Sequential Learning Algorithms Samplers

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- $X_t \in \{0,1\}$ is a binary sequence for $t = 1, 2, \ldots$
- Let $\{X_t\}$ be i.i.d. Bernouilli(p) sequence.
- Let \hat{X}_t be the prediction from the algorithm.
- Consider a randomized prediction, i.e., X_t is a Bernoulli(α) random variable.
- Claim 1: If p were known, the 'best' algorithm, the one that minimizes the expected number of mistakes upto to time T, would be, for

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- i.e., a fixed prediction of 1 if $p \ge 0.5$ and 0 otherwise.
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■ Claim 2: More generally, let \hat{X}_t be the prediction using the data $X_1, \dots X_{t-1}$, from any algorithm. Then,

$$\Pr(\hat{X}_t \neq X_t \mid X_1, \dots, X_{t-1}) \ge \min\{(p, (1-p))\}$$

i.e., a fixed prediction conditioned on the data has the minimum expected error.

Proof

■ The last claim leads us to suggest the following *historical majority* algorithm (HMA) for prediction of a Benroulli sequence when *p* is unknown.

$$\hat{X}_t = \text{majority from } \{X_1, X_2, \dots, X_{t-1}\}$$

 \blacksquare The expected loss for this algorithm upto time T is

$$\sum_{t=1}^T \mathsf{Pr}ig(\hat{X}_t
eq X_tig) - \min\{p, (1-p)\}$$

- The best error that any algorithm could have had is with the knowledge of p. It is the difference that is the loss.
- The number of mistakes that the HM algorithm would make up to time T is the number of times that X_t was not equal to the majority of $\{X_1, X_2, \dots, X_{t-1}\}$ for $t = 1, \dots, T$.



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Worksheet

- Without loss of generality, assume p > 0.5.

$$L_t := \Pr(\hat{X}_t \neq 1)$$

$$L_t = \Pr\left(\sum_{k=1}^{t-1} X_k \le \frac{t-1}{2}\right)$$

- Without loss of generality, assume $p \ge 0.5$.
- Let L_t be the probability that the prediction for step t is different from the "full knowledge case", i.e.,

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■ Clearly, L_t is the loss at step t.

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$$L_{t} = \Pr\left(\sum_{k=1}^{t-1} X_{k} \le \frac{t-1}{2}\right)$$

$$= \Pr\left(\sum_{k=1}^{t-1} (X_{k}) - (t-1)p \le \frac{t-1}{2} - (t-1)p\right)$$

$$\leq \exp \frac{-2\left(\frac{t-1}{2}-(t-1)p\right)^2}{t-1}$$
 Applying Hoeffding Inequality

$$= \exp{-2(t-1)(0.5-p)^2}$$

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$$\le \exp\frac{-2(\frac{t-1}{2} - (t-1)p)^{2}}{t-1} \quad \text{Applying Hoeffding Inequality}$$

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$$\begin{split} L_t &= \Pr \left(\sum_{k=1}^{t-1} X_k \leq \frac{t-1}{2} \right) \\ &= \Pr \left(\sum_{k=1}^{t-1} (X_k) - (t-1)p \leq \frac{t-1}{2} - (t-1)p \right) \\ &\leq \exp \frac{-2\left(\frac{t-1}{2} - (t-1)p\right)^2}{t-1} \quad \text{Applying Hoeffding Inequality} \\ &= \exp -2(t-1)(0.5-p)^2 \end{split}$$

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 \blacksquare The cumulative regret in T steps is

$$\sum_{t=1}^{T} L_{t} \leq \sum_{t=1}^{T-1} \exp{-2(t-1)(0.5-p)^{2}}$$

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- Consider the following shooting game: You are shooting into a unit square and you have to "predict" the fixed point, say \hat{X}_t for the *t*-th attempt.
- Let X_t be the actual location of the hit.
- For every \hat{X}_t there is a cost $||\hat{X}_t X_t||^2$.
- First consider the offline problem, or the one shot problem: the sequence X_t for t = 1, ..., T, is available to you and you have to choose the *best* fixed point

$$X_T^* := \arg\min_{x \in [0,1]^2} \sum_{t=1}^T ||X_t - x||^2$$

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 $\hat{X}_t = X_T^*$ minimizes the total cost.



- Now convert this to an online question: Predict the point after every shot such that the cumulative error is minimized.
- Specifically

$$\min \sum_{t=1} \|\hat{X}_t - X_t\|^2$$

with X_t to be determined before the t-th shot and can be done using the information X_1, \ldots, X_{t-1} .

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■ How does this compare with the "offline" algorithm above.

Specifically, what is the cumulative *loss* in the online algorithm compared to the full information offline case, i.e.,



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$$\sum_{t=1}^{T} \|\hat{X}_{t} - X_{t}\|^{2} - \sum_{t=1}^{T} \|\hat{X}_{t} - X_{t}\|^{2} = \sum_{t=1}^{T} \|X_{t-1}^{*} - X_{t}\|^{2} + \sum_{t=1}^{T} \|X_{t-1}^{*} - X_{t}\|^{2}$$

Lemma 1

$$\sum_{t=1}^{T} \|X_t^* - X_t\|^2 \le \sum_{t=1}^{T} \|X_T^* - X_t\|^2$$

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Proof

- Proof is by induction;
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$$\sum_{t=1}^{T} \|X_{t}^{*} - X_{t}\|^{2} \leq \sum_{t=1}^{T} \|X_{T}^{*} - X_{t}\|^{2}$$

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We now prove the preceding inequality

$$\sum_{t=1}^{T-1} \|X_t^* - X_t\|^2 \leq \sum_{t=1}^{T-1} \|X_{T-1}^* - X_t\|^2 \quad \text{from induction hypothesis}$$

$$\leq \sum_{t=1}^{T-1} \|X_T^* - X_t\|^2 \quad \text{because } X_T^* \text{ is not the minimizer} \quad \mathbb{R}$$

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$$L_T \leq \sum_{t=1}^{T} \|X_{t-1}^* - X_t^*\| \cdot \|X_{t-1}^* + X_t^* - 2X_t\|$$

using Cauchy-Schwartz inequality $|\langle x, y \rangle| \le ||x|| \cdot ||y||$.

$$\leq \sum_{t=1}^{T} \|X_{t-1}^* - X_t^*\| \cdot (\|X_{t-1}^*\| + \|X_t^*\| + \|2X_t\|)$$
 using triangle inequality

$$\leq \sum_{t=1}^{T} 4 \|X_{t-1}^* - X_t^*\|$$
 using the upper bound for $\|\cdot\|$ in 2nd term

Use the formula for x_i^* , triangle inequality and upper bound for $\|\cdot\|$

$$||X_{t-1}^* - X_t^*|| = ||X_{t-1}^* - \frac{(t-1)X_{t-1}^* + X_t}{t}|| = \frac{1}{t}||X_{t-1}^* - X_t||$$
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$$R_T \le 4 \sum_{t=1}^{T} \frac{8}{T} \le 8(1 + \ln(T))$$

Theorem For any sequence $X_1, ... X_T$, in the unit circle the FTL algorithm has a loss of at most $(1 + 8 \ln(T))$.

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