

①

Assignment 4

i) a_n = 4a_{n-1} a₀ = 1 n > 1

$$a_n = 4a_{n-1}$$

$$(\text{substitute}) = 4[4a_{n-2}] = 4^2 a_{n-2}$$

$$(\text{substitute}) = 4[4^2 a_{n-3}] = 4^3 a_{n-3}$$

$$= 4^k a_{n-k}$$

where k is some integer, & we must choose an appropriate value for k

$$\text{let } k = n$$

$$\Rightarrow a_{n-n} = a_0 = 1 \quad [\text{Given}]$$

$$\Rightarrow 4^n a_{n-n} = 4^n a_0 = \boxed{4^n \forall n \geq 0}$$

b) a_n = ca_{n-1} a₀ = 1 n > 1

Given that, a_n = ca_{n-1}

$$(\text{substitute}) = c(c a_{n-2}) = c^2 a_{n-2}$$

$$(\text{substitute}) = c[c^2 a_{n-3}] = c^3 a_{n-3}$$

for some value of k = c^k a_{n-k}

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where k is some integer, and we must choose an appropriate value of k

Suppose that, $k = n$

$$\Rightarrow c^n a_{n-n}$$

$$= c^n a_0$$

$\Rightarrow c^n \neq 0$ and $c > 1$

c) $b_n = b_{n-1} + \frac{n}{2}$

$b_0 = 0 \quad \forall n \geq 1$

Given that, $b_n = b_{n-1} + \frac{n}{2}$

$$= [b_{n-2} + \left(\frac{n}{2} - 1\right) + \frac{n}{2}]$$

$$= [b_{n-3} + \left(\frac{n}{2} - 2\right)] + \left(\frac{n}{2} - 1\right) + \frac{n}{2}$$

$$= [b_{n-(k+1)} + \left(\frac{n}{2} - k\right)] + \dots + \left(\frac{n}{2} - 2\right) + \left(\frac{n}{2} - 1\right) + \frac{n}{2}$$

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$$= b_{n-(k+1)} + \left(\frac{n}{2} - k\right) + \dots + \left(\frac{n}{2} - 2\right) + \left(\frac{n}{2} - 1\right) + \frac{n}{2}$$

We want to stop at b_0 , we know that $b_0 = 0$

$$\text{So, let } \frac{n}{2} - (k+1) = 0$$

$$\Rightarrow \frac{n}{2} - k - 1 = 0$$

$$\Rightarrow \frac{n}{2} = k + 1$$

$$\Rightarrow n = 2(k+1) = 2k+2$$

$$\therefore \frac{n}{2} - (k+1) = 0$$

$$\frac{n}{2} - k = 1$$

$$\frac{n}{2} - (k-1) = 2 \text{ etc.}$$

So, we get

$$b_0^{\leq 0} + 1 + 2 + 3 + \dots + \left(\frac{n}{2} - 1\right) + \frac{n}{2}$$

$$\Rightarrow 1 + 2 + 3 + \dots + \left(\frac{n}{2} - 1\right) + \frac{n}{2}$$

$$= \frac{n/2}{2} \left[1 + \frac{n}{2} \right] = \frac{n}{4} \left(\frac{2+n}{2} \right)$$

$$= \boxed{\frac{n(n+2)}{8}}$$

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$$d) c_n = 4c_{n-1} + n$$

$$c_0 = 0$$

$$+ n > 1$$

Given that, $c_n = 4c_{n-1} + n$

$$= 4[4c_{n-2} + (n-1)] + n$$

$$= 4^2 c_{n-2} + 4(n-1) + n$$

$$= 4^2 [4c_{n-3} + (n-2)] + 4(n-1) + n$$

$$= 4^3 c_{n-3} + 4^2(n-2) + 4(n-1) + n$$

$$= 4^k c_{n-k} + 4^{k-1}(n-(k-1)) + \dots + 4(n-1) + 4^0 n$$

when $n=k$, we get

$$= 4^n c_0 + 4^{n-1} \cdot 1 + 4^{n-2} \cdot 2 + 4^{n-3} \cdot 3 + \dots + 4^1(n-1) + n$$

$$= \sum_{j=0}^{n-1} 4^{n-j} \cdot j$$

$$c_n = \sum_{j=0}^{n-1} 4^{n-j} \cdot j$$

$$= 0 \cdot 4^n + 1 \cdot 4^{n-1} + 2 \cdot 4^{n-2} + \dots + (n-1)4 + n$$

$$= \sum_{j=0}^{n-1} (n-j) 4^j = n \sum_{j=0}^{n-1} 4^j - \sum_{j=0}^{n-1} j \cdot 4^j$$

