHCP server look at it and drop it.
DHCP offer:	when server gets message it then send offer to host. host take the first offer it receive if more then one is given.
DHCP req:	Host says I take it send re to server
DHCP ACK:	Server sends the ip address to host along with subnet mask ,default gw, dns server server: record the all list of ip address, MAC and there lease time(expiration time) UDP PORT client:68 Server:67

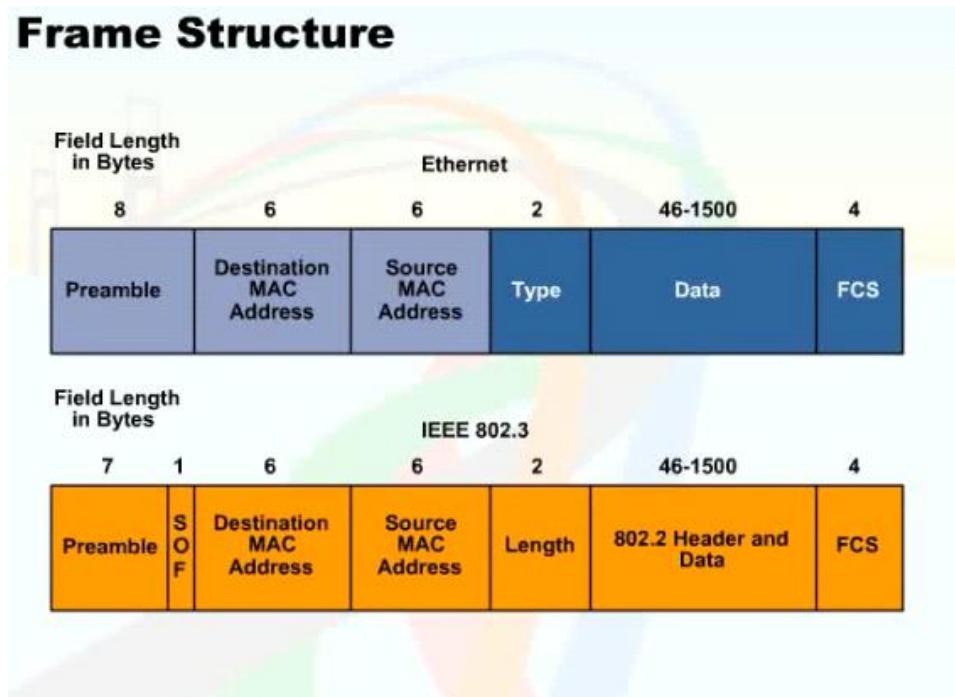
MAC

21 May 2018 00:38

destination is first component
preamble is not count in frame size

omkar hills
turahalli

Frame Structure



frame	72bytes (data min)
preamble	seq of known bit, used for sync and channel estimation
sof	start of frame
Destination add	group addressing (multicast) higher bit order is one for broadcasting all bits are made one
length/type	actual content
data and padding	
FCS	extra error-detecting code

etherenet: connect no of com system to from a local area network
10 base 5

mapping of ip to mac

static mapping	dynamic mapping
table is created with logical and physical add	Each time a machine knows the logical address of the other machine, it can use a protocol to find physical address

ARP:

ARP associates an ip add with the physical address.

Logical address-- ---mapping via ARP--->Physical address

ARP request	broadcast (ip of m1,physical m1 ,intended m2 ip address)
ARP replay	unicast (m2 ip and m2 physical)

Word Offset	Byte 0	Byte 1	Byte 2	Byte 3		
0x0000	Hardware Type (0x01)		Protocol Type (0x80)			
0x0010	HLEN (0x06)	PLEN (0x04)	Operation			
0x0020	Sender Hardware Address					
0x0030			Sender Protocol Address			
0x0040						
0x0050	Target Hardware Address					
0x0060			Target Protocol Address			
0x0070						

hw type	type of h/w
protocol type	ipv4,ipv6
h/w len	length of physical add
protocol len	len of logical add
operation	request(1),replay(2)
target h/w address	in a request not filled, fill via replay

RARP:

RARP maps physical to logical(ip)

