

<i>Z, len</i>	integer
<i>f</i>	float
<i>n</i>	int
<i>ident, id</i>	identifier
<i>label, l</i>	label
<i>dcls</i>	global variable declarations
<i>fndefns</i>	function declarations
<i>opt_tid</i>	optional thread id
<i>ef_sig</i>	external signature
<i>p</i>	pointer
<i>typ</i>	
<i>fundef, fd</i>	
<i>fn_body</i>	
<i>fn</i>	
<i>ge</i>	
δ	

<i>signedness</i>	<pre> ::= Signed Unsigned </pre>	Signedness
<i>intsize</i>	<pre> ::= I8 I16 I32 </pre>	Integer sizes
<i>floatsize</i>	<pre> ::= F32 F64 </pre>	Float sizes
<i>type, ty</i>	<pre> ::= void int(<i>intsize</i>, <i>signedness</i>) float(<i>floatsize</i>) pointer(<i>ty</i>) array(<i>ty</i>, <i>len</i>) function(<i>ty</i>*, <i>ty</i>) struct(<i>id</i>, <i>φ</i>) union(<i>id</i>, <i>φ</i>) comp_pointer(<i>id</i>) (<i>ty</i>) </pre>	Types the void type integer types floating-point types pointer types (*ty) array types (ty[<i>len</i>]) function types struct types union types pointer to named struct or union S
<i>typelist, ty</i> *	<pre> ::= nil <i>ty</i>:::<i>ty</i>* </pre>	Type list
<i>fieldlist, φ</i>	<pre> ::= nil (<i>id</i>, <i>ty</i>):::<i>φ</i> </pre>	Field lists
<i>unary_operation, op</i> ₁	<pre> ::= ! ~ - </pre>	unary Boolean negation Integer complement opposite
<i>binary_operation, op</i> ₂	<pre> ::= + - * / % & ^ << >> </pre>	binary addition subtraction multiplication division modulo bitwise and bitwise or bitwise xor left shift right shift

		==	equality
		!=	not equal
		<	less than
		>	greater than
		<=	less than equal
		>=	greater than equal
<i>expr, e</i>	::=		typed expression
		<i>a</i> ^{<i>ty</i>}	expression
<i>expr_descr, a</i>	::=		basic expressions
		<i>n</i>	integer literal
		<i>f</i>	float literal
		<i>id</i>	variable
		* <i>e</i>	unary pointer dereference
		& <i>e</i>	address-of
		<i>op</i> ₁ <i>e</i>	unary operation
		<i>e</i> ₁ <i>op</i> ₂ <i>e</i> ₂	binary operation
		(<i>ty</i>) <i>e</i>	type cast
		<i>e</i> ₁ ? <i>e</i> ₂ : <i>e</i> ₃	conditional
		<i>e</i> ₁ && <i>e</i> ₂	sequential and
		<i>e</i> ₁ <i>e</i> ₂	sequential or
		sizeof (<i>ty</i>)	size of a type
		<i>e</i> . <i>id</i>	access to a member of a struct or union
<i>opt_lhs</i>	::=		optional lhs expression
		(<i>id</i> : <i>ty</i>)=	
<i>opt_e</i>	::=		optional expression
		<i>e</i>	
<i>e</i> *	::=		expression list
<i>atomic_statement, astmt</i>	::=		atomic
		CAS	compare and swap
		ATOMIC_INC	locked inc
<i>statement, s</i>	::=		statements
		skip	do nothing
		<i>e</i> ₁ = <i>e</i> ₂	assignment [lvalue = rvalue]
		<i>opt_lhs e'</i> (<i>e</i> *	function or procedure call
		<i>s</i> ₁ ; <i>s</i> ₂	sequence
		if (<i>e</i> ₁) then <i>s</i> ₁ else <i>s</i> ₂	conditional
		while (<i>e</i>) do <i>s</i>	while
		do <i>s</i> while (<i>e</i>)	do while
		for (<i>s</i> ₁ ; <i>e</i> ₂ ; <i>s</i> ₃) <i>s</i>	for loop
		break	break

	 	continue return <i>opt_e</i> switch (<i>e</i>) <i>ls</i> <i>l</i> : <i>s</i> goto <i>l</i> thread_create(<i>e</i> ₁ , <i>e</i> ₂) <i>opt_lhs astmt</i> (<i>e</i> [*]) mfence	continue return switch labelled statement goto thread creation atomic operation mfence
<i>labeled_statements, ls</i>	::= 	default : <i>s</i> case <i>n</i> : <i>s</i> ; <i>ls</i>	labeled statements default labeled case
<i>arglist</i>	::= 	<i>ty id</i> <i>ty id, arglist</i>	Argument lists
<i>varlist</i>	::= 	<i>ty id</i> ; <i>varlist</i>	Local variable lists
<i>fndefn_internal</i>	::= 	<i>ty id</i> (<i>arglist</i>) { <i>varlist s</i> }	function definition
<i>program</i>	::= 	<i>dcls fndefns</i> main = <i>id</i>	programs
<i>val, v</i>	::= 	<i>n</i> <i>f</i> <i>p</i> undef	untyped values integer value floating point value pointer undef
<i>extval, evl</i>	::= 	extint <i>n</i> extfloat <i>f</i>	external values external integer value external floating point value
<i>vs</i>	::= 	nil <i>v</i> :: <i>vs</i> <i>vs</i> @[<i>v</i>]	value list
<i>arg_list, args</i>	::= 	nil <i>id</i> ^{<i>ty</i>} :: <i>args</i>	argument lists
<i>evs</i>	::= 	nil	eventval list

