

Q1:

(a)

Source	Port	Destination	Port	Protocol	Action
*	*	64.37.180.23	80	TCP	Allow
*	*	64.37.180.23	443	TCP	Allow
*	*	64.37.160.63	22	TCP	Allow
*	*	*	*	*	Deny

(b)

Source	Port	Destination	Port	Protocol	Action	Notes
10.20.2.63	*	*	*	*	Deny	Blocking TCP/UDP Int_Web
10.100.X.X	*	*	53	TCP	Allow	Allow TCP internal Network
10.100.X.X	*	*	8080	TCP	Allow	Allow HTTP
10.100.X.X	*	*	8443	TCP	Allow	Allow HTTPS
10.20.2.22	*	*	53	TCP	Allow	Allow Remote SSH

(c)

www.companyxyz.com should resolve to the external IP for the Ext_Web which is 64.37.180.23. This IP is external so external networks can communicate with this IP address. If we used the internal IP 10.20.2.63 only those in the internal network can connect to this IP address.

Remote.companyxyz.com should resolve to the external IP of remote access which is 64.37.160.63. This external IP can be connected to by external networks and internal networks. If the internal IP was used 10.20.2.22 only those in the internal network can connect to this IP address.

(d)

The company must port forward the internal IP and the ports, The IP 10.10.1.20 and the ports 8080 and 8443 must be port forwarded (using TCP as the service). Port forwarding allows external networks to connect to a specific service in an internal network.

(e)

Source	Port	Destination	Port	Protocol	Action
10.20.2.22	*	10.100.X.X	22	TCP	Allow
10.20.2.22	*	10.100.X.X	*	*	Deny

(f)

Yes, as EXT_FW allows access to the EXT_WEB server, if appropriate SSH is setup on the EXT_WEB it can be directly accessed by SSH assuming port forwarding has also been done.

(g) Int_Web is only allowed to respond to requests from the internal network. The user cannot directly connect to this network from home. The only way to connect to this server is to connect to a machine on the internal network by ssh into RemoteAccess (64.37.160.63:22) and to then ssh into Int_Web. Thus, a direct connection from home directly cannot work.

(h)

Source	Port	Destination	Port	Protocol	Action
10.100.X.X	*	10.20.2.22	22	TCP	Allow
10.100.X.X	*	10.20.2.63	8080	TCP	Allow
10.100.X.X	*	10.20.2.63	8443	TCP	Allow
10.100.X.X	*	*	53	TCP	Allow
*	*	*	*	*	Deny

(i)

The Guest_Wifi and Home_Eth networks are two independent subnetworks with their own IP ranges. These two networks cannot communicate with each other.

A segregated Wifi network will allow you to prevent guests to access your private Samba Share folders and access to the printer. This segregated network ensures that those that want to see your private family documents must be physically connected to your server.

Give the laptop a static IP address and create a firewall rules that allows connection from the allocated IP address to the subnetwork of Home_WAP.

The external IP (80.245.7.8) will show on the webservers log. This is because when the user is sending a request to the webserver the home router will send the request to the webserver using the external IP address and provide it a port number. This process is done by using a method known as Network Address Translation (NAT), this maps a single IP address space into another in this case we are mapping the internal IP into the external IP. NAT modifies the IP header whilst it is in transit from the router.

Question 2:

(a) $L/(R1) + L/(R2)$

(b) (i) 2 Users

(ii) As there is a 2Mbps link and each user is continuously transmitting 1Mbps there can be up to two maximum simultaneous transmissions without packet delay as the total 2Mbps link is being used. If more than two users are using this link for instance three users, a link of 3Mbps is required which is more than the available bandwidth therefore, there will be a packet queue formed at the router as packets will be queuing faster than the packets that are being transported across the network.

(iii) 0.2

(iv) $0.2^3 = 0.008$

- (c) (i) 500Kbps
(ii) $(4*8*10^6)/(500*1000) = 64$ seconds
(iii) 100Kbps, 320 seconds
- (d) The first packet will arrive at N(L/R) time. The second packet will arrive at N(L/R) + L/R time. Where P is the last packet to reach the destination the overall equation must be $N(L/R) + (P-1)(L/R) = (N + P - 1)(L/R)$.
- (e) The IP address and port number of a process of the other host.
- (f) TCP has a reliable data transfer ensuring a connection using three way TCP.
Security is not provided by either this must be done in the application layer
A guarantee of an arrival time not provided by either
Guarantee that a certain throughput is maintained is not provided by either.
- (g) A TCP server needs two sockets as this protocol work by request and response. TCP sends data to the client and requests a response from the client, since data is being transported between both machines it requires two sockets. A TCP server only needs two sockets even with n simultaneous connections because server hosts can support many simultaneous TCP connections. TCP requires the IP and port of both the source and destination, when the packet is received at the destination machine, they can be demultiplexed into different sockets.
- (h) (i)

One RTT for handshake + (One RTT for handshake + One RTT for page request + page data) + 5(One RTT for handshake + One RTT for handshake + One image request)

$$(0.1 + 0.1 + \frac{1*1000*8}{1*10^6}) + 5*(0.1 + 0.1 + \frac{50*1000*8}{1*10^6}) = 3.208s$$

(ii)

Getting the first webpage is required to get other data.

