# 876. Middle of the Linked List

| : <b>≡</b> Tags |                     |
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| Property        | @September 12, 2022 |

## Question

#### 原文:

Given the **nead** of a singly linked list, return the middle node of the linked list.

If there are two middle nodes, return **the second middle** node.

#### 我的理解:

給定一串鏈結陣列,返回中間的node的數值,如果node個數是偶數個有返回後面那個 node數值

#### 翻譯:

给出一个单链表的 "头",返回*链表的中间节点*。 如果有两个中间节点,返回**第二个中间**节点。

自評翻譯正確性:70% (要返回的是包含該node與其後方的整串)

Word Memory :

### Code

```
/**
 * Definition for singly-linked list.
 * struct ListNode {
 * int val;
 * ListNode *next;
 * ListNode() : val(0), next(nullptr) {}
 * ListNode(int x) : val(x), next(nullptr) {}
 * ListNode(int x, ListNode *next) : val(x), next(next) {}
 * };
 */
class Solution {
```

```
public:
    ListNode* middleNode(ListNode* head) {
        ListNode * pre=head;
        ListNode * pro=head;
        int i=0, j=0;
        if(head->next==NULL){
            return head;
        while(pre!=NULL){
            pre=pre->next;
            i++;
        }
        i=i/2;
        while(j<i){
            pro=pro->next;
            j++;
        return pro;
    }
};
```

思路:先開一個 pre 去遍歷抓整個list長度,紀錄為i,然後i除2, 之後再用一個 pro 遍歷並用 j 記步,一直到j==i就是一半的位置

#### Success Details >

Runtime: 0~ms, faster than 100.00% of C++ online submissions for Middle of the Linked List.

Memory Usage:  $7.1\,$  MB, less than 71.98% of C++ online submissions for Middle of the Linked List.

Next challenges:

Delete the Middle Node of a Linked List

Maximum Twin Sum of a Linked List

Show off your acceptance:







| Time Submitted   | Status   | Runtime | Memory | Language |
|------------------|----------|---------|--------|----------|
| 09/11/2022 22:51 | Accepted | 0 ms    | 7.1 MB | срр      |

### 優良code參考

```
class Solution {
public:
    ListNode* middleNode(ListNode* head) {
        ListNode *slow = head, *fast = head;
        while (fast && fast->next) {
            slow = slow->next;
            fast = fast->next->next;
        }
        return slow;
    }
};
```

思路:分成兩個pointer下去跑,slow & fast,fast跑的是slow的兩倍快,所以如果fast已經抵達終點或者說已經不能再往下跑了停止時,slow跑的剛好是fast的一半,也就剛好是在一半的位置。 (great solution)