

# Information Systems

German-Russian Institute of Advanced Technology  
(GRIAT)

Summer term 2015

Rita Schindler | TU Ilmenau, Germany  
[www.tu-ilmenau.de/dbis](http://www.tu-ilmenau.de/dbis)

## Goal of this Lecture

- ▶ Knowledge about database systems, their foundations
  - ▶ modeling & design, query formulation
- ▶ ... their implementation
  - ▶ storage & index structures, query optimization & processing, transaction processing
- ▶ ... and application
  - ▶ particularly, in data warehousing
- ▶ Prerequisites:
  - ▶ Basic understanding of relational databases, SQL

▶ 2

Information Systems | R. Schindler | GRIAT 22.04.2015

## Overview

- ▶ Introduction & Conceptual Design
- ▶ Relational Database Theory
- ▶ SQL
- ▶ Storage and Index Structures
- ▶ Query Processing
- ▶ Transaction Processing & Recovery
- ▶ Architecture & Design of Data Warehouses
- ▶ Querying Data Warehouses
- ▶ Indexing & Materialized Views

▶ 3

Information Systems | R. Schindler | GRIAT 22.04.2015

## Course Material

- ▶ Handouts
- ▶ Exercises, tasks and questions at the end of each chapter
- ▶ Recommended readings:
  - ▶ Garcia-Molina, Ullman, Widom  
Database Systems – The Complete Book  
Pearson/Prentice Hall, 2. ed., 2009
  - ▶ ...
  - ▶ C.J. Date  
Introduction to Database Systems  
Addison-Wesley, 8. edition, 2003



▶ 4

Information Systems | R. Schindler | GRIAT 22.04.2015

## Список Литературы

Издательская группа "Диалектика-Вильямс"

<http://www.dialektika.com>

Подразделение:

**Издательский дом "Вильямс"**

<http://www.williamspublishing.com/>

Гектор Гарсиа-Молина, Джеффри Д. Ульман, Дженнифер Уидом

**Системы баз данных. Полный курс**

1088 стр.; 2004; Вильямс

ISBN 5-8459-0384-X, 0-1303-1995-3;

К. Дж. Дейт

**Введение в системы баз данных. Восьмое издание**

1328 стр.; 2008; Вильямс.

ISBN 978-5-8459-0788-2, 0-321-19784-4;

► 5

Information Systems | R. Schindler | GRIAT 22.04.2015

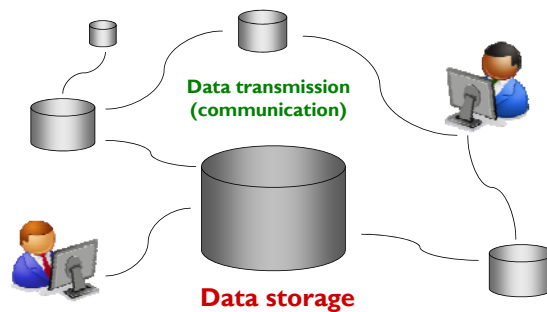
## Information Systems

### Chapter 1: Foundations & Conceptual Database Design

Rita Schindler | TU Ilmenau, Germany  
[www.tu-ilmenau.de/dbis](http://www.tu-ilmenau.de/dbis)

## Information Systems

Information exchange: Overcoming of space and time



► 7

Information Systems | R. Schindler | GRIAT 22.04.2015

## Information Systems (cont.)

- Communication and data storage – **core components** of the most complex networked systems of computer science

### Example **Amazon**

- Online-Marketplace
- Technology:
  - Linux,
  - Oracle DB-Software,
  - three databases with capacities of 7.8 TB, 18.5 TB, 24.7 TB
- Only their combination provides the most service offerings
- Our focus: **data storage**

► 8

Information Systems | R. Schindler | GRIAT 22.04.2015

## Databases

Databases today are essential to every business.

Whenever you visit a major Web site –

- ▶ Google, Yahoo!, Amazon.com, ...
- ▶ or thousands of smaller sites that provide information –

there is a **database** behind the scene serving up the information you request

The power of databases

- ▶ comes from a body of knowledge and technology that has developed over several decades
- ▶ is embodied in a specialized software package

▶ 9

Information Systems | R. Schindler | GRIAT 22.04.2015

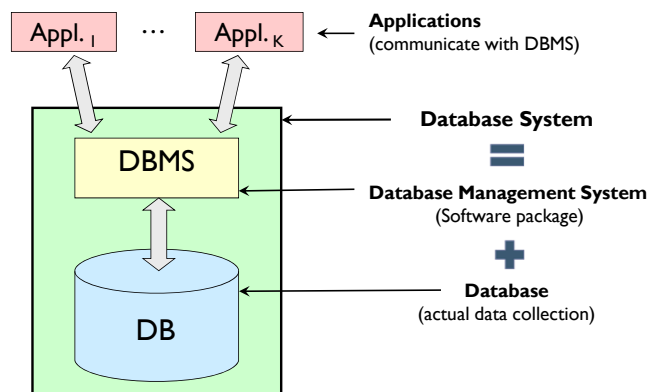
## What is a Database System?

- ▶ A **Database System** (DBS) consists of a **Database** (DB) and a **Database Management System** (DBMS)
- ▶ A Database is a (typically very large) integrated collection of interrelated data which are stored in files.
  - ▶ Data describe information and activities about one or more related organizations (portion of the real world).  
For example, an university database might contain information about
    - ▶ entities (e.g., students, courses, faculties, ...)
    - ▶ relationships among entities (e.g., 'Bill is taking Database Systems course')
- ▶ A **Database Management System**
  - ▶ is a collection of software packages designed to store, access, and manage databases.
  - ▶ It provides users and applications with an environment that is convenient and efficient to use.

▶ 10

Information Systems | R. Schindler | GRIAT 22.04.2015

## What is a Database System?



▶ 11

Information Systems | R. Schindler | GRIAT 22.04.2015

## Purpose of a Database System

- ▶ **Data Integration**: All data are uniformly managed.
- ▶ **Efficient Data Access**: Database languages are provided to store, access and manage data.
- ▶ **Data Dictionary**: Contains all data about objects and structures of the database (metadata)
- ▶ **User/Application-Views**: Different views for different users and applications
- ▶ **Integrity Constraints**: are enforced by the DBMS.
- ▶ **Security Mechanisms**: to protect data from security threats
- ▶ **Transactions**: combine sets of operations on data into logical units
- ▶ **Synchronization**: of concurrent user transactions
- ▶ **Recovery**: of data after system crash
- ▶ **ad-hoc queries**, report generation, interfaces to other database systems, interfaces for application programming, ...

