

Topic V15

Array Indexing vs Pointers

Reading: (Section 2.14)

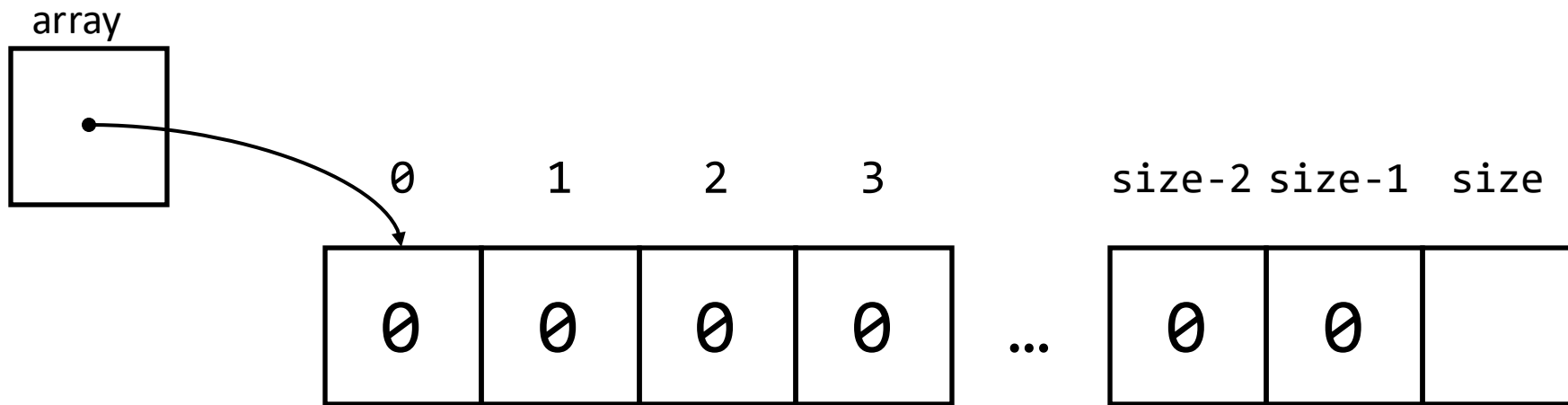
Example of a Task

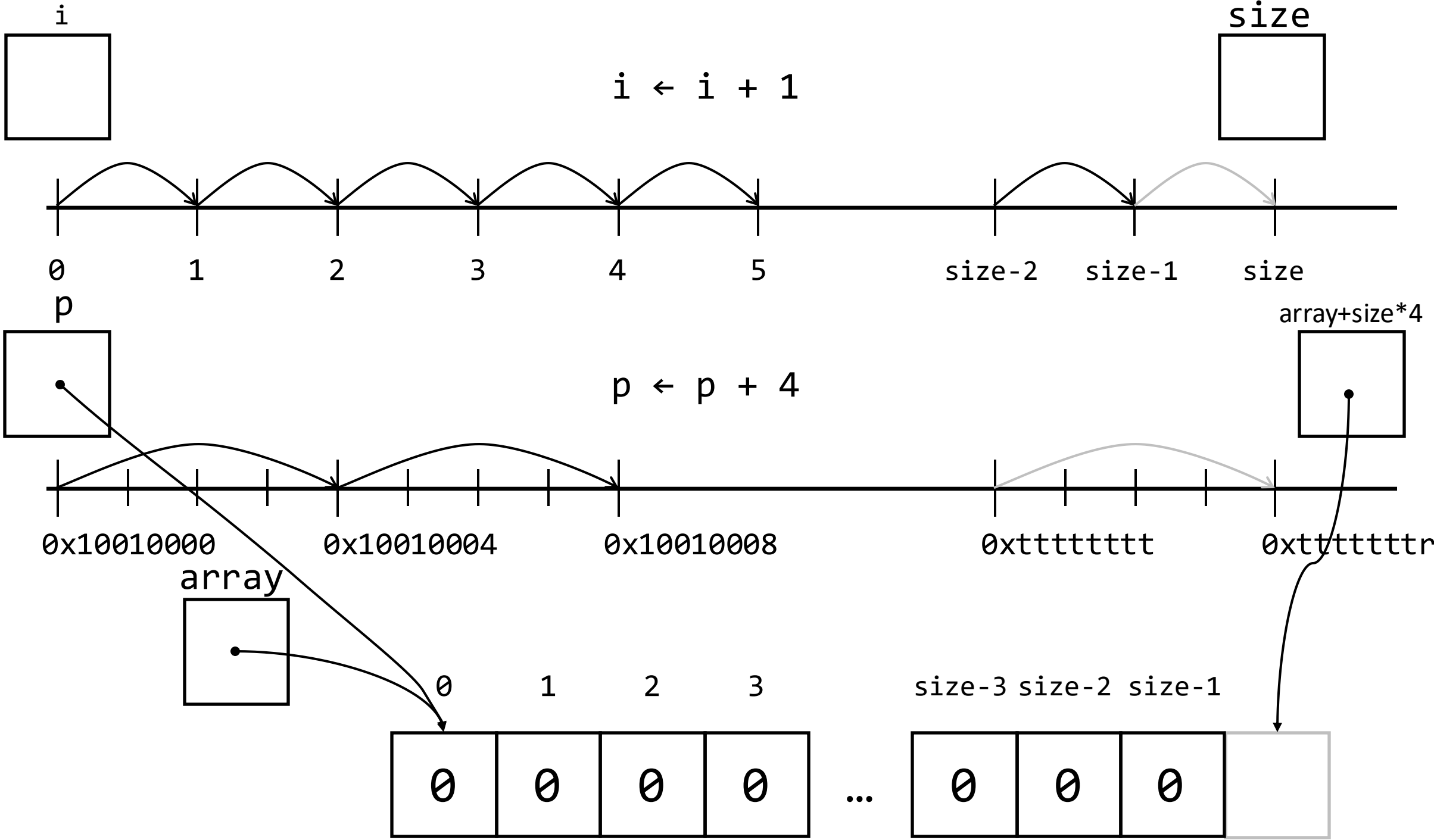
`clear`: writes zero in all elements of an integer array

