# Operating Systems: Midterm Review

CSC-4320/6320 -Summer 2014

# What to Study?

- Slides, notes, and reading assignments up to and including "Deadlocks"
- Past homework assignments
- Past programming assignments
- There should be very few surprises
- Type of Questions:
  - general "what do you think?" questions
  - homework-like problems
- Exam will be available for 24 hours starting 6:00 p.m.
   Thursday.
  - You are expected to work alone
  - You may use your book, notes, google
    - I am more interested in seeing that you can find the answers and have an understanding of the topic than I am in your ability to memorize.

 What is the difference between an Interrupt and an Exception/Trap? Give two examples of each

- What is the difference between an Interrupt and an Exception/Trap? Give two examples of each
  - Both are events that the OS must react to
  - An Interrupt is due to some external event
    - e.g., keyboard input
    - e.g., disk operation completion
  - A trap/exception is caused by an instruction
    - e.g., illegal memory access
    - e.g., system call instruction

 What are the advantages/disadvantages of User vs. Kernel Threads?

- What are the advantages/disadvantages of User vs. Kernel Threads?
  - User-level
    - good: very low overhead for creation/maintenance because the kernel's not involved
    - bad: cannot take advantage of multi-core architectures
    - bad: if one thread blocks, they all block
  - Kernel-level
    - bad: high overhead because the kernel is involved
    - removes the two "bad" of user threads

Explain what happens during a context-switch

- Explain what happens during a context-switch
  - Assume A is running and we are switching to B because A's time quantum has expired
  - Save A's "state" (registers) into its PCB structure in the kernel
  - Move A's PCB to the "Ready" queue
  - Restore B's "state" (registers) into the CPU by reading it from B's PCB
  - Move B's PCB to the "Running" queue
  - Tell the CPU to start a new fetch-decode-execute cycle
    - The PC register contains the address of the next instruction that B must execute

- What is a zombie?
- What is an orphan?

- What is a zombie?
  - A process that has terminated (i.e., called "exit()"), but whose parent has not acknowledged the death (i.e., not called wait() or waitpid()).
  - It is kept around so that later the parent can find out its return value (e.g., exit(42));
- What is an orphan?
  - A process whose parent has died
  - It is adopted by process PID=1 (meaning that its PPID=1)

- What is priority inversion?
- Why should context-switching overhead be low?
- What happens if at the end of a time quantum of a running process the ready queue is empty?
- etc. etc.

- There may be a general question about the programming assignments
- Just to make sure that you've done them and understood concepts in there
  - Example: What system program is used to trace system calls on a Linux system? strace
- Example: What can be done to force the init process to adopt a child process you create? Create a child process that creates a child and then calls exit, thus the grandchild will be adopted by init.

### Preparing for General Questions

- Most of the lectures consist of topics that beg for a "how does this work?" question
- Each such topic corresponds to an OS concept
  - There are not that many such concepts in each lecture
- Go through general questions in homework assignments, questions in the chapter practice exercises may be a good place to review your understanding of the chapter

### **Problems**

- All problems will resemble homework problems with maybe a few more creative questions added
  - Given a program with calls to fork(), what happens?
  - Given a set of jobs with arrival time, burst time, and priority, what do various scheduling algorithms do?
  - Given a multi-threaded program, what are the possible outputs?
  - Given a multi-threaded program, how would you add locks?
  - Given some scenario, is there a possible deadlock?

### **Problems**

- Some problems will be: "look at that code, and tell me what it does"
  - either pseudo-code or C++/Java code

### Concurrency

- Since we didn't have a concurrency
   Homework assignment, we can do a few "similar" concurrency examples..
- Spinning locks vs. blocking locks
- Interleaved output
- Locks
- Deadlocks

# Spinning vs. Blocking Locks

#### • Spinning:

- burn CPU cycles checking the lock continuously, which is wasteful
- But as soon as the lock is released by whoever had it you grab it, which is good

#### Blocking lock:

- Go to "sleep" asking for somebody to wake up when the key's ready, which avoids being a useless CPU hog
- But this requires much more work as the OS is now involved to put you to sleep and wake you up
  - Moving your PCB from various queues, etc.