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$$\Rightarrow c_n = \frac{n(1-4^n)}{1-4} - \sum_{j=0}^{n-1} j \cdot 4^j$$

$$= \frac{n}{2} (4^n - 1) - \sum_{j=0}^{n-1} j \cdot 4^j$$

$$S_n = \sum_{j=0}^{n-1} j \cdot 4^j = ?$$

$$\text{Now, } 4^{k+1} - 4^k = 4^k(4-1) = 4^k \cdot 3 = 4^k$$

$$S_n = \sum_{j=0}^{n-1} j \cdot \frac{4^{j+1} - 4^j}{2} = \frac{1}{2} \sum_{j=0}^{n-1} j (4^{j+1} - 4^j)$$

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$$e) t_n = 10t_{n-1} + 21t_{n-2} \quad t_0 = 0 \quad t_1 = 1 \quad n \geq 2$$

$$t_2 = 10t_1 + 21t_0$$

$$= 10(0) + 21(1) = 21$$

$$t_3 = 10t_2 + 21t_1 = 10 + 21^2$$

$$t_4 = 10t_3 + 21t_2 = 10(10 + 21^2) + 21^3$$

Also, here $A = 10$, $B = -21$

Characteristics equation is $x^2 - 10x + 21 = 0$

$$\Rightarrow x^2 - 10x + 21 = 0$$

$$x^2 - 3x - 7x + 21 = 0$$

$$x(x-3) - 7(x-3) = 0$$

$$(x-7)(x-3) = 0$$

$$x = 3, 7$$

∴ Solution for $t_n = ax_1^n + bx_2^n = a(3)^n + b(7)^n$
where a and b are constants

Given conditions are $t_0 = 0$ and $t_1 = 1$

$$t_0 = a(3)^0 + b(7)^0 \Rightarrow a+b = 0 \quad \dots \textcircled{1}$$

$$t_1 = a(3)^1 + b(7)^1$$

$$\Rightarrow 3a + 7b = 1 \quad \dots \textcircled{2}$$

⑦

from eqn ① and ②

$$\begin{array}{r} 3a + 3b = 0 \\ 3a + 7b = 1 \\ \hline -4b = -1 \\ b = \frac{1}{4} \end{array}$$

$$a = -\frac{1}{4}$$

The solutions of recurrence is:

$$\boxed{-\frac{1}{4}(3)^n + \frac{1}{4}(7)^n = t_n}$$

$$f) t_n = -5t_{n-1} + 14t_{n-2} + 3^n \quad t_0 = 0, \quad t_1 = 1 \quad \forall n \geq 2$$

$$t_n + 5t_{n-1} - 14t_{n-2} = \boxed{3^n}$$

$$b = 3 \quad \uparrow \quad P(n) = 1$$

∴ degree $d=0$

Characteristic equation :

$$(x-2)(x+7) = \text{not } 0$$

$$(x-2)(x+7)(x-3) = 0$$

$$a_n = c_1 2^n - c_2 7^n + c_3 3^n$$

from recurrence \rightarrow boundary values

$$t_0 = 0 \quad \text{---} \textcircled{1}$$

$$t_1 = 1 \quad \text{---} \textcircled{2}$$

$$t_2 = -5t_1 + 14t_0 + 3^2$$

$$= 4 \quad \text{---} \textcircled{3}$$

So, we get 3 equations:

$$\underline{n=0} \quad c_1 - c_2 + c_3 = 0 \quad \text{---} \textcircled{A}$$

$$\underline{n=1} \quad 2c_1 - 7c_2 + 3c_3 = 1 \quad \text{---} \textcircled{B}$$

$$\underline{n=2} \quad 4c_1 - 49c_2 + 9c_3 = 4 \quad \text{---} \textcircled{C}$$

$$\boxed{c_1 = -\frac{6}{5}}, \quad \boxed{c_2 = \frac{1}{20}}, \quad \boxed{c_3 = \frac{5}{4}}$$

$$\Rightarrow \boxed{t_n = -\frac{6}{5}2^n - \frac{1}{20}7^n + \frac{5}{4}3^n}$$

⑨

g)