▶ 12

Information Systems | R. Schindler | GRIAT 22.04.2015

## History of Relational Database Systems

- ▶ 1970: Ted Codd (IBM) → relational model as conceptual foundation of relational DBS
- ▶ 1974: System R (IBM) → first prototype of a RDBMS
  - ▶ only two modules: RDS (optimizing SQL processor), RSS (access method); approx. 80.000 LOC (PL/I, PL/S, Assembler), approx. 1,2 MB code size
  - ▶ query language SEQUEL
  - ▶ first installation 1977
- ▶ 1975: University of California at Berkeley (UCB) → Ingres
  - ▶ query language QUEL
  - ▶ predecessor of Postgres, Sybase, ...
- ▶ 1979: Oracle Version 2
- ▶ Today: 24 Bill. \$ market

▶ 13

Information Systems | R. Schindler | GRIAT 22.04.2015

## Some of the Largest Database Systems

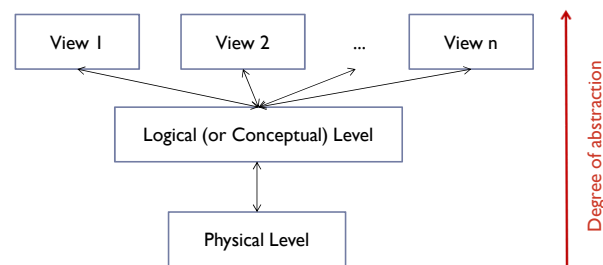
- ▶ eBay Data Warehouse: 6,5 PB (= 6,5 · 10<sup>15</sup> Bytes)
  - ▶ Teradata DBMS, >100 servers, 17 trillion records, 50 TB/day
- ▶ WalMart Data Warehouse: 2,5 PB
  - ▶ Teradata DBMS, NCR MPP hardware;
  - ▶ Product information (sales etc.) from 2.900 stores; 50.000 queries/week
- ▶ Facebook 400 TB
  - ▶ Hadoop/Hive, 610 nodes, 15 TB/day
- ▶ US Library of Congress 10-20 TB
  - ▶ not fully digitized

▶ 14

Information Systems | R. Schindler | GRIAT 22.04.2015

## Different Views of Data

- ▶ A major purpose of a DBMS is to provide users with an abstract view of data,
- ▶ i.e. it hides details of how data are stored and maintained on a computer.



▶ 15

Information Systems | R. Schindler | GRIAT 22.04.2015

## Different Views of Data (cont.)

- ▶ **Physical Level** describes how data records are actually stored and how files and indexes are organized and used
- ▶ **Logical Level** (sometimes also called Conceptual Level) describes what data are stored in the DBS in terms of entities and relationships; emphasis on logical structure of the database
- ▶ **View Level** describes how users and applications see the data
- ▶ Abstraction is achieved by describing each level in terms of a **database schema**, which, in turn, is based on a data model

▶ 16

Information Systems | R. Schindler | GRIAT 22.04.2015

## Data Models, Schemas, and Instances

- ▶ A Data Model is a collection of concepts for describing
  - ▶ data and relationships among data
  - ▶ data semantics and data constraints
- ▶ Object-Based logical Models
  - ▶ Entity-Relationship (ER) Model
  - ▶ Object-Oriented (OO) Model
- ▶ Record-Based logical Models
  - ▶ Relational Model
  - ▶ Network Model
  - ▶ Hierarchical Model

▶ 17

Information Systems | R. Schindler | GRIAT 22.04.2015

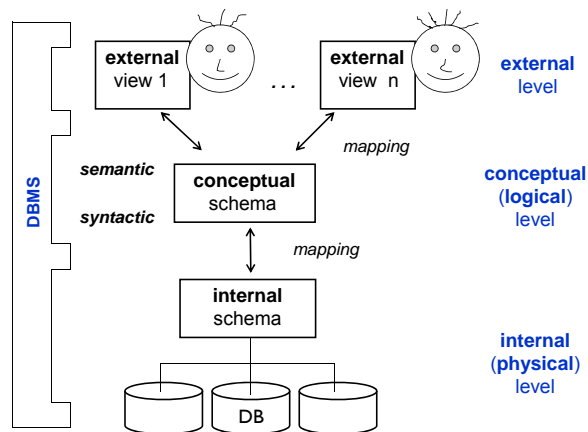
## Data Models, Schemas, and Instances (cont.)

- ▶ A database **schema** is a description of a particular collection of data, using a given data **model**.
- ▶ An **instance** of a database schema is the actual content of the database at a particular point in time.
- ▶ Schemas exist at different levels of abstraction
  - ▶ **Conceptual** (or Logical) Schema: typically builds the basis for designing a database (⇒ **main focus in this course**)
  - ▶ View (or **External**) **Schemas**: typically determined during requirements analysis (often require integration into one conceptual schema)
  - ▶ **Physical Schema**: storage structures associated with relations

▶ 18

Information Systems | R. Schindler | GRIAT 22.04.2015

## Three Level Schema Architecture



▶ 19

Information Systems | R. Schindler | GRIAT 22.04.2015

## Data Independence

- ▶ Ability to modify definition of schema at one level without affecting a schema definition at a higher level
- ▶ Achieved through the use of three levels of data abstraction (also called **three level schema architecture**)
  - ▶ **Logical Data Independence**:
    - ▶ Ability to modify logical schema without causing application programs to be rewritten
  - ▶ **Physical Data Independence**:
    - ▶ Ability to modify physical schema without causing logical schema or applications to be rewritten (occasionally necessary to improve performance)

▶ 20

Information Systems | R. Schindler | GRIAT 22.04.2015

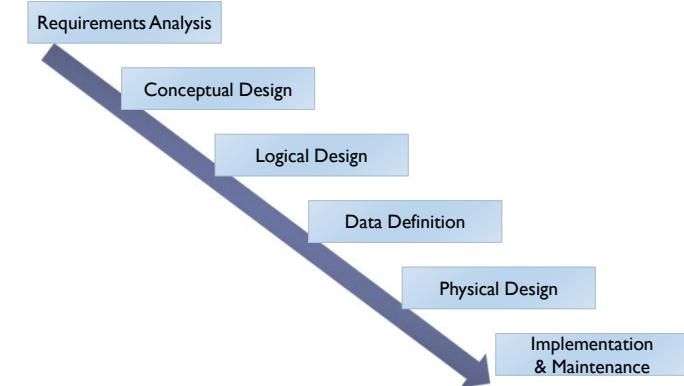
## Database Languages

- ▶ A Database Management System offers two different types of languages (for the user)
  - ▶ Data Definition Language (DDL)
    - ▶ Specification language (notation) for defining a database schema; includes syntax and semantics
    - ▶ DDL compiler generates set of tables stored in the DBMS's data dictionary (contains metadata, i.e. data about data)
    - ▶ Data storage and definition language – special type of DDL in which storage structures and access methods used by the DBS are specified
  - ▶ Data Manipulation Language (DML)
    - ▶ Language for accessing and manipulating the data that is organized according to underlying data model
    - ▶ Two classes of languages
      - Procedural – user specifies how required data is retrieved
      - Declarative – user specifies what data is required without specifying how to get those data
    - ▶ (declarative  $\triangleq$  non-procedural)

