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## Memory Allocators

#### Section 1

## Introduction

Applicable for

LINUX AND C

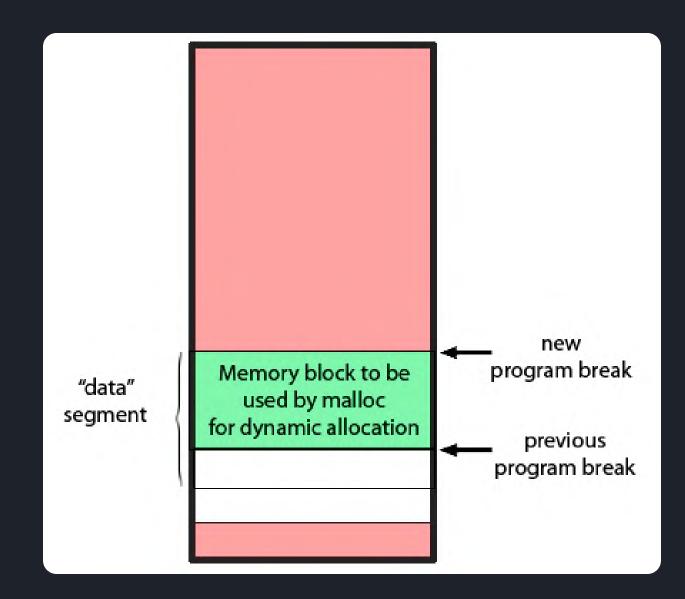
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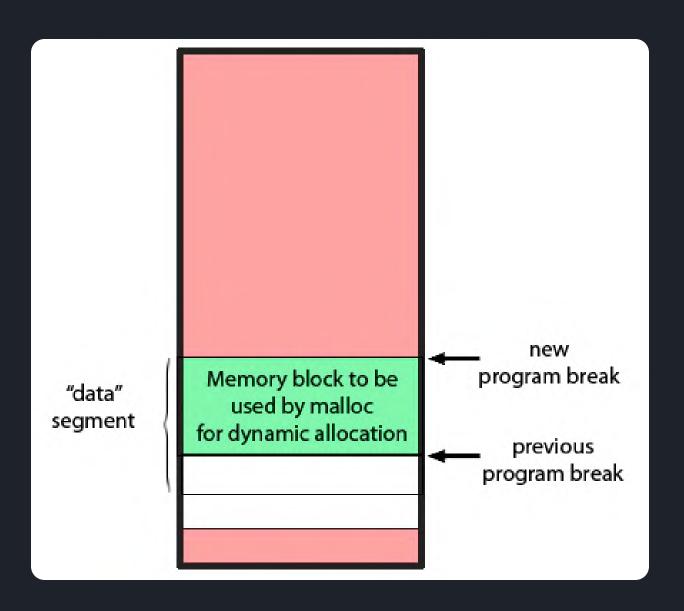
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```
void *malloc(size_t size)
{
    void *block;
    block = sbrk(size);
    if (block == (void*) -1)
        return NULL;
    return block;
}
```

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  do not give any way to use the fragmented memory and is thus wasted.
- **Copy on write (COW)** Unix based operating systems generally employ copy-on-write mechanism for virtual memory management, specially in cases like fork(). So it may be possible that the memory requested has been marked as *not present* and accessing that memory can cause page fault issues.

## Dynamic Storage Allocation: A Survey and Critical Review \* \*\*

Paul R. Wilson, Mark S. Johnstone, Michael Neely, and David Boles\*\*\*

Department of Computer Sciences
University of Texas at Austin
Austin, Texas, 78751, USA
(wilson|markj|neely@cs.utexas.edu)

Abstract. Dynamic memory allocation has been a fundamental part of most com-

