#### Propositions and Predicates

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In this class, we shall present how Coq's type system allows us to express properties of programs and/or mathematical objects. We will try to show the great expressive power of this formalism, mostly by examples.

Let e and e' be two expressions of the same type. We can build a proposition which expresses the equality between e and e'.

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Check 1+1 = 2.

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$$1+1=2$$
: Prop

Check 
$$2 = 3$$
.

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Check negb (negb true) = true.

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## Building Propositions from Predicates

Check Zlt.

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 $Zlt: Z \rightarrow Z \rightarrow Prop$ 

Check Zlt 2 3.

2 < 3 : Prop

### **Building Propositions from Predicates**

```
Check Zlt.

Zlt: Z -> Z -> Prop

Check Zlt 2 3.

2 < 3: Prop

Check le.

le: nat -> nat -> Prop

Check le 0%nat 6%nat.
```

Check Zlt.

## **Building Propositions from Predicates**

```
Zlt: Z \rightarrow Z \rightarrow Prop
Check Zlt 2 3.
2 < 3: Prop
Check le.
le: nat \rightarrow nat \rightarrow Prop
Check le 0%nat 6%nat.
(0 <= 6)%nat: Prop
```

#### Don't be mistaken!

Check Zlt\_bool 2 3.

Zlt\_bool 2 3 : bool

Definition Zmax n p := if n < p then p else n.

#### Don't be mistaken!

Check Zlt\_bool 2 3. Zlt\_bool 2 3: bool

```
Definition Zmax n p := if n n < p" has type "Prop" ... *)
```

#### Don't be mistaken!

Check Zlt bool 2 3.

```
Zlt_bool 2 3 : bool
Definition Zmax n p := if n < p then p else n.
(* Error : the term " n < p " has type "Prop" ... *)
Definition Zmax n p := if Zlt_bool n p then p else n.
Notice that the following examples are well formed propositions:
Ztl bool 2 3 = true
Zlt bool 2 3 = false
Zeq_bool (6*6) (9*4) = true
6*6=9*4
45 \le 7 \text{max} 34 45
```

#### Quantifiers and Connectives

The following are well-formed propositions:

```
forall n:Z, 0 <= n * n
```

(\* The square of any integer is greater or equal than 0 \*)

(\* There exists at least some integer whose square is 4 \*) exists n:Z, n \* n = 4

(\* Z is unbounded \*)
forall n:Z, exists p:Z, n < p.</pre>

(\* A well-formed, unprovable proposition \*)
forall n : Z, n ^ 2 <= 2 ^ n.</pre>

There exists some useful notations for nested quantifiers, which we shall present in further examples.

```
(* Zlt is irreflexive *)
```

Check Zlt\_irrefl.

```
(* Zlt is irreflexive *)
```

Check Zlt\_irrefl.

 $Zlt\_irrefl: for all \ n: Z, \sim n < n$ 

```
(* Zlt is irreflexive *)
```

Check Zlt\_irrefl.  $Zlt_irrefl: forall \ n: Z, \sim n < n$ 

Check forall n : Z,  $\sim$  n < n.

```
(* Zlt is irreflexive *)
Check Zlt irrefl.
Zlt_irrefl : forall n : Z. \sim n < n
Check forall n : Z, \sim n < n.
forall n: Z_n \sim n < n: Prop
(* There is no integer square root of 2 *)
Check \sim(exists n:Z, n*n = 2).
Require Import List.
(* No number in the empty list of integers ! *)
forall z:Z, \sim In z nil.
\sim (exists z:Z. In z nil).
```

## Implication (->)

```
(* Zle_trans *)
forall n m p : Z, n <= m -> m <= p -> n <= p.

(* Zlt_asym *)
forall n p:Z, n < p -> ~ p < n.</pre>
```

## Disjunction (or)

```
forall n:Z, 0 <= n \lor n < 0.

forall n p : Z, n < p \lor p <= n.

forall n p : Z, n < p \lor p = n \lor p < n.

(forall n : nat, n = 0 \lor exists p:nat, p < n)\( n \)

forall l:list Z,
    l = nil \lor exists a, exists l', l = a::l'.</pre>
```

# Conjonction (and)

```
let (q,r) := Zdiv_{eucl} 456 37 in
456 = 37 * q + r \land
0 <= r < 37. (* 0 <= r \land r < 37 *)
forall a b q r: Z, 0 < b ->
a = b * q + r ->
0 <= r < b ->
q = a / b \land r = a mod b.
```

# Logical Equivalence (iff)

```
Coq.ZArith.Zbool.Zle_bool :
forall n m : Z, n <= m <-> Zle_bool n m = true

forall l1 l2 : list Z,
    (forall z:Z, In z (l1 ++ l2) <->
        In z l1 ∨ In z l2).
```

## **Building new Predicates**

#### **Building new Predicates**

```
(* number of occurences of n in 1 *)
Fixpoint multiplicity (n:Z)(1:list Z) : nat :=
  match 1 with
    nil => 0%nat
  | a::1' => if Zeq_bool n a
             then S (multiplicity n 1')
             else multiplicity n 1'
  end.
(* 1' is a permutation of 1 *)
Definition is_perm (1 1':list Z) :=
    forall n, multiplicity n l = multiplicity n l'.
```

### Specifying a merge function

```
(* The binary function m preserves
   the elements' multiplicity *)

Definition preserves_multiplicity
      (m : list Z -> list Z -> list Z) :=
  forall l l' n,
      multiplicity n (m l l') =
      (multiplicity n l + multiplicity n l')%nat.
```