▶ 21

Information Systems | R. Schindler | GRIAT 22.04.2015

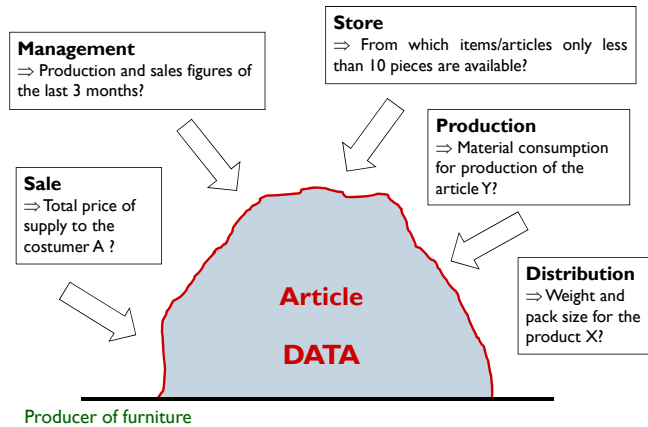
## Database Design



▶ 22

Information Systems | R. Schindler | GRIAT 22.04.2015

## Design steps by example (1.)

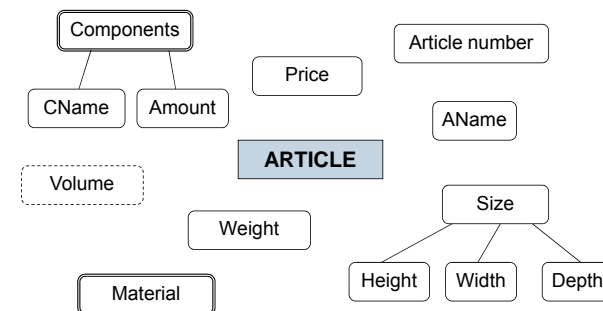


▶ 23

Information Systems | R. Schindler | GRIAT 22.04.2015

## Design steps by example (2.)

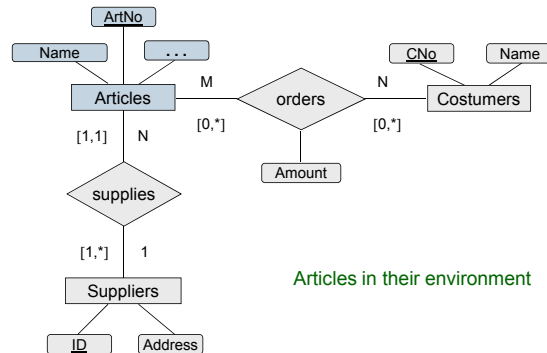
Features of interest (abstraction) of the articles (Notation according to [EN02])



▶ 24

Information Systems | R. Schindler | GRIAT 22.04.2015

### Design steps by example (3.)



► 25

Information Systems | R. Schindler | GRIAT 22.04.2015

### Design steps by example (4.)

Articles	ArtNo	Name	Price	...	Store	Quantity	...
	1357	Table	74.90		A13	13	
	2468	Bed	111.00		C06	7	
	1298	Shelf	89.90		B01	56	

Orders	CNo	ArtNo	Amount
	34567	1357	1
	31246	2468	2
	34567	1298	2



Total price for the orders of each customer?

```
SELECT CNo, SUM ( Amount * Price )
FROM Articles A, Orders O
WHERE A . ArtNo = O . ArtNo
GROUP BY CNo ;
```

► 26

Information Systems | R. Schindler | GRIAT 22.04.2015

### Exercises / Questions

- We can store (very) large collections of data in a traditional file system or in a database system. What are the distinctions?
- A database system consists of two parts. Can you characterize these?

► 27

Information Systems | R. Schindler | GRIAT 22.04.2015

### Conceptual Database Design

- The first step is the abstract representation of the structure of a database.
- Questions that are addressed during conceptual design:
  - What are the **entities and relationships** of interest (mini-world)?
  - What **information about entities and relationships** among entities needs to be stored in the database?
  - What are the **constraints** (or business rules) that (must) hold for the entities and relationships?
- Note: Design is independent of all physical considerations (DBMS, OS, ...).

► 28

Information Systems | R. Schindler | GRIAT 22.04.2015

## Entity-Relationship Data Model

- ▶ The most common model for this phase of design is the entity-relationship model.
- ▶ **Entity-Relationship model (ER model)**
  - ▶ 1976 P. P. Chen
  - ▶ "The Entity-Relationship Model – Toward a Unified View of Data" ACM Transactions on Database Systems, Vol. 1, March 1976, Pages 9-36
  - ▶ Today there are many extensions of the model.
- ▶ A database schema in the ER model can be represented pictorially (**Entity-Relationship diagram**)

## Example

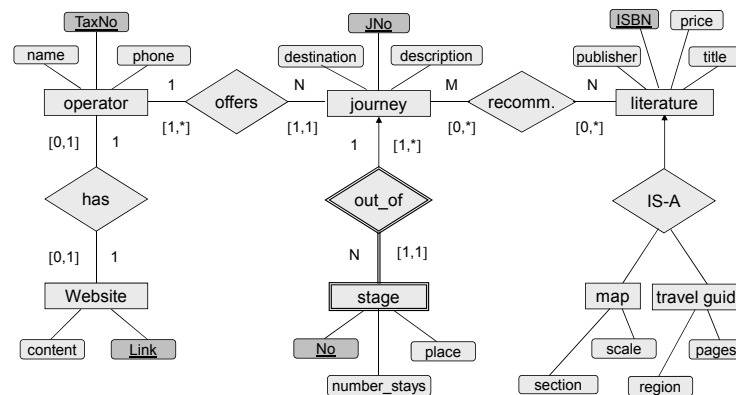
The tourist industry wants to set up a supraregional data collection containing tour operators, offered journeys, and other information which may be of interest to clients. For the **journeys** the destination and a short description are of importance. Moreover detailed information about the **stages** of a journey, i.e. concerning the places and the number of overnight stays there are meaningful.

Naturally details – like the name and the telephone number – of the **tour operators** and their **Web appearances** are needed. Of interest are the link and the content itself. Also travel **literature** – deviced into **travel guides** and **maps** shall be contained in the data collection. Out of these literature recommendations for the different journeys will be derived.

