

Optimization Solving Linear Equations

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system of linear equations

$$Ax = b$$

EXAMPLE 1 Determine if the following homogeneous system has a nontrivial solution. Then describe the solution set.

$$3x_1 + 5x_2 - 4x_3 = 0$$
$$-3x_1 - 2x_2 + 4x_3 = 0$$
$$6x_1 + x_2 - 8x_3 = 0$$

Solution Let A be the matrix of coefficients of the system and row reduce the augmented matrix $\begin{bmatrix} A & \mathbf{0} \end{bmatrix}$ to echelon form:

$$\begin{bmatrix} 3 & 5 & -4 & 0 \\ -3 & -2 & 4 & 0 \\ 6 & 1 & -8 & 0 \end{bmatrix} \sim \begin{bmatrix} 3 & 5 & -4 & 0 \\ 0 & 3 & 0 & 0 \\ 0 & -9 & 0 & 0 \end{bmatrix} \sim \begin{bmatrix} 3 & 5 & -4 & 0 \\ 0 & 3 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

Example from: D.C. Lay, "Linear Algebra"



$$\begin{bmatrix} 1 & 0 & -\frac{4}{3} & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix} \qquad \begin{aligned} x_1 & -\frac{4}{3}x_3 &= 0 \\ x_2 & = 0 \\ 0 &= 0 \end{aligned}$$
$$\mathbf{x} = \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} \frac{4}{3}x_3 \\ 0 \\ x_3 \end{bmatrix} = x_3 \begin{bmatrix} \frac{4}{3} \\ 0 \\ 1 \end{bmatrix} = x_3 \mathbf{v}, \quad \text{where } \mathbf{v} = \begin{bmatrix} \frac{4}{3} \\ 0 \\ 1 \end{bmatrix}$$

 $\mathbf{x}' = (2, 1)^T$

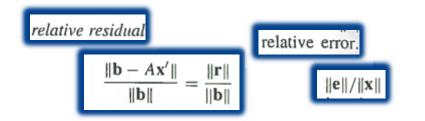


Solving Linear Equations

Solve the following system:

$$2.0000x_1 + 2.0000x_2 = 6.0000$$

 $2.0000x_1 + 2.0005x_2 = 6.0010$ $\mathbf{x} = (1, 2)^T$ $2.000x_1 + 2.000x_2 = 6.000$
 $2.000x_1 + 2.001x_2 = 6.001$



$$\frac{\|\mathbf{r}\|_{\infty}}{\|\mathbf{b}\|_{\infty}} = \frac{0.0005}{6.0010} \approx 0.000083 \qquad \frac{\|\mathbf{e}\|_{\infty}}{\|\mathbf{x}\|_{\infty}} = 0.5$$

Can the relative error be esitimated by the residual error?

$$\frac{\|\mathbf{r}\|}{\|A\|} \le \|\mathbf{e}\| \le \|A^{-1}\| \|\mathbf{r}\|$$

$$\frac{\|\mathbf{b}\|}{\|A\|} \le \|\mathbf{x}\| \le \|A^{-1}\| \|\mathbf{b}\|$$

$$\frac{1}{\|A\| \|A^{-1}\|} \frac{\|\mathbf{r}\|}{\|\mathbf{b}\|} \le \frac{\|\mathbf{e}\|}{\|\mathbf{x}\|} \le \|A\| \|A^{-1}\| \frac{\|\mathbf{r}\|}{\|\mathbf{b}\|}$$

The number $||A|| ||A^{-1}||$ is called the *condition number* of A and will be denoted by cond(A). Thus

(10)
$$\frac{1}{\text{cond}(A)} \frac{\|\mathbf{r}\|}{\|\mathbf{b}\|} \le \frac{\|\mathbf{e}\|}{\|\mathbf{x}\|} \le \text{cond}(A) \frac{\|\mathbf{r}\|}{\|\mathbf{b}\|}$$

$$||A||_{\infty} = \max_{1 \le i \le m} \left(\sum_{j=1}^{n} |a_{ij}| \right)$$

Cond(A)~16000



Try this:

- Calculate the relative residual for: $A = \begin{bmatrix} 3 & 2 \\ 1 & 1 \end{bmatrix}$ and $b = \begin{bmatrix} 5 \\ 2 \end{bmatrix}$ when the computed solution is $x' = \begin{bmatrix} 1.1 \\ 0.88 \end{bmatrix}$
- Calculate $cond_{\infty}(A)$
- The correct solution $\mathbf{x'} = \begin{bmatrix} \mathbf{1} \\ \mathbf{1} \end{bmatrix}$ does this comply with the bounds on the relative error: $\frac{\|e\|}{\|x\|}$



Solving Linear Equations

Let x^* be a vector that minimizes $||Ax - b||^2$; that is, for all $x \in \mathbb{R}^n$,

$$||Ax - b||^2 \ge ||Ax^* - b||^2$$
.