#### Rules of thumb:

- Spinning for a long time is wasteful
- Blocking if the lock is never taken for a long time is wasteful

- Two threads, two global variables a=1 and b=1
- Thread #1: b++; a=b-a
- Thread #2: a++
- What are the possible outputs?
- 3 "clean" interleavings:
  - a++; b++; a=b-a (0,2)
  - b++; a++; a=b-a (0,2)
  - b++; a=b-a; a++ (2,2)
- There can also be bad interleaving of the non-atomic "a++" and "a=b-a" "instructions" since they both write to the same variable
- These can happen only after b++ has executed and then we're left with both "a++" and "a=b-a" to execute

- What are the "bad" interleavings for "a++" and "a=ba"?
  - Remember that at this point b=2

#### • Case #1:

- "a++" load the value of a, which is 1, computes 2, and is about to write back to a, but there is a context switch
- "a=b-a" loads the value of b (2) and of a(1), computes ba=1, and is about to write back to a, but there is a context switch
- "a++" resumes and writes back value 2 to a
- "a=b-a" resumes and overwrites value 2 by value 1
- So in the end we have (1,2)

- What are the "bad" interleavings for "a++" and "a=ba"?
  - Remember that at this point b=2
- Case #2
  - "a=b-a" loads the value of b(2) and of a (1), computes ba=1, and is about to write back to a, but there is a context switch
  - "a++" loads the value of a, which is 1, computes 2, and is about to write back to a, but there is a context switch
  - "a=b-a" resumes and writes back value 1 to a
  - "a++" resumes and overwrites value 1 by value 2
  - So in the end we have (2,2)

- In the end we found these possible outputs:
  - Clean #1: (0,2)
  - Clean #2: (0,2)
  - Clean #3: (2,2)
  - Unclean #1: (1,2)
  - Unclean #2: (2,2)
- So, in total, we have 3 possible outputs: (0,2),
   (2,2), and (1,2)

- When writing concurrent code in which multiple threads update the same (global) variables, one must make sure to avoid all "unclean" interleaving
- For this purpose we create critical sections
- Each critical section (which can have multiple zones in the code!) is defined by a lock, with a lock(A)/unlock(A) pair of statments

 Say our code has the following statements and that many threads will execute these statements

```
x++;
...
y++;
x++;
...
z=2;
```

 Say our code has the following statements and that many threads will execute these statements

```
....

X++;

...

y++;

X++;

...

z=2;

...
```

A brute force approach is to use one lock
...

```
lock(A);
X++;
unlock(A);
lock(A);
y++;
X++;
unlock(A);
lock(A);
z=2;
unlock(A);
```

A brute force approach is to use

Say our code has the following statements and that many threads will execute these statements

```
one lock
                                        lock(A);
                NOT good because a thread cannot do x++
                                        X++
X++;
                 while another one is doing y++.
                 Loss of concurrency, and thus performance.
y++;
                  If you want everything to be sequential,
X++;
                  Why use threads in the first place
                                         ятпоск(A);
z=2;
                                        lock(A);
                                        z=2;
                                        unlock(A);
```

 Say our code has the following statements and that many threads will execute these statements

```
x++;
...
y++;
x++;
...
z=2;
```

- First Question to answer, how many locks?
- Should there be a lock for protecting updates of x?
- Should there be a lock for protecting updates of y?
- Should there be a lock for protecting updates of z?

Say our code has the following statements and that many threads will execute these statements
 ....
 x++;
 ...
 y++;
 x++;
 ...
 z=2;
 ...

```
Solution:
lock(A);
X++;
unlock(A);
lock(B);
y++;
unlock(B);
lock(A);
X++;
unlock(A);
lock(A);
z=2;
unlock(A);
```

### Deadlocks

- We haven't seen much on this topic
- Make sure you understand the resource allocation graphs
  - If the "gray boxes" have more than "one dot": if there is a cycle, there may be a deadlock
  - If the "gray boxes" have only "one dot": if there is a cycle, there is a deadlock