Concerning literature ISBN, title, publisher, and price shall be the characteristic features. For Maps still the section and the scale should be added.

Concerning travel guides the city or region described and the number of pages (hints for content and weight in the luggage) should be added.

## Example



## Entities



journey

values	properties
33557	number
Norway	destination
journey by mailboat	description

### Entity $e_i$

- ▶ real-world object or thing with an independent existence and which is distinguishable from other objects.
- ▶ Examples are a person, car, book, ... , examen, order, ...



## Attributes

### Attribute $a_i$

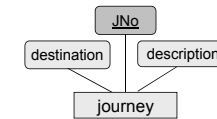


- ▶ An entity is represented by a set of attributes (its descriptive **properties**), e.g., name, age, salary, price etc.
- ▶ Attribute **values** that describe each entity become a **major part** of the data eventually stored in a **database**.
- ▶ Attribute Types:
  - ▶ simple (atomic) or composite
  - ▶ single-valued or multi-valued
  - ▶ stored or derived
  - ▶ each of the above can be an **optional** attribute, in case an entity may not have an applicable value for an attribute. → NULL
- ▶ With each attribute a **domain** is associated, i.e., a set of permitted values for an attribute.
- ▶ Possible domains are *number*, *string*, *date* etc.

▶ 33

Information Systems | R. Schindler | GRIAT 22.04.2015

## Entity Types and Sets



- ▶ **Entity Type  $E$** 
  - ▶ Collection of entities that all have the same attributes, e.g., persons, cars, customers etc.
- ▶ **Entity Set  $E$** 
  - ▶ Collection of entities of a particular entity type at any point in time; entity set is typically referred to using the same name as entity type.

▶ 34

Information Systems | R. Schindler | GRIAT 22.04.2015

## Key attributes

- ▶ Entities of an entity type need to be distinguishable.
- ▶ A **key** of an entity type  $E$  is a set  $K$  of one or more attributes whose values uniquely determine each entity in an entity set.
  - ▶ For given any two distinct entities  $e_1$  and  $e_2$  in  $E$ ,  $e_1$  and  $e_2$  cannot have identical values for each of the attributes in the key  $K$ .
  - ▶ It is possible for  $e_1$  and  $e_2$  to agree in some of this attributes, but never in all attributes.
- ▶ There also can be **more than one possible** key for an entity set. Then it is customary to pick one key as the „preferred key“, and to act as if that were the only key.
- ▶ If there are no attributes with the key property, we define an **artificial** key.

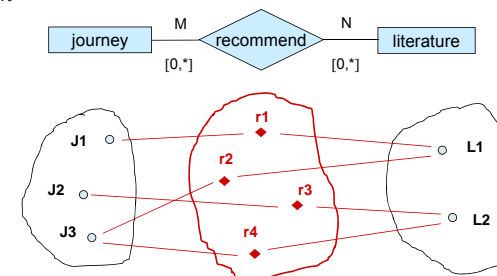
▶ 35

Information Systems | R. Schindler | GRIAT 22.04.2015

## Relationship Types and Sets

### Relationship Type $R$

- ▶ describes a set of similar relationships
- ▶ An  $n$ -ary relationship type  $R$  links  $n$  entity types  $E_1, \dots, E_n$ .



▶ 36

Information Systems | R. Schindler | GRIAT 22.04.2015

## Relationships

### Relationship (instance)

- association among two or more entities, e.g., “customer ‘Bill’ orders product ‘SkyPhone’ ”
- Each relationship in a relationship set  $R$  of relationship type involves entities  $e_1 \in E_1, \dots, e_n \in E_n$

$$R \subseteq \{ (e_1, \dots, e_n) \mid e_1 \in E_1, \dots, e_n \in E_n \},$$

where  $(e_1, \dots, e_n)$  is a relationship.



▶ 37

Information Systems | R. Schindler | GRIAT 22.04.2015

## Relationships, Types, Sets

- Degree of a relationship:** refers to the number of entity types that participate in the relationship type (binary, ternary, ...).
- Roles:** The same entity type can participate more than once in a relationship type.



- Role labels clarify semantics of a relationship, i.e., the way in which an entity participates in a relationship.  
→ **recursive relationship**.

▶ 38

Information Systems | R. Schindler | GRIAT 22.04.2015

## Relationships, Types, Sets (cont.)

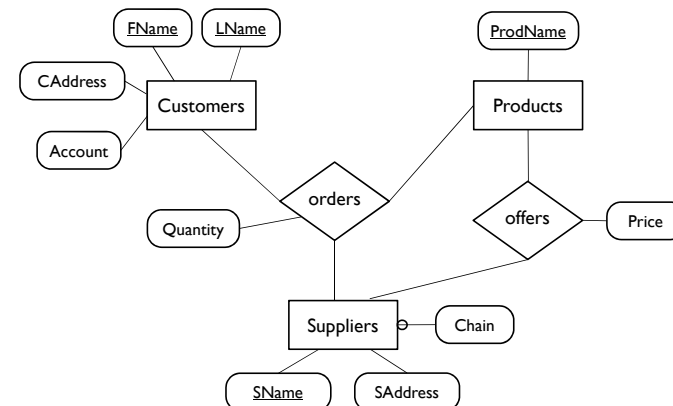
### Relationship Attributes

- A relationship type can have attributes describing properties of a relationship.
  - “customer ‘Bob’ ordered product ‘SkyPhone’ on October 15, 2014, for \$650”.
  - These are attributes that cannot be associated with participating entities only, i.e., they make only sense in the context of a relationship.
- Note that a relationship does not have key attributes! The identification of a particular relationship in a relationship set occurs through the keys of participating entities.

▶ 39

Information Systems | R. Schindler | GRIAT 22.04.2015

## Example of an Entity-Relationship Diagram



▶ 40

Information Systems | R. Schindler | GRIAT 22.04.2015

### Example of an Entity-Relationship Diagram (cont.)

- ▶ **Rectangles** represent entity types
- ▶ **Ellipses** represent attributes
- ▶ **Diamonds** represent relationship types
- ▶ **Lines** link attributes to entity types and entity types to relationship types
- ▶ **Key** attributes are underlined
- ▶ **Empty Circle** at the end of a line linking an attribute to an entity type represents an optional (null) attribute (not mentioned in textbook)
- ▶ Not in the above diagram: **Double Ellipses** represent multi-valued attributes

▶ 41

Information Systems | R. Schindler | GRIAT 22.04.2015

### Constraints on Relationship Types

- ▶ Limit the number of possible combinations of entities that may participate in a relationship set.
- ▶ There are two types of constraints:
  - ▶ **cardinality ratio** and
  - ▶ **participation constraints**
- ▶ Very useful concept in describing binary relationship types.