#### 1 Introduction

Paul Wilson and his colleagues have written an excellent survey paper on allocation techniques that discusses some of the goals in more detail - "Dynamic Storage Allocation: A Survey and Critical Review" in *International Workshop on Memory Management*, September 1995.

**Conclusion:** They discuss, minimizing space by minimizing wastage (generally due to fragmentation) must be the primary goal in any allocator.

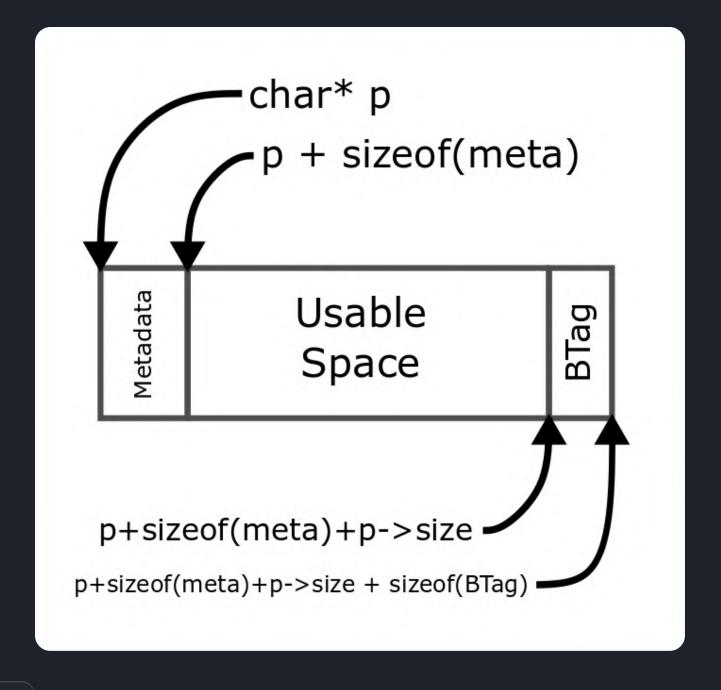
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# Implementations

A MEMORY ALLOCATOR BY **DOUG LEA** 

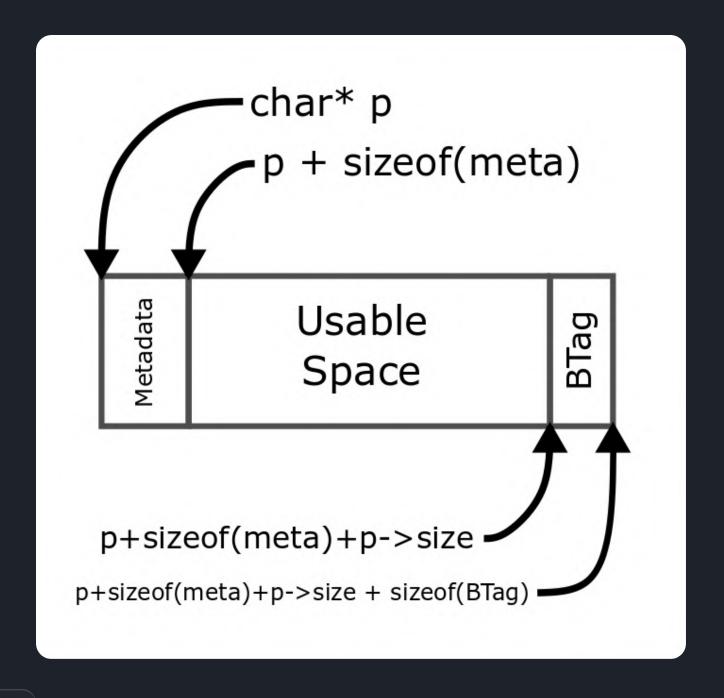
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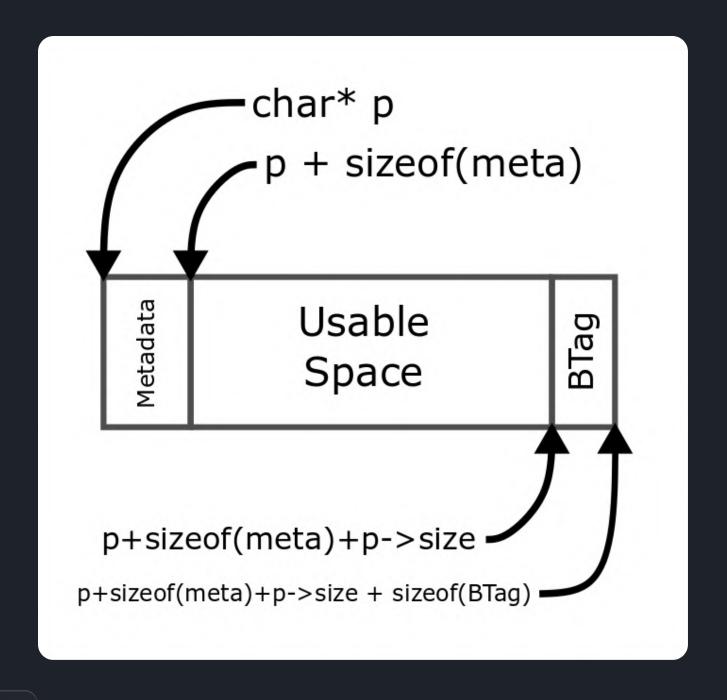
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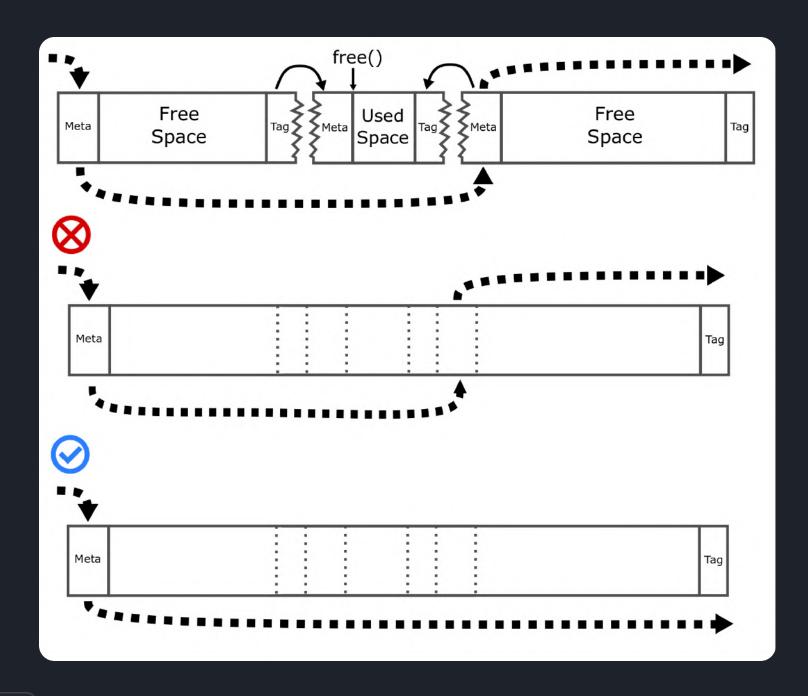


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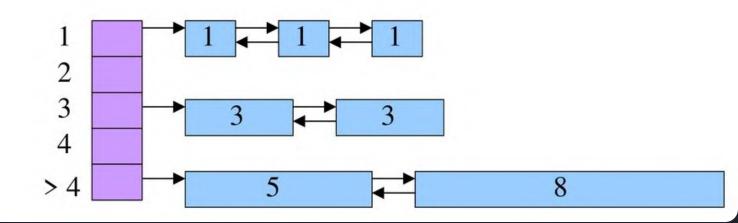
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#### This allows for two important capabilities:

- Two bordering unused chunks can be coalesced into one larger chunk. This minimizes the number of unusable small chunks.
- All chunks can be traversed starting from any known chunk in either a forward or backward direction.

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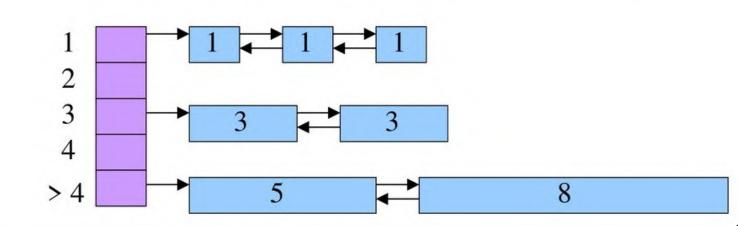
- Have a bin for each block size, up to a limit
  Advantages: no search for requests up to that size
  Disadvantages: many bins, each storing a pointer