Getting first page: One RTT for handshake + (One RTT for handshake + One RTT for page request + page data)

$$= (0.1 + 0.1 + \frac{1*1000*8}{1*10^6}) = 0.208$$

Now the Images can be done in parallel

(One RTT for handshake + One RTT for handshake + One image request)

$$= *(0.1 + 0.1 + \frac{50*1000*8}{1*10^6}) = 0.6$$

Total time taken = $0.208 + 0.6 = 0.808s$

(iii)

One RTT for handshake + (one RTT + document) + 5 (one RTT + image)

$$= (0.1 + 0.1 + \frac{1*1000*8}{1*10^6}) + 5*(0.1 + \frac{50*1000*8}{1*10^6}) = 2.708s$$

- (i) **(i)** Source: 33000, Destination: 80
(ii) Source: 80, Destination 33000
(iii) No, as HTTP is a TCP protocol so UDP cannot be used to contact the server. HTTP relies on the three-way TCP handshake.

(j) 01011101 11110010
00110010 00001101

UDP Checksum requires it to be the 1's complement, in this case there are two bits that are not correct. This segment is not correctly received. The receiver must request the packet again.

Question 3

- (a) (i) gaia.cs.umass.edu/cs453/index.html
(ii) HTTP 1.1.
(iii) Persistent because of Connection:keep-alive
(iv) Mozilla/5.0 Browser The browser type is needed as different browsers handle HTTP messages differently.
- (b) (i) Server was successfully able to find the document as a 200 OK response was returned.
Tue, 07 Mar 2008 12:39:45 GMT
(ii) Sat, 10 Dec 2005 18:27:46 GMT
(iii) 3874 bytes
(iv) The first 5 bytes are "<!doc". Yes the server has agreed to keep the persistent connection as there is the "Connection: Keep-Alive" in the response.
- (c) (i) 71.192.34.104
(ii) Source: 192.168.1.100, Source Port: 4335
Destination: 63.233.169.104, Destination Port: 80
- (d) First TCP SYN segment sent used by the GET at 7.109267? The first SYN packet was sent at 7.075657000
Acknowledgment is received at 7.108986
- (e) (i) 6.069168s
Source: 71.192.34.104, Source Port: 4335, Destination IP: 64.233.169.104,
Destination Port: 80
The fields that are still the same are the source port, destination port, destination IP address.
(ii) The source IP address got changed. This is because the source IP in the home side was an internal IP address, this had to be translated to an external IP address using NAT so external networks can communicate with this computer. The TCP checksum is also modified as this checksum is formed with the data in the header, since the source IP changed the checksum must also change.

Question 4

The screenshot shows two terminal windows side-by-side. Both windows are titled 'Activities' and have tabs for 'Terminal'. The left window is titled 'nebulakvm02.eps.surrey.ac.uk:6999 (on heron010)' and the right window is titled 'nebulakvm05.eps.surrey.ac.uk:7001 (on heron010)'. The time at the top center is 14:45.

Terminal Content (Left Window):

```
# UDPPingClient.py
# We will need the following module to generate randomized lost packets
import random
from socket import *
import time
print("Client Starting")
# Create a UDP socket
clientSocket = socket(AF_INET, SOCK_DGRAM)
# Notice the use of SOCK_DGRAM for UDP packets
# Assign IP address and port number to socket
clientSocket.bind(('', 1024))
rec = ("10.0.0.1", 1024)
while True:
    # Generate random number in the range of 0 to 10
    rand = random.randint(0, 10)
    # Recieve message
    message = clientSocket.recvfrom(1024)
    # Print the message received
    print("Message Received: " + message[0].decode("utf-8"))
    if rand < 5:
        # Set timeout
        clientSocket.settimeout(1.0)
        # Bind the socket
        clientSocket.bind(('', 1024))
        # Send the message
        clientSocket.sendto("Ping 10", rec)
        # Try to receive
        try:
            # Keep asking
            message, address = clientSocket.recvfrom(1024)
            # Get the received message
            recieveTime = time.time()
            # Print RTT
            print("RTT = " + str(recieveTime - recieveTime))
            print("Message Received: " + message.decode("utf-8"))
        except timeout:
            print("Packet is most likely lost if there is a timeout. Timeout has been set above of 1 sec")
            except timeout:
```