$$t_n = 3t_{n-1} + n + 3^n \quad t_0 = 0 \quad \forall n \geq 1$$

$$\underbrace{t_n - 3t_{n-1}}_{(x-3)} = n + 3^n$$

$\uparrow \quad \uparrow$

$$b_1 = 1 \quad b_2 = 3$$

$$p_1(n) = n \quad p_2(n) = 1$$

$$d_1 = 1 \quad d_2 = 0$$

Characteristic equation:

$$(x-3)(x-1)^2(x+3) = 0$$

$$\Rightarrow (x-3)^2(x+1)^2 = 0$$

$$x = \underbrace{1, 3}$$

multiplicity = 2

General solution: $t_n = c_1 1^n + c_2 1^n \cdot n + c_3 \cdot 3^n + c_4 n \cdot 3^n$ From recurrence \rightarrow boundary values

$$t_0 = 0 \quad \text{--- } ①$$

$$t_1 = 3t_0 + 1 + 3^1 = 4 \quad \text{--- } ②$$

$$t_2 = 3t_1 + 2 + 3^2 = 23 \quad \text{--- } ③$$

$$t_3 = 3t_2 + 3 + 3^3 = 98 \quad \text{--- } ④$$

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So, we get

$$\underline{n=0} \quad c_1 + c_2 + 0c_3 + 0c_4 = 0$$

$$\underline{n=1} \quad c_1 + c_2 + 3c_3 + 3c_4 = 4$$

$$\underline{n=2} \quad c_1 + 2c_2 + 9c_3 + 18c_4 = 23$$

$$\underline{n=3} \quad c_1 + 3c_2 + 27c_3 + 81c_4 = 98$$

$$\therefore c_1 = -1, \quad c_2 = -\frac{3}{4}; \quad c_3 = 1, \quad c_4 = \frac{11}{12}$$

$$\Rightarrow t_n = -1 - \frac{3}{4}n + 3^n + \frac{11}{12}n \cdot 3^n$$

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$$2) a) T(n) = 4T\left(\frac{n}{2}\right) + n \quad \forall n > 1$$

By Induction:

let a = first term

r = common ratio

m = total number of n terms

⇒ we can simplify it, as the sum of indeterminate size is a G.P

$$S_n = a + ar + ar^2 + ar^3 + ar^4 + \dots + ar^{m-1}$$

$$= \frac{a(r^m - 1)}{r - 1}$$

for $n = 2^k$

$$\therefore S_n = 2^n + 2^{n+1} + 2^{n+2} + \dots + 2^{2^k - 1}$$

$$= \frac{2^n(2^n - 1)}{2 - 1} = 4^n - 2^n$$

∴ The general formula for a is:

$$a(n) = 4^n - 2^n + 4^n \cdot m \\ = 4^n(m+1) - 2^n$$

$$\text{Hence, } T(k) = a(2^k) \\ \Rightarrow 4^{2^k}(m+1) - 2^{2^k}$$

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By Big-oh of $T(n)$

(say we don't know the boundary conditions)

Let $n = 2^k$ for $k \in \mathbb{I}^+$

$$\Rightarrow k \log_2^2 = \log_2 n \Rightarrow k = \log_2 n = \lg n \quad (\lg \equiv \log_2)$$

$$\text{rec } T(n) = 4T\left(\frac{n}{2}\right) + n$$

$$T(2^k) = 4T(2^{k-1}) + 2^k$$

$$\Rightarrow t_k = 4t_{k-1} + 2^k$$

$$\text{where } t_k \triangleq T(2^k) = T(n)$$

$$\Rightarrow \underbrace{t_k - 4t_{k-1}}_{\substack{\uparrow \\ (x-4)}} = 2^k$$

$$(x-4)(x-2)^4 = 0$$

$$\text{roots} = 4, 2$$

$$\Rightarrow t_k = c_1 4^k + c_2 2^k$$

$$\Rightarrow T(2^k) = c_1 \underbrace{4^k}_{n=2^k} + c_2 \underbrace{2^k}_{n=2^k}$$

$$n = 2^k \quad k = \lg n$$

$$\Rightarrow T(n) = c_1 4^{\lg n} + c_2 n$$

$$\Rightarrow T(n) \in O(4^{\lg n}) \quad |n \text{ is power of 2)}$$

(2)

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$$6) T(n) = 5T\left(\frac{n}{2}\right) + n^2 \quad \forall n > 1$$