We refer to the vector x^* as a least-squares solution to Ax = b.

Theorem 12.1 The unique vector \mathbf{x}^* that minimizes $\|\mathbf{A}\mathbf{x} - \mathbf{b}\|^2$ is given by the solution to the equation $\mathbf{A}^{\mathsf{T}} \mathbf{A} \mathbf{x} = \mathbf{A}^{\mathsf{T}} \mathbf{b}$; that is, $\mathbf{x}^* = (\mathbf{A}^{\mathsf{T}} \mathbf{A})^{-1} \mathbf{A}^{\mathsf{T}} \mathbf{b}$.

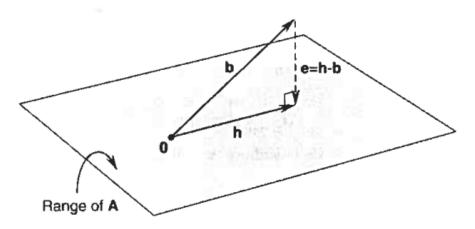


Figure 12.1 Orthogonal projection of b on the subspace $\mathcal{R}(A)$.



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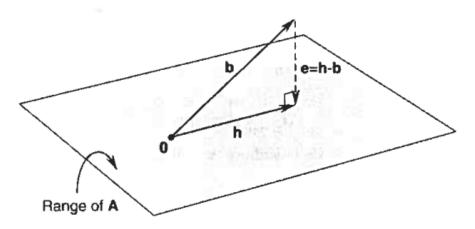


Figure 12.1 Orthogonal projection of b on the subspace $\mathcal{R}(A)$.



An alternative method of arriving at the least-squares solution is to proceed as follows. First, we write

$$f(x) = \|Ax - b\|^{2}$$

$$= (Ax - b)^{\mathsf{T}} (Ax - b)$$

$$= \frac{1}{2} x^{\mathsf{T}} (2A^{\mathsf{T}} A) x - x^{\mathsf{T}} (2A^{\mathsf{T}} b) + b^{\mathsf{T}} b.$$

Therefore, f is a quadratic function. The quadratic term is positive definite because rank A = n. Thus, the unique minimizer of f is obtained by solving the FONC (see Exercise 6.31); that is,

$$\nabla f(\boldsymbol{x}) = 2\boldsymbol{A}^{\mathsf{T}} \boldsymbol{A} \boldsymbol{x} - 2\boldsymbol{A}^{\mathsf{T}} \boldsymbol{b} = \boldsymbol{0}.$$

The only solution to the equation $\nabla f(x) = 0$ is $x^* = (A^T A)^{-1} A^T b$.



Example 12.1 Suppose that you are given two different types of concrete. The first type contains 30% cement, 40% gravel, and 30% sand (all percentages of weight). The second type contains 10% cement, 20% gravel, and 70% sand. How many pounds of each type of concrete should you mix together so that you get a concrete mixture that has as close as possible to a total of 5 pounds of cement, 3 pounds of gravel, and 4 pounds of sand?

The problem can be formulated as a least-squares problem with



Example 12.2 Line Fitting. Suppose that a process has a single input $t \in \mathbb{R}$ and a single output $y \in \mathbb{R}$. Suppose that we perform an experiment on the Table 12.1 Experimental Data for Example 12.2.

process, resulting in a number of measurements, as displayed in Table 12.1.

$$y = mt + c$$

$$\|Ax - b\|^2 = \sum_{i=0}^{2} (mt_i + c - y_i)^2.$$



Example 12.3 Attenuation Estimation. A wireless transmitter sends a discrete-time signal $\{s_0, s_1, s_2\}$ (of duration 3) to a receiver, as shown in Figure 12.3. The real number s_i is the value of the signal at time i.

The transmitted signal takes two paths to the receiver: a direct path, with delay 10 and attenuation factor a_1 , and an indirect (reflected) path, with delay 12 and attenuation factor a_2 . The received signal is the sum of the signals from these two paths, with their respective delays and attenuation factors.

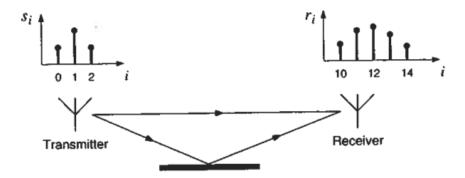


Figure 12.3 Wireless transmission in Example 12.3.