		$evl :: evs$	
$memory_chunk, c$::=	Mint32	
$mobject_kind$::=	MObjStack	
$rmw_instr, rmwi$::=	ADD v CAS $v v'$ SET v	
mem_event, me	::=	write $p \text{ } memory_chunk \text{ } v$ read $p \text{ } memory_chunk \text{ } v$ alloc $p \text{ } n \text{ } mobject_kind$ free $p \text{ } mobject_kind$ rmw $p \text{ } memory_chunk \text{ } v \text{ } rmwi$ fence	
$event, ev$::=	call $id \text{ } evs$ return $typ \text{ } evl$ exit n fail	
$thread_event, te$::=	ext $event$ mem mem_event exit start $p \text{ } vs$	thread events externally observable event memory event thread-local event normal exit thread start (bootstrap)
opt_pty	::=		optional pointer/type pair
opt_v	::=		optional value
ps	::=		pointer list
ρ, ρ', ρ''	::=		environment
$cont, \kappa_S$::=	stop $[- ; s_2] \cdot \kappa_S$ $[\text{while } (e) \text{ do } s] \cdot \kappa_S$ $[\text{do } s \text{ while } (e)] \cdot \kappa_S$ $[\text{for } (; e_2 ; \diamond s_3) s] \cdot \kappa_S$ $[\text{for } (; \diamond e_2 ; s_3) s] \cdot \kappa_S$	statement continuation sequence while do while for loop, pending increment for loop, pending condition evaluation

	$\begin{array}{ l} [opt_lhs \text{ fd } (_) \mid \rho] \cdot \kappa_s \\ [\text{switch } \kappa_s] \\ [\text{free } ps; \text{return } opt_v] \cdot \kappa_s \end{array}$	call awaiting args switch protecting break
$expr_cont, \kappa_e$	$::=$	expression continuations
	$[op_1^{ty} _] \cdot \kappa_e$	unary operation
	$[- \ op_2^{ty_1 * ty_2 \rightarrow ty} \ e_2] \cdot \kappa_e$	binary operation
	$[v \ op_2^{ty_1 * ty_2 \rightarrow ty} _] \cdot \kappa_e$	binary operation
	$[(ty) _^{ty'}] \cdot \kappa_e$	
	$[- \ ^{ty} ? e_2 : e_3] \cdot \kappa_e$	
	$[- \cdot \delta] \cdot \kappa_e$	access to member of struct
	$[* _^{ty}] \cdot \kappa_e$	load
	$[- \ ^{ty} = e_2] \cdot \kappa_s$	assignment
	$[v \ ^{ty} = _] \cdot \kappa_s$	assignment
	$[opt_lhs _ \ ^{ty} (e^*)] \cdot \kappa_s$	call function
	$[opt_lhs \ v \ ^{ty} (vs, e^*)] \cdot \kappa_s$	call args
	$[opt_lhs \ astmt (vs, e^*)] \cdot \kappa_s$	atomic args
	$[\text{if } (_^{ty}) \text{ then } s_1 \text{ else } s_2] \cdot \kappa_s$	if
	$[\text{while } (_e) \text{ do } s] \cdot \kappa_s$	while
	$[\text{do } s \text{ while } (_e)] \cdot \kappa_s$	dowhile
	$[\text{for } (; _e_2; s_3) s] \cdot \kappa_s$	for loop, pending test
	$[\text{return } _] \cdot \kappa_s$	function return
	$[\text{switch } (_) \text{ ls}] \cdot \kappa_s$	switch
	$[\text{thread_create } (_, e_2)] \cdot \kappa_s$	thread creation
	$[\text{thread_create } (p, _)] \cdot \kappa_s$	thread creation
$state, \sigma$	$::=$	states
	$\text{lval } (e) \cdot \kappa_e \mid \rho$	
	$e \cdot \kappa_e \mid \rho$	
	$v \cdot \kappa_e \mid \rho$	
	$s \cdot \kappa_s \mid \rho$	
	$vs \cdot \kappa_s$	
	$\text{bind } (fn, vs, args) \cdot \kappa_s \mid \rho$	
	$\text{alloc } (vs, args) \cdot \kappa_s \mid \rho$	
	$opt_lhs \ \text{ext } (_^{tyP}) \cdot \kappa_s \mid \rho$	
	$opt_lhs \ v \cdot \kappa_s \mid \rho$	

$\sigma \xrightarrow{te} \sigma'$ Labelled Transitions (parameterised over ge)

$$\begin{array}{c}
 \frac{}{n^{ty} \cdot \kappa_e \mid \rho \longrightarrow n \cdot \kappa_e \mid \rho} \text{ STEPCONSTINT} \\
 \frac{}{f^{ty} \cdot \kappa_e \mid \rho \longrightarrow f \cdot \kappa_e \mid \rho} \text{ STEPCONSTFLOAT} \\
 \frac{}{id^{ty} \cdot \kappa_e \mid \rho \longrightarrow \text{lval } (id^{ty}) \cdot [* _^{ty}] \cdot \kappa_e \mid \rho} \text{ STEPVAREXPRBYVALUE} \\
 \frac{\rho ! id = \text{Some } p}{\text{lval } (id^{ty}) \cdot \kappa_e \mid \rho \longrightarrow p \cdot \kappa_e \mid \rho} \text{ STEPVARLOCAL} \\
 \frac{\rho ! id = \text{None} \quad \text{Genv.find_symbol ge } id = \text{Some } p}{\text{lval } (id^{ty}) \cdot \kappa_e \mid \rho \longrightarrow p \cdot \kappa_e \mid \rho} \text{ STEPVARGLOBAL}
 \end{array}$$

$$\begin{array}{c}
\text{access_mode } \textcolor{blue}{ty}' = \text{By_value } c \\
\text{typ} = \text{type_of_chunk } c \\
