

Term Project - SIC/XE Assembler Phase (1&2)

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pass 1

○ Requirements specifications.

1. The pass1 is to execute by entering
pass1 <source-file-name>
2. The source file for the main program for this phase is to be named pass1.c
3. You should build a parser that is capable of handling source lines that are instructions, storage declaration, comments, and assembler directives (a directive that is not implemented should be ignored possibly with a warning)
 1. For instructions, the parser is to minimally be capable of decoding 2, 3 and 4-byte instructions as follows:
 - a. 2-byte with 1 or 2 symbolic register reference (e.g., TIXR A, ADDR S,A)
 - b. RSUB (ignoring any operand or perhaps issuing a warning)
 - c. 3-byte PC-relative with symbolic operand to include immediate, indirect, and indexed addressing
 - d. 3-byte absolute with non-symbolic operand to include immediate, indirect, and indexed addressing
 - e. 4-byte absolute with symbolic or non-symbolic operand to include immediate, indirect, and indexed addressing
 2. The parser is to handle all storage directives (BYTE, WORD, RESW, and RESB).
4. The output of this phase should contain (at least):
 1. The symbol table.
 2. The source program in a format similar to the listing file described in your text book except that the object code is not generated as shown below. A meaningful error message is printed below the line in which the error occurred.

pass 2 :

Specifications :

- a) The assembler is to execute by entering

assemble <source-file-name>
- b) The source file for the main program for this phase is to be named assemble.cpp
- c) The output of the assembler should include (at least):
 1. Object-code file whose format is the same as the one described in the text book in section 2.1.1 and 2.3.5.

2. A report at the end of pass2. Pass1 and Pass2 errors should be included as part of the assembler report, exhibiting both the offending line of source code and the error.

d) The assembler should support:

1. EQU and ORG statements.

2. Simple expression evaluation. A simple expression includes simple (A <op> B) operand arithmetic, where <op> is one of +, -, *, / and no spaces surround the operation, eg. A+B.

Bonus

1. General expression evaluation.

2. Literals (Including LTORG)

=C'<ASCII-TEXT>', =X'HEX-TEXT', =<DECIMAL-TEXT> forms.

3. Control sections

◦ Design.

Very Basic Design.

Two Secondary Object classes

First Class is the ID class which encloses information about SIC/XE operations. The op-code and the format whether it's 1, 2, or 3/4.

another class symOb, which is the object pushed in the hash map it contains the added address and a Boolean variable to know if this address is relocatable or absolute.

another LitOb class that stores objects of the litTab as an integer addresses.

◦ Main data structures :

- HashMaps : C++ built in hash maps of String and objects of written classes which holds information (example : OpTab < String operation , ID (int Opcode, int Format) >)
 - String arrays that represent the possible field in a SIC/XE program.
-

◦ Algorithms Description for pass 1:

Main Process Algorithm:

Pseudo code

```
- // Read first line & check opcode
- If (opcode = START) THEN
-   Locctr = #operand
- End IF
- // loop over the code till END
- While (opcode != END) do
-   If there's a symbol in LABEL field , THEN
-     Search for it in the Symbol Table
-     If (found)THEN
-       Duplicate error => true
-     Else
-       Insert (Label, Locctr) into the Syml Table
-     END IF
- // increment the LOCCTR
-   Search for the opcode in the OPTABLE
-   If (found)
-     Add number of bytes of opcode(2,3,4) to the LOCCTR
-   Else if(opcode = WORD)
-     Add 3 to LOCCTR
-   Else if (opcode = RESW)
-     Add 3 * #operand to LOCCTR
-   ELSE if (opcode = RESW)
-     Add #operand to LOCCTR
-   ELSE if(opcode = BYTE)
```

- Add length of the constant to the LOCCTR
- ELSE
- Invalid operation code => true
- Read Next Line
- END WHILE
- Get Length of the program = LOCCTR – Starting Address

How it works:

Basic idea

First, the program reads the first line of the assembly code and checks that its opcode is called “START”, and assign the LOCCTR to the starting address which is found in the operand field of the first line.

Then the program loops as long as the opcode isn’t equal to “END”, and checks if there’s a symbol in the LABEL field and add it in the symbol table iff it doesn’t already exists in it.

After that, the program increments the LOCCTR in one of three ways, either the opcode is found in optable therefore add number of bytes of opcode, or the opcode equals WORD therefore add 3, or opcode equals Byte therefore add length of the constant, or the opcode equals RESW therefore add 3* number in the operand, or the opcode equals RESB therefore add number in the operands, or the opcode doesn’t satisfy any of these conditions and therefore there will be an error.

Finally get the size of the program.