# Specifying a merge function (2)

```
(* let's assume the following predicate "to be sorted"
  is defined *)
Parameter sorted_Zle : list Z -> Prop.

Definition preserves_sort
      (m : list Z -> list Z -> list Z) :=
  forall l l', sorted_Zle l -> sorted_Zle l' ->
            sorted_Zle (m l l').
```

Definition merge\_spec (m : list Z -> list Z -> list Z):=

preserves\_sort m ∧ preserves\_multiplicity m.

forall P Q : Prop,  $\sim$  (P  $\vee$  Q) ->  $\sim$  P  $\wedge$   $\sim$  Q.

```
forall P Q : Prop, \sim (P \lor Q) -> \sim P \land \sim Q.
```

forall P : Prop,  $\sim$  P <-> P -> False.

```
forall P Q : Prop, \sim (P \lor Q) -> \sim P \land \sim Q.
```

forall P : Prop, 
$$\sim$$
 P <-> P -> False.

forall P Q R:Prop, 
$$(P \land Q \rightarrow R) \longleftrightarrow (P \rightarrow Q \rightarrow R)$$
.

```
forall P Q : Prop, \sim (P \vee Q) -> \sim P \wedge \sim Q. forall P : Prop, \sim P <-> P -> False. forall P Q R:Prop, (P \wedge Q -> R) <-> (P -> Q -> R). forall P Q, P \vee Q -> Q \vee P.
```

False\_ind: forall P : Prop, False -> P

```
forall P Q : Prop, \sim (P \vee Q) -> \sim P \wedge \sim Q. forall P : Prop, \sim P <-> P -> False. forall P Q R:Prop, (P \wedge Q -> R) <-> (P -> Q -> R). forall P Q, P \vee Q -> Q \vee P.
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forall P Q : Prop, \sim (P \vee Q) -> \sim P \wedge \sim Q. forall P : Prop, \sim P <-> P -> False. forall P Q R:Prop, (P \wedge Q -> R) <-> (P -> Q -> R). forall P Q, P \vee Q -> Q \vee P.
```

absurd: forall A C : Prop, A -> ~ A -> C

False\_ind: forall P : Prop, False -> P

forall P : nat -> Prop,  $\sim$  (exists n, P n) -> forall n,  $\sim$  P n.

```
forall P : nat \rightarrow Prop, \sim (exists n, P n) \rightarrow
                              forall n, \sim P n.
nat_ind: forall P : nat -> Prop,
 P N ->
 (forall n:nat, P n \rightarrow P (S n)) ->
 forall n:nat, P n.
(forall P:Prop, P \lor \sim P) <->
(forall P:Prop, \sim \sim P \rightarrow P).
```

Definition or\_ex (P Q : Prop) := P  $\lor$  Q  $\land$  ( $\sim$ P  $\land$   $\sim$  Q).

Quantifying over propositions and predicates

Definition or\_ex (P Q : Prop) := P  $\vee$  Q  $\wedge$  ( $\sim$ P  $\wedge$   $\sim$  Q).

Lemma or\_ex\_not\_iff : forall P Q, or\_ex P Q ->  $\sim$  (P <-> Q).

Quantifying over propositions and predicates

```
SearchRewrite (rev (rev _)).
rev_involutive:
   forall (A : Type) (1 : list A), rev (rev 1) = 1
```

```
SearchRewrite (rev (rev )).
rev involutive:
   forall (A : Type) (1 : list A), rev (rev 1) = 1
forall (A:Type)(P:A->Prop), ~(exists x, P x ) ->
                              forall x. ~ P x.
forall (A:Type)(x y z:A), x = y \rightarrow y = z \rightarrow x = z.
forall (A B:Type)(a:A)(b:B), fst (a,b) = a.
forall (A B : Type)(p:A*B), p = (fst p, snd p).
```

## A Little Case Study

```
boolean function *)
Definition decides (A:Type)(P:A->Prop)(p : A -> bool) :=
```

Compatibility between a predicate and a

forall a:A, P a <-> (p a)=true.

# A Little Case Study

```
boolean function *)

Definition decides (A:Type)(P:A->Prop)(p : A -> bool) :=
  forall a:A, P a <-> (p a)=true.

Definition decides2
     (A:Type)(P:A->A->Prop)(p : A -> A-> bool) :=
  forall a b :A , P a b <-> p a b = true.

Check decides2 _ Zle Zle_bool.

decides2 Z Zle Zle_bool : Prop
```

Compatibility between a predicate and a

```
Require Import Relations.
Print order.
Record order (A : Type) (R : relation A) : Prop :=
Build order
  { ord_refl : reflexive A R;
     ord_trans : transitive A R;
     ord_antisym : antisymmetric A R }
Print antisymmetric.
antisymmetric =
fun (A : Type) (R : relation A) =>
   forall x y : A, R x y \rightarrow R y x \rightarrow x = y
: forall A : Type, relation A -> Prop
```

```
Section sort_spec.
Parameter sorted :
  forall (A:Type), relation A -> list A -> Prop.
```

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Parameter sorted :
  forall (A:Type), relation A -> list A -> Prop.

Variable s:
  forall A:Type,(A->A->bool) -> list A -> list A.
```

```
Section sort_spec.
Parameter sorted:
  forall (A:Type), relation A -> list A -> Prop.
Variable s:
  forall A:Type, (A \rightarrow A \rightarrow bool) \rightarrow list A \rightarrow list A.
Definition sort_correct := (* to improve ! *)
    forall (A:Type)
              (R : relation A)
              (r : A \rightarrow A \rightarrow bool),
                    order A R -> decides2 A R r ->
                   forall 1, let 1' := s A r l in
                       sorted A R 1' A
                        forall a, In a 1 <-> In a 1'.
```