Topic V09

Conditional Instructions and Loops

Reading: (Section 2.7)

Conditional Operations (recap)

Branch to a labeled instruction if a condition is true Otherwise, continue sequentially

```
beq rs1, rs2, L1
if (rs1 == rs2) branch to instruction labeled L1
```

bne rs1, rs2, L2 if (rs1 != rs2) branch to instruction labeled L2

jal zero, L3 unconditional jump to instruction labeled L3

While Loops

How can we improve this code?

```
C code:
while (save[i] == k)
    i = i + j;
...
```

RISC-V code:

```
Loop: slli t1, s3, 2
add t1, t1, s6
lw t0, 0(t1)
bne t0, s5, DoneLoop
add s3, s3, s4
jal zero, Loop

DoneLoop:
```

Assumption

```
i \leftrightarrow s3

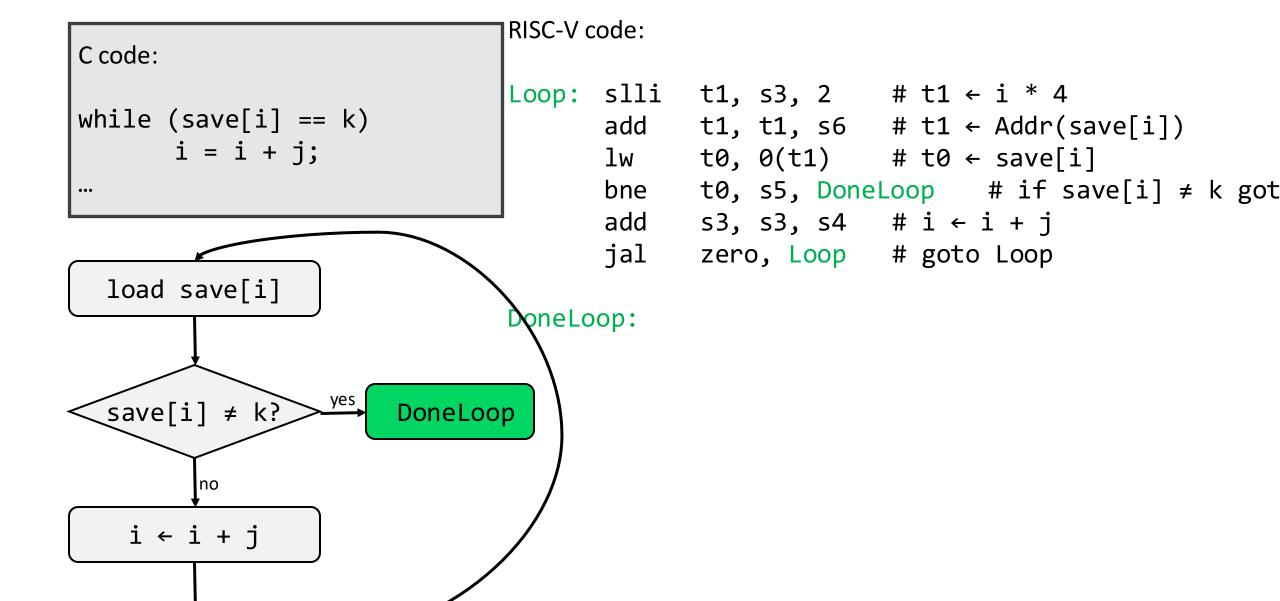
j \leftrightarrow s4

k \leftrightarrow s5

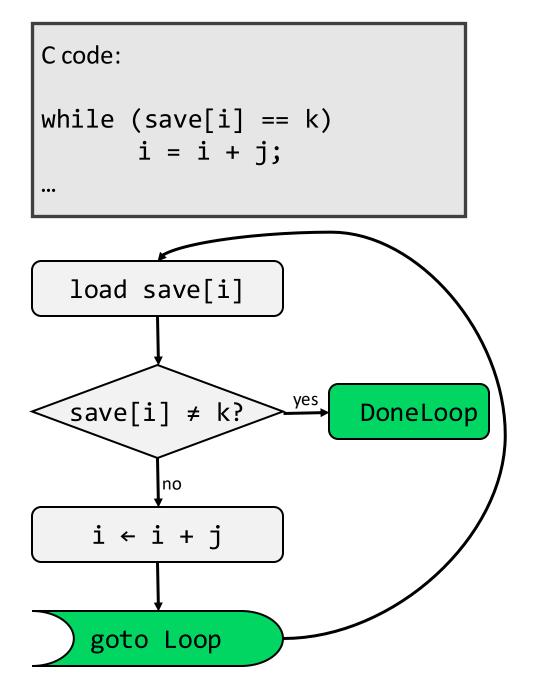
base of save[] \leftrightarrow s6
```