TCP

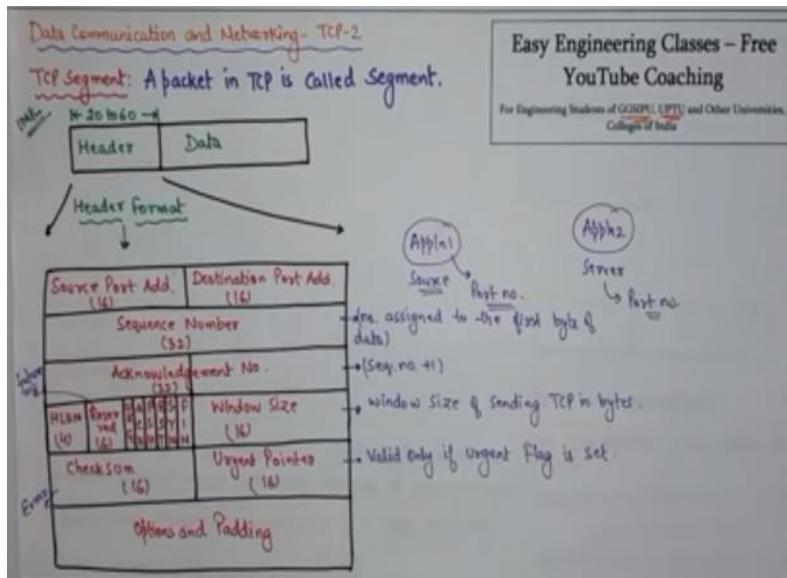
21 May 2018 01:47

Tcp is most widely used for data transmission in communication network such as internet.
TCP provide process to process communication

Feature:

1. process to process
2. port no
3. numbering system
4. byte no
5. seq no {segment no}
6. ack no
7. flow control
8. error control
9. congestion control

sync	seq no(6000) client
syn+ack	syn(1000)server ,ack(6001),receiving window(5000)server
ack	sync (6000),ack (1001),receiving window of client(6000)