Terminal Content (Right Window):

```
# UDPPingServer.py
# We will need the following module to generate randomized lost packets
import random
from socket import *
print("Server Starting")
# Create a UDP socket
serverSocket = socket(AF_INET, SOCK_DGRAM)
# Notice the use of SOCK_DGRAM for UDP packets
# Assign IP address and port number to socket
serverSocket.bind(('', 1024))
rec = ("10.0.0.1", 1024)
while True:
    # Generate random number in the range of 0 to 10
    rand = random.randint(0, 10)
    # Recieve message
    message = serverSocket.recvfrom(1024)
    # Print the message received
    print("Message Received: COMPUTER NETWORKING!")
    if rand < 5:
        # Set timeout
        serverSocket.settimeout(1.0)
        # Bind the socket
        serverSocket.bind(('', 1024))
        # Send the message
        serverSocket.sendto("Ping 10", rec)
        # Try to receive
        try:
            # Keep asking
            message, address = serverSocket.recvfrom(1024)
            # Get the received message
            recieveTime = time.time()
            # Print RTT
            print("RTT = " + str(recieveTime - recieveTime))
            print("Message Received: COMPUTER NETWORKING!")
        except timeout:
            print("Packet is most likely lost if there is a timeout. Timeout has been set above of 1 sec")
            except timeout:
```

The screenshot shows two terminal windows side-by-side. The left window is titled 'nebulakvm02.eps.surrey.ac.uk:6999 (on heron010)' and the right window is titled 'nebulakvm05.eps.surrey.ac.uk:7001 (on heron010)'. Both windows are running on an Ubuntu system.

Terminal 1 (Client Side):

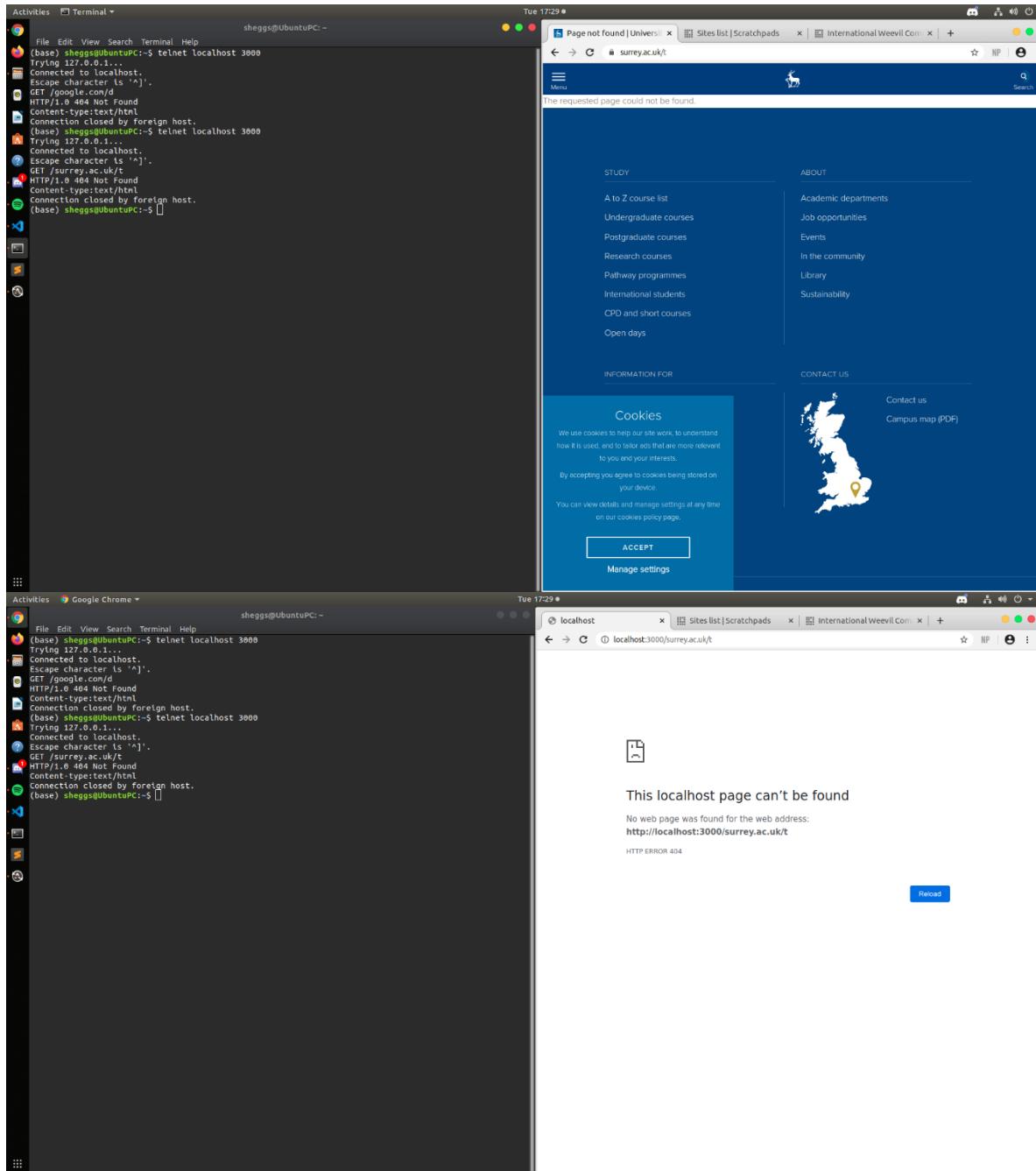
```
# UDPPingClient.py
# We will need the following module to generate randomized lost packets
import random
from socket import *
import time
print("Client Starting")
# Create a UDP socket
# Notice the use of SOCK_DGRAM for UDP packets
clientSocket = socket(AF_INET, SOCK_DGRAM)
# Assign IP address and port number to socket
clientSocket.bind(("", 1024))
# Getting the current time
messageSendTime = time.time()
# Creating client
clientSocket.settimeout(1)
# Binding the socket
clientSocket.bind(("", 1024))
# Sending the message
clientSocket.sendto(message, (destination, port))
# Try to receive
try:
    # Keep asking
    message, addr = clientSocket.recvfrom(1024)
    # Get the received time
    recieveTime = time.time()
    # Print RTT
    print("DEBUG: " + str(recieveTime - messageSendTime))
    print("Message Recieved: " + message.decode("utf-8") + " RTT = " + str(recieveTime - messageSendTime))
    # Packet is most likely lost if there is a timeout. Timeout has been set above of 1 sec
except timeout:
    print("No answer from server")
```