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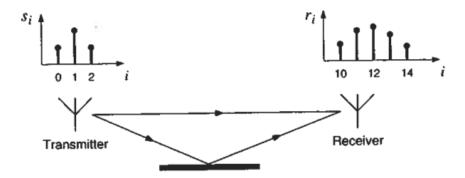


Figure 12.3 Wireless transmission in Example 12.3.



Suppose that the received signal is measured from times 10 through 14 as $r_{10}, r_{11}, \ldots, r_{14}$, as shown in the figure. We wish to compute the least-squares estimates of a_1 and a_2 , based on the following values:

The problem can be posed as a least-squares problem with

$$m{A} = egin{bmatrix} s_0 & 0 \ s_1 & 0 \ s_2 & s_0 \ 0 & s_1 \ 0 & s_2 \end{bmatrix}, \qquad m{x} = egin{bmatrix} a_1 \ a_2 \end{bmatrix}, \qquad m{b} = egin{bmatrix} r_{10} \ r_{11} \ r_{12} \ r_{13} \ r_{14} \end{bmatrix}.$$



The recursive least-mean sqaure

$$\left\|egin{bmatrix} m{A_0} \ m{A_1} \end{bmatrix}m{x} = egin{bmatrix} m{b^{(0)}} \ m{b^{(1)}} \end{bmatrix}
ight\|^2.$$

$$\boldsymbol{x}^{(1)} = \boldsymbol{G}_1^{-1} \begin{bmatrix} \boldsymbol{A}_0 \\ \boldsymbol{A}_1 \end{bmatrix}^{\mathsf{T}} \begin{bmatrix} \boldsymbol{b}^{(0)} \\ \boldsymbol{b}^{(1)} \end{bmatrix},$$

$$G_1 = \begin{bmatrix} A_0 \\ A_1 \end{bmatrix}^\mathsf{T} \begin{bmatrix} A_0 \\ A_1 \end{bmatrix}.$$

$$\begin{aligned} \boldsymbol{x}^{(1)} &= \boldsymbol{G}_{1}^{-1} \begin{bmatrix} \boldsymbol{A}_{0} \\ \boldsymbol{A}_{1} \end{bmatrix}^{\mathsf{T}} \begin{bmatrix} \boldsymbol{b}^{(0)} \\ \boldsymbol{b}^{(1)} \end{bmatrix} \\ &= \boldsymbol{G}_{1}^{-1} \left(\boldsymbol{G}_{1} \boldsymbol{x}^{(0)} - \boldsymbol{A}_{1}^{\mathsf{T}} \boldsymbol{A}_{1} \boldsymbol{x}^{(0)} + \boldsymbol{A}_{1}^{\mathsf{T}} \boldsymbol{b}^{(1)} \right) \\ &= \boldsymbol{x}^{(0)} + \boldsymbol{G}_{1}^{-1} \boldsymbol{A}_{1}^{\mathsf{T}} \left(\boldsymbol{b}^{(1)} - \boldsymbol{A}_{1} \boldsymbol{x}^{(0)} \right), \end{aligned}$$

where G_1 can be calculated using

$$\boldsymbol{G}_1 = \boldsymbol{G}_0 + \boldsymbol{A}_1^\mathsf{T} \boldsymbol{A}_1.$$



The recursive least-mean sqaure

$$G_{k+1} = G_k + A_{k+1}^{\top} A_{k+1}$$

 $x^{(k+1)} = x^{(k)} + G_{k+1}^{-1} A_{k+1}^{\top} \left(b^{(k+1)} - A_{k+1} x^{(k)} \right).$

In the special case where the new data at each step are such that A_{k+1} is a matrix consisting of a single row, $A_{k+1} = a_{k+1}^{\mathsf{T}}$, and $b^{(k+1)}$ is a scalar,

$$P_{k+1} = P_k - \frac{P_k a_{k+1} a_{k+1}^{\mathsf{T}} P_k}{1 + a_{k+1}^{\mathsf{T}} P_k a_{k+1}},$$

 $x^{(k+1)} = x^{(k)} + P_{k+1} a_{k+1} \left(b_{k+1} - a_{k+1}^{\mathsf{T}} x^{(k)} \right).$



The recursive least-mean sqaure

Example 12.6 Let

$$A_0 = \begin{bmatrix} 1 & 0 \\ 0 & 1 \\ 1 & 1 \end{bmatrix} b^{(0)} = \begin{bmatrix} 1 \\ 1 \\ 1 \end{bmatrix},$$

$$A_1 = a_1^{\mathsf{T}} = \begin{bmatrix} 2 & 1 \end{bmatrix} b^{(1)} = b_1 = \begin{bmatrix} 3 \end{bmatrix},$$

$$A_2 = a_2^{\mathsf{T}} = \begin{bmatrix} 3 & 1 \end{bmatrix} b^{(2)} = b_2 = \begin{bmatrix} 4 \end{bmatrix}.$$