\text{Val.has_type } v \text{ typ} \\
\hline
p \cdot [\textcolor{blue}{*}_{\textcolor{blue}{ty}'}] \cdot \kappa_e \mid_\rho \xrightarrow{\text{mem (read } p \text{ } c \text{ } v)} v \cdot \kappa_e \mid_\rho \quad \text{STEPLoadByValue} \\
\\
\text{access_mode } \textcolor{blue}{ty}' = \text{By_reference} \setminus / \text{ access_mode } \textcolor{blue}{ty}' = \text{By_nothing} \\
\hline
p \cdot [\textcolor{blue}{*}_{\textcolor{blue}{ty}'}] \cdot \kappa_e \mid_\rho \longrightarrow p \cdot \kappa_e \mid_\rho \quad \text{STEPLoadNotByValue} \\
\\
\hline
\textcolor{blue}{\&e}^{\textcolor{blue}{ty}} \cdot \kappa_e \mid_\rho \longrightarrow \text{lval}(e) \cdot \kappa_e \mid_\rho \quad \text{STEPAddr} \\
\\
\hline
e_1 ? e_2 : e_3^{\textcolor{blue}{ty}} \cdot \kappa_e \mid_\rho \longrightarrow e_1 \cdot [\textcolor{blue}{_}^{\text{typeof } e_1 ? e_2 : e_3}] \cdot \kappa_e \mid_\rho \quad \text{STEPCondition} \\
\\
\hline
\text{is_true } v \text{ ty} \\
\hline
v \cdot [\textcolor{blue}{_}^{\textcolor{blue}{ty}} ? e_2 : e_3] \cdot \kappa_e \mid_\rho \longrightarrow e_2 \cdot \kappa_e \mid_\rho \quad \text{STEPConditionTrue} \\
\\
\hline
\text{is_false } v \text{ ty} \\
\hline
v \cdot [\textcolor{blue}{_}^{\textcolor{blue}{ty}} ? e_2 : e_3] \cdot \kappa_e \mid_\rho \longrightarrow e_3 \cdot \kappa_e \mid_\rho \quad \text{STEPConditionFalse} \\
\\
\hline
\textcolor{blue}{*}e^{\textcolor{blue}{ty}} \cdot \kappa_e \mid_\rho \longrightarrow e \cdot [\textcolor{blue}{*}_{\textcolor{blue}{ty}}] \cdot \kappa_e \mid_\rho \quad \text{STEPDeref} \\
\\
\hline
\text{lval}(\textcolor{blue}{*}e^{\textcolor{blue}{ty}}) \cdot \kappa_e \mid_\rho \longrightarrow e \cdot \kappa_e \mid_\rho \quad \text{STEPDerefLval} \\
\\
\hline
e.\textcolor{blue}{id}^{\textcolor{blue}{ty}} \cdot \kappa_e \mid_\rho \longrightarrow \text{lval}(e.\textcolor{blue}{id}^{\textcolor{blue}{ty}}) \cdot [\textcolor{blue}{*}_{\textcolor{blue}{ty}}] \cdot \kappa_e \mid_\rho \quad \text{STEPField} \\
\\
\hline
\text{typeof } e = \text{struct}(id', \phi) \\
\text{field_offset } id \ \phi = \text{OK } \delta \\
\hline
\text{lval}(e.\textcolor{blue}{id}^{\textcolor{blue}{ty}}) \cdot \kappa_e \mid_\rho \longrightarrow \text{lval}(e) \cdot [\textcolor{blue}{_} \cdot \delta] \cdot \kappa_e \mid_\rho \quad \text{STEPFStruct1} \\
\\
\hline
p' = \text{MPtr.add } p \ (\text{Int.repr } \delta) \\
\hline
p \cdot [\textcolor{blue}{_} \cdot \delta] \cdot \kappa_e \mid_\rho \longrightarrow p' \cdot \kappa_e \mid_\rho \quad \text{STEPFStruct2} \\
\\
\hline
\text{typeof } e = \text{union}(id', \phi) \\
\hline
\text{lval}(e.\textcolor{blue}{id}^{\textcolor{blue}{ty}}) \cdot \kappa_e \mid_\rho \longrightarrow \text{lval}(e) \cdot \kappa_e \mid_\rho \quad \text{STEPFUnion} \\
\\
\hline
v = \text{Vint}(\text{Int.repr}(\text{sizeof } \textcolor{blue}{ty}')) \\
\hline
\text{sizeof}(\textcolor{blue}{ty}')^{\textcolor{blue}{ty}} \cdot \kappa_e \mid_\rho \longrightarrow v \cdot \kappa_e \mid_\rho \quad \text{STEPSizeOf} \\
\\
\hline
op_1 e^{\textcolor{blue}{ty}} \cdot \kappa_e \mid_\rho \longrightarrow e \cdot [op_1^{\text{typeof } e} \textcolor{blue}{_}] \cdot \kappa_e \mid_\rho \quad \text{STEPUnOp1} \\
\\
\hline
\text{sem_unary_operation } op_1 \ v \ \textcolor{blue}{ty} = \text{Some } v' \\
\hline
v \cdot [op_1^{\textcolor{blue}{ty}} \textcolor{blue}{_}] \cdot \kappa_e \mid_\rho \longrightarrow v' \cdot \kappa_e \mid_\rho \quad \text{STEPUnOp} \\
\\
\hline
(e_1 \ op_2 \ e_2)^{\textcolor{blue}{ty}} \cdot \kappa_e \mid_\rho \longrightarrow e_1 \cdot [\textcolor{blue}{_} \ op_2^{\text{typeof } e_1 * \text{typeof } e_2 \rightarrow \textcolor{blue}{ty}} \textcolor{blue}{e}_2] \cdot \kappa_e \mid_\rho \quad \text{STEPBinOp1} \\
\\
\hline
\text{valid_arg } op_2 \ \textcolor{blue}{ty}_1 \ \textcolor{blue}{ty}_2 \ v = \text{true} \\
\hline