diameter

23 May 2018 08:14

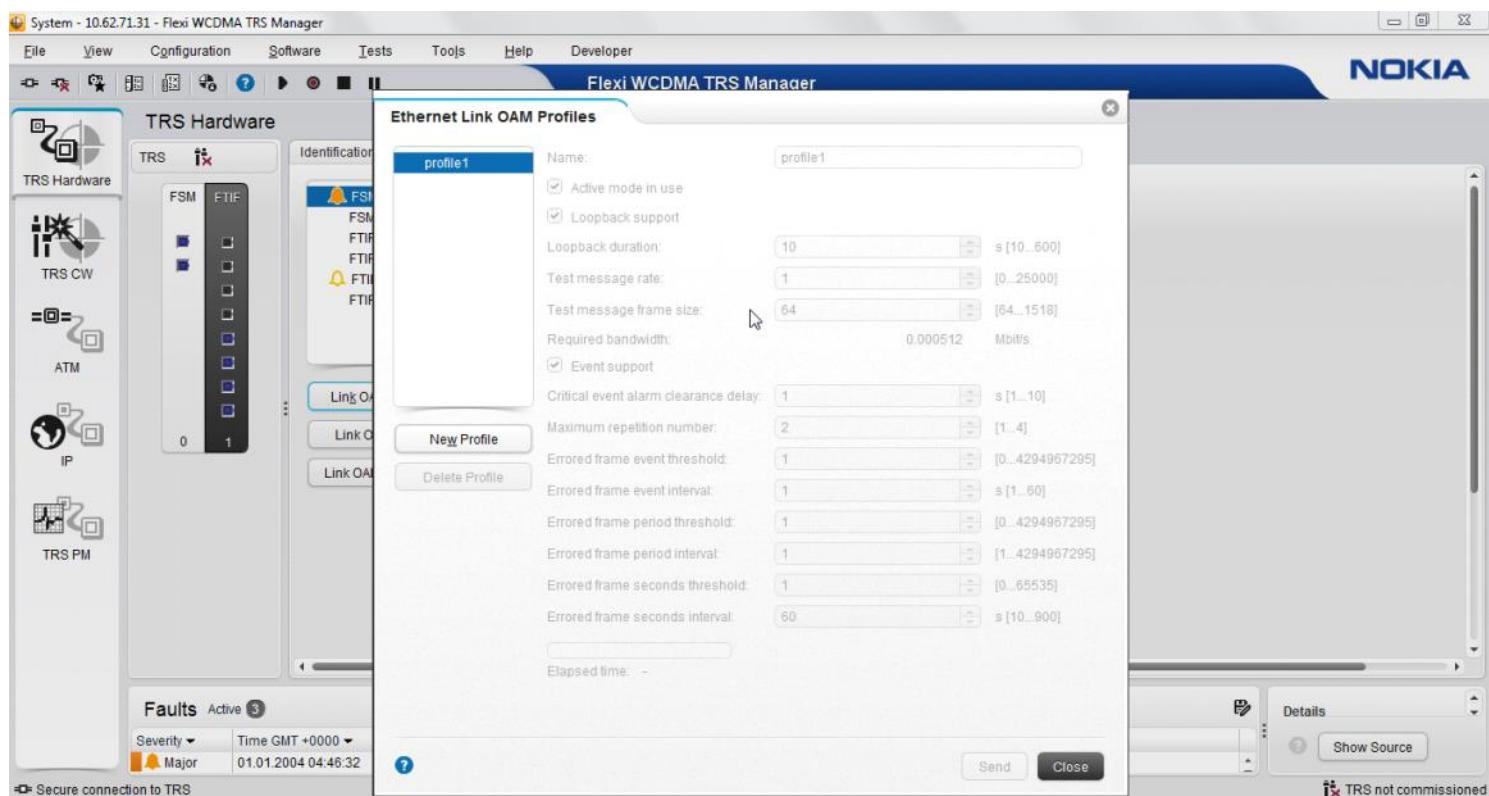
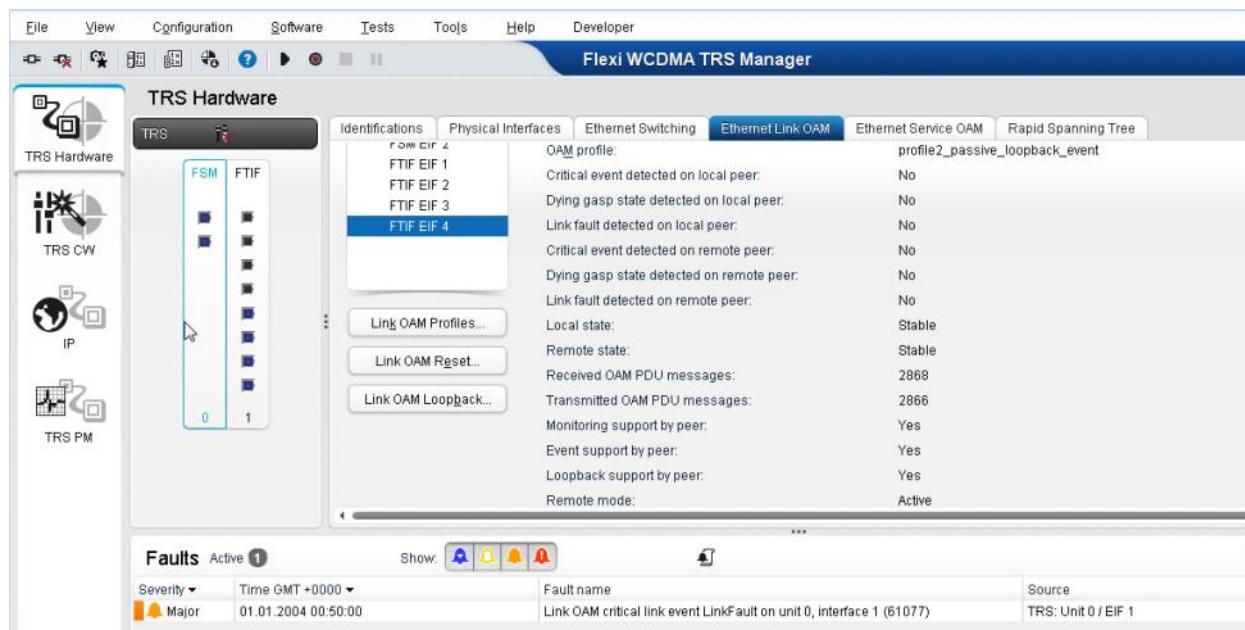
one protocol:
itf dia top of that they adopted

dia enable service provider
qos, smart charging, load balancing

dia (data for ip n/w) is diff for ss7(vioce network pre ip)

