

Review of Chapters 3-6

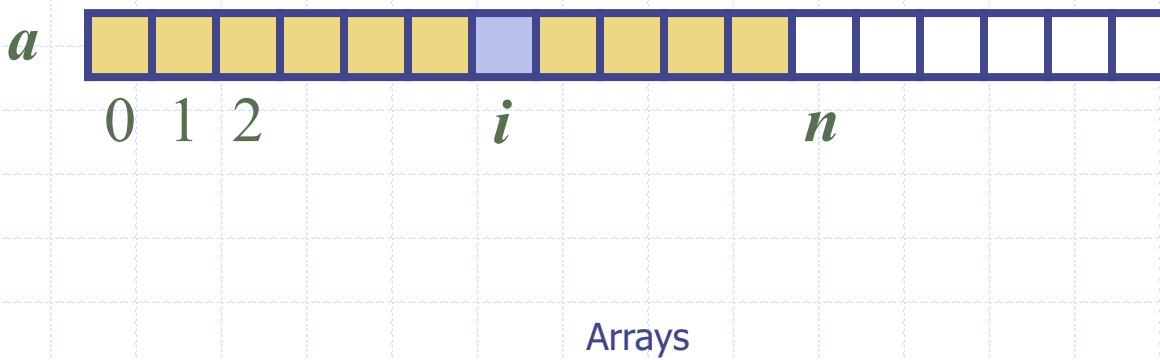
Feb. 14, 2018

Chapter 3 Fundamental Data Structures

- Arrays
- Linked lists
 - Singly linked lists
 - Doubly linked lists
 - Circularly Linked List

Arrays

- What is array?
- Capacity of array
- Array operations (access, add, delete elements) and efficiency.
- When to use array?



Sorting an Array

□ The Insertion-Sort Algorithm

Algorithm InsertionSort(A):

Input: An array A of n comparable elements

Output: The array A with elements rearranged in nondecreasing order

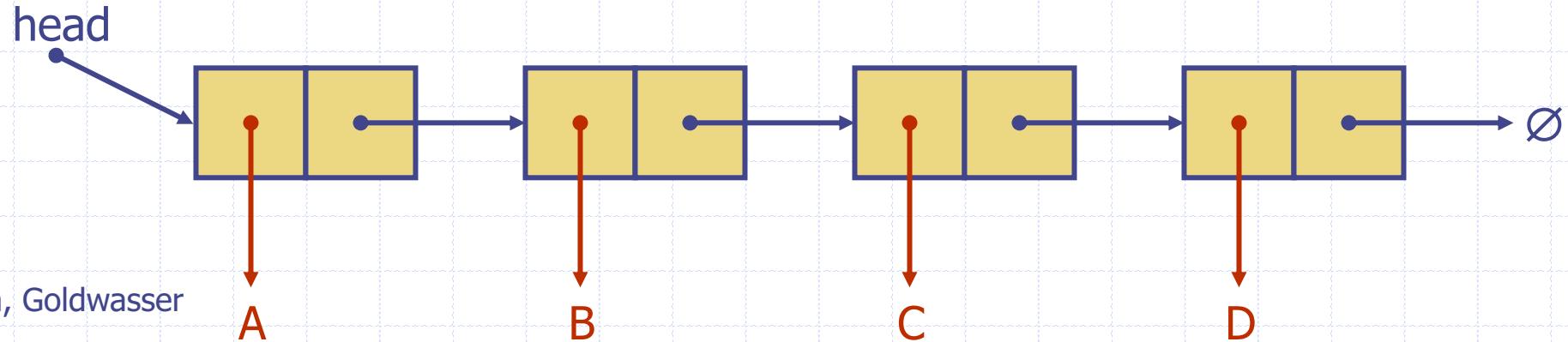
for k from 1 to $n - 1$ **do**

 Insert $A[k]$ at its proper location within $A[0], A[1], \dots, A[k]$.

Code Fragment 3.5: High-level description of the insertion-sort algorithm.

Singly Linked Lists

- What is singly linked list?
- Capacity of a singly linked list
- Linked list operations and efficiency.
 - Find an element in the list
 - Inserting/removing at the Head / tail
- When to use singly linked list?