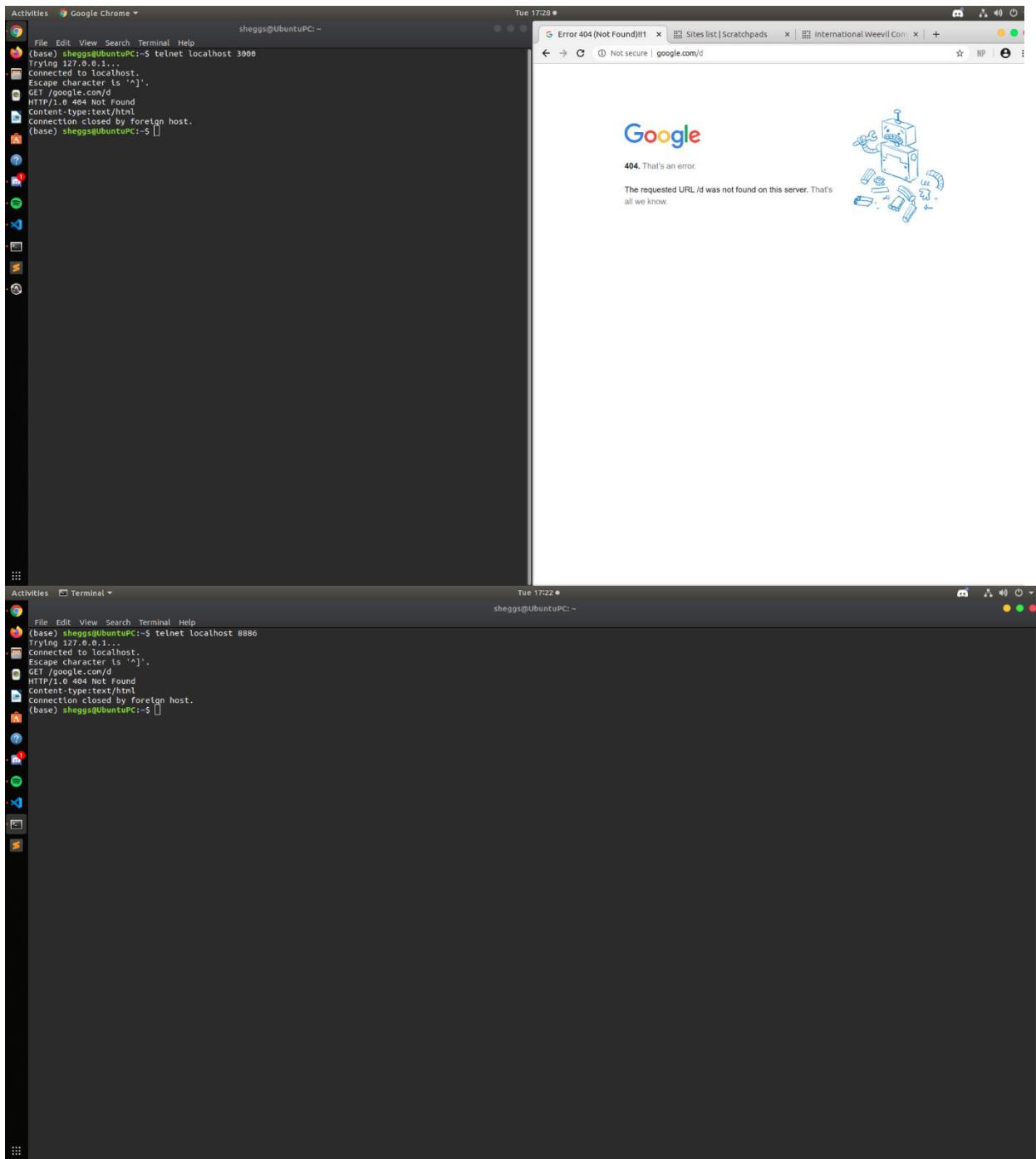
Terminal 2 (Server Side):

```
# UDPPingServer.py
# We will need the following module to generate randomized lost packets
import random
from socket import *
print("Server Starting")
# Create a UDP socket
# Notice the use of SOCK_DGRAM for UDP packets
serverSocket = socket(AF_INET, SOCK_DGRAM)
# Assign IP address and port number to socket
serverSocket.bind(("", 12000))
rec = ("10.0.0.1", 1024)
while True:
    # Generate random number in the range of 0 to 10
    rand = random.randint(0, 10)
    # Receive
    message, address = serverSocket.recvfrom(1024)
    # If
    if rand == 0:
        print("Message Recieved. COMPUTER NETWORKING!")
    else:
        print("Message Recieved. COMPUTER NETWORKING!")
    # Other
    else:
        print("Message Recieved. COMPUTER NETWORKING!")
```

