

lecture 9

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No, the family does not have a SDR.

The largest number of sets in the family with an SDR is 5.

With $n=6$, $k=1$, $\min_{i=1,2,\dots,6} |A_i| + 6 - 1 = 6$;

With $n=6$, $k=2$, $\min_{i_1, i_2=1,2,\dots,6} |A_{i_1} \cup A_{i_2}| + 6 - 2 = 6$

With $n=6$, $k=3$, $\min_{i_1, i_2, i_3=1,2,\dots,6} |A_{i_1} \cup A_{i_2} \cup A_{i_3}| + 6 - 3 = 6$

With $n=6$, $k=4$, $\min_{i_1, i_2, i_3, i_4=1,2,\dots,6} |A_{i_1} \cup A_{i_2} \cup A_{i_3} \cup A_{i_4}| + 6 - 4 = 5$

With $n=6$, $k=5$, $\min_{i_1, i_2, i_3, i_4, i_5=1,2,\dots,6} |A_{i_1} \cup A_{i_2} \cup A_{i_3} \cup A_{i_4} \cup A_{i_5}| + 6 - 5 = 5$

With $n=6$, $k=6$, $\min_{i_1, i_2, i_3, i_4, i_5, i_6=1,2,\dots,6} |A_{i_1} \cup A_{i_2} \cup A_{i_3} \cup A_{i_4} \cup A_{i_5} \cup A_{i_6}| + 6 - 6 = 5$

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We can convert this problem to a matching problem that one can sit on any seat but the latest one. So, we can choose any number from the set, but the set do not have the number i . So the number of the SDR is the permutations of the matching problem, the answer is D_n ;

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1. $A \rightarrow a, B \rightarrow a, C \rightarrow b, D \rightarrow d$, a reject B ;

2. $B \rightarrow d$, d reject D ;

3. $D \rightarrow b$, b reject C ;

4. $C \rightarrow a$, a reject A ;

5. $A \rightarrow b$, b reject A ;

6. $A \rightarrow c$;

$A \rightarrow c$, $B \rightarrow d$, $C \rightarrow a$, $D \rightarrow b$;

1. $a \rightarrow D$, $b \rightarrow B$, $c \rightarrow D$, $d \rightarrow C$, D reject a ;

2. $a \rightarrow C$, C reject d ;

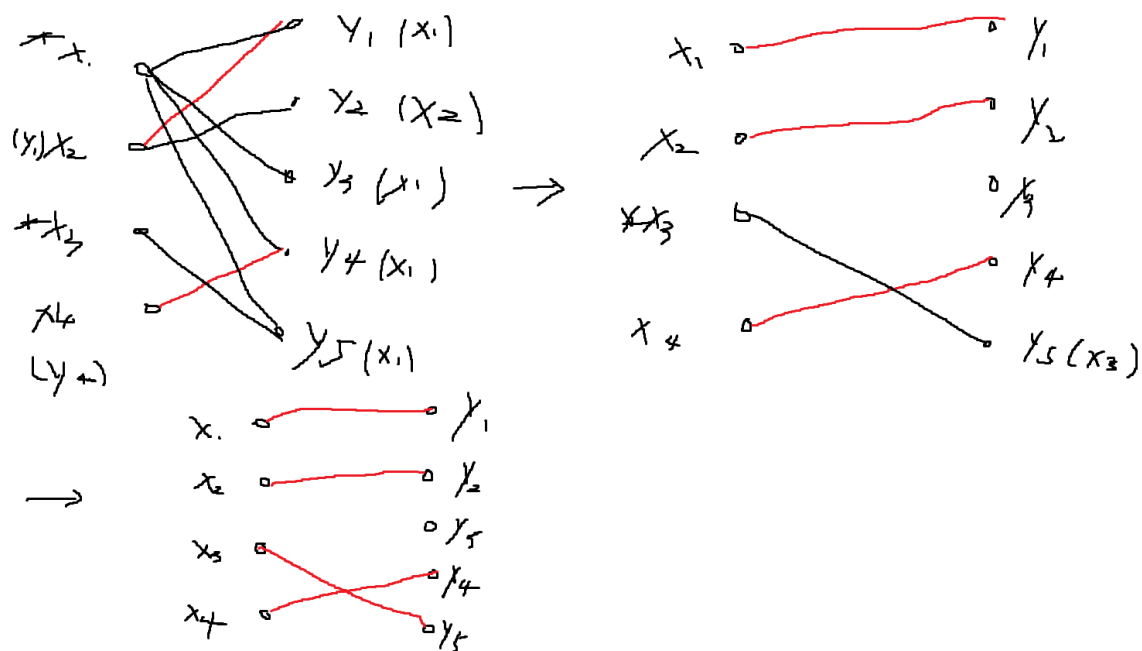
3. $d \rightarrow B$, B reject b ;

4. $b \rightarrow D$, D reject c ;

5. $c \rightarrow A$;

$A \rightarrow c$, $B \rightarrow d$, $C \rightarrow a$, $D \rightarrow b$;

CONT'D



Here, we get the max-matching is $\{(x_1, y_1), (x_2, y_2), (x_3, y_5), (x_4, y_4)\}$;

min-cover $\{x_1, x_2, x_4, y_5\}$;

We get a max-matching $M = \{(x_3, y_5), (x_1, y_1), (x_2, y_2), (x_4, y_4)\}$, but we can find the vertex y_3 is still uncovered. And we should construct a subgraph composed of edges incident to the vertex y_3 . Then we find the max-matching of the subgraph, and add it to the max-matching of M , then we get a minimum edge cover. The max-matching of the subgraph is $\{(x_4, y_3)\}$, so we get the minimum edge cover for the graph is $M = \{(x_1, y_1), (x_2, y_2), (x_3, y_5), (x_4, y_3), (x_4, y_4)\}$

