# Theorems for long division and root extraction

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#### 1 Notations

$$\mathbb{N} = \{0, 1, 2, \ldots\}.$$

#### Three-by-two theorem $\mathbf{2}$

This theorem is extremely useful for basecase division.

**Theorem 1** (three-by-two). Fix  $B \in \mathbb{N}$ , B > 1, the base of our positional number system. Fix also numbers  $x \in \mathbb{R}$ ,  $y \in \mathbb{R}$  such that:

- $0 < x < B^3$ :
- $B \le y \le B^2$ ;
- x/y < B.

Define:

- $q = \lfloor \frac{x}{y} \rfloor$ , the true quotient;
- $q_e = \lfloor \frac{\lfloor x \rfloor}{\lfloor y \rfloor} \rfloor$ , our estimate of the quotient.

Then either  $(q = q_e)$  or  $(q = q_e - 1)$ .

Proof.

Lemma 1.  $q \leq q_e$ .

Proof. Define  $\delta = x - \lfloor x \rfloor$ ; note that  $0 \le \delta < 1$ . We have  $\frac{x}{y} \le \frac{x}{\lfloor y \rfloor} = \frac{\lfloor x \rfloor + \delta}{\lfloor y \rfloor}$ . Then  $q = \lfloor \frac{x}{y} \rfloor \le \lfloor \frac{\lfloor x \rfloor + \delta}{\lfloor y \rfloor} \rfloor$ . We now want to prove  $\lfloor \frac{\lfloor x \rfloor + \delta}{\lfloor y \rfloor} \rfloor = \lfloor \frac{\lfloor x \rfloor}{\lfloor y \rfloor} \rfloor$ . Since  $\delta < 1$ , for any integers M, N, K, the following holds:  $(M < KN) \Longrightarrow ((M + \delta) < KN)$ . Substitute  $M = \lfloor x \rfloor$ ,  $N = \lfloor y \rfloor$ ,  $K = \lfloor \frac{M}{N} \rfloor$ .

Substitute 
$$M = \lfloor x \rfloor$$
,  $N = \lfloor y \rfloor$ ,  $K = \lfloor \frac{M}{N} \rfloor$ .

**Lemma 2.** |x| < B(|y| + 1).

*Proof.* 
$$|x| \le x < By < B(|y| + 1).$$

Define now the following values:

- $u = \lfloor x \rfloor;$
- v = |y|;
- $q_{\max} = \frac{u+1}{v}$ ;
- $q_{\min} = \frac{u}{v+1}$ .

Lemma 3. The following bounds hold:

- 1.  $q_{\text{max}} \frac{u}{v} \le \frac{1}{B}$ ;
- 2.  $\frac{u}{v} q_{\min} < 1$ .

*Proof.* 1.  $\frac{u+1}{v} - \frac{u}{v} = \frac{1}{v} \le \frac{1}{B}$ ;

2.  $\frac{u}{v} - \frac{u}{v+1} = \frac{u}{v(v+1)}$ . By lemma 2, u < B(v+1), so  $\frac{u}{v} - q_{\min} < \frac{B(v+1)}{v(v+1)} = B/v \le 1.$ 

**Lemma 4.**  $q \ge q_e - 1$ .

*Proof.* We have:

- $q_{\min} < \frac{x}{y} < q_{\max};$
- $q_{\min} < \frac{u}{v} < q_{\max}$ .

Taking floor of both sides of these inequalities, we get:

- $\lfloor q_{\min} \rfloor \leq q \leq \lfloor q_{\max} \rfloor$ ;
- $\lfloor q_{\min} \rfloor \le q_e \le \lfloor q_{\max} \rfloor$ .

By lemma 3,  $q_{\text{max}} - q_{\text{min}} < 1 + \frac{1}{B} < 2$ . This means that  $\lfloor q_{\text{max}} \rfloor - \lfloor q_{\text{min}} \rfloor$  is either 0, 1, or 2. We are only interested in the case of it being 2, as in other cases  $(q \geq q_e - 1)$  holds automatically.

Suppose  $\lfloor q_{\text{max}} \rfloor - \lfloor q_{\text{min}} \rfloor = 2$  and  $q_e - q = 2$ . Then,

- $\lfloor q_{\max} \rfloor = q_e = \lfloor \frac{u}{v} \rfloor$ , which implies  $\frac{u}{v} \ge q_e = q + 2$ ;
- $\lfloor q_{\min} \rfloor = q$ , which implies  $q_{\min} < q + 1$ .

Together, these statements imply  $\frac{u}{v} - q_{\min} > 1$ , contradicting lemma 3.  $\square$ 

# 3 Approximation of inverse theorem

This theorem is useful for calculating the inverse of a number with Netwon's method; namely, it tells us how to find the initial approximation of the inverse.