Proxy Server

404 Images



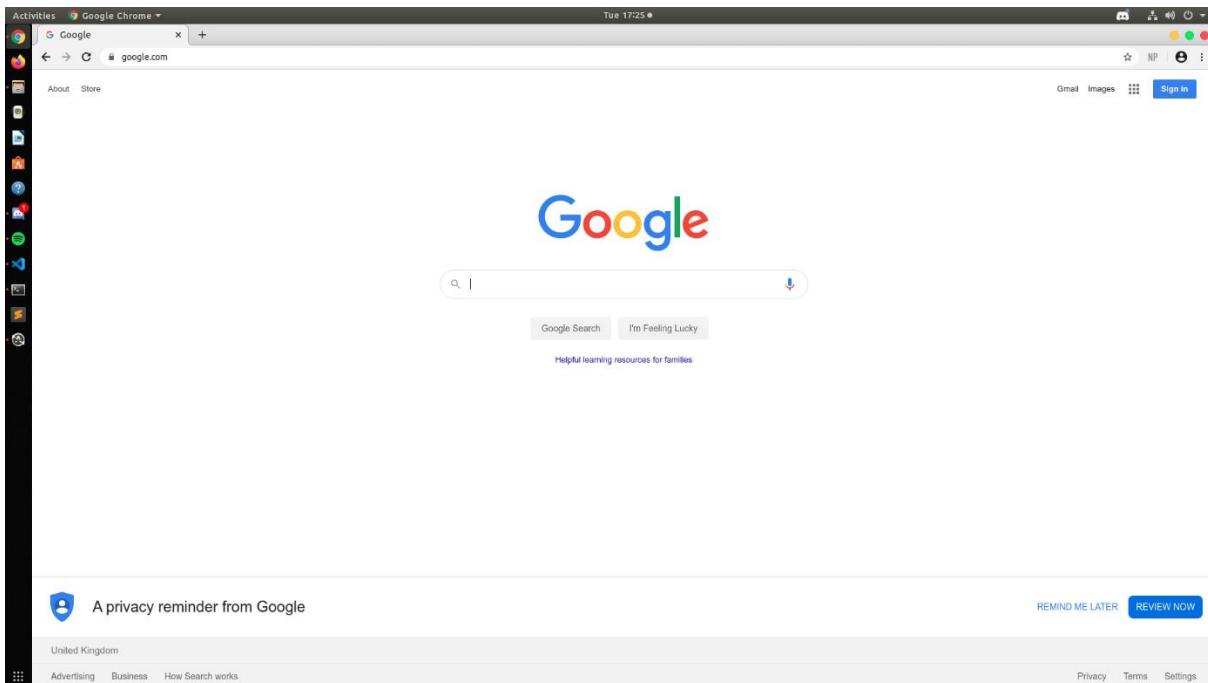


301 Moved

When it is cached (200 request returned but 301 message is shown on page)