General Algorithms -Utilities-

Read line by line from input file

-reads a line from a txt file till the end of file

```
string fgetline(FILE* file)
{
    char str[77];
    if ( fgets (str , 77 , infile) == NULL )
    {
        eFile = -1; // END OF FILE( eFile : GLOBAL )
    }
    string d(str);
    unsigned pos = d.find("\n");
    d =d.substr (0,pos); // IGNORE THE RUBISH
    return uppercaseString(d);
}
```

Convert letters into UpperCase

-loops on every character and converts them using toupper(char) function.

```
string uppercaseString(string s)
{
    char c;
    int i=0;
    while (s[i])
    {
        c=s[i];
        s[i] = (char) toupper(c);
        i++;
    }
    return s;
}
```

Print in output file

```
void print(FILE* pFile,string s)
{
    fputs((uppercaseString(s)+"\n").c_str(),pFile);
}
```

Print algorithm with specific format in output file

- takes a string array and lin number and location and formats a specific string to a specific format then calls the output function above.

```
string format( int line , int loc , string s [] )
{
    string out=""; // the output string
    string lineStr = toString(line);// line number in string
    string locStr = inc(loc);// loc counter in string
    out += lineStr+" "; // add line value to first for out string
    // for loop to reserve the space in various length for string
    for ( int i = 0 ; i < (6 - locStr.size()) ; i++)
    {
        out+='0';
    }
    out += locStr+" ";
    out += s[0]; // Label value
}
```

```

for ( int i = 0 ; i < (8- s[0].size()); i++ )
{
    out += ' ';
}
out += ' ';
out += s[1]; // Mnemonic value
for ( int i = 0 ; i < (6- s[1].size()); i++ )
{
    out += ' ';
}
out += "  ";
out += s[2]; // Operand value
for ( int i = 0 ; i < (30 - s[2].size()) ; i++)
{
    out += ' ';
}
out += s[3]; // comment value
out += "\n";
return out;
}

```

Open file with dynamic or static path

```

//Create Folder if the default folder not exist

void CreateFolder(const char * path)
{
    if(!CreateDirectory(path ,NULL))
    {
        return;
    }
}

// initFiles to handle dynamic and static path for open file form any folder
void initFiles(int argc, char *argv[]){
CreateFolder("C:\\\\AssemblerFiles\\\\");// create default folder
    string filename(argv[1]);
    if ( argc > 2 )//if argument have path
    {
        string Input = argv[2]; // path string ( dynamic path )
        infile = fopen(argv[2],"r"); // open the file
        Input = Input.substr(0,Input.size()-4)+"OUT.txt";// out file path
        outfile = fopen(Input.c_str(),"w"); // open output file
    }
    else
    {
        // static path ( if files exist in default folder )
        infile = fopen(("C:\\\\AssemblerFiles\\\\"+filename+".txt").c_str(),"r");
        outfile = fopen (("C:\\\\AssemblerFiles\\\\"+filename+"OUT.txt").c_str(),"w");
    }
// check file opened or not
    if(infile == NULL)
    {
        cout << " File Can't Open !! ";
    }
    if(outfile == NULL)
    {
        cout << "\n  OUT FILE ERROR !! ";
    }
}

```

```

        outfile = fopen (( "C:\\AssemblerFiles\\"+filename+"OUT.txt").c_str(), "w");
        cout << "\n    BUT THE OUT FILE IN ( AssemblerFiles ) FOLDER " << endl;
    }
    cout << "\n\n    successfully loaded ( File : " << filename << " ) \n\n" << endl;
}

```

Print SYMTAB in OUTPUT FILE

- a method that loops on the symTab hashmap that is created at run-time to print at the end of pass one in the output file.

```

void printSymTab()
{
    map<string,symOb*>::iterator it;
    for (map<string,symOb*>::iterator it=symTab.begin();it!=symTab.end();++it)
    {
        symOb* symTemp = it->second;
        string name =it->first;
        int temp1 = symTemp->getAddress();
        int temp2 = symTemp->getRelocatable();
        print(outfile,rowTableFormat(name,temp1,temp2));
    }
}

string rowTableFormat(string name,int address , int r)
{
    // Get format of row's to print
    string out="";
    for ( int i = 0 ; i < 5 ; i++)
    {
        out+="' ';
    }
    string addstr = hexaChange(address);
    string reAbs = (r == 1)? "Relocatable" : "Absolute";
    out += name;
    for ( int i = 0 ; i < (12 - name.length()) ; i++)
    {
        out+="' ';
    }
    out+=addstr;
    for ( int i = 0 ; i < (12 - addstr.size()) ; i++)
    {
        out+="' ';
    }
    out += reAbs;
    out+="\n";
    return out;
}

```

CheckWord():

this method is called when a certain algorithm wants to check if this operand for the word directive is valid or not.


```

bool checkWord(string word)
{
    int i = 0;

    if(word[i] == '-')
    {
        i++;
    }

    for( ; i < word.size() ; i++)
    {
        if(word[i] > '9' || word[i] < '0')
        {
            return false;
        }
    }
    return true;
}