v \cdot [\textcolor{blue}{_} \ op_2^{\textcolor{blue}{ty}_1 * \textcolor{blue}{ty}_2 \rightarrow \textcolor{blue}{ty}} \textcolor{blue}{e}_2] \cdot \kappa_e \mid_\rho \longrightarrow e_2 \cdot [v \ op_2^{\textcolor{blue}{ty}_1 * \textcolor{blue}{ty}_2 \rightarrow \textcolor{blue}{ty}} \textcolor{blue}{_}] \cdot \kappa_e \mid_\rho \quad \text{STEPBinOp2} \\
\\
\hline
\text{sem_binary_operation } op_2 \ v_1 \ \textcolor{blue}{ty}_1 \ v_2 \ \textcolor{blue}{ty}_2 = \text{Some } v \\
\hline
v_2 \cdot [v_1 \ op_2^{\textcolor{blue}{ty}_1 * \textcolor{blue}{ty}_2 \rightarrow \textcolor{blue}{ty}} \textcolor{blue}{_}] \cdot \kappa_e \mid_\rho \longrightarrow v \cdot \kappa_e \mid_\rho \quad \text{STEPBinOp} \\
\\
\hline
(\textcolor{blue}{ty}) e^{\textcolor{blue}{ty}'} \cdot \kappa_e \mid_\rho \longrightarrow e \cdot [(\textcolor{blue}{ty}) \textcolor{blue}{_}^{\text{typeof } e}] \cdot \kappa_e \mid_\rho \quad \text{STEPCast1} \\
\\
\hline
\text{cast } v \ \textcolor{blue}{ty}' \ \textcolor{blue}{ty} \ v' \\
\hline
v \cdot [(\textcolor{blue}{ty}) \textcolor{blue}{_}^{\textcolor{blue}{ty}'}] \cdot \kappa_e \mid_\rho \longrightarrow v' \cdot \kappa_e \mid_\rho \quad \text{STEPCast2}
\end{array}$$

$$\begin{array}{c}
\frac{n_0 = \text{Int.repr } 0 \quad n_1 = \text{Int.repr } 1}{e_1 \&\& e_2^{ty} \cdot \kappa_e \mid_\rho \longrightarrow e_1 ? (e_2 ? (n_1^{ty}) : (n_0^{ty})^{ty}) : n_0^{ty ty} \cdot \kappa_e \mid_\rho} \text{STEPANDBOOL} \\
\\
\frac{n_0 = \text{Int.repr } 0 \quad n_1 = \text{Int.repr } 1}{e_1 \mid \mid e_2^{ty} \cdot \kappa_e \mid_\rho \longrightarrow e_1 ? (n_1^{ty}) : (e_2 ? (n_1^{ty}) : (n_0^{ty})^{ty})^{ty} \cdot \kappa_e \mid_\rho} \text{STEPORBOOL} \\
\\
\frac{}{\text{thread_create}(e_1, e_2) \cdot \kappa_s \mid_\rho \longrightarrow e_1 \cdot [\text{thread_create}(_, e_2)] \cdot \kappa_s \mid_\rho} \text{STEPTHREAD} \\
\\
\frac{}{p \cdot [\text{thread_create}(_, e_2)] \cdot \kappa_s \mid_\rho \longrightarrow e_2 \cdot [\text{thread_create}(p, _)] \cdot \kappa_s \mid_\rho} \text{STEPTHREADFN} \\
\\
\frac{}{v \cdot [\text{thread_create}(p, _)] \cdot \kappa_s \mid_\rho \xrightarrow{\text{start } p \text{ v} :: \text{nil}} \text{skip} \cdot \kappa_s \mid_\rho} \text{STEPTHREADEVT} \\
\\
\frac{}{e_1 = e_2 \cdot \kappa_s \mid_\rho \longrightarrow \text{lval}(e_1) \cdot [_ \text{typeof } e_1 = e_2] \cdot \kappa_s \mid_\rho} \text{STEPASSIGN1} \\
\\
\frac{}{v_1 \cdot [_^{ty} = e_2] \cdot \kappa_s \mid_\rho \longrightarrow e_2 \cdot [v_1^{ty} = _] \cdot \kappa_s \mid_\rho} \text{STEPASSIGN2} \\
\\
\frac{\text{type_to_chunk } ty_1 = \text{Some } c \quad \text{cast_value_to_chunk } c \ v_1 = v_2}{v_1 \cdot [p_1^{ty_1} = _] \cdot \kappa_s \mid_\rho \xrightarrow{\text{mem(write } p_1 \ c \ v_2)} \text{skip} \cdot \kappa_s \mid_\rho} \text{STEPASSIGN} \\
\\
\frac{}{s_1 ; s_2 \cdot \kappa_s \mid_\rho \longrightarrow s_1 \cdot [_ ; s_2] \cdot \kappa_s \mid_\rho} \text{STEPSEQ} \\
\\
\frac{}{\text{opt_lhs } e' (e^*) \cdot \kappa_s \mid_\rho \longrightarrow e' \cdot [\text{opt_lhs } _ \text{typeof } e' (e^*)] \cdot \kappa_s \mid_\rho} \text{STEPCALL} \\
\\
\frac{\text{Genv.find_funct ge } v = \text{Some } fd \quad \text{type_of_fundef } fd = ty}{v \cdot [\text{opt_lhs } _^{ty} (\text{nil})] \cdot \kappa_s \mid_\rho \longrightarrow \text{nil} \cdot [\text{opt_lhs } fd(_) \mid_\rho] \cdot \kappa_s} \text{STEPCALLARGSNONE} \\
\\
\frac{}{v \cdot [\text{opt_lhs } _^{ty} (e :: e^*)] \cdot \kappa_s \mid_\rho \longrightarrow e \cdot [\text{opt_lhs } v^{ty} (\text{nil}, e^*)] \cdot \kappa_s \mid_\rho} \text{STEPCALLARGS1} \\
\\
\frac{}{v_1 \cdot [\text{opt_lhs } v^{ty} (vs, e :: e^*)] \cdot \kappa_s \mid_\rho \longrightarrow e \cdot [\text{opt_lhs } v^{ty} (vs @ [v_1], e^*)] \cdot \kappa_s \mid_\rho} \text{STEPCALLARGS2} \\
\\
\frac{\text{Genv.find_funct ge } v = \text{Some } fd \quad \text{type_of_fundef } fd = ty}{v' \cdot [\text{opt_lhs } v^{ty} (vs, \text{nil})] \cdot \kappa_s \mid_\rho \longrightarrow vs @ [v'] \cdot [\text{opt_lhs } fd(_) \mid_\rho] \cdot \kappa_s} \text{STEPCALLFINISH} \\