Let $n = 2^k$ for some $k \in \mathbb{I}_+$

$$\Rightarrow T(2^k) = 5T(2^{k-1}) + \underbrace{(2^k)^2}_{(2^k)^2} \\ (2^k)^2 = 2^{2k} = 4^k$$

$$\Rightarrow T(2^k) = 5T(2^{k-1}) + 4^k$$

Let $t_k \triangleq T(2^k)$. Then :

$$t_k = 5t_{k-1} + 4^k$$

$$\Rightarrow t_k - 5t_{k-1} = 4^k \\ \underbrace{(x-5)}_{(x-5)(x-4)} \cdot \underbrace{(x-4)}_{=0} = 0$$

$$\Rightarrow t_k = c_1 5^k + c_2 k \cdot 4^k$$

Putting n back

$$T(n) = c_1 \underbrace{n^2}_{n=2^k} + c_2 n^2 \underbrace{\lg n}_{2^k=n} \\ \text{so, } n^2 = 4^k \Rightarrow k = \lg n$$

$$\Rightarrow T(n) = O(n^2 \lg n) \quad |n \text{ is a power of 2}|$$

3)

Tower of Hanoi

(14)

$t_n \triangleq$ No. of steps required to move n discs from 1 to 3

$$t_n = 2t_{n-1} + 1 \quad \forall n \geq 2$$

$$t_1 = 1 \quad \leftarrow \text{Hanoi recurrence}$$

pattern: $t_n = 2^n - 1 \quad \forall n \geq 1$
(FORMULA)

PROOF : By induction

CLAIM: Solution to $t_n = 2t_{n-1} + 1 \quad \forall n \geq 2$
with $t_1 = 1$ is the formula

$$t_n = 2^n - 1 \quad \forall n \geq 1$$

Value given by recurrence $\stackrel{?}{=}$ value given by formula?

Check base case(s), & then
use induction hypothesis.

Base case

recursion

formula

$$n=2$$

$$3$$

$$3$$

$$n=3$$

$$7$$

$$7$$

Induction hypothesis : Assume the formula solves the recurrence
for $n = 1, 2, 3, \dots, k$

We need to prove that $n = k+1$

$$\text{Recursion: } t_{k+1} = \underbrace{2t_k + 1}_{\downarrow}$$

$$\begin{aligned} \text{We have to show that the formula for } n=k \\ = \underbrace{2(2^k - 1) + 1}_{\downarrow} \end{aligned}$$

t_k from formula

$$\begin{aligned} &= 2^{k+1} - 2 + 1 \\ &= 2^{k+1} - 1 \end{aligned}$$

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which is exactly the same formula, but k is replaced by $k+1$

Solving, by characteristic eqn method:

$$t_n = 3t_{n-1} - 4t_{n-2} = 0 \text{ etc.}$$

$$x^2 - 3x - 4 = 0$$

3 or more t_n values

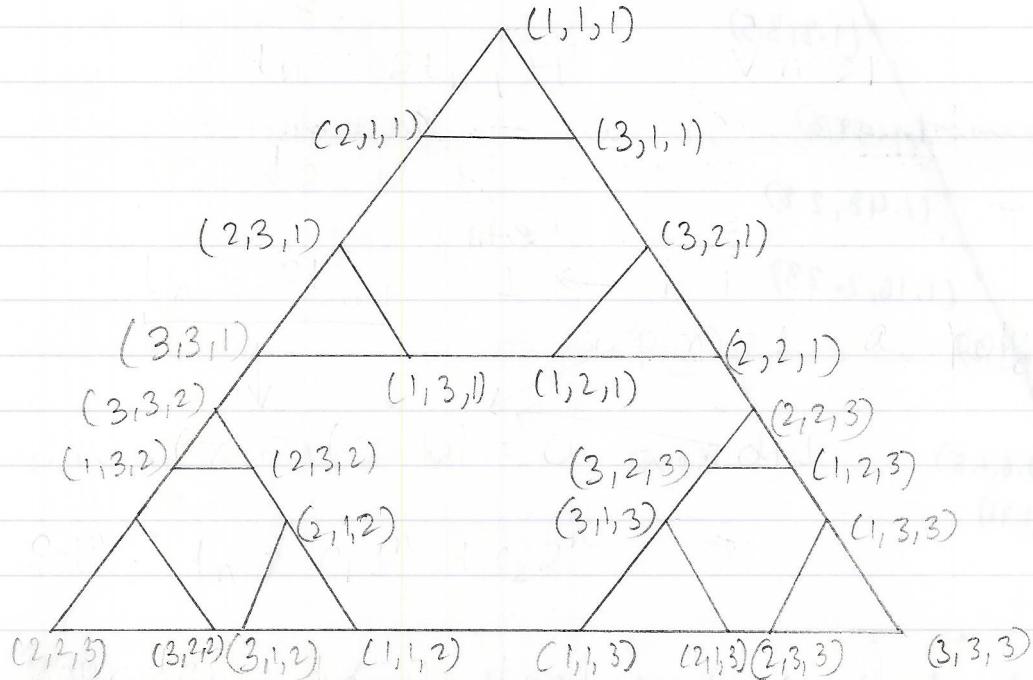
Degree

3 \Rightarrow quadratic eqn

4 \Rightarrow cubic eqn

5 \Rightarrow quartic eqn ($x^4 + \dots$)

Now, look at the Hanoi recursion:



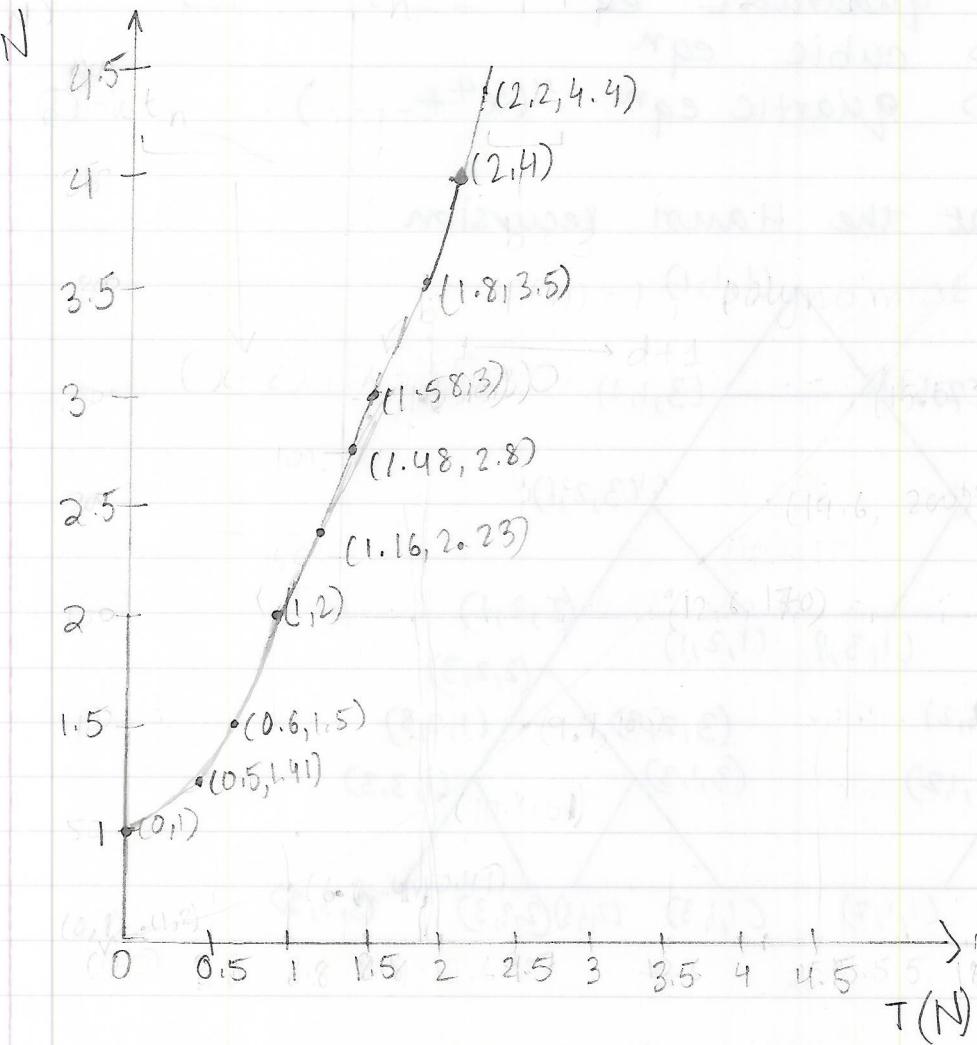
(17)

$$T(N) = \sum_{i=0}^{N-1} 2^i = \sum_{i=0}^{N-1} 2^i = 2^N - 1$$

$$S_n = a \left(\frac{r^n - 1}{r - 1} \right)$$

The complexity of the algorithm is exponential
It grows very fast, as a power of 2

$$\therefore T(N) = 2^N$$



4) Binary-search

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Initial condition: $T(1) = O(1) = 1$

