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# Lab Report — Symmetric Encryption & Hashing

## Tools Used:

Ubuntu Bash Terminal, OpenSSL, GHex (Hex Editor)

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## Objective

To perform experiments with symmetric encryption (AES in multiple modes), hashing algorithms, and keyed hashing (HMAC). The aims were to explore encryption/decryption behavior, visualize the effect of block cipher modes on images, observe the impact of ciphertext corruption, understand padding requirements, and compare digest outputs across algorithms.

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## Environment & Setup

- Ubuntu environment running under Windows subsystem
  - OpenSSL pre-installed (`openssl` command available)
  - GHex installed for manual byte-level editing
  - Workspace contained: plaintext files, encrypted binary files, BMP images, hash outputs, and HMAC outputs
  - All commands were executed directly from the terminal
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# Task 1 — AES Encryption With Different Modes

## Goal

Encrypt and decrypt a text file using at least three AES modes (CBC, ECB, CFB) and verify correctness.

## Files Used

`plain.txt` — a simple text file.

## Commands Executed (Examples)

### AES-128-CBC

```
openssl enc -aes-128-cbc -e -in plain.txt -out aes128_cbc.bin \  
-K 00112233445566778899aabbccddeeff \  
-iv 112233445566778899aabbccddeeff00  
  
openssl enc -aes-128-cbc -d -in aes128_cbc.bin -out decrypted_cbc.txt \  
-K 00112233445566778899aabbccddeeff \  
-iv 112233445566778899aabbccddeeff00
```

### AES-128-ECB

```
openssl enc -aes-128-ecb -e -in plain.txt -out aes128_ecb.bin \  
-K 00112233445566778899aabbccddeeff
```

### AES-128-CFB

```
openssl enc -aes-128-cfb -e -in plain.txt -out aes128_cfb.bin \  
-K 00112233445566778899aabbccddeeff \  
-iv 112233445566778899aabbccddeeff00
```

## Observations

- Decryption succeeded for all modes; output matched the original plaintext.

- CBC and CFB required IVs; ECB ignored IV.
  - Stream-like modes behave differently in terms of padding and internal processing.
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## Task 2 — ECB vs CBC Encryption on an Image

### Goal

Encrypt a BMP image in ECB and CBC mode, patch the header, and observe visual differences.

### Procedure

1. Encrypt the file using AES-128-ECB and AES-128-CBC.
2. Open the encrypted binaries using GHex.
3. Copy the first 54 bytes (BMP header) from the original file into each encrypted output.
4. Save modified files as viewable BMP images.

### Observations

- **ECB Image:** Structural patterns from the original were still visible because ECB encrypts each block independently. Repeated blocks in the plaintext produce repeated blocks in ciphertext.
- **CBC Image:** Appeared as random noise. No recognizable shapes. The chaining mechanism hides repeated patterns.

### Conclusion

CBC provides better protection of visual structure than ECB.

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# Task 3 — Corrupted Ciphertext Behavior

## Goal

Modify a single byte in ciphertext and observe how corruption spreads across different AES modes.

## Steps

1. Create a long text file (`long_plain.txt`) with >64 bytes.
2. Encrypt using modes ECB, CBC, CFB, OFB.
3. Open ciphertext in GHex.
4. Flip one bit at byte index 29 and save as a corrupted file.
5. Decrypt using same key/IV.

## Results

- **ECB:** Only the block containing the corrupted byte was affected. Others decrypted normally.
  - **CBC:** The corrupted block was completely garbled, and the corresponding byte in the next block flipped. Remaining blocks unaffected.
  - **CFB:** Only a limited region was corrupted. Resembles stream cipher behavior.
  - **OFB:** Only the specific byte position was corrupted; no propagation occurred.
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# Task 4 — Padding Requirements

## Goal

Determine which AES modes require padding.

## Experiment

Encrypted a short text file (<16 bytes) with CBC, ECB, CFB, and OFB.

## Findings

- **Modes requiring padding:**
  - **ECB, CBC** (because they operate in fixed 16-byte blocks)
- **Modes not requiring padding:**
  - **CFB, OFB** (work byte-by-byte)

OpenSSL automatically added padding for CBC and ECB.

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# Task 5 — Generating Message Digests

## Goal

Generate digests of a file using MD5, SHA-1, and SHA-256.

## Commands

```
openssl dgst -md5 file.txt  
openssl dgst -sha1 file.txt  
openssl dgst -sha256 file.txt
```

## Observations

- Each algorithm produced a different fixed-length hash.
  - SHA-256's 256-bit digest is longest and provides the strongest security.
  - Even 1-bit change in the input caused drastically different hash outputs.
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# Task 6 — Keyed Hashing (HMAC)

## Goal

Generate HMAC values for a file using different keys and algorithms.

## Commands

```
openssl dgst -sha256 -hmac "mykey123" file.txt
```

```
openssl dgst -sha1 -hmac "mykey123" file.txt
```

```
openssl dgst -md5 -hmac "mykey123" file.txt
```

## Observations

- HMAC outputs changed whenever the key changed.
- HMAC accepts keys of any length; long keys are internally hashed or padded.
- SHA-256 based HMAC produced the longest and strongest output

# Task 7 — Avalanche Effect Test

## Goal

The objective of this task is to demonstrate the **Avalanche Effect**, a fundamental property of cryptographic hash functions. The Avalanche Effect states that even if a single bit in the input changes, more than 50% of the output hash bits will change, making the new hash appear completely unrelated to the original.

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## Step 1 — Creating the Input File and Generating $H_1$

### A. Creating the Input File

A text file named `input_avalanche.txt` was created using the following command:

```
gedit input_avalanche.txt
```

A few lines of text (3–4 lines) were manually added and the file was saved.

## B. Generating the First Hash Value ( $H_1$ )

The initial hash values were generated using MD5 and SHA256:

```
# MD5 H1
```

```
openssl dgst -md5 input_avalanche.txt
```

```
# SHA256 H1
```

```
openssl dgst -sha256 input_avalanche.txt
```

These two outputs were recorded as  **$H_1(\text{MD5})$**  and  **$H_1(\text{SHA256})$**  in the report.

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## Step 2 — Modifying One Bit and Generating $H_2$

### A. Flipping One Bit Using GHex

The file was opened in the hex editor:

```
ghex input_avalanche.txt &
```

Inside GHex, one byte was modified to change exactly one bit.

For example:

A byte with value **3A** was changed to **3B**, which differs by only 1 bit.

The file was then saved.

### B. Generating the Second Hash Value ( $H_2$ )

```
# MD5 H2
```

```
openssl dgst -md5 input_avalanche.txt
```

```
# SHA256 H2
```

```
openssl dgst -sha256 input_avalanche.txt
```

These outputs were recorded as  $H_2(\text{MD5})$  and  $H_2(\text{SHA256})$ .

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## Step 3 — Observation

### Difference Between $H_1$ and $H_2$

A clear and significant difference was observed:

- The  $H_1$  and  $H_2$  hash values looked completely different.
- Even though only a single bit in the input was changed, both MD5 and SHA256 produced totally different output hashes.
- It appears as if the hashes came from two entirely different files.

### Reason — Avalanche Effect

This behavior confirms the **Avalanche Effect**:

- A minor change (even a 1-bit modification) in the input results in a drastically different hash value.
  - For secure hash functions, **more than 50% of the output bits typically change**.
  - This property ensures that hash functions are highly sensitive to input changes and prevents attackers from predicting output patterns.
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