For binary relationships, the cardinality ratio must be one of the following types:

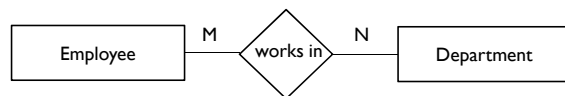
▶ 42

Information Systems | R. Schindler | GRIAT 22.04.2015

### Constraints on Relationship Types

For binary relationships, the cardinality ratio must be one of the following types:

- ▶ **Many-To-Many, N:M (default)**



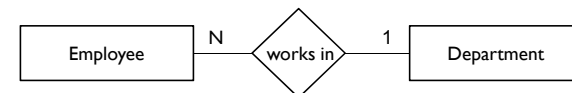
- ▶ **Meaning:** An employee can work in many departments ( $\geq 0$ ), and a department can have several employees ( $N, M \equiv \text{any number}$ ).

▶ 43

Information Systems | R. Schindler | GRIAT 22.04.2015

### Constraints on Relationship Types (cont.)

- ▶ **Many-To-One, N:1**



- ▶ **Meaning:** An employee can work in at most one department ( $\leq 1$ ), and a department can have several employees.

- ▶ **One-To-Many, 1:N**



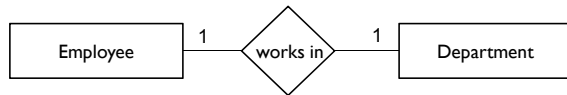
- ▶ **Meaning:** An employee can work in many departments ( $\geq 0$ ), but a department can have at most one employee.

▶ 44

Information Systems | R. Schindler | GRIAT 22.04.2015

### Constraints on Relationship Types (cont.)

#### ► One-To-One, 1:1



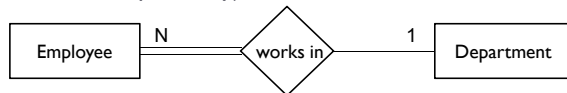
- Meaning: An employee can work in at most one department, and a department can have at most one employee.

### Constraints on Relationship Types (cont.)

- A 1:N relationship type (and the counterpart N:1) is also often called a **functional relationship**.
- Cardinality ratio of a relationship can affect the placement of a relationship attribute. E.g., in case of a N:1 relationship type, one can place a relationship attribute at a participating entity type.

### Constraints on Relationship Types (cont.)

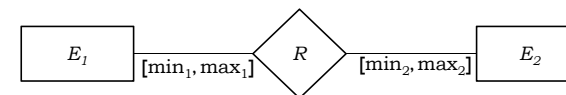
- Participation constraint: specifies whether the existence of an entity  $e \in E$  depends on being related to another entity via the relationship type  $R$ .
- **total**: each entity  $e \in E$  **must** participate in a relationship, it cannot exist without that participation (total participation aka existence dependency).



- **partial**: default; each entity  $e \in E$  **can** participate in a relationship

### Constraints on Relationship Types (cont.)

- Instead of a cardinality ratio or participation constraint, more precise cardinalities can be associated with relationship types:



- Each entity  $e_1 \in E_1$  must participate in relationship set  $R$  **at least**  $\min_1$  and **at most**  $\max_1$  times (analogous for  $e_2 \in E_2$ ).

## Frequently used cardinalities

Relationship	$[min_1, max_1]$	$[min_2, max_2]$	pictorial notation
N:M	$[0, *]$	$[0, *]$	
N:1	$[0, 1]$	$[0, *]$	
1:1	$[0, 1]$	$[0, 1]$	

► 49

Information Systems | R. Schindler | GRIAT 22.04.2015

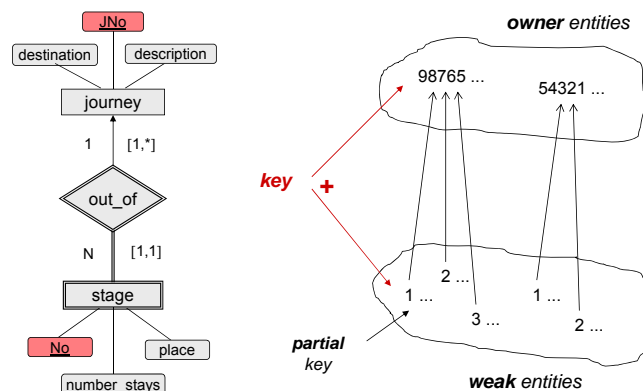
## Weak Entity Sets

- Entity type, at whose identification (key) is involved a relationship (N:1) to a owner entity type
- Existence of the weak entities depends from the existence of their owner entity
- partial key
  - attribute(s) for identification of the weak entities relating to one owner entity
- key
  - partial key of the weak entities **plus** key of the owner entity set

► 50

Information Systems | R. Schindler | GRIAT 22.04.2015

## Example

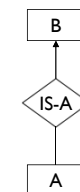


► 51

Information Systems | R. Schindler | GRIAT 22.04.2015

## IS-A Hierarchies

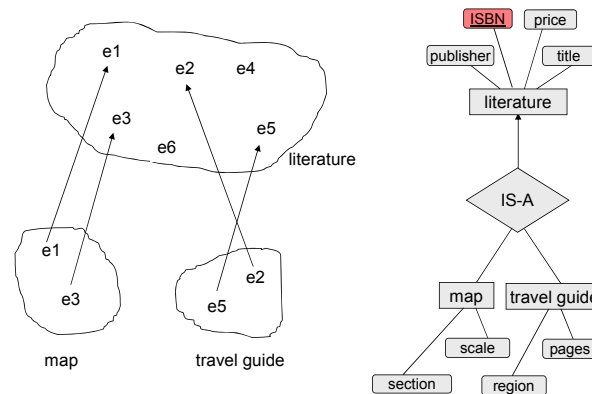
- Two entity sets A and B are in the relationship IS-A, if entity set B is a generalisation of entity set A, or equivalently, A is a special kind of B.
- Entity set A is a **subset** of entity set B.
- A inherits the attributes of B, but also has additional attributes that don't make sense for those members of B that are not also members of A.
- Technically, each entity a in set A is related to exactly one entity b in set B, such that a and b are really the same entity.
- The **key** attributes for entity set A are actually attributes of entity set B. (A inherit the from B.)