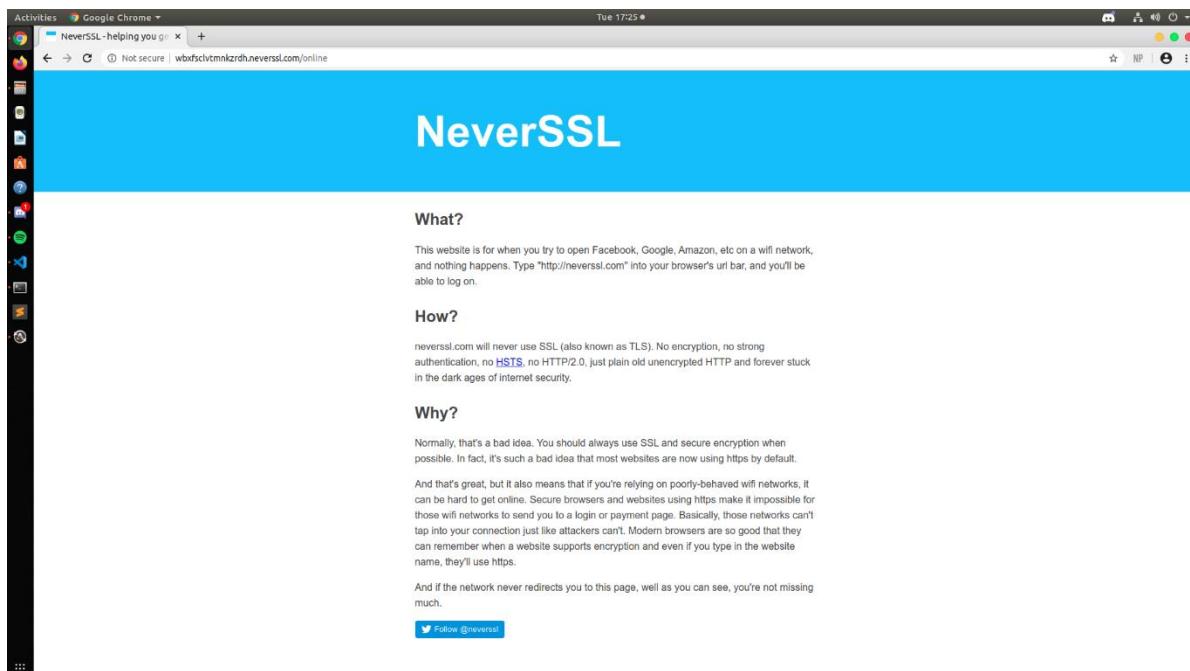
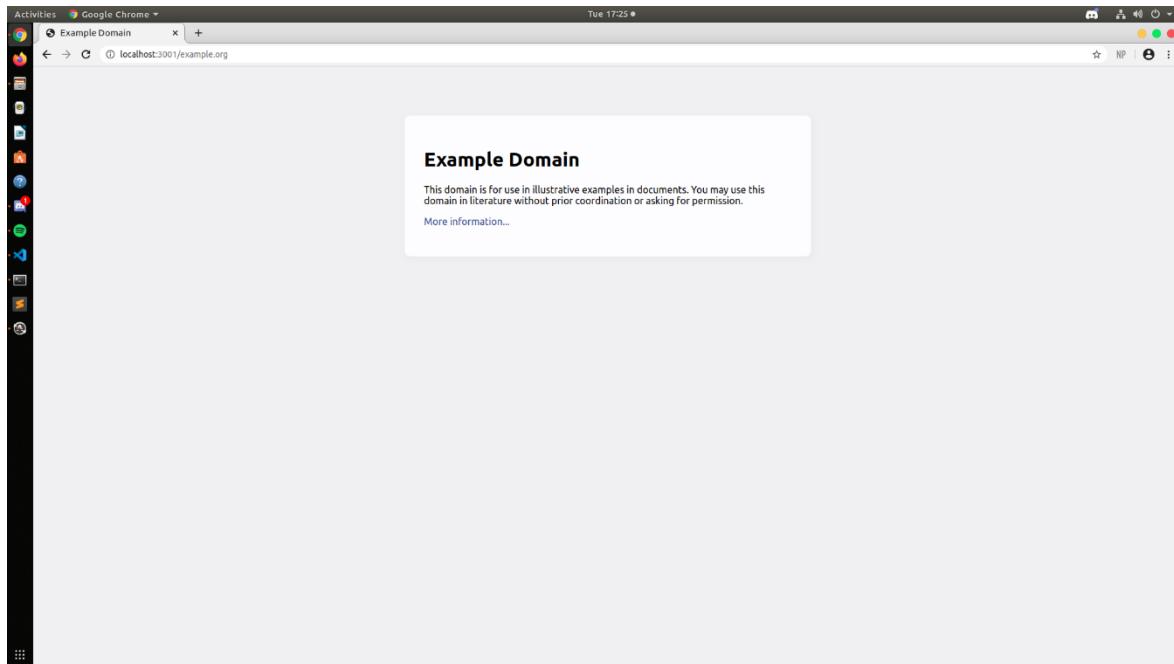
```
Activities Visual Studio Code Tue 17:21 ● sheggs@UbuntuPC:~  
File Edit View Search Terminal Help  
sheggs@UbuntuPC:~$ telnet localhost:8886  
telnet: could not resolve localhost:8886/telnet: Name or service not known  
[base] sheggs@UbuntuPC:~$ telnet localhost 8886  
Trying 127.0.0.1...  
Connected to localhost.  
Escape character is '^]'.  
GET /google.com  
HTTP/1.0 200 OK  
Content-Type:text/html  
Content-Length:219  
Location: http://www.google.com/  
Content-Type: text/html; charset=UTF-8  
Date: Tue, 28 Apr 2026 00:28:39 GMT  
Expires: Thu, 29 May 2026 00:28:39 GMT  
Cache-Control: public, max-age=2592000  
Content-Type: qwe  
Content-Length: 219  
X-XSS-Protection: 0  
X-Frame-Options: SAMEORIGIN  
<HTML><HEAD><meta http-equiv="content-type" content="text/html; charset=utf-8">  
<TITLE>301 Moved</TITLE><BODY>  
<H1>301 Moved</H1>  
<P>The document has moved  
<A HREF="http://www.google.com/">here</A>.  
</BODY></HTML>
```

