

188. Best Time to Buy and Sell Stock IV

Hard 2157 135 Add to List Share

You are given an integer array `prices` where `prices[i]` is the price of a given stock on the i^{th} day.

Design an algorithm to find the maximum profit. You may complete at most k transactions.

Notice that you may not engage in multiple transactions simultaneously (i.e., you must sell the stock before you buy again).

Example 1:

Input: $k = 2$, `prices = [2,4,1]`

Output: 2

Explanation: Buy on day 1 (price = 2) and sell on day 2 (price = 4), profit = 4-2 = 2.

Example 2:

Input: $k = 2$, `prices = [3,2,6,5,0,3]`

Output: 7

Explanation: Buy on day 2 (price = 2) and sell on day 3 (price = 6), profit = 6-2 = 4. Then buy on day 5 (price = 0) and sell on day 6 (price = 3), profit = 3-0 = 3.

Constraints:

- $0 \leq k \leq 100$
- $0 \leq \text{prices.length} \leq 1000$
- $0 \leq \text{prices}[i] \leq 1000$

Notation:

`prices`: the stock price array with length n

i : the i^{th} day (i will go from 0 to $n-1$)

k : the maximum number of transactions allowed to complete

$T[i][k]$: the maximum profit that could be gained at the end of i^{th} day with at most k transactions.

Base case:

$$T[-1][k] = T[i][0] = 0$$

no stock or no transaction yield no profit.

Actions:

buy: there should be 0 stock held in our hand.

sell: there should be 1 stock held in our hand.

rest

Definition:

$$T[i][k][0]:$$

the maximum profit at the end of i^{th} day with at most k transactions and with 0 stock in our hand after taking the action.

$$T[i][k][1]:$$

the maximum profit at the end of i^{th} day with at most k transactions and with 1 stock in our hand after taking the action.

Base case:

$$T[-1][k][0] = 0$$

Profit is 0 if there is no stock available.

$$T[-1][k][1] = -\text{Infinity}$$

It's impossible to have 1 stock in hand if there is no stock.

$$T[i][0][0] = 0$$

Profit is 0 if there is no transaction are allowed.

$$T[i][0][1] = -\text{Infinity}$$

It's impossible to have 1 stock in hand if there is no transaction.

Recurrence relations:

$$T[i][k][0] = \max(\text{rest}, \text{sell}) = \max(T[i-1][k][0], T[i-1][k][1] + \text{price}[i])$$

$$T[i][k][1] = \max(\text{rest}, \text{buy}) = \max(T[i-1][k][1], T[i-1][k-1][0] - \text{price}[i])$$


```

1 class Solution {
2 public:
3     int maxProfit(int k, vector<int>& prices) {
4         if (k <= 0) {
5             return 0;
6         }
7
8         // base case
9         vector<int> t_i_k_0(k, 0);
10        vector<int> t_i_k_1(k, INT_MIN); // T[-1][1][1]
11
12        // recurrence
13        for (int price : prices) {
14            vector<int> t_i_k_0_temp(t_i_k_0);
15            // T[i][0][0] = max(T[i-1][0][0], T[i-1][0][1] + prices[i])
16            t_i_k_0[0] = max(t_i_k_0_temp[0], t_i_k_1[0] + price);
17            // T[i][0][1] = max(T[i-1][0][1], -prices[i])
18            t_i_k_1[0] = max(t_i_k_1[0], -price);
19            for (int i = 1; i < k; i++) {
20                // T[i][k][0] = max(T[i-1][k][0], T[i-1][k][1] + prices[i])
21                t_i_k_0[i] = max(t_i_k_0_temp[i], t_i_k_1[i] + price);
22                // T[i][k][1] = max(T[i-1][k][1], T[i-1][k-1][0] - prices[i])
23                t_i_k_1[i] = max(t_i_k_1[i], t_i_k_0_temp[i - 1] - price);
24            }
25        }
26
27        return *max_element(t_i_k_0.begin(), t_i_k_0.end());
28    }
29 };

```