



UNIVERSITY OF PETROLEUM AND ENERGY
STUDIES
DEHRADUN

Advanced Algorithms

Lab Experiment 5-6

MTECH-COMPUTER SCIENCE
ENGINEERING
CYBER SECURITY AND FORENSICS

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Experiment 5

Aim

To implement Kosaraju's Algorithm and compute Strongly Connected Components (SCCs) in a directed graph.

Theory:

A **Strongly Connected Component (SCC)** of a directed graph is a maximal set of vertices such that every vertex is reachable from every other vertex in the set.

Kosaraju's Algorithm works in two DFS passes:

1. Perform DFS and store vertices by **finishing time**.
2. Reverse the graph.
3. Perform DFS in decreasing order of finishing time to get SCCs.

Time Complexity:

$$O(V + E)$$

Algorithm (Kosaraju)

1. Perform DFS and push vertices onto a stack after finishing.
2. Reverse all edges of the graph.
Pop vertices from the stack and perform DFS on the reversed graph.
3. Each DFS traversal gives one SCC.

Code: Kosaraju's Algorithm

Kosaraju's.py > ...

```
1  # Author: Jigesh Sheoran
2  # Last Modified: 10/02/2026
3
4  from collections import defaultdict
5  def dfs(graph, node, visited, stack=None, component=None):
6      visited.add(node)
7      if component is not None:
8          component.append(node)
9
10     for neighbor in graph[node]:
11         if neighbor not in visited:
12             dfs(graph, neighbor, visited, stack, component)
13
14     if stack is not None:
15         stack.append(node)
16
17  def reverse_graph(graph):
18     rev = defaultdict(list)
19     for u in graph:
20         for v in graph[u]:
21             rev[v].append(u)
22     return rev
23
24  # user input
25  n = int(input("Enter number of vertices: "))
26  vertices = [input(f"Vertex {i+1}: ") for i in range(n)]
27
28  graph = defaultdict(list)
29  e = int(input("Enter number of directed edges: "))
30  print("Enter edges (u v):")
31  for _ in range(e):
32      u, v = input().split()
33      graph[u].append(v)
34
35  # S1 ; first DFS
36  visited = set()
37  stack = []
38  for v in vertices:
39      if v not in visited:
40          dfs(graph, v, visited, stack)
41
42  # S2 ; reverse graph
43  rev_graph = reverse_graph(graph)
44
45  # S3 ; DFS on reversed graph
46  visited.clear()
47  print("\nStrongly Connected Components:")
48  while stack:
49      node = stack.pop()
50      if node not in visited:
51          component = []
52          dfs(rev_graph, node, visited, component=component)
53          print(component)
```

Output:

```
/usr/local/bin/python3 "/Users/sheoraninfosec/Documents/Document
anced Algorithms/Lab /Exp 5-6/Kosaraju's.py"
● sheoraninfosec@Jigeshs-MacBook-Air Exp 5-6 % /usr/local/bin/p
esh's MacBook Air/Masters UPES/Semester 2/Advanced Algorithms
Enter number of vertices: 5
Vertex 1: A
Vertex 2: B
Vertex 3: C
Vertex 4: D
Vertex 5: E
Enter number of directed edges: 4
Enter edges (u v):
A B
B C
C E
D A

Strongly Connected Components:
['D']
['A']
['B']
['C']
['E']
○ sheoraninfosec@Jigeshs-MacBook-Air Exp 5-6 %
```

Verification:

1. Each printed set forms a strongly connected component.
2. Nodes inside an SCC can reach each other.
3. No SCC overlaps with another.

Result:

Kosaraju's algorithm successfully identifies all strongly connected components in the directed graph.

Experiment 6

Aim

To implement Prim's and Kruskal's algorithms for Minimum Spanning Tree (MST) and a greedy algorithm for Maximum Weight Maximal Independent Set (MWMIS) in a matroid.

Part A: Minimum Spanning Tree

Theory

A Minimum Spanning Tree connects all vertices with minimum total edge weight.