Redirects the client to google.com (On first try, other attempts the cached result shows the error)



Real 200 responses

CSS hrefs have been fixed but for some websites it can't get the required data.



Activities Google Chrome ▾

Tue 17:26 •

localhost:3001/open-up.eu

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OPEN UP! Opening Up the Natural History Heritage for Europeana

Search form

Search Search... All Taxonomy

Main menu

- Home
- About Us
- News and Events
- Partners
- Logistics Information
- How to join
- Contact/Privacy

You are here

Home > Page not found

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The OpenUp! Network involves partner institutions from 13 European countries, contributing more than 8.5 million records of multimedia objects from the natural history domain to Europeana by now.

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europaea aggregator open-up.eu/en/ RTNER

Activities Terminal ▾

Tue 17:24 • sheggs@UbuntuPC: ~

```
</html>
Connection closed by foreign host.
(base) sheggs@UbuntuPC:~$ telnet localhost 3001
Trying 127.0.0.1...
Connected to localhost.
Escape character is '^'.
GET /neverssl.com
HTTP/1.0 200 OK
Content-Type: text/html
Content-Length: 1416
Connection: close
Last-Modified: Thu, 06 Jun 2019 04:19:28 GMT
Date: Tue, 26 May 2020 02:04:49 GMT
ETag: "92bd819fcf2865f0be341e0239d23bf"
Vary: Accept-Encoding
Server: Apache/2.4.34 (Ubuntu)
X-Amz-CF-Pop: LHR3-C2
X-Amz-CF-Id: g-lEYtjgXL_WOIDNKLARhep6U1998SEiwUlwZHu-IVAbipCBGbb9A==
Age: 51596

<html>
  <head>
    <title>NeverSSL - Connecting ...</title>
  </head>
  <body>
    font-family: Montserrat, helvetica, arial, sans-serif;
    font-size: 1em;
    color: #444444;
    margin: 0;
  </body>
  <h2>
    font-weight: 700;
    font-size: 1.6em;
    margin-top: 30px;
  </h2>
  <p>
    line-height: 1.6em;
  </p>
</html>
```

Activities Terminal ▾

Tue 17:24 • sheggs@UbuntuPC: ~

```
(base) sheggs@UbuntuPC:~$ telnet localhost 3001
Trying 127.0.0.1...
Connected to localhost.
Escape character is '^'.
GET /example.org
HTTP/1.0 200 OK
Age: 475139
Content-Size: max-age=604800
Content-Type: text/html; charset=UTF-8
Date: Tue, 26 May 2020 16:24:11 GMT
ETag: "3147526947+lendl"
Expires: Tue, 02 Jun 2020 16:24:11 GMT
Last-Modified: Tue, 26 May 2020 07:18:26 GMT
Server: ECS (nby/D1D4)
Vary: Accept-Encoding
X-Cache: HIT
Content-Length: 1256
Connection: close

<!doctype html>
<html>
<head>
<title>Example Domain</title>
<meta charset="utf-8" />
<meta http-equiv="Content-Type" content="text/html; charset=utf-8" />
<meta name="viewport" content="width=device-width, initial-scale=1" />
<style type="text/css">
body {
  background-color: #f0f0f2;
  margin: 0;
  padding: 0;
  font-family: -apple-system, system-ui, BlinkMacSystemFont, "Segoe UI", "Open Sans", "Helvetica Neue", Arial, sans-serif;
}
div {
  width: 600px;
  margin: 0 auto;
  padding: 2em;
  background-color: #fdfdff;
  border-radius: 0.5em;
  box-shadow: 2px 3px 7px 2px rgba(0,0,0,0.02);
}
a:link, a:visited {
  color: #38488f;
  text-decoration: none;
}
@media (max-width: 700px) {
  div {
    margin: 0 auto;
    width: auto;
  }
}
</style>
</head>
```