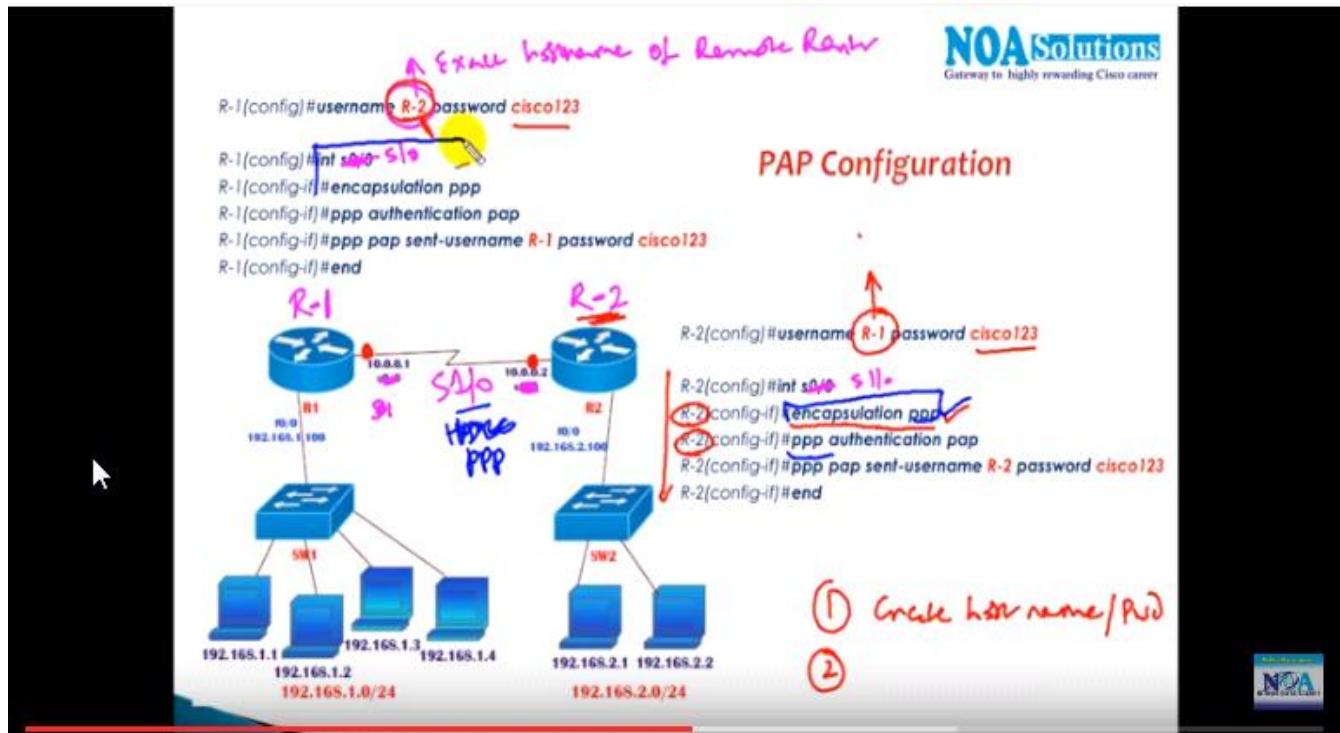
LOAM

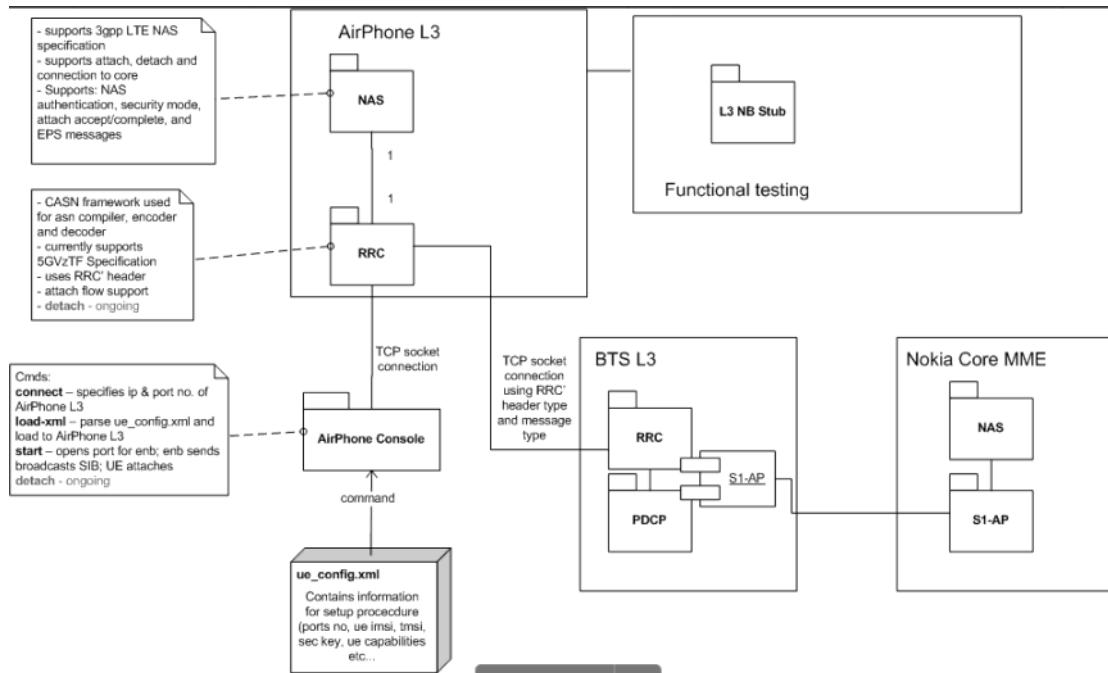
04 June 2018 01:16



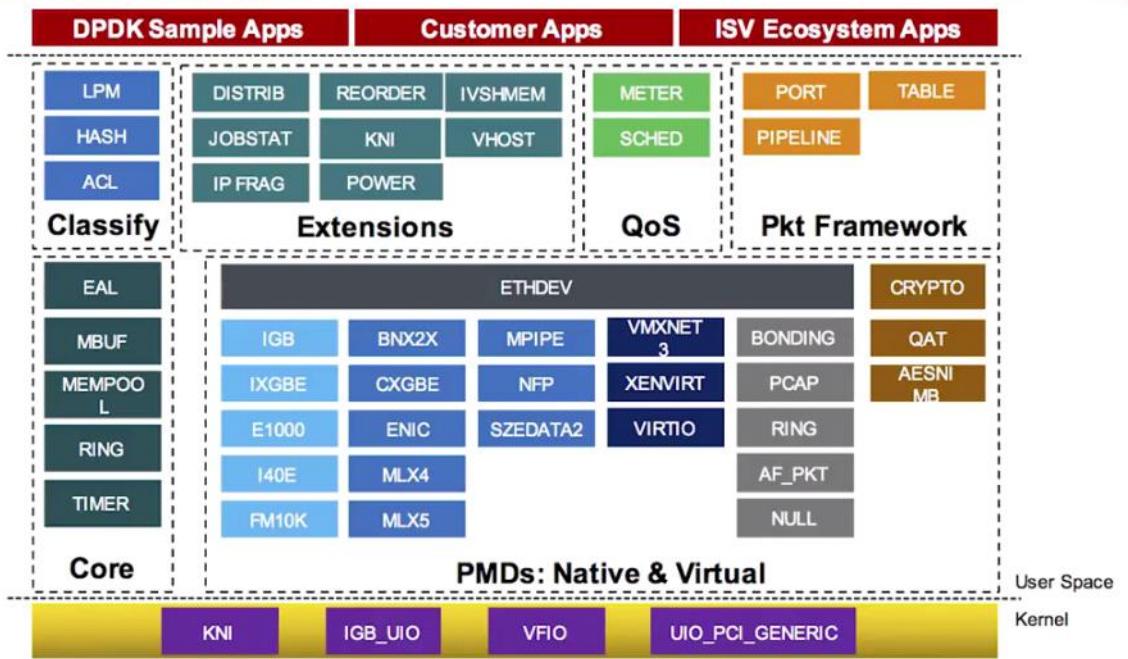
MLPPP

04 June 2018 01:38





DPDK Architecture



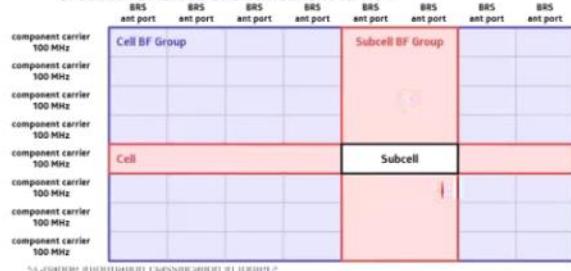
Terminology

Cell

- The cell is defined as the cell as seen by the UE with a single physical-layer cell identity and a single global cell ID.
- According to V5G211: "There are 504 unique physical-layer cell identities". LTE 36.311: "The IE CellIdentity is used to unambiguously identify a cell within a PLMN".
- Cell maps to one component carrier.