- Except for a final bin for all larger free blocks
  - For allocating larger amounts of memory
  - For splitting to create smaller blocks, when needed



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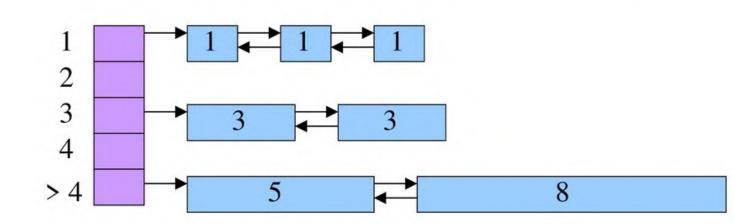


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**Best fit strategy** (searching if the required size block exists in the bin else breaking a large block into that size) is used to allocate memory for new allocation requests.

These two algorithms are used to allocate and free memory blocks in malloc().

SOURCE BLOG: http://gee.cs.oswego.edu/dl/html/malloc.html

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A big disadvantage: dlmalloc works only for single threaded programs. It fails for multi-threaded programs.

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THE GLIBC IMPLEMENTATION

ptmalloc - **pthreads malloc** is the current glibc implementation for malloc. It is a derived version of **dlmalloc** except that it has a **pthreads** wrapper which helps malloc support multithreaded programs for memory allocation.

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IMPLEMENTATION DETAILS : https://sourceware.org/glibc/wiki/MallocInternals
SOURCE CODE : https://github.com/hustfisher/ptmalloc

#### THANK YOU - PRESENTATION BY

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