

## 7.1 Background

We have discussed Abstract Lists with explicit linear orders

- Arrays, linked lists, strings

We saw three cases which restricted the operations:

- Stacks, queues, dequeues

Following this, we looked at search trees for storing implicit linear orders: Abstract Sorted Lists

- Run times were generally  $\Theta(\ln(n))$

We will now look at a restriction on an implicit linear ordering:

- Priority queues

## 7.1.1 Definition

With queues

- The order may be summarized by *first in, first out*

If each object is associated with a priority, we may wish to pop that object which has highest priority

With each pushed object, we will associate a nonnegative integer (0, 1, 2, ...) where:

- The value 0 has the *highest* priority, and
- The higher the number, the lower the priority

## 7.1.2 Operations

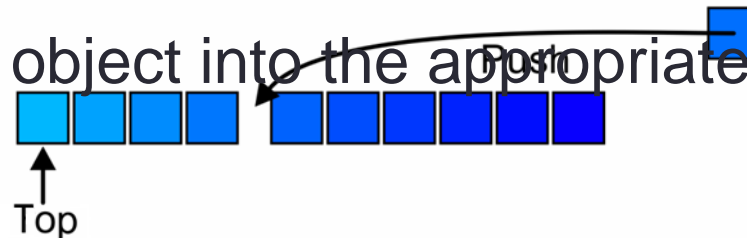
The top of a priority queue is the object with highest priority



Popping from a priority queue removes the current highest priority object:



Push places a new object into the appropriate place



## 7.1<sup>3</sup> Lexicographical Priority

Priority may also depend on multiple variables:

- Two values specify a priority:  $(a, b)$
- A pair  $(a, b)$  has higher priority than  $(c, d)$  if:
  - $a < c$ , or
  - $a = c$  and  $b < d$

For example,

- $(5, 19)$ ,  $(13, 1)$ ,  $(13, 24)$ , and  $(15, 0)$  all have *higher* priority than  $(15, 7)$

## 7.1.4 Process Priority in Unix

This is the scheme used by Unix, e.g.,

```
% nice +15 ./a.out
```

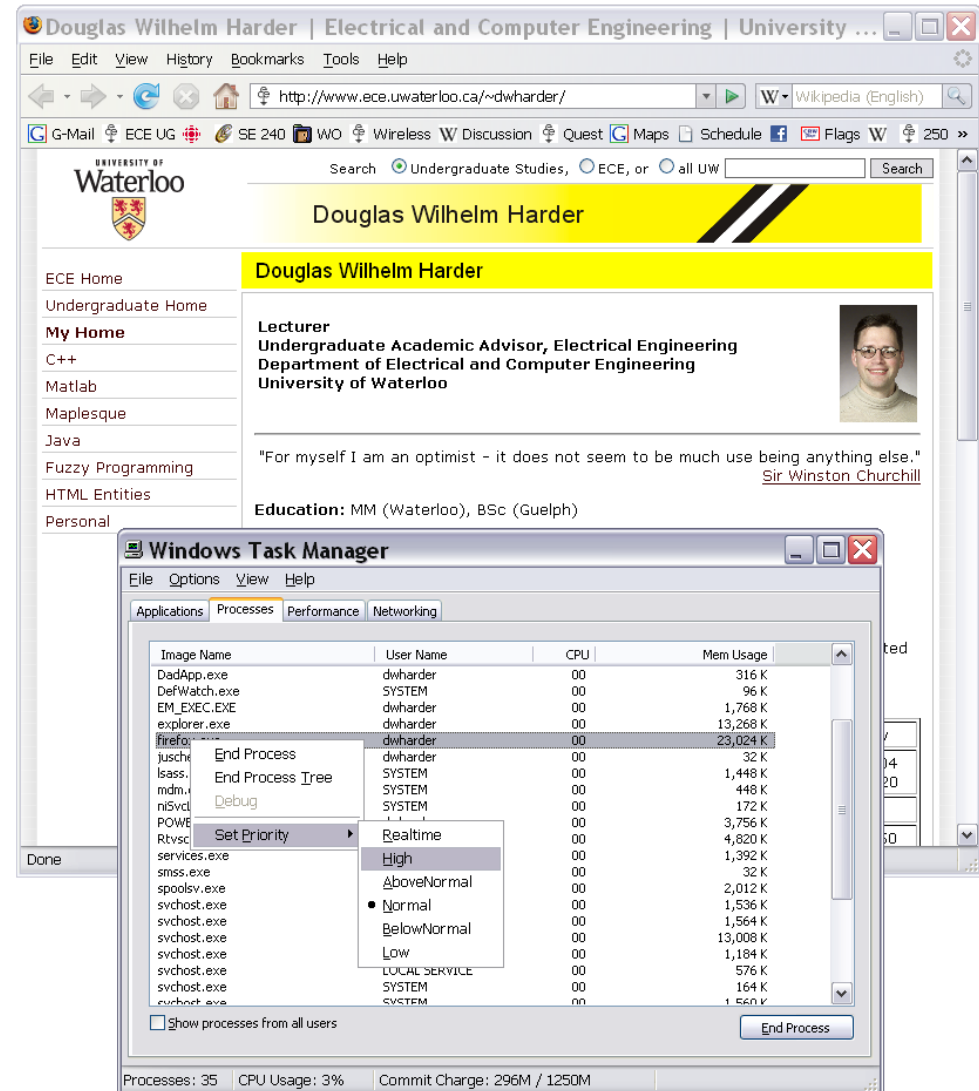
reduces the priority of the execution of the routine `a.out` by 15

This allows the processor to be used by interactive programs

- This does not significantly affect the run-time of CPU-bound processes

# 7.1.4 Process Priority in Windows

The priority of processes in Windows may be set in the *Windows Task Manager*



## 7.1.5 Implementations

Our goal is to make the run time of each operation as close to  $\Theta(1)$  as possible

We will look at two implementations using data structures we already know:

- Multiple queues—one for each priority
- An AVL tree

The next topic will be a more appropriate data structure: the heap

7.1.5.1

## Multiple Queues

Assume there is a fixed number of priorities, say  $M$

- Create an array of  $M$  queues
- Push a new object onto the queue corresponding to the priority
- Top and pop find the first empty queue with highest priority



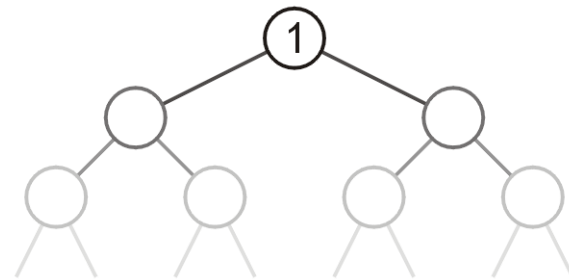
## 7.1.5.3 Heaps

Can we do better?

- That is, can we reduce some (or all) of the operations down to  $\Theta(1)$ ?

The next topic defines a *heap*

- A tree with the top object at the root
- We will look at binary heaps
- Numerous other heaps exists:
  - $d$ -ary heaps
  - Leftist heaps
  - Skew heaps
  - Binomial heaps
  - Fibonacci heaps
  - Bi-parental heaps



# References

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