\\
\frac{}{\text{opt_lhs } astmt(e :: e^*) \cdot \kappa_s \mid_\rho \longrightarrow e \cdot [\text{opt_lhs } astmt(\text{nil}, e^*)] \cdot \kappa_s \mid_\rho} \text{STEPATOMIC} \\
\\
\frac{}{v \cdot [\text{opt_lhs } astmt(vs, e :: e^*)] \cdot \kappa_s \mid_\rho \longrightarrow e \cdot [\text{opt_lhs } astmt(vs @ [v], e^*)] \cdot \kappa_s \mid_\rho} \text{STEPATOMICARGS} \\
\\
\frac{\text{sem_atomic_statement } astmt \ (\text{vs} ++ v :: \text{nil}) = \text{Some } (p, rmwi) \quad \text{Val.has_type } v' \ (\text{type_of_chunk } \text{Mint32})}{v \cdot [astmt(vs, \text{nil})] \cdot \kappa_s \mid_\rho \xrightarrow{\text{mem(rmw } p \text{ Mint32 } v' \ rmwi)} \text{skip} \cdot \kappa_s \mid_\rho} \text{STEPATOMICFINISHNONE} \\
\\
\frac{\text{sem_atomic_statement } astmt \ (\text{vs} ++ v :: \text{nil}) = \text{Some } (p, rmwi) \quad \text{Val.has_type } v' \ (\text{type_of_chunk } \text{Mint32})}{v \cdot [(id : ty) = astmt(vs, \text{nil})] \cdot \kappa_s \mid_\rho \xrightarrow{\text{mem(rmw } p \text{ Mint32 } v' \ rmwi)} (id : ty) = v' \cdot \kappa_s \mid_\rho} \text{STEPATOMICFINISHSOME} \\
\\
\frac{}{\text{mfence} \cdot \kappa_s \mid_\rho \xrightarrow{\text{mem fence}} \text{skip} \cdot \kappa_s \mid_\rho} \text{STEPFENCE}
\end{array}$$

$$\begin{array}{c}
\frac{}{\text{continue} \cdot [-; s] \cdot \kappa_s \mid_\rho \longrightarrow \text{continue} \cdot \kappa_s \mid_\rho} \text{STEPCONTINUE} \\
\frac{}{\text{break} \cdot [-; s] \cdot \kappa_s \mid_\rho \longrightarrow \text{break} \cdot \kappa_s \mid_\rho} \text{STEPSBREAK} \\
\frac{}{\text{if } (e) \text{ then } s_1 \text{ else } s_2 \cdot \kappa_s \mid_\rho \longrightarrow e \cdot [\text{if } (\text{typeof } e) \text{ then } s_1 \text{ else } s_2] \cdot \kappa_s \mid_\rho} \text{STEPIFTHEELSE} \\
\frac{\text{is_true } v \text{ } ty}{v \cdot [\text{if } (\text{_}ty) \text{ then } s_1 \text{ else } s_2] \cdot \kappa_s \mid_\rho \longrightarrow s_1 \cdot \kappa_s \mid_\rho} \text{STEPIFTHEELSETRUE} \\
\frac{\text{is_false } v \text{ } ty}{v \cdot [\text{if } (\text{_}ty) \text{ then } s_1 \text{ else } s_2] \cdot \kappa_s \mid_\rho \longrightarrow s_2 \cdot \kappa_s \mid_\rho} \text{STEPIFTHEELSEFALSE} \\
\frac{}{\text{while } (e) \text{ do } s \cdot \kappa_s \mid_\rho \longrightarrow e \cdot [\text{while } (\text{_}e) \text{ do } s] \cdot \kappa_s \mid_\rho} \text{STEPWHILE} \\
\frac{\text{is_true } v \text{ } (\text{typeof } e)}{v \cdot [\text{while } (\text{_}e) \text{ do } s] \cdot \kappa_s \mid_\rho \longrightarrow s \cdot [\text{while } (e) \text{ do } s] \cdot \kappa_s \mid_\rho} \text{STEPWHILETRUE} \\
\frac{\text{is_false } v \text{ } (\text{typeof } e)}{v \cdot [\text{while } (\text{_}e) \text{ do } s] \cdot \kappa_s \mid_\rho \longrightarrow \text{skip} \cdot \kappa_s \mid_\rho} \text{STEPWHILEFALSE} \\
\frac{}{\text{continue} \cdot [\text{while } (e) \text{ do } s] \cdot \kappa_s \mid_\rho \longrightarrow \text{while } (e) \text{ do } s \cdot \kappa_s \mid_\rho} \text{STEPCONTINUEWHILE} \\
\frac{}{\text{break} \cdot [\text{while } (e) \text{ do } s] \cdot \kappa_s \mid_\rho \longrightarrow \text{skip} \cdot \kappa_s \mid_\rho} \text{STEPSBREAKWHILE} \\
\frac{}{\text{do } s \text{ while } (e) \cdot \kappa_s \mid_\rho \longrightarrow s \cdot [\text{do } s \text{ while } (e)] \cdot \kappa_s \mid_\rho} \text{STEPEDOWHILE} \\
\frac{\text{is_true } v \text{ } (\text{typeof } e)}{v \cdot [\text{do } s \text{ while } (\text{_}e)] \cdot \kappa_s \mid_\rho \longrightarrow \text{do } s \text{ while } (e) \cdot \kappa_s \mid_\rho} \text{STEPEDOWHILETRUE} \\
\frac{\text{is_false } v \text{ } (\text{typeof } e)}{v \cdot [\text{do } s \text{ while } (\text{_}e)] \cdot \kappa_s \mid_\rho \longrightarrow \text{skip} \cdot \kappa_s \mid_\rho} \text{STEPEDOWHILEFALSE} \\