**Theorem 2** (approximation of inverse). Fix  $B \in \mathbb{N}$ , B > 1, the base of our positional number system. Fix then  $n \in \mathbb{N}$ , n > 0, the number of words in our initial approximation. Fix a number  $x \in \mathbb{R}$  such that  $B^n \leq x < B^{n+1}$ . Define:

- $r = \frac{B^{2n}}{x}$ , the true inverse (scaled up by 2n places);
- $r_e = \lfloor \frac{B^{2n}}{|x|+1} \rfloor$ , our estimate of the scaled-up inverse.

Then:

- $r 2 < r_e < r$ ;
- $\bullet \ B^{n-1} \le r_e < B^n.$

*Proof.*  $r_e < r$  is trivial: we increased the denominator  $(\lfloor x \rfloor + 1 > x)$  and then took floor of the fraction.

Define now the following values:

- $\bullet u = |x|;$
- $\bullet \ r' = \frac{B^{2n}}{n+1}.$

Note that  $r_e = \lfloor r' \rfloor$ . Then  $r - r' = \frac{B^{2n}}{u(u+1)} \leq \frac{B^{2n}}{B^{2n} + B^n} < 1$ . Now we have

$$r - r_e = (r - r') + (r' - \lfloor r' \rfloor) < 1 + 1 = 2.$$

We can prove  $B^{n-1} \leq r_e < B^n$  by substituting the maximum and minimum possible values of  $\lfloor x \rfloor + 1$  into  $r_e = \lfloor \frac{B^{2n}}{\lfloor x \rfloor + 1} \rfloor$ . The maximum possible value,  $B^{n+1}$ , gives us  $r_e \geq B^{n-1}$ ; and the minumum possible value,  $B^n + 1$ , gives us  $r_e \leq B^n - 1$ .

## 4 Root extraction

We are given  $d \in \mathbb{N}$  and root order  $n \in \mathbb{N}, n \geq 2$ . We need to calculate  $\lfloor \sqrt[n]{d} \rfloor$ . Define the "true" root  $\xi = \sqrt[n]{d}$ . We are going to iterate  $y \mapsto \Phi(y)$ , where

$$\Phi(y) = \frac{1}{n} \left( \frac{d}{y^{n-1}} + (n-1) \cdot y \right).$$

If  $y = \xi \cdot \delta$ , then  $\Phi(y) = \xi \cdot \varphi(\delta)$ , where

$$\varphi(x) = \frac{1 + (n-1)x^n}{nx^{n-1}}.$$

**Theorem 3** (icky).  $1 < \varphi(x) < x \text{ for } x > 1$ .

*Proof.* We have

$$\varphi(x) < x \Leftrightarrow 1 + (n-1)x^n < nx^n \Leftrightarrow 1 - x^n < 0.$$

Now we will prove  $\varphi(x) > 1$ .

$$\frac{1+(n-1)x^n}{nx^{n-1}} > 1 \Leftrightarrow (1+\varepsilon)^{n-1}(\varepsilon(n-1)-1) > -1, \text{ where } \varepsilon = x-1 > 0.$$
 Substituting  $\lambda = \varepsilon(n-1) > 0$  and  $m = n-1$ , we get

$$(1 + \frac{\lambda}{m})^m (\lambda - 1) > -1.$$

The sequence  $E_m = (1 + \frac{\lambda}{m})^m$  increases monotonically for  $\lambda > 0$ , and  $\lim_{m \to \infty} E_m = e^{\lambda}$ . This means  $0 < (1 + \frac{\lambda}{m})^m < e$ ; we are now going to prove

$$e^{\lambda}(\lambda - 1) > -1$$

for  $\lambda > 0$ .

The derivative  $\frac{d}{d\lambda}e^{\lambda}(\lambda-1)=e^{\lambda}\cdot\lambda$  is positive for  $\lambda>0$ ; and  $e^{\lambda}(\lambda-1)=-1$  for  $\lambda=0$ .

**Theorem 4** (root extraction). Consider now the following process: we start with an arbitrary integer  $y_0 \ge \xi$ , and then, while  $y_i > \xi$ , put  $y_{i+1} = \lfloor \Phi(y_i) \rfloor$ . This process will terminate at some finite step  $k \ge 0$  with  $y_k = \lfloor \xi \rfloor$ .

*Proof.* Note that  $\Phi(y) = \xi \varphi(y/\xi)$ .

**Lemma 5.**  $\lfloor \Phi(y_i) \rfloor < y_i \text{ for any integer } y_i > \xi.$ 

Proof. 
$$\lfloor \Phi(y_i) \rfloor \leq \Phi(y_i) < y_i$$
.

**Lemma 6.** If, for some integer  $y_i$ , we have  $y_i > \xi$  and  $y_{i+1} = \lfloor \Phi(y_i) \rfloor \leq \xi$ , then  $y_{i+1} = \lfloor \xi \rfloor$ .

*Proof.* We have 
$$y_{i+1} = \lfloor \Phi(y_i) \rfloor \leq \xi < \Phi(y_i)$$
.

Note that  $(y > \xi) \Leftrightarrow (y^n > d)$ , and

$$\lfloor \Phi(y) \rfloor = \lfloor (\lfloor d/y^{n-1} \rfloor + (n-1) \cdot y)/n \rfloor.$$