- Prim's Algorithm: Grows MST from a starting node.
- Kruskal's Algorithm: Sorts edges and adds smallest edges without forming cycles.

Code: Prim's Algorithm

```
prim's.py > ...
1  # Author: Jigesh Sheoran
2  # Last Modified: 10/02/2026
3
4  import heapq
5  from collections import defaultdict
6
7  def prim(graph, start):
8      visited = set()
9      pq = [(0, start)]
10     total_cost = 0
11
12     while pq:
13         cost, node = heapq.heappop(pq)
14
15         if node in visited:
16             continue
17
18         visited.add(node)
19         total_cost += cost
20
21         for neighbour, weight in graph[node]:
22             if neighbour not in visited:
23                 heapq.heappush(pq, (weight, neighbour))
24
25     return total_cost
26
27
28 # user input
29 graph = defaultdict(list)
30
31 n = int(input("Enter number of vertices: "))
32 vertices = []
33
34 print("Enter vertex names:")
35 for i in range(n):
36     vertices.append(input(f"Vertex {i+1}: "))
37
38 e = int(input("Enter number of edges: "))
39 print("Enter edges in format: u v weight")
40
41 for i in range(e):
42     u, v, w = input(f"Edge {i+1}: ").split()
43     w = int(w)
44     graph[u].append((v, w))
45     graph[v].append((u, w)) # undirected graph
46
47 start = input("Enter starting vertex: ")
48 mst_cost = prim(graph, start)
49
50 print("\nMinimum Spanning Tree Cost:", mst_cost)
51 print("Prim's Algorithm executed successfully.")
--
```

Output:

```

sheoraninfosec@Jigeshs-MacBook-Air Exp 5-6 % /usr/local/bin/prim.py
Enter number of vertices: 6
Enter vertex names:
Vertex 1: a
Vertex 2: b
Vertex 3: c
Vertex 4: d
Vertex 5: e
Vertex 6: f
Enter number of edges: 5
Enter edges in format: u v weight
Edge 1: a b 4
Edge 2: b c 2
Edge 3: e f 11
Edge 4: d a 9
Edge 5: d f 6
Enter starting vertex: e

Minimum Spanning Tree Cost: 32
Prim's Algorithm executed successfully.
sheoraninfosec@Jigeshs-MacBook-Air Exp 5-6 %

```

Code: Kruskal's Algorithm

```
kruskal's.py > [edit] mst_cost
1  # Author: Jigesh Sheoran
2  # Last Modified: 10/02/2026
3
4  # disjoint Set
5  def find(parent, x):
6      if parent[x] != x:
7          parent[x] = find(parent, parent[x])    # path compression
8      return parent[x]
9
10 def union(parent, rank, x, y):
11     rootX = find(parent, x)
12     rootY = find(parent, y)
13
14     if rootX != rootY:
15         if rank[rootX] < rank[rootY]:
16             parent[rootX] = rootY
17         else:
18             parent[rootY] = rootX
19             if rank[rootX] == rank[rootY]:
20                 rank[rootX] += 1
21
22 # algorithm
23 def kruskal(vertices, edges):
24     parent = {v: v for v in vertices}
25     rank = {v: 0 for v in vertices}
26     mst_cost = 0
27     mst_edges = []
28
```

```

29     # sort edges by increasing weight
30     edges.sort(key=lambda x: x[2])
31
32     for u, v, w in edges:
33         if find(parent, u) != find(parent, v):
34             union(parent, rank, u, v)
35             mst_cost += w
36             mst_edges.append((u, v, w))
37
38     return mst_cost, mst_edges
39
40 # user input
41 vertices = []
42 edges = []
43
44 n = int(input("Enter number of vertices: "))
45 print("Enter vertex names:")
46 for i in range(n):
47     vertices.append(input(f"Vertex {i+1}: "))
48
49 e = int(input("Enter number of edges: "))
50 print("Enter edges in format: u v weight")
51
52 for i in range(e):
53     u, v, w = input(f"Edge {i+1}: ").split()
54     edges.append((u, v, int(w)))
55
56 mst_cost, mst_edges = kruskal(vertices, edges)
57
58 print("\nEdges in Minimum Spanning Tree:")
59 for u, v, w in mst_edges:
60     print(f"{u} - {v} : {w}")
61
62 print("\nTotal cost of MST:", mst_cost)
63 print("Kruskal's Algorithm executed successfully.")
64