\frac{}{\text{continue} \cdot [\text{do } s \text{ while } (e)] \cdot \kappa_s \mid_\rho \longrightarrow e \cdot [\text{do } s \text{ while } (\text{_}e)] \cdot \kappa_s \mid_\rho} \text{STEPEDOCONTINUEWHILE} \\
\frac{}{\text{break} \cdot [\text{do } s \text{ while } (e)] \cdot \kappa_s \mid_\rho \longrightarrow \text{skip} \cdot \kappa_s \mid_\rho} \text{STEPEDOBREAKWHILE} \\
\frac{}{\text{for } (s_1; e_2; s_3) s \cdot \kappa_s \mid_\rho \longrightarrow s_1 \cdot [\text{for } (; \diamond e_2; s_3) s] \cdot \kappa_s \mid_\rho} \text{STEPFORINIT} \\
\frac{}{\text{skip} \cdot [\text{for } (; \diamond e_2; s_3) s] \cdot \kappa_s \mid_\rho \longrightarrow e_2 \cdot [\text{for } (; \text{_}e_2; s_3) s] \cdot \kappa_s \mid_\rho} \text{STEPFORCOND} \\
\frac{\text{is_true } v \text{ } (\text{typeof } e_2)}{v \cdot [\text{for } (; \text{_}e_2; s_3) s] \cdot \kappa_s \mid_\rho \longrightarrow s \cdot [\text{for } (; e_2; \diamond s_3) s] \cdot \kappa_s \mid_\rho} \text{STEPFORTRUE} \\
\frac{\text{is_false } v \text{ } (\text{typeof } e_2)}{v \cdot [\text{for } (; \text{_}e_2; s_3) s] \cdot \kappa_s \mid_\rho \longrightarrow \text{skip} \cdot \kappa_s \mid_\rho} \text{STEPFORFALSE} \\
\frac{}{\text{skip} \cdot [\text{for } (; e_2; \diamond s_3) s] \cdot \kappa_s \mid_\rho \longrightarrow s_3 \cdot [\text{for } (; \diamond e_2; s_3) s] \cdot \kappa_s \mid_\rho} \text{STEPFORINCR} \\
\frac{}{\text{break} \cdot [\text{for } (; e_2; \diamond s_3) s] \cdot \kappa_s \mid_\rho \longrightarrow \text{skip} \cdot \kappa_s \mid_\rho} \text{STEPFORBREAK} \\
\frac{}{\text{continue} \cdot [\text{for } (; e_2; \diamond s_3) s] \cdot \kappa_s \mid_\rho \longrightarrow s_3 \cdot [\text{for } (; \diamond e_2; s_3) s] \cdot \kappa_s \mid_\rho} \text{STEPFORCONTINUE}
\end{array}$$

$$\begin{array}{c}
\text{call_cont } \kappa_s = (\text{Kcall None (Internal } \textcolor{blue}{fn}) \ \rho'' \ \kappa'_s) \\
\textcolor{blue}{fn}.\text{(fn_return)} = \text{Tvoid} \\
\textcolor{blue}{ps} = \text{sorted_pointers_of_env } \rho' \\
\hline
\text{return} \cdot \kappa_s \mid_{\rho'} \longrightarrow \text{skip} \cdot [\text{free } \textcolor{blue}{ps}; \text{return None}] \cdot \kappa_s \mid_{\rho'} \quad \text{STEPRETURNNONE} \\
\hline
\text{skip} \cdot [\text{free } \textcolor{blue}{p}::\textcolor{blue}{ps}; \text{return } \textcolor{blue}{opt_v}] \cdot \kappa_s \mid_{\rho} \xrightarrow{\text{mem (free } \textcolor{blue}{p} \text{ MObjStack)}} \text{skip} \cdot [\text{free } \textcolor{blue}{ps}; \text{return } \textcolor{blue}{opt_v}] \cdot \kappa_s \mid_{\rho} \quad \text{STEPRETURN} \\
\hline
\begin{array}{c}
\text{call_cont } \kappa_s = \kappa'_s \\
\text{get_fundef } \kappa'_s = \text{Some (Internal } \textcolor{blue}{fn}) \\
\textcolor{blue}{fn}.\text{(fn_return)} <> \text{Tvoid} \\
\hline
\text{return } \textcolor{blue}{e} \cdot \kappa_s \mid_{\rho'} \longrightarrow \textcolor{blue}{e} \cdot [\text{return } _] \cdot \kappa_s \mid_{\rho'} \quad \text{STEPRETURNSOME}
\end{array} \\
\hline
\frac{\textcolor{blue}{ps} = \text{sorted_pointers_of_env } \rho}{\textcolor{blue}{v} \cdot [\text{return } _] \cdot \kappa_s \mid_{\rho} \longrightarrow \text{skip} \cdot [\text{free } \textcolor{blue}{ps}; \text{return (Some } \textcolor{blue}{v})] \cdot \kappa_s \mid_{\rho}} \quad \text{STEPRETURNSOME1} \\
\hline
\frac{}{\text{switch } (\textcolor{blue}{e}) \textcolor{blue}{ls} \cdot \kappa_s \mid_{\rho} \longrightarrow \textcolor{blue}{e} \cdot [\text{switch } (_) \textcolor{blue}{ls}] \cdot \kappa_s \mid_{\rho}} \quad \text{STEPSWITCH} \\
\hline
\frac{\textcolor{blue}{s} = \text{seq_of_labeled_statement (select_switch } \textcolor{blue}{n} \textcolor{blue}{ls})}{\textcolor{blue}{n} \cdot [\text{switch } (_) \textcolor{blue}{ls}] \cdot \kappa_s \mid_{\rho} \longrightarrow \textcolor{blue}{s} \cdot [\text{switch } \kappa_s] \mid_{\rho}} \quad \text{STEPSELECTSWITCH} \\
\hline