Parameters:

`array`: address of first array element

`size`: number of elements in array





Array Indexing vs. Pointers

Array Indexing:

```
void clear1(int array[], int size){  
    int i;  
    for (i = 0; i < size; i += 1)  
        array[i] = 0;  
}
```

```
t0 ← 0  
while (t0 < size){  
    t2 ← array+4*t0  
    M[t2] ← zero  
    t0 ← t0 + 1  
}
```

a0 array

t0 i

Array Indexing vs. Pointers

Array Indexing:

```
void clear1(int array[], int size){  
    int i;  
    for (i = 0; i < size; i += 1)  
        array[i] = 0;  
}
```

```
t0 ← 0  
while (t0 < size){  
    t1 ← 4*t0  
    t2 ← array+t1  
    M[t2] ← zero  
    t0 ← t0 + 1  
}
```

Must multiply by 4

and add to base

in every loop iteration

a0

array

t0

i

Pointer Indexing:

```
void clear2(int *array, int size){  
    int *p;  
    for (p = &array[0]; p < &array[size]; p += 1)  
        *p = 0;  
}
```

```
t0 ← array  
t2 ← &array[size]  
while (t0 < t2){  
    M[t0] ← zero  
    t0 ← t0 + 4  
}
```

Pointer corresponds
to memory address

a0

array

t0

p

```

void clear1(int array[], int size){
    int i;
    for (i = 0; i < size; i += 1)
        array[i] = 0;
}

```

```

        ble    a1, zero, done # if n <= 0
        mv     t0, zero      # i ← 0
loop1:  slli   t1, t0, 2      # t1 ← i * 4
        add    t2, a0, t1    # t2 ← &array[i]
        sw     zero, 0(t2)   # array[i] ← 0
        addi   t0, t0, 1     # i ← i + 1
        blt    t0, a1, loop1 # if i < size
                                goto loop1

```

```

void clear2(int *array, int size){
    int *p;
    for (p = &array[0]; p < &array[size]; p += 1)
        *p = 0;
}

```

```

        ble    a1, zero, done # if n <= 0
        mv     t0, a0        # p ← &array[0]
        slli   t1, a1, 2     # t1 ← size * 4
        add    t2, a0, t1    # t2 ← &array[size]
loop2:  sw     zero, 0(t0)   # M[p] ← 0
        addi   t0, t0, 4     # p ← p + 4
        blt    t0, t2, loop2 # if p < &array[size]
                                goto loop2

```

add, addi, slli: 1 cycle; sw: 5 cycles; blt, ble: 2 cycles

Assume size is very large. Compute CPI. Which is faster? By how much?

<div># of instructions = 5 * size + 2</div> <div># of cycles = 1 + 2 + size * (1 + 1 + 5 + 1 + 2)</div> <div>= 10 * size + 3</div> <div>CPI = $\frac{10 * size + 3}{5 * size + 2} \approx 2$</div>					<div># of instructions = 3 * size + 4</div> <div># of cycles = 5 + size * (5 + 1 + 2)</div> <div>= 8 * size + 5</div> <div>CPI = $\frac{8 * size + 5}{3 * size + 4} \approx 2.67$</div>				
ble	a1, zero, done	# if n <= 0	2		ble	a1, zero, done	# if n <= 0	2	
mv	t0, zero	# i ← 0	1		mv	t0, a0	# p ← &array[0]	1	
loop1: slli	t1, t0, 2	# t1 ← i * 4	1		slli	t1, a1, 2	# t1 ← size * 4	1	
add	t2, a0, t1	# t2 ← &array[i]	1		add	t2, a0, t1	# t2 ← &array[size]	1	
sw	zero, 0(t2)	# array[i] ← 0	5		loop2: sw	zero, 0(t0)	# M[p] ← 0	5	
addi	t0, t0, 1	# i ← i + 1	1		addi	t0, t0, 4	# p ← p + 4	1	
blt	t0, a1, loop1	# if i < size	2		blt	t0, t2, loop2	# if p < &array[size]	2	
		goto loop1					goto loop2		

of cycles(index) = 10 * size + 3

of cycles(ptr) = 8 * size + 5

Speedup = $\frac{10 * size + 3}{8 * size + 5} \approx 1.25$

add, addi, slli: 1 cycle; sw: 5 cycles; blt, ble: 2 cycles

Assume size is very large. Compute CPI. Which is faster? By how much?

Comparison of Array Indexing vs. Pointers

Array indexing requires shift to be inside loop

It is part of index calculation for incremented i

Recalculate address in each iteration

Compiler can achieve same effect as manual use of pointers

Induction variable elimination

Better to make program clearer and safer