► 52

Information Systems | R. Schindler | GRIAT 22.04.2015

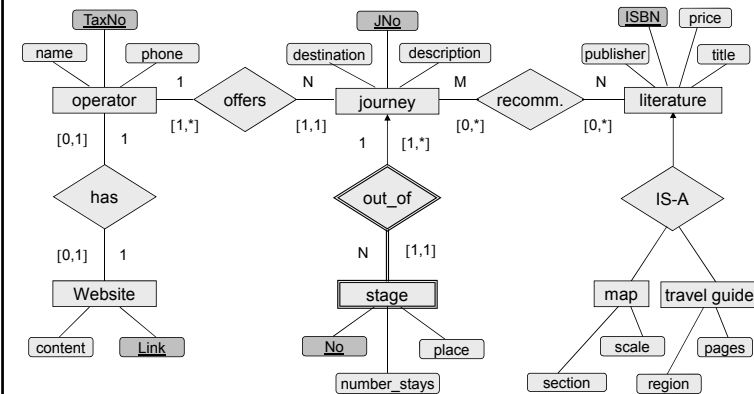
## Example



► 53

Information Systems | R. Schindler | GRIAT 22.04.2015

## Example



► 54

Information Systems | R. Schindler | GRIAT 22.04.2015

## Design Choices for ER Conceptual Design

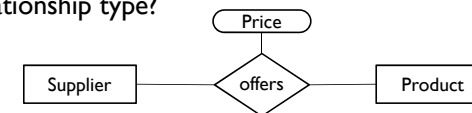
- It is possible to define entities and their relationships in a number of different ways (in the same model!).
- Should a real-world concept be modeled as an entity type, attribute, or relationship type?
- Is "Address" an attribute or an entity type? Decision depends upon the use one wants to make of address information. If one is interested in the structure, e.g., (City, Street, Zip-Code), Address must be modeled as an entity type or complex attribute.

► 55

Information Systems | R. Schindler | GRIAT 22.04.2015

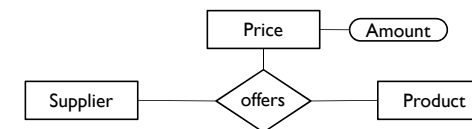
## Design Choices for ER Conceptual Design

- Should a concept be modeled as an entity type or relationship type?



- Here a supplier cannot offer the same product for different prices! Why?

- Modeling price as an entity type resolves this problem:



► 56

Information Systems | R. Schindler | GRIAT 22.04.2015

## Steps in Designing an Entity-Relationship Schema

1. Identify entity types (entity type vs. attribute)
2. Identify relationship types
3. Identify and associate attributes with entity and relationship types
4. Determine attribute domains
5. Determine key attributes for entity types
6. Associate (refined) cardinality ratio(s) with relationship types

## Exercises / Questions

## Exercises: Problem 1

In the following table, E1 and E2 stand for entity types and R stands for a relationship type between two respective entity types. Characterize each relationship type by determining:

1. the cardinality ratio
2. the participation constraints for entities of E1 in R
3. the participation constraints for entities of E2 in R

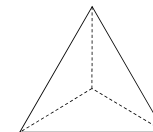
Discuss possible semantic alternatives (what might this relationship mean?) as well as necessary additional semantic information (what else do we need to know?).

Model these relationship types in ER diagrams and add attributes where necessary.

E1	R	E2	Comments
Manufacturers	supply	Parts	Price is specified
Employees	works_in	Rooms	No more than 3 in one room
Students	to-be-friends	Students	with attribut <i>since</i>
Students	are	Persons	
Supermarkets	have	Chain stores	

## Exercises: Problem 2

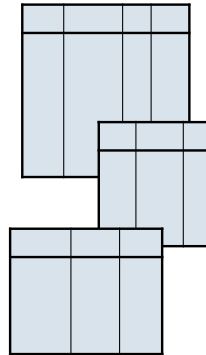
A **Polyhedron** is defined by the hull of its bounding **Faces**. These faces are defined by their bounding **Edges**. An edge is defined by its start and end **Points** in three-dimensional space. (You can imagine a tetrahedron as the simplest form of polyhedrons.)



Design a suitable ER-diagram and specify the respective functionalities and cardinalities. (Attributes can be ignored.)

## The relational data model

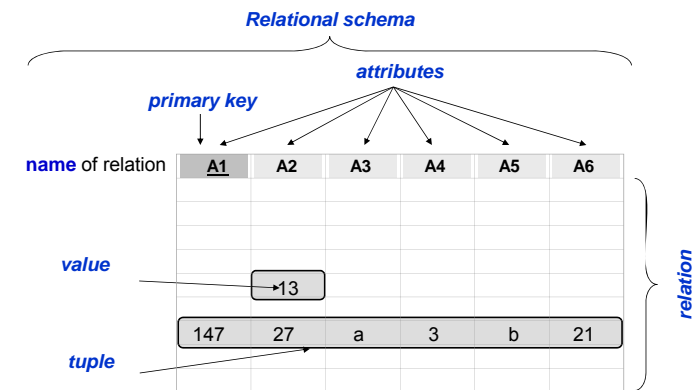
- ▶ Logical data model
- ▶ E. F. Codd (IBM) – 1970
- ▶ Represent data as two-dimensional tables called **relations**.
- ▶ Each row (tuple) of a table describes with the values in the columns an entity or a relationship



▶ 61

Information Systems | R. Schindler | GRIAT 22.04.2015

## Relations



▶ 62

Information Systems | R. Schindler | GRIAT 22.04.2015

## Basics of the Relational Model

### Domains $D_i$

- ▶ set of permitted (**atomic**) values for an attribute  $A_i$

### Schema $R (A_1, A_2, \dots, A_n)$

- ▶ descriptive attributes
  - ▶  $R$  – name of relation
  - ▶  $n$  – degree of relation (count of attributes)

### Relation $r (R)$

- ▶ finite set of  $n$ -tuples  $r = \{t_1, t_2, \dots, t_m\}$
- ▶  $m$  – number of tuples

### $n$ -tuple $t$

- ▶ list (ordered set) of  $n$  values from the domains corresponding to attributes of the schema respectively
- ▶  $t_k = \langle v_1, v_2, \dots, v_n \rangle$  with  $v_i \in D_i$ ,  $i = 1, \dots, n$ ,  $k = 1, \dots, m$

▶ 63

Information Systems | R. Schindler | GRIAT 22.04.2015

## Keys

### Key

- ▶ a set of attributes, whose values can identify the tuples (rows), but no subset has this property → a key must be minimal
- ▶ For a relation several candidate keys may exist.

### Primary Key (PK)

- ▶ If there are more than one candidate keys, one of those is selected to be the **primary key** of the relation.

### Foreign Key (FK)

- ▶ An attribute or attributes of one relation, referencing some key of a second („foreign“) relation
- ▶ Values of the foreign key appearing in the first relation must also appear in the referenced attributes of the second relation.