\frac{}{\text{break} \cdot [\text{switch } \kappa_s] \mid_{\rho} \longrightarrow \text{skip} \cdot \kappa_s \mid_{\rho}} \quad \text{STEPBREAKSWITCH} \\
\hline
\frac{}{\text{continue} \cdot [\text{switch } \kappa_s] \mid_{\rho} \longrightarrow \text{continue} \cdot \kappa_s \mid_{\rho}} \quad \text{STEPCONTINUESWITCH} \\
\hline
\frac{}{\textcolor{blue}{l}::\textcolor{blue}{s} \cdot \kappa_s \mid_{\rho} \longrightarrow \textcolor{blue}{s} \cdot \kappa_s \mid_{\rho}} \quad \text{STEPLABEL} \\
\hline
\begin{array}{c}
\text{call_cont } \kappa_s = \kappa'_s \\
\text{get_fundef } \kappa'_s = (\text{Some (Internal } \textcolor{blue}{fn})) \\
\text{find_label } \textcolor{blue}{l} \textcolor{blue}{fn}.\text{(fn_body)} \ \kappa'_s = \text{Some } (\textcolor{blue}{s}', \ \kappa''_s) \\
\hline
\text{goto } \textcolor{blue}{l} \cdot \kappa_s \mid_{\rho} \longrightarrow \textcolor{blue}{s}' \cdot \kappa''_s \mid_{\rho} \quad \text{STEPGOTO}
\end{array} \\
\hline
\begin{array}{c}
\textcolor{blue}{args} = \textcolor{blue}{fn}.\text{(fn_params)} ++ \textcolor{blue}{fn}.\text{(fn_vars)} \\
\textcolor{blue}{fd} = \text{Internal } \textcolor{blue}{fn} \\
\hline
\textcolor{blue}{vs} \cdot [\textcolor{blue}{opt_lhs } \textcolor{blue}{fd}(_) \mid_{\rho}] \cdot \kappa_s \longrightarrow \text{alloc } (\textcolor{blue}{vs}, \textcolor{blue}{args}) \cdot [\textcolor{blue}{opt_lhs } \textcolor{blue}{fd}(_) \mid_{\rho}] \cdot \kappa_s \mid_{\rho_{\text{empty}}} \quad \text{STEPFUNCTIONINTERNAL}
\end{array} \\
\hline
\frac{\textcolor{blue}{n} = \text{Int.repr(sizeof } \textcolor{blue}{ty})}{\text{alloc } (\textcolor{blue}{vs}, \textcolor{blue}{id}^{\textcolor{blue}{ty}}::\textcolor{blue}{args}) \cdot \kappa_s \mid_{\rho} \xrightarrow{\text{mem (alloc } \textcolor{blue}{p} \textcolor{blue}{n} \text{ MObjStack)}} \text{alloc } (\textcolor{blue}{vs}, \textcolor{blue}{args}) \cdot \kappa_s \mid_{\rho \oplus (\textcolor{blue}{id} \mapsto \textcolor{blue}{p})}} \quad \text{STEPALLOCLOCAL} \\
\hline
\begin{array}{c}
\textcolor{blue}{args} = \textcolor{blue}{fn}.\text{(fn_params)} \\
\textcolor{blue}{fd} = (\text{Internal } \textcolor{blue}{fn}) \\
\hline
\text{alloc } (\textcolor{blue}{vs}, \text{nil}) \cdot [\textcolor{blue}{opt_lhs } \textcolor{blue}{fd}(_) \mid_{\rho'}] \cdot \kappa_s \mid_{\rho''} \longrightarrow \text{bind } (\textcolor{blue}{fn}, \textcolor{blue}{vs}, \textcolor{blue}{args}) \cdot [\textcolor{blue}{opt_lhs } \textcolor{blue}{fd}(_) \mid_{\rho'}] \cdot \kappa_s \mid_{\rho''} \quad \text{STEPBINDARGS}
\end{array} \\
\hline
\begin{array}{c}
\rho! \textcolor{blue}{id} = \text{Some } \textcolor{blue}{p} \\
\text{type_to_chunk } \textcolor{blue}{ty} = (\text{Some } \textcolor{blue}{c}) \\
\text{cast_value_to_chunk } \textcolor{blue}{c} \ \textcolor{blue}{v}_1 = \textcolor{blue}{v}_2 \\
\hline
\text{bind } (\textcolor{blue}{fn}, \textcolor{blue}{v}_1::\textcolor{blue}{vs}, \textcolor{blue}{id}^{\textcolor{blue}{ty}}::\textcolor{blue}{args}) \cdot \kappa_s \mid_{\rho} \xrightarrow{\text{mem (write } \textcolor{blue}{p} \textcolor{blue}{c} \textcolor{blue}{v}_2)}} \text{bind } (\textcolor{blue}{fn}, \textcolor{blue}{vs}, \textcolor{blue}{args}) \cdot \kappa_s \mid_{\rho} \quad \text{STEPBINDARGS}
\end{array} \\
\hline
\frac{\textcolor{blue}{s} = \textcolor{blue}{fn}.\text{(fn_body)}}{\text{bind } (\textcolor{blue}{fn}, \text{nil}, \text{nil}) \cdot \kappa_s \mid_{\rho} \longrightarrow \textcolor{blue}{s} \cdot \kappa_s \mid_{\rho}} \quad \text{STEPTRANSFERFUN} \\
\hline
\begin{array}{c}
\text{true (* event_match (external_function id targs ty) vs t vres -> *)} \\
\textcolor{blue}{fd} = \text{External } \textcolor{blue}{id} \ \textcolor{blue}{ty}^* \ \textcolor{blue}{ty} \\
\textcolor{blue}{vs} = \text{map val_of_eval } \textcolor{blue}{evs} \\
\hline
\textcolor{blue}{vs} \cdot [\textcolor{blue}{opt_lhs } \textcolor{blue}{fd}(_) \mid_{\rho}] \cdot \kappa_s \xrightarrow{\text{ext (call } \textcolor{blue}{id} \textcolor{blue}{evs})} \textcolor{blue}{opt_lhs} \ \text{ext}(_ \text{typ}) \cdot \kappa_s \mid_{\rho} \quad \text{STEPEXTERNALCALL}
