

3.3) Draw a clear resource allocation graph based on the information below.

(The sets P - processes, R - resources, E- edges as follows)

$$P = \{P_1, P_2, P_3, P_4\}$$

$$R = \{R_1, R_2, R_3, R_4, R_5\}$$

$$E = \{(P_1, R_5), (P_1, R_2), (R_3, P_1), (R_5, P_1), (P_2, R_3), (R_2, P_2), (P_3, R_3), (P_3, R_1), (R_2, P_3), (R_4, P_3), (P_4, R_3), (P_4, R_1), (R_4, P_4), (P_4, R_5)\}$$

The number of resource instances are

R1 has 1 instance

R2 has 2 instances

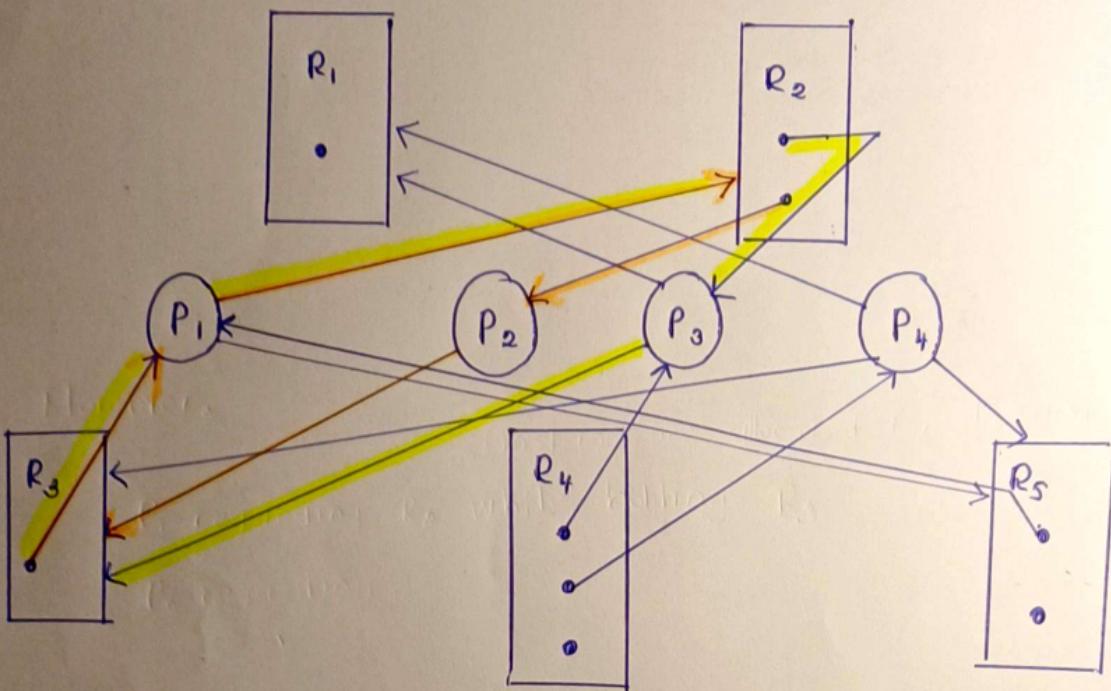
R3 has 1 instance

R4 has 3 instances

R5 has 2 instances

3.4) Using the resource graph that you have drawn in section 3.3, identify deadlock sequences (if any). Justify your answer.

Resource Allocation Graph:



This is the RAG we can draw based on the details provided in the question

To answer question 3.4, we need to identify whether there is a deadlock sequence or not. If a deadlock exists, we must justify the answer.

For that, we can use resource allocation graph reduction.

draw the resource allocation graph reduction diagram.
Therefore, let's try to ~~reduce the graph~~,

* Let's start with process that hold resources and don't request any resources.

Unfortunately, there is no any processes, as I explain below:

- P_1 Holds R_3 and R_5

BUT

P_1 Requesting R_2 and R_5 [Yes, R_5 is already allocated to P_1 , but still P_1 requesting R_2 & instances are allocated to $P_2 + P_3$]

- P_2 Holds R_2 BUT Requesting R_3 [P_2 can not be reduced yet.]
Since R_3 only has one instance & that is already allocated to P_1

- P_3 Holds R_2 and R_4 BUT Requesting R_1 and R_3 .

[It also can't reduced since R_3 only has one instance + that is already allocated to P_1]

- P_4 Holds R_4 . BUT Requesting R_1, R_3 and R_5 .

[It also can't reduced since R_3 only has one instance + that is already allocated to P_1]

Therefore, \rightarrow (instances are allocated for P_2 and P_3)

P_1 requesting R_2 while holding R_3

P_2 requesting R_3 while holding R_2

\Rightarrow (One instance allocated for P_1)

∴ there is a circular wait: $P_1 \rightarrow R_2 \rightarrow P_2 \rightarrow R_3 \rightarrow P_1$

P_1 requesting R_2 , but R_2 allocated for P_2 , but P_2 requesting R_3 , but that is allocated for P_1 .

of

P_1 not releasing R_3 , since to complete the entire task, P_1 need R_2 .

However, R_2 allocated for P_2 , but P_2 can't complete its execution without R_3 . Therefore P_1 & P_2 never releasing the resources.

Moreover,

- P_1 requesting R_2 , while holding R_3
 P_3 requesting R_3 , while holding R_2
- (Instances are allocated for P_2 and P_3)
→ (There is only one instance which is already allocated for P_1)

∴ there is a circular wait : $P_1 \rightarrow R_2 \rightarrow P_3 \rightarrow R_3 \rightarrow P_1$

P_1 requesting R_3 , but R_2 allocated for P_3 , but P_3 requesting R_3 , but that is allocated for P_1 .

P_1 not releasing R_3 , since to complete the entire task of P_1 , need R_2 . However, R_2 allocated for P_3 , but P_3 can't complete its execution without R_3 . Therefor $P_1 + P_3$ also never releasing the resources.

R_2 only has 2 instances
 R_3 only has 1 instance

- given for P_2
- given for P_3
- given for P_1

Therefore, these resources don't have free instances available for allocation when other processes are requesting them.

The system is in a deadlock state because there is a circular wait. Each process in the cycle is waiting for a resource held by the next process in the cycle. Therefore, no process can proceed.