Other Linked Lists

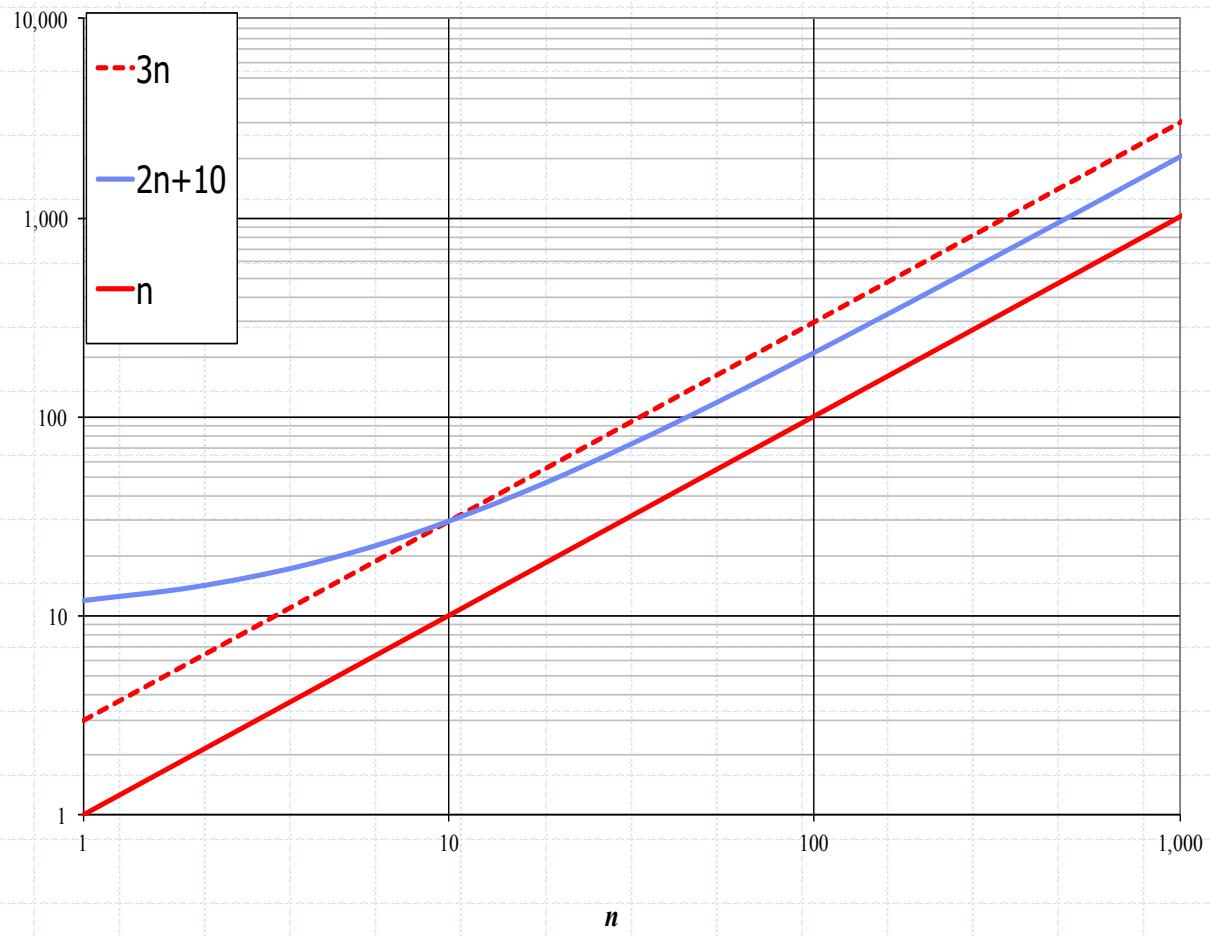
- ❑ What is doubly linked list ?
- ❑ Circularly Linked List

CS 2233-01 Data Structures and Algorithms

Chapter 4 Algorithm Analysis

Asymptotic Notation (Big-Oh Notation)

- **Algorithm analysis:** determine how running time increases as the *size* of the problem increases.
- **Big Oh-notation:** Given functions $f(n)$ and $g(n)$, we say that $f(n)$ is $O(g(n))$ if there are positive constants c and n_0 such that
$$f(n) \leq cg(n) \text{ for } n \geq n_0$$



Basic asymptotic efficiency classes

1	constant
$\log n$	logarithmic
n	linear
$n \log n$	$n\text{-log-}n$
n^2	quadratic
n^3	cubic
2^n	exponential
$n!$	factorial

Big-Oh Rules

- If $f(n)$ is a polynomial of degree d , then $f(n)$ is $O(n^d)$, i.e.,
 1. Drop lower-order terms
 2. Drop constant factors, e.g. $n^4 + 100n^2 + 10n + 50$
- Use the smallest possible class of functions
 - Say “ $2n$ is $O(n)$ ” instead of “ $2n$ is $O(n^2)$ ”
- Use the simplest expression of the class
 - Say “ $3n + 5$ is $O(n)$ ” instead of “ $3n + 5$ is $O(3n)$ ”