\end{array}
\end{array}$$

$ \begin{aligned} &\text{Val.has_type } v \text{ } \textit{typ} \\ &\textit{fd} = \text{External } \textit{id} \text{ } \textit{ty}^* \text{ } \textit{ty} \\ &\textit{typ} = \text{match } (\text{opttyp_of_type } \textit{ty}) \text{ with } \text{Some } x \Rightarrow x \mid \text{None} \Rightarrow \text{Ast.Tint} \text{ end} \\ &\textit{v} = \text{val_of_eval } \textit{evl} \end{aligned} $	STEPEXTERNALRET
$ \frac{}{\textit{opt_lhs} \text{ ext } (_ \text{ } \textit{typ}) \cdot \kappa_s \mid_\rho \xrightarrow{\text{ext } (\text{return } \textit{typ } \textit{evl})} \textit{opt_lhs } v \cdot \kappa_s \mid_\rho} $	
$ \frac{ \begin{aligned} &\rho! \textit{id} = \text{Some } p \\ &\text{type_to_chunk } \textit{ty} = \text{Some } c \\ &\text{cast_value_to_chunk } c \text{ } v_1 = v_2 \end{aligned} }{ (\textit{id}:\textit{ty}) = v_1 \cdot \kappa_s \mid_\rho \xrightarrow{\text{mem } (\text{write } p \text{ } c \text{ } v_2)} \text{skip} \cdot \kappa_s \mid_\rho } $	STEPEXTERNALSTORESOMELOCAL
$ \frac{ \begin{aligned} &\rho! \textit{id} = \text{None} \\ &\text{Genv.find_symbol } \text{ge } \textit{id} = \text{Some } p \\ &\text{type_to_chunk } \textit{ty} = \text{Some } c \\ &\text{cast_value_to_chunk } c \text{ } v_1 = v_2 \end{aligned} }{ (\textit{id}:\textit{ty}) = v_1 \cdot \kappa_s \mid_\rho \xrightarrow{\text{mem } (\text{write } p \text{ } c \text{ } v_2)} \text{skip} \cdot \kappa_s \mid_\rho } $	STEPEXTERNALSTORESOMEGLOBAL
$ \frac{}{v \cdot \kappa_s \mid_\rho \longrightarrow \text{skip} \cdot \kappa_s \mid_\rho} $	STEPEXTERNALSTORENONE
$ \frac{}{\text{skip} \cdot [-; s_2] \cdot \kappa_s \mid_\rho \longrightarrow s_2 \cdot \kappa_s \mid_\rho} $	STEPSKIP
$ \frac{}{\text{skip} \cdot [\text{while } (e) \text{ do } s] \cdot \kappa_s \mid_\rho \longrightarrow \text{while } (e) \text{ do } s \cdot \kappa_s \mid_\rho} $	STEPWHILELOOP
$ \frac{}{\text{skip} \cdot [\text{do } s \text{ while } (e)] \cdot \kappa_s \mid_\rho \longrightarrow e \cdot [\text{do } s \text{ while } (-e)] \cdot \kappa_s \mid_\rho} $	STEPDOWHILELOOP
$ \frac{}{\text{skip} \cdot [\text{switch } \kappa_s] \mid_\rho \longrightarrow \text{skip} \cdot \kappa_s \mid_\rho} $	STEPSKIPSWITCH
$ \frac{ \text{call_cont } \kappa_s = [\textit{fd}(_) \mid_{\rho'}] \cdot \kappa'_s }{ \text{skip} \cdot [\text{free nil; return } \textit{opt_v}] \cdot \kappa_s \mid_{\rho''} \longrightarrow \text{skip} \cdot \kappa'_s \mid_{\rho'} } $	STEPRETURNNONEFINISH
$ \frac{ \begin{aligned} &\text{type_to_chunk } \textit{ty} = (\text{Some } c) \\ &\rho'! \textit{id} = \text{Some } p \\ &\text{call_cont } \kappa_s = [(\textit{id}:\textit{ty}) = \textit{fd}(_) \mid_{\rho'}] \cdot \kappa'_s \\ &\text{cast_value_to_chunk } c \text{ } v_1 = v_2 \end{aligned} }{ \text{skip} \cdot [\text{free nil; return } (\text{Some } v_1)] \cdot \kappa_s \mid_{\rho''} \xrightarrow{\text{mem } (\text{write } p \text{ } c \text{ } v_2)} \text{skip} \cdot \kappa'_s \mid_{\rho'} } $	STEPRETURN SOMEFINISHLOCAL
$ \frac{ \begin{aligned} &\text{type_to_chunk } \textit{ty} = (\text{Some } c) \\ &\rho'! \textit{id} = \text{None} \\ &\text{Genv.find_symbol } \text{ge } \textit{id} = \text{Some } p \\ &\text{call_cont } \kappa_s = [(\textit{id}:\textit{ty}) = \textit{fd}(_) \mid_{\rho'}] \cdot \kappa'_s \\ &\text{cast_value_to_chunk } c \text{ } v_1 = v_2 \end{aligned} }{ \text{skip} \cdot [\text{free nil; return } (\text{Some } v_1)] \cdot \kappa_s \mid_{\rho''} \xrightarrow{\text{mem } (\text{write } p \text{ } c \text{ } v_2)} \text{skip} \cdot \kappa'_s \mid_{\rho'} } $	STEPRETURN SOMEFINISHGLOBAL
$ \frac{}{\text{skip} \cdot \text{stop} \mid_\rho \xrightarrow{\text{exit}} \text{skip} \cdot \text{stop} \mid_\rho} $	STEPSTOP

Definition rules: 94 good 0 bad
Definition rule clauses: 178 good 0 bad