Question 5

- (a) (i) MTA = Mail Transfer Agent
(ii)
(A) “Too many recipients” has an error code of 452 and not 552 so this error code is incorrect.
(B) Error 552 means too much mail data.
(iii)
(A) 220 message means service ready in the RFC documentation. This message is an initial message, a receiver will send this reply when the connection is completed.
(B) 220 message timeouts are 5 minutes. Length of all objects is not compulsory, the RFC states “implementation techniques that impose no limits on the lengths of these objects should be used”. However, maximum total lengths for local-part, domain, path, command line, reply line, message buffer and recipients’ buffer has been stated and appropriate error respond codes have been allocated 500,501,452 and 522.
(C) A TCP connection must be established for a 220 message to be sent. This message can be delayed until the system allows for more mail to be processed. A TCP connection is needed to be started for any SMTP messages to be sent.
(iv) SMTP messages must be queued to be sent from the SMTP client to the SMTP server. SMTP clients will periodically attempt to send messages to a SMTP server, mail must be periodically retried by the sender as it cannot respond with an error message under any circumstances. Failure of a message being sent can only be determined after a timeout of a few minutes.
- (b) SRV maps hostname and port numbers of services with specified services. This allows the use of subdomains for different services. SRV record allows admins to use several servers for a single domain. Allowing services to move between hosts with little problems.

Question 6

- (a) (i)
 $n = (11)(19) = 209$
 $z = (11-1)(19-1) = 180$
(ii)
 $e = 17$ is an acceptable choice as e is less than the value of n and has no common factor with z .
(iii)
 $de = 1 \pmod{z}$
 $d(17) = 1 \pmod{180}$
 $d = 53$
(iv)
Encryption
 $m = 7$
 $c = m^e \pmod{n}$
 $c = 7^{17} \pmod{209}$
 $c = 182$
Decryption
 $m = c^d \pmod{n}$
 $m = 182^{53} \pmod{209}$
 $m = 7$
- (b) Alice: Secret Key S_A , Public Key $(T_A) = (g^S_A) \pmod{p}$, Shared Key $T_B^S_A \pmod{p}$

Bob: Secret Key S_B , Public Key $(T_B) = (g^S_B) \text{ mod } p$, Shared Key $T_A^S_B \text{ mod } p$

(i)

$$S = (T_B^S_A \text{ mod } p) = (g^S_B)^S_A \text{ mod } p = g^{S_A} \text{ mod } p \wedge S_B = T_A^S_B \text{ mod } p = S'$$

(ii)

$$p = 11, g = 2, S_A = 5, S_B = 12$$

$$(T_A) = (g^S_A) \text{ mod } p = 2^5 \text{ mod } 11 = 10$$

$$(T_B) = (g^S_B) \text{ mod } p = 2^{12} \text{ mod } 11 = 4$$

(iii)

$$S = T_B^S_A \text{ mod } p \quad S' = T_A^S_B \text{ mod } p$$

$$S = 4^5 \text{ mod } 11 = 1$$

$$S' = 10^{12} \text{ mod } 11 = 1$$

(iv) Eve will first communicate with Alice and receive the public key T_A and then communicate with Bob and get a public key of T_B . Both Alice and Bob will make a shared key with Eve and then Eve will be the “man-in-the-middle” having access the plain data being sent across the networking and having the ability to send any message to either parties.

- (c) Client = 128.238.38.162 Server = 216.75.194.220

(i) Client

(ii) IP = 216.75.194.220, Port = 443

(iii) 283

(iv) Three

(v) Master Secret

(vi) first is bc, last is 29

- (d) No longer will work as the MAC that Bob is appending is unique and not shared between Bob and Alice. Alice won't be able to decrypt the message.

Question 7

- (a) A host associated with a base station is operating in infrastructure mode. If it has no infrastructure to connect with it is called an ad hoc network. Ad hoc networks have no infrastructure thus the hosts must provide routing, DNS, address assignment and other features. Whilst infrastructure mode must connect to a base station and the host is connected via this base station, so all traditional services are already provided.
- (b) A user that is mobile has to switch from one base station to another if they exceed the range of the currently connected base station in a process known as handoff. A user that walks around her house with a laptop is not considered mobile as it doesn't switch between base stations.
- (c) Beacon frames send data regarding an AP which consists of its SSID and MAC address. This will be used to aid the wireless station to select an AP to connect to. Beacon frames are transmitted across one of the 11 channels.
- (d) SSID, MAC Address
- (e) Ethernet has no packet loss whilst 802.11 is more likely to lose packets and higher risk of collisions. Ethernet has a physical link which makes it much more reliable.
- (f) Active scanning is by constantly broadcasting a probe frame that will be received by all access points in the range of the wireless host. Passive scanning is scanning channels and waiting for beacon frames.
- (g) (i)
No, it will not break down. Both can work parallel over the same channel. If both APs transmitted at the same time, there will be collisions. Maximum transmission rate is 11Mbps.

- (ii) Now they are across different channels there are no collisions and the maximum transmission rate is 22Mbps. 11Mbps per channel.
- (iii) Host will send a DHCP discovery message via the AP in order to obtain an IP address. Once this IP address is allocated this is now connected to the subnet.

References

[1] <https://tools.ietf.org/html/rfc5321#section-4.5.3.1>