First compute the vector $x^{(0)}$ minimizing $||A_0x - b^{(0)}||^2$. Then, use the RLS algorithm to find $x^{(2)}$ minimizing

$$\begin{bmatrix} A_0 \\ A_1 \\ A_2 \end{bmatrix} x - \begin{bmatrix} b^{(0)} \\ b^{(1)} \\ b^{(2)} \end{bmatrix} \end{bmatrix}^2.$$

We have

$$P_0 = (A_0^{\mathsf{T}} A_0)^{-1} = \begin{bmatrix} 2/3 & -1/3 \\ -1/3 & 2/3 \end{bmatrix},$$

$$x^{(0)} = P_0 A_0^{\mathsf{T}} b^{(0)} = \begin{bmatrix} 2/3 \\ 2/3 \end{bmatrix}.$$

Applying the RLS algorithm twice, we get:

$$P_{1} = P_{0} - \frac{P_{0}a_{1}a_{1}^{\mathsf{T}}P_{0}}{1 + a_{1}^{\mathsf{T}}P_{0}a_{1}} = \begin{bmatrix} 1/3 & -1/3 \\ -1/3 & 2/3 \end{bmatrix},$$

$$x^{(1)} = x^{(0)} + P_{1}a_{1} \left(b_{1} - a_{1}^{\mathsf{T}}x^{(0)}\right) = \begin{bmatrix} 1 \\ 2/3 \end{bmatrix},$$

$$P_{2} = P_{1} - \frac{P_{1}a_{2}a_{2}^{\mathsf{T}}P_{1}}{1 + a_{2}^{\mathsf{T}}P_{1}a_{2}} = \begin{bmatrix} 1/6 & -1/4 \\ -1/4 & 5/8 \end{bmatrix},$$

$$x^{(2)} = x^{(1)} + P_{2}a_{2} \left(b_{2} - a_{2}^{\mathsf{T}}x^{(1)}\right) = \begin{bmatrix} 13/12 \\ 5/8 \end{bmatrix}.$$



Solution with minimum norm

minimize $\|x\|$ subject to Ax = b.

Theorem 12.2 The unique solution x^* to Ax = b that minimizes the norm ||x|| is given by

$$\boldsymbol{x}^* = \boldsymbol{A}^{\top} (\boldsymbol{A} \boldsymbol{A}^{\top})^{-1} \boldsymbol{b}.$$

KACZMARZ'S ALGORITHM

Let a_j^{T} denote the jth row of A, and b_j the jth component of b, and μ a positive scalar, $0 < \mu < 2$. With this notation, Kaczmarz's algorithm is:

- 1. Set i := 0, initial condition $x^{(0)}$.
- 2. For $j = 1, \dots, m$, set $\boldsymbol{x}^{(im+j)} = \boldsymbol{x}^{(im+j-1)} + \mu \left(b_j \boldsymbol{a}_j^\top \boldsymbol{x}^{(im+j-1)} \right) \frac{\boldsymbol{a}_j}{\boldsymbol{a}_j^\top \boldsymbol{a}_j}.$
- 3. Set i := i + 1; go to step 2.

Theorem 12.3 In Kaczmarz's algorithm, if $\mathbf{x}^{(0)} = \mathbf{0}$, then $\mathbf{x}^{(k)} \to \mathbf{x}^* = \mathbf{A}^{\top} (\mathbf{A} \mathbf{A}^{\top})^{-1} \mathbf{b}$ as $k \to \infty$.



Moore-Penrose invers

Theorem 12.6 Consider a system of linear equations Ax = b, $A \in \mathbb{R}^{m \times n}$, rank A = r. The vector $x^* = A^{\dagger}b$ minimizes $||Ax - b||^2$ on \mathbb{R}^n . Furthermore, among all vectors in \mathbb{R}^n that minimize $||Ax - b||^2$, the vector $x^* = A^{\dagger}b$ is the unique vector with minimal norm.

For the case in which a matrix $A \in \mathbb{R}^{m \times n}$ with $m \geq n$ and rank A = n, we can easily check that the following is a pseudoinverse of A:

$$A^{\dagger} = (A^{\top}A)^{-1}A^{\top}.$$

For the case in which a matrix $A \in \mathbb{R}^{m \times n}$ with $m \leq n$ and rank A = m, we can easily check, as we did in the previous case, that the following is a pseudoinverse of A:

 $\mathbf{A}^{\dagger} = \mathbf{A}^{\mathsf{T}} (\mathbf{A} \mathbf{A}^{\mathsf{T}})^{-1}.$