▶ 64

Information Systems | R. Schindler | GRIAT 22.04.2015



## Translation of ER Schema into Tables

- ▶ An **ER schema** can be represented by a **collection of tables** which represent contents of the database (instance).
- ▶ **Primary keys** allow entity types and relationship types to be expressed uniformly as tables.
- ▶ For each entity and relationship type, a **unique table** can be derived which is assigned the name of the corresponding entity or relationship type.
- ▶ Each table has a **number of columns that correspond to the (atomic) attributes** and which have unique names. An attribute of a table has the same domain as the attribute in the ER schema.
- ▶ Translating an ER schema into a collection of tables is the basis for deriving a relational database schema from an ER diagram.

▶ 65

Information Systems | R. Schindler | GRIAT 22.04.2015

## Translating Entity Types into Tables

- ▶ Given an entity type  $E_i$  with (atomic) attributes  $A_1, \dots, A_n$  and associated domains  $D_1, \dots, D_n$ .
- ▶ Translation of the entity type CUSTOMER into a table

FName	LName	CAddress	Account
Bob	Smith	Ilmenau	5,000
Michael	Gordon	Erfurt	2,000
...	...	...	...

```
create table Customers (
  FName varchar(40), LName varchar(40),
  CAddress varchar(100), Account real);
```

- ▶ A row in such a table corresponds to an entity from the entity set.

▶ 66

Information Systems | R. Schindler | GRIAT 22.04.2015

## Translating Relationship Types into Tables

- ▶ A M:N relationship type is represented as a table with columns for the primary key attributes of the participating entity types, and any descriptive attributes of the relationship type.

- ▶ Example: Relationship type offers

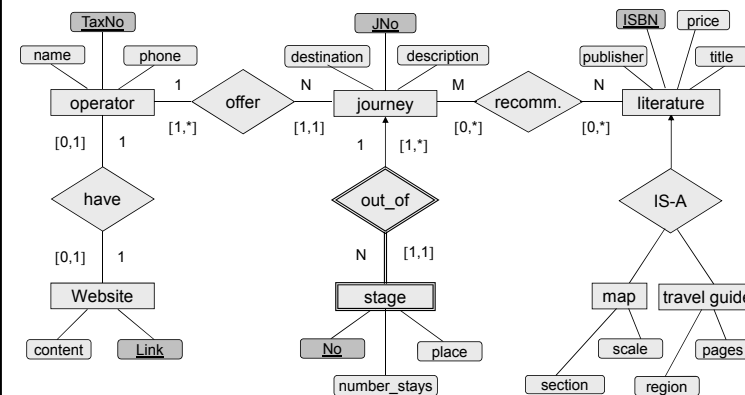
ProdName	SName	Price
SkyPhone	HaI-Mart	650
Super Tablet	K+B	1,200
...	...	...

- ▶ ProdName and SName are the primary key attributes of the entity types PRODUCTS and SUPPLIERS respectively.
- ▶ Translation of 1:N and N:1 (functional) and 1:1 relationship types into tables sometimes can be optimized  $\rightarrow$  no table for relationship type necessary!

▶ 67

Information Systems | R. Schindler | GRIAT 22.04.2015

## Example Tourism



▶ 68

Information Systems | R. Schindler | GRIAT 22.04.2015

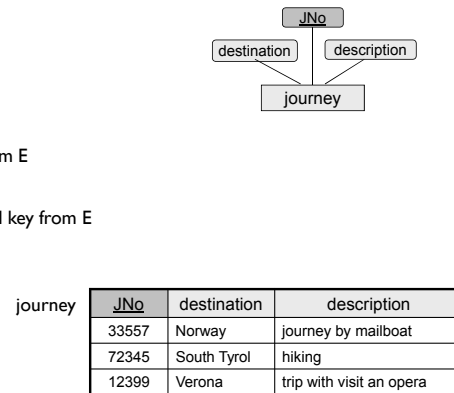
## Example: entity types

### Entity Type E



### Relation

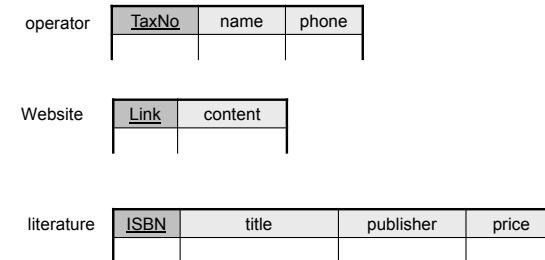
- ▶ Attributes  
attributes from E
- ▶ Primary key  
(one) minimal key from E



▶ 69

Information Systems | R. Schindler | GRIAT 22.04.2015

## Example



▶ 70

Information Systems | R. Schindler | GRIAT 22.04.2015

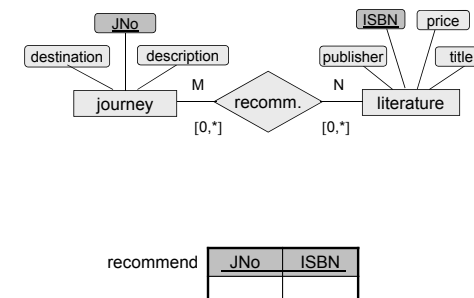
## Example: relationship types

### Relationship Type R



### Relation (in addition to the relations for the entity types)

- ▶ Attributes  
primary key attributes of the participating entity types (in function of foreign keys)  
**plus**  
attributes of the relationship itself
- ▶ Primary key  
 $M:N \Rightarrow$  common primary key consisting of primary keys of both sides  
 $1:N \Rightarrow$  primary key of the N-Side  
 $1:1 \Rightarrow$  one primary key (both are candidate keys !)



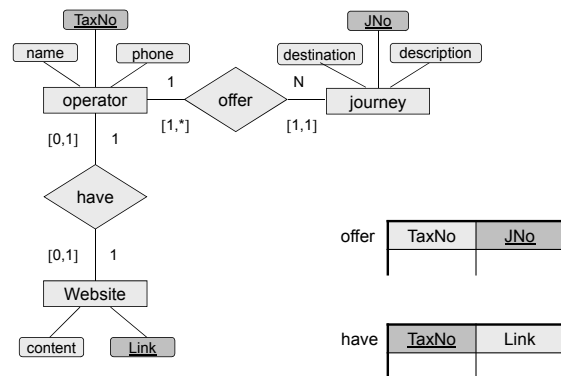
▶ 71

Information Systems | R. Schindler | GRIAT 22.04.2015

▶ 72

Information Systems | R. Schindler | GRIAT 22.04.2015

## Example



73

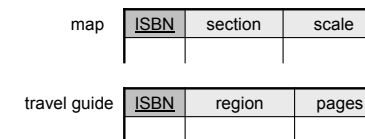
Information Systems | R. Schindler | GRIAT 22.04.2015

## Example: specialized entity types

### Specialized entity type

### Relation (for specialized entity type and IS-A relationship)

- Attributes
  - attributes of specialized entity type
  - plus** key attribute(s) of the supertype
- Primary key
  - primary key of the supertype



