# **Asymptotic Analysis - countPositiveElements Function**

### Function Overview:

The function loops through an array A of size n and counts how many elements are positive.

```
Python Function:

def countPositiveElements(A):
    count = 0  # Line 5
    for x in A:  # Line 6
        if x > 0:  # Line 7
            count += 1 # Line 8
    return count  # Line 9
```

# **Cost Table Recap:**

| Line | Operation    | Cost | How many times it runs       |
|------|--------------|------|------------------------------|
| 5    | count = 0    | c1   | 1                            |
| 6    | for x in A   | c2   | n                            |
| 7    | if x > 0     | с3   | n                            |
| 8    | count += 1   | c4   | n (worst case: all positive) |
| 9    | return count | с5   | 1                            |

#### Effort Calculation (Line-by-Line Explanation)

#### Step 1:

$$T(n) = c1 + n*c2 + n*c3 + n*c4 + c5$$

We sum all the individual operations:

- Line 5: 1 time -> c1
- Line 6, 7, 8: each runs n times -> n\*c2 + n\*c3 + n\*c4
- Line 9: 1 time -> c5

## Step 2:

$$T(n) = c1 + c5 + n*(c2 + c3 + c4)$$

Combine constants c1 and c5 and factor out n from the rest.

Step 3:

$$T(n) = c6 + n*c7$$

Rename c1 + c5 as c6, and c2 + c3 + c4 as c7 for simplicity.

Step 4:

$$T(n) \le n*c6 + n*c7$$

Upper bound idea: if constants are positive, c6 + n\*c7 is less than or equal to n\*c6 + n\*c7

Step 5:

$$T(n) \le n^*(c6 + c7)$$

Factor out n.

Final Step:

$$T(n) <= n*c8$$

Combine all constants into one constant c8. Since asymptotic analysis ignores constants, this simplifies the notation.

Final Form: T(n) = O(n)

This shows the time complexity is linear because the effort increases proportionally to n.

Why 
$$T(2n) = 2T(n)$$
?

If the array size doubles, then:

$$T(2n) = 2n * c8 = 2 * (n * c8) = 2 * T(n)$$

So yes, linear growth confirms this.