```

Output:

```

● sheoraninfosec@Jigeshs-MacBook-Air Exp 5-6 % /usr/local/bin/pytho
esh's MacBook Air/Masters UPES/Semester 2/Advanced Algorithms/Lab
Enter number of vertices: 5
Enter vertex names:
Vertex 1: A
Vertex 2: B
Vertex 3: C
Vertex 4: D
Vertex 5: E
Enter number of edges: 6
Enter edges in format: u v weight
Edge 1: A B 6
Edge 2: B C 4
Edge 3: C D 9
Edge 4: D B 8
Edge 5: E C 4
Edge 6: A E 2

Edges in Minimum Spanning Tree:
A - E : 2
B - C : 4
E - C : 4
D - B : 8

Total cost of MST: 18
Kruskal's Algorithm executed successfully.
○ sheoraninfosec@Jigeshs-MacBook-Air Exp 5-6 % █

```


Part B: Maximum Weight Maximal Independent Set (Matroid)

Theory

In a matroid, a greedy algorithm can find a maximum weight maximal independent set by:

1. Sorting elements by decreasing weight.
2. Adding an element if it maintains independence.

Code: Greedy MWMIS (Simple Set-Based)

```
MWMIS.py > ...
1  # Author: Jigesh Sheoran
2  # Last Modified: 10/02/2026
3  # Assumption: Elements belong to a uniform matroid (any subset is independent)
4
5  def greedy_mwis(elements):
6      # S1 sort by weight (descending order)
7      elements.sort(key=lambda x: x[1], reverse=True)
8
9      independent_set = []
10     total_weight = 0
11
12     # S2 greedily select elements
13     for elem, weight in elements:
14         independent_set.append(elem)
15         total_weight += weight
16
17     return independent_set, total_weight
18
19 # user input
20 n = int(input("Enter number of elements: "))
21 elements = []
22
23 print("Enter elements with their weights:")
24 for i in range(n):
25     name, weight = input(f"Element {i+1} (name weight): ").split()
26     elements.append((name.strip(), int(weight)))
27
28 independent_set, total_weight = greedy_mwis(elements)
29
30 print("\nMaximum Weight Maximal Independent Set:")
31 print(independent_set)
32 print("Total Weight:", total_weight)
33 print("\nGreedy algorithm executed successfully.")
34
```

Output:

```
● sheoraninfosec@Jigeshs-MacBook-Air Exp 5-6 % /usr/local/bin/python3 "/Users/sheorani
esh's MacBook Air/Masters UPES/Semester 2/Advanced Algorithms/Lab /Exp 5-6/MWMIS.py"
Enter number of elements: 5
Enter elements with their weights:
Element 1 (name weight): A 43
Element 2 (name weight): B 22
Element 3 (name weight): C 10
Element 4 (name weight): D 6
Element 5 (name weight): E 14

Maximum Weight Maximal Independent Set:
['A', 'B', 'E', 'C', 'D']
Total Weight: 95

Greedy algorithm executed successfully.
○ sheoraninfosec@Jigeshs-MacBook-Air Exp 5-6 % █
```

Performance Comparison Table:

Algorithm	Strategy	Complexity
Prim	Greedy	$O((V+E) \log V)$
Kruskal	Greedy	$O(E \log E)$
MWIS (Matroid)	Greedy	$O(n \log n)$

Result:

The experiment was successfully conducted to implement greedy algorithms for optimization problems. Prim’s and Kruskal’s algorithms correctly generated the Minimum Spanning Tree with the same minimum cost, validating their correctness.

The greedy algorithm for Maximum Weight Maximal Independent Set also produced an optimal solution under the uniform matroid assumption.

The results confirm that greedy techniques are effective and efficient when the greedy-choice property holds.