save is an array of 4-byte integers

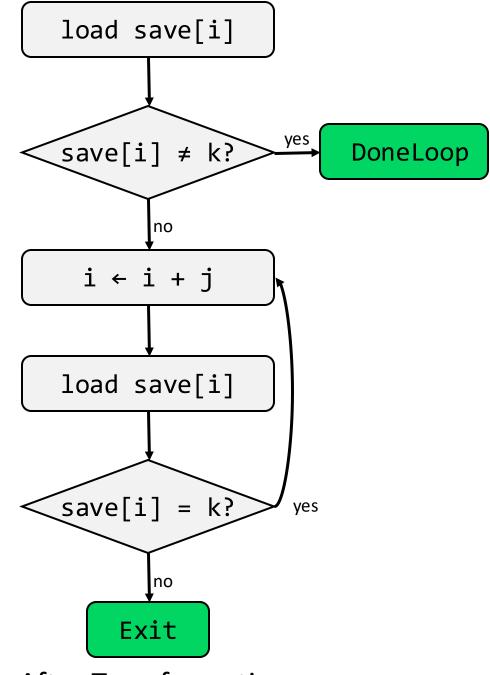
```
# t1 ← i * 4
# t1 ← Addr(save[i])
# t0 ← save[i]
# if save[i] ≠ k goto DoneLoop
# i ← i + j
# goto Loop
```



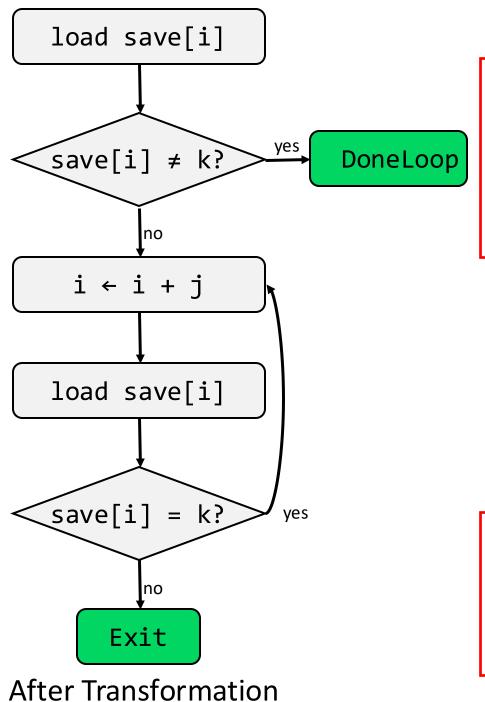
goto Loop



Before Transformation



After Transformation



Before Transformation:

```
Loop: slli t1, s3, 2
add t1, t1, s6
lw t0, 0(t1)
bne t0, s5, DoneLoop
add s3, s3, s4
jal zero, Loop
```

slli t1, s3, 2

```
# t1 ← i * 4
# t1 ← Addr(save[i])
# t0 ← save[i]
# if save[i] ≠ k goto
# i ← i + j
# goto Loop
```

DoneLoop:

After Transformation:

```
add t1, t1, s6
lw t0, 0(t1)
bne t0, s5, DoneLoop

Loop: add s3, s3, s4
slli t1, s3, 2
add t1, t1, s6
lw t0, 0(t1)
beq t0, s5, Loop
```

t1 ← Addr(save[i])
t0 ← save[i]
if save[i] ≠ k goto

 $# t1 \leftarrow i * 4$

i ← i + j
t1 ← i * 4
t1 ← Addr(save[i])

t0 ← save[i]
if save[i] ≠ k goto

DoneLoop:

Non-Branch Comparison

Set result to 1 if a condition is true Otherwise, set to 0

```
Set less than slt rd, rs1, rs2
```

Set less than immediate slti rd, rs, immediate

```
if (rs1 < rs2)
rd ← 1;
else
rd ← 0;
```

```
if (rs1 < imm)
rd ← 1;
else
rd ← 0;
```

```
int signum(int x)
{
  if(x > 0) return 1;
  if(x < 0) return -1;
  return 0;
}</pre>
```

Example: slt avoids branches

```
RISC-V Assembly: signum: slt t1, a0, zero #t1 \leftarrow 1 if x < 0 slt a0, zero, a0 # a0 \leftarrow 1 is x > 0 sub a0, a0, t1 # 1 is x>0; 0 if x==0, -1 if x<0 ret
```

Other Branches

Branch on less than blt rs1, rs2, L1

```
if (rs1 < rs2)
        PC ← PC + {imm,1b'0}
else
        PC ← PC + 4;</pre>
```

Branch on greater than or equal bge rs1, rs2, L1

```
if (rs1 ≥ rs2)
          PC ← PC + {imm,1b'0}
else
          PC ← PC + 4;
```

What about the other types of comparisons; why are there only instructions for **blt** and bge?

Other Branches

Suppose you have two registers, s1 and s2, that you wish to compare:

Branch on greater than or equal bge s2, s1, L1

How do you branch on s1 less than or equal s2?

Branch on less than blt s2, s1, L1

How do you branch on s1 greater than s2?

Other Branches (Pseudo Instructions)

```
Branch on equal zero

beqz s1, L1

Equivalent to:

beq s1, zero, L1
```

```
if (rs1 = 0)
        PC ← PC + {imm,1b'0}
else
        PC ← PC + 4;
```

```
Branch on not equal zero
bnez s1, L1
Equivalent to:
bne s1, zero, L1
```

```
if (rs1 ≠ 0)
        PC ← PC + {imm,1b'0}
else
        PC ← PC + 4;
```

Signed vs. Unsigned

Signed comparison: slt, slti

Unsigned comparison: sltu, sltiu

```
slt t0, s0, s1 #signed
t0←1 because -1 < +1
```

```
sltu t0, s0, s1 #unsigned
```

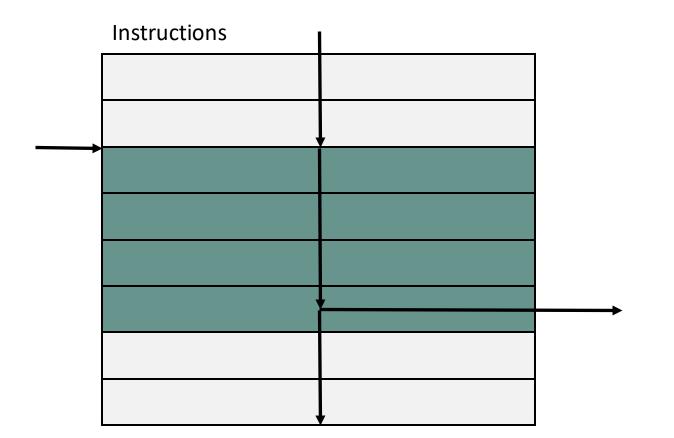
```
t0←0 because +4,294,967,295 > +1
```

Basic Blocks

A basic block is a sequence of instructions with:

No branch targets (except at beginning)

No embedded branches (except at end)



A compiler identifies basic blocks for optimization

An advanced processor can accelerate execution of basic blocks

t1, s3, 2 slt slli t1, t1, s6 add t0, 0(t1) blt lw sltu t0, s5, Exit bne slli Loop: t1, s3, 2 Loop: add s3, s3, s4 bge t1, t1, s6 add t1, s3, 2 slli t0, 0(t1) lw beq add t1, t1, s6 t0, s5, Exit bne t0, 0(t1) lw s3, s3, s4 add bnez t0, s5, Loop beq jal zero, Loop Exit: Exit: beqz

Recap