Subcell

- Implementation-specific concept defined as the intersection of subcell BF group and component carrier.



Cell BF Group

- The set of cells on different carrier frequencies using the same physical antennas and the same analog beam forming hardware
- In 5G17 there is one Cell BF Group per RAU. In the future, there will be more due to more antennas and/or multiple frequency bands.

Subcell BF Group

- Implementation-specific concept consisting of a subset of BRS antenna ports of a Cell BF Group controlled by the same packet scheduler instance.
- MU-MIMO is possible between subcell BF Groups.
- SU-MIMO cannot cross subcell BF Group boundary.
- In 5G17 there are always 2 BRS antenna ports per subcell BF Group. In the future the number of antenna ports might be different (especially if the number of component carriers is smaller).

BRS Antenna Port

- The BRS antenna port is defined by V5G211: "Beam reference signals are transmitted on one or several of antenna ports 0 to 7".
- In case of analog beam forming, the BRS antenna ports with the same number on different component carriers are assumed to be mapped to the same analog beam direction and same physical antenna.

NOKIA

Carrier aggregation

08 June 2018 21:00

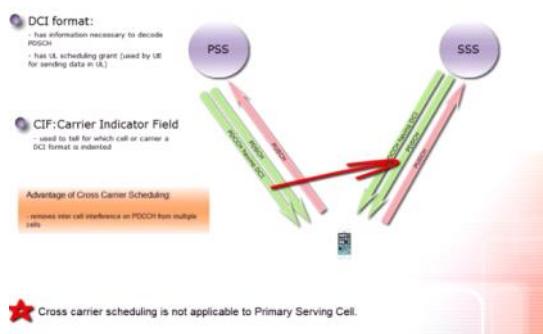
PCC :primary control carrier	1. add and remove secondary component carrier dynamically 2. dynamically activate and deactivate secondary cell 3. handles all RRC and NAS procedures 4. Receive measurement reports and control mobility of UE 5. primary cell can only be changed at time of handover.	
SCC carrier component secondary	A ue can aggregate maximum up to 5 carriers i.i 1pcc and 4 SCC max BW 10mz(20,20,20,20,20) actual no if SCC depends on UE capability Not possible to configure UE with more uplink component carrier then downlink	

Carrier allocation schemes in CA(36.101)

Intra band contiguous(CA_X)	pcc and scc are conti and same band
Intra band noncontiguous (CA_X-X)	pcc and scc are allocated from same band but not contiguous
inter band noncontiguous (CA_X-Y)	pcc and scc are in two different freq band

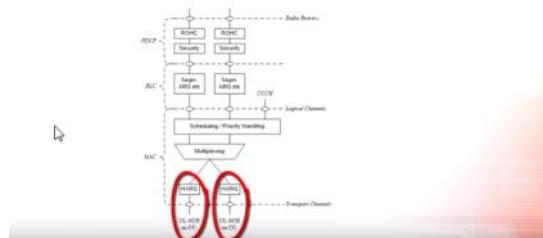
Precondition for CA:ue capability

UE category	
supporting band combination	
CA BW class	
BW set	



Role of MAC Layer in CA

Configures one HARQ entity for each Secondary Cell



Role of Physical Layer in CA

Two new PUCCH formats were defined:

- Format 1b (HARQ ACK/NACK for up-to 2 component carriers)
- Format 3 (HARQ ACK/NACK for up-to 5 component carriers)