74

Information Systems | R. Schindler | GRIAT 22.04.2015

## Example: weak entity types

### Weak entity type W

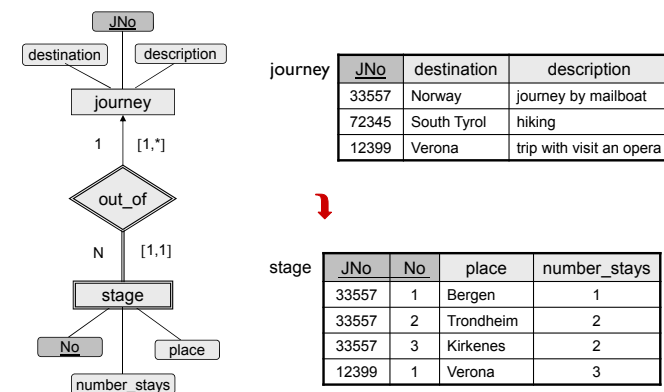
### Relation (for weak entity type and its relationship to owner)

- Attributes
  - attributes of the weak entity type
  - plus** primary key attribute(s) of the owner entity type
- Primary key
  - common primary key consisting of partial key and the owner's primary key

75

Information Systems | R. Schindler | GRIAT 22.04.2015

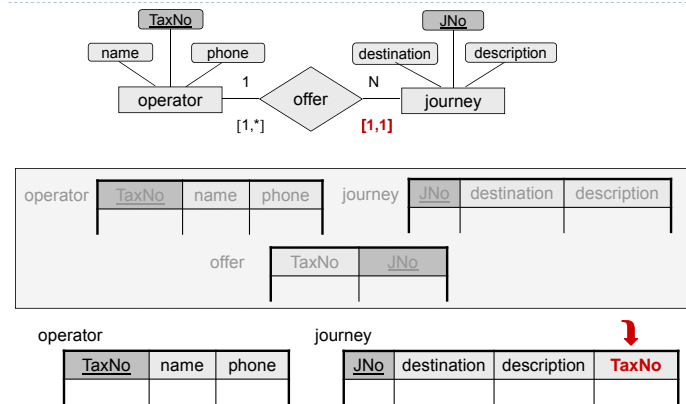
## Example



76

Information Systems | R. Schindler | GRIAT 22.04.2015

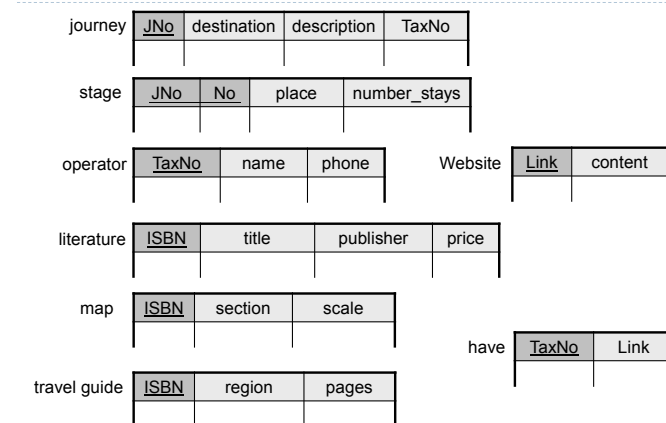
### Example: cardinality constraint [1,1]



▶ 77

Information Systems | R. Schindler | GRIAT 22.04.2015

### Example: Database Schema



▶ 78

Information Systems | R. Schindler | GRIAT 22.04.2015

### Summary of Conceptual Design

- ▶ Conceptual design follows requirements analysis, yields a high level description of data to be stored (conceptual level).
- ▶ ER model is a popular model for conceptual design, constructs are expressive, close to the way people think about applications; supported by many CASE tools.
- ▶ Basic constructs are entities, relationships, and attributes. Some additional constructs: IS-A hierarchies, cardinality ratios, ...
- ▶ There are many variations on ER model constructs.

▶ 79

Information Systems | R. Schindler | GRIAT 22.04.2015

### Summary of Conceptual Design (cont.)

- ▶ Several kinds of integrity constraints can be expressed in the ER model: key constraints, structural constraints, constraints on specializations
  - ▶ Some of them can be expressed in SQL when translating entity and relationship types into tables
  - ▶ Not all constraints can be expressed in the ER model
  - ▶ Constraints play an important role in determining a good database design for an application domain.

▶ 80

Information Systems | R. Schindler | GRIAT 22.04.2015

### Summary of Conceptual Design (cont.)

- ▶ ER design is subjective: There are many ways to model a given scenario! Analyzing alternative schemas is important!
- ▶ Entity type vs. attribute, entity type vs. relationship type, binary vs. n-ary relationship type, use of IS-A, generalization and specialization, . . .
- ▶ Ensuring a good database design includes analyzing and further refining relational schema obtained through translating ER schema.

▶ 81

Information Systems | R. Schindler | GRIAT 22.04.2015

### Summary

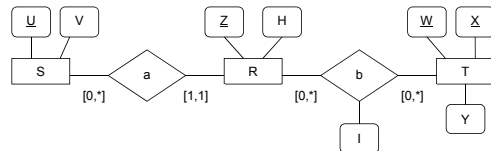
- ▶ databases, DBMS, data model, and schema
- ▶ data independence by three level schema architecture
- ▶ database design process
- ▶ conceptual design using the ER model
- ▶ translating conceptual schema into relational schema (tables)

▶ 82

Information Systems | R. Schindler | GRIAT 22.04.2015

### Exercises (1)

- ▶ Design a relational database schema for the following ER Diagram, i.e. the necessary relational schemas with their necessary attributes and their primary keys underlined.
- ▶ Consider possible optimizations to the database schema (reduction of tables).

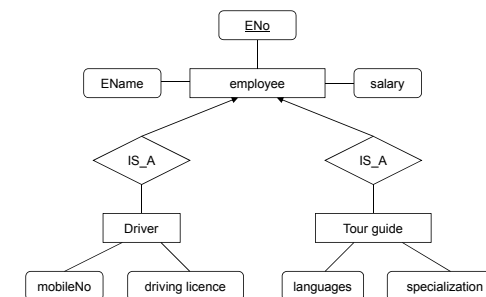


▶ 83

Information Systems | R. Schindler | GRIAT 22.04.2015

### Exercises (2)

- ▶ Design a relational database schema for the following ER Diagram, i.e. the necessary relational schemas with their necessary attributes and their primary keys underlined.



▶ 84

Information Systems | R. Schindler | GRIAT 22.04.2015

### Exercises (3)

- Design a relational database schema for the following ER Diagram.