Relatives of Big-Oh

big-Omega

- $f(n)$ is $\Omega(g(n))$ if there is a constant $c > 0$ and an integer constant $n_0 \geq 1$ such that

$$f(n) \geq c g(n) \text{ for } n \geq n_0$$

big-Theta

- $f(n)$ is $\Theta(g(n))$ if there are constants $c_1 > 0$ and $c_2 > 0$ and an integer constant $n_0 \geq 1$ such that

$$c_1 g(n) \leq f(n) \leq c_2 g(n) \text{ for } n \geq n_0$$

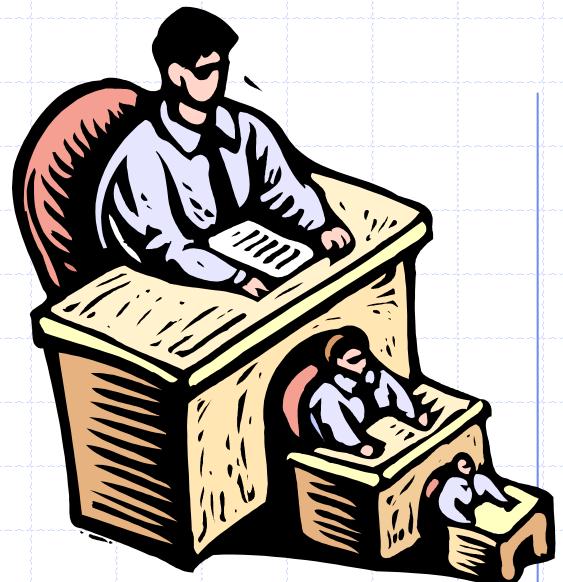
Questions in Mid-term exam

- Compare different efficiency classes
- Estimate running time of a piece of code (in big-Oh term)
- Choose data structures/algorithms to solve a problem, under constraints of time or space

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Data Structures and Algorithms

Chapter 5 Recursion



Content of a Recursive Method

- Base case(s)

- Values of the input variables for which we perform no recursive calls are called **base cases** (there should be at least one base case).
- Every possible chain of recursive calls **must** eventually reach a base case.

- Recursive calls

- Calls to the current method.
- Each recursive call should be defined so that it makes progress towards a base case.

A Classic Recursive Algorithm: Binary Search

- ❑ A Classic Recursive Algorithm: *Binary search*
- ❑ What is **Linear Recursion** (Binary Recursion, Multiple Recursion)?
- ❑ What is **Tail Recursion**?
- ❑ Convert Tail recursion to non-recursive methods?

Chapter 6 Stacks, Queues, and Deques

- ❑ Presentation for use with the textbook Data Structures and Algorithms in Java, 6th edition, by M. T. Goodrich, R. Tamassia, and M. H. Goldwasser, Wiley, 2014

Stacks

- Definition & ADT
- Array-based stacks
- Performance and Limitations of Array-based Stacks
- Linked list-based stacks
- Applications of stacks



```
public interface Stack<E> {  
    int size();  
    boolean isEmpty();  
    E top();  
    void push(E element);  
    E pop();  
}
```

Applications of stacks

- Text parsing
- Parentheses matching
- Evaluating Arithmetic Expressions
- Evaluate postfix notation

Queues

- Definition & ADT
- Queues V.S. Stacks
- A Circular Queue

```
public interface Queue<E> {  
    int size();  
    boolean isEmpty();  
    E first();  
    void enqueue(E e);  
    E dequeue();  
}
```

Deque (Double-Ended Queues) ADT

addFirst(*e*): Insert a new element *e* at the front of the deque.

addLast(*e*): Insert a new element *e* at the back of the deque.

removeFirst(): Remove and return the first element of the deque
(or `null` if the deque is empty).

removeLast(): Remove and return the last element of the deque
(or `null` if the deque is empty).

Deque ADT (cont.)

`first()`: Returns the first element of the deque, without removing it (or null if the deque is empty).

`last()`: Returns the last element of the deque, without removing it (or null if the deque is empty).

`size()`: Returns the number of elements in the deque.

`isEmpty()`: Returns a boolean indicating whether the deque is empty.

Implementation

- ❑ Circular Array
- ❑ Doubly Linked List
- ❑ Performance of the Deque Operations

Method	Running Time
size, isEmpty	$O(1)$
first, last	$O(1)$
addFirst, addLast	$O(1)$
removeFirst, removeLast	$O(1)$

Midterm will be on Friday, Feb. 16.