Physical-layer-processing chain is repeated for every UL Serving Cell



CA Bandwidth Class

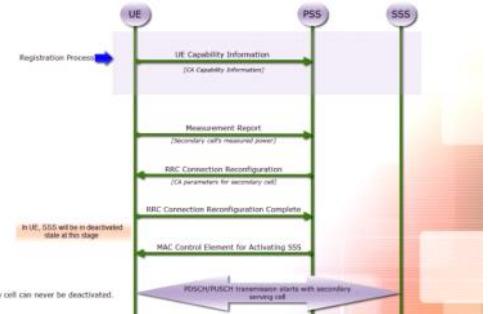
CA Bandwidth Class	Aggregated Transmission Bandwidth Configuration	Number of contiguous CC
A	$N_{\text{RB}} \leq 100$	1
B	$200 < N_{\text{RB}} \leq 300$	2
C	$100 < N_{\text{RB}} \leq 200$	3
D	$200 < N_{\text{RB}} \leq 300$	3
E	$300 < N_{\text{RB}} \leq 400$	4

CA_3C:
 - two contiguous carriers in frequency band 3
 - transmission bandwidth from 100 RB to 200 RB

CA_25A-25A:
 - two non-contiguous carriers in frequency band 25
 - maximum transmission bandwidth of 100 RB



★ CA_25A X



GPRS

15 June 2018 18:25

TDMA based scheduling

reduced mac scheduling ----> Uplink (2 time data)
 resource allocator(bsc)
 pcu (PACKET CONTROL UNIT)

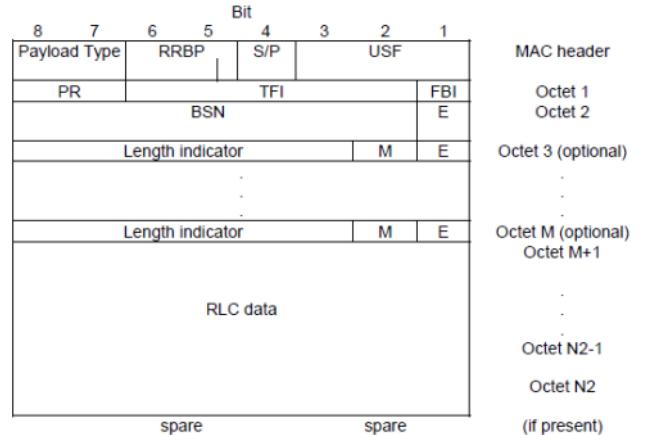
MAC:

rlc:

SIGNALING	dl UI
control	dl ul

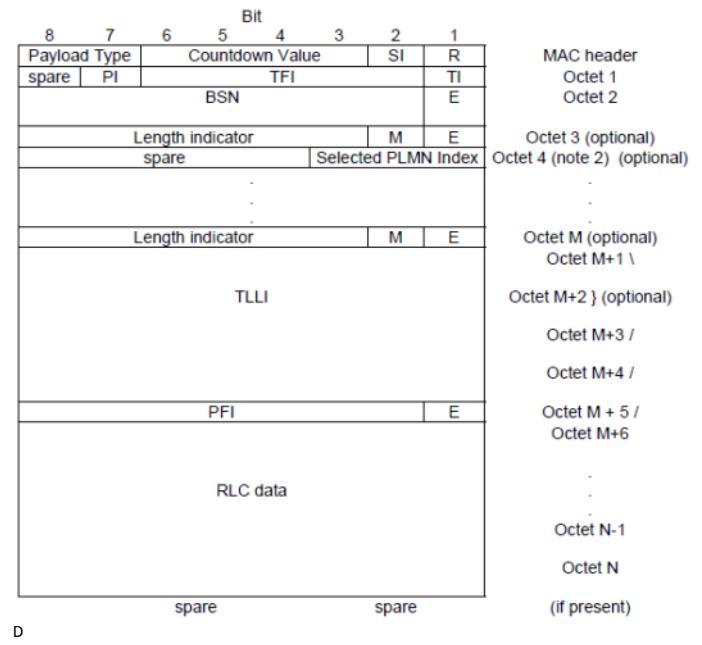
one rlc-->one radio block

DL data block:



RRBP	relative receive block preoid
TFI	I2 link between UE and BTS ,every user have its TBF-->TFI
FBI	final block identity,enables on last data
USF	
s/p	supplementry polling

UL data block:



D

RLC Block Structure - Uplink RLC/MAC control block

< 44.060 - Figure 10.3.2.1: Uplink RLC/MAC control block together with its MAC header >