Recurrence relation: The time to search in an array of N elements = time to search in an array of $N/2$ elements + 1 comparison

$$T(N) = T(N/2) + 1$$

Next, we perform telescoping :

$$T(N) = T(N/2) + 1$$

$$T(N/2) = T(N/4) + 1$$

$$T(N/4) = T(N/8) + 1$$

$$\dots$$

$$T(4) = T(2) + 1$$

$$T(2) = T(1) + 1$$

Summing up the left & right sides

$$T(N) + T(N/2) + T(N/4) + T(N/8) + \dots + T(2)$$

$$= T(N/2) + T(N/4) + T(N/8) + \dots + T(2) + T(1) + (1 + 1 + \dots + 1)$$

The no. of 1's on the right side is $\log N$

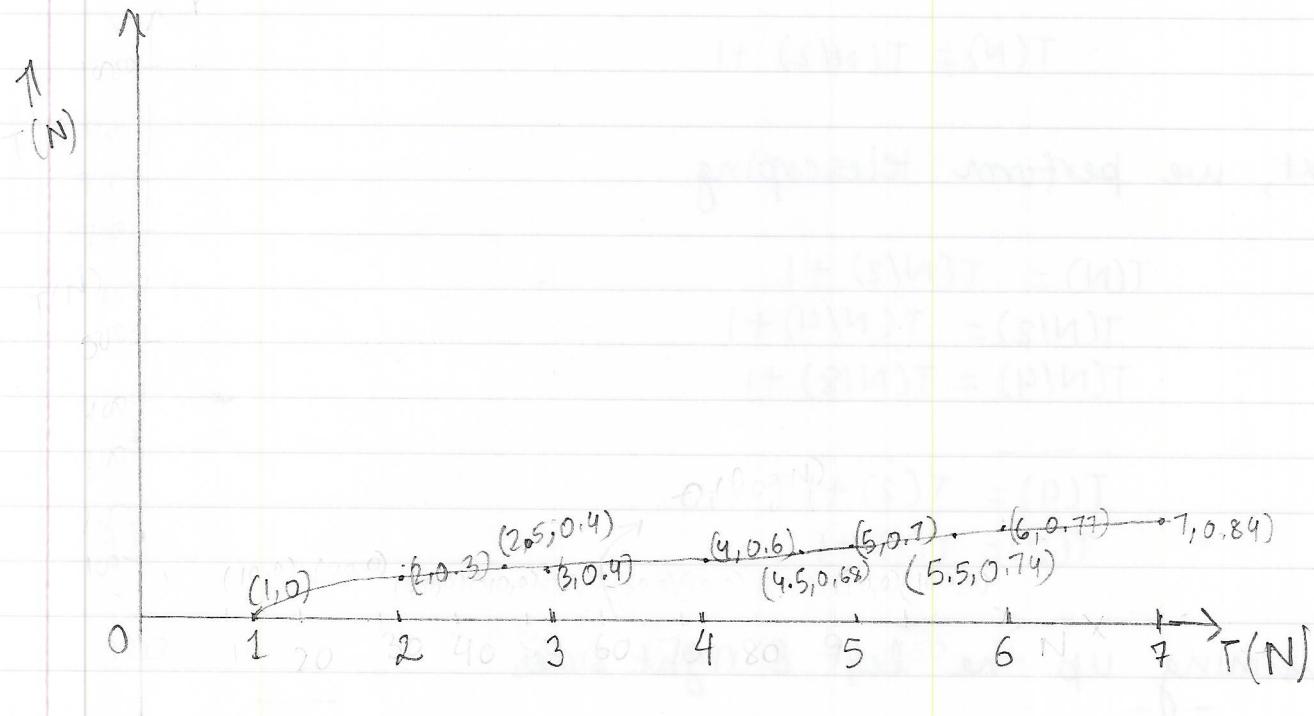
(19)

$$\Rightarrow T(N) = T(1) + \log N$$

$$\therefore T(N) = 1 + \log N$$

\therefore The runtime of binary search is

$$T(N) = O(\log N)$$



The binary search algorithm reduces the search for an element in a search sequence of n to size of $n/2$, when n is even. Hence, if $f(n)$ is the no. of comparisons required to search

$$\text{Then, } f(n) = f(n/2) + 2 \text{ when } n \text{ is even.}$$