```

ChecByte():

this method is called when a certain algorithm wants to check if this operand for the Byte directive is valid or not. returns -1 if not, returns the actual number in the operand field if yes.

```

int checkByte(string byte)
{
    if(byte[0] != 'C' && byte[0] != 'c' && byte[0] != 'X' && byte[0] != 'x')
    {
        cout << "NO X OR C" << endl;
        return -1;
    }

    if(byte[1] != ',' )
    {
        cout << "first comma" << endl;
        return -1;
    }

    if(byte[byte.size() - 1] != ',' )
    {
        cout << "last comma" << endl;
        return -1;
    }

    // a character input
    if(byte[0] == 'C' || byte[0] == 'c')
    {
        // check
        cout << "number of chars" << endl;
        return byte.size()-3;
    }
}

```

```

    }

// a hex word input
int numberOfHalfBytes = 0 ;
if(byte[0] == 'X' || byte[0] == 'x')
{
    for(int i = 2 ; i < byte.size()-1 ; i++)
    {
        if(((byte[i] >= '0') && (byte[i] <= '9')) || ((byte[i] >= 'a') && (byte[i] <=
'f')) || ((byte[i] >= 'A') && (byte[i] <= 'F'))))
        {
            numberOfHalfBytes++;
        }
        else
        {
            cout << "hex out of range " << byte[i] << endl;
            return -1;
        }
    }
    if(numberOfHalfBytes % 2 == 1)
    {
        cout << "odd number of hex half bytes" << endl;
        return -1;
    }
    return numberOfHalfBytes/2;
}
cout << "main return" << endl;
return -1;
}

```

CheckResrve():

this method is called when a certain algorithm wants to check if this operand for the RESW/RESB directive is valid or not. returns -1 if not, returns the actual number in the operand field if yes.

```

int checkReserve(string res)
{
    int i = 0;
    if(res[i] == '-')
    {
        i++;
        return -1;
    }

    for ( ; i < res.size() ; i++)
    {
        if(res[i] < '0' || res[i] > '9' )
        {
            return -1;
        }
    }
}

```

```
int temp = atoi (res.c_str());  
return temp;  
}
```

Parsing Algorithms:

-There are two main parsing Methods

PrimaryParsing()

it parses the line into chunks according to empty white spaces or tabs.

it ignores any space after 3 words and understands them as a comment, which is later used for the second parsing method which handles the main parsing and is like the brain that understands where each word should be in an array of length four [LABEL, OPERATION, OPERAND, COMMENT] that the main processing algorithm uses.

```
int primaryParsing(string inputLine, string * line){  
  
    bool spaceFound = false;  
    bool commentField = false;  
    int arrayPointer = 0;  
    int actualArrayPointer = -1;  
    bool actualWrite = false;  
    int i=0;  
  
    if(inputLine.empty())  
    {  
        return -1;  
    }  
  
    if(inputLine.at(0) == ' ' )  
    {  
        spaceFound = true;  
        actualWrite = false;  
        i++;  
    }  
  
    for(; i<inputLine.size() ; i++)  
    {  
  
        if(inputLine.at(i) == ' ' && !spaceFound && !commentField)  
        {  
            spaceFound = true;  
            arrayPointer++;  
            actualWrite = false;  
        }  
        else if(inputLine.at(i) == ' ' && spaceFound && !commentField)  
        {  
            actualWrite = false;  
            continue;  
        }  
        else if(inputLine.at(i) == '\t')  
        {  

```

```

        continue;
    }
    else
    {
//      concatenate on current array pointer;

        if(inputLine.at(i) == '.')
        {
//      put rest of the string in line[3]
            line[3] = inputLine.substr (i,inputLine.size()-i);
            return 0;
        }

        if(arrayPointer == 3)
        {
            commentField = true;
        }
        if(!actualWrite)
        {
            actualWrite = true;
            actualArrayPointer++;
        }
        line[arrayPointer] += inputLine.at(i);
        spaceFound = false;
    }
}
return actualArrayPointer;
}

```

the second method in the parsing algorithms

Secondary Parsing()

This method is responsible for performing checks that every string in each of the array slots is as stated. The Algorithm is as follow :

if the array is full and the first slot has "operation code" then we should ask if the operation is of format 1 then the rest of the slots are comments and must be moved to slot 4, else then concatenate slot 3 and 4 where both of them are comment and set them in slot 4 , then shift the operation and the operand to the right slots.

this method uses a returned value from primary parsing, which is the number of slots that actually have words written in them. referred to as SplitSize a variable of which the main algorithm depends on.

The Map Method :

An initializer method which initializes the hash map with the opcodes and formats to each instruction.

this method was built with a clever trick. we obtained the appendix of the SIC/XE in a written format and then wrote a java parser program that reads that file line by line and then outputs in the java console a code written in C++ which is designed to be copy pasted directly without filling the hashmap by hand one by one.

Assumptions pass 1 :

-It is assumed that the user cannot use a label name which is the same as the name of an instruction or a directive. Further information about this part in the bonus section.

Bonus :

- Format #4 instructions (instructions preceded by “+” operator), it checks first if the instruction contains “+” operator then remove from the string and sets a Boolean named formatFour to true, after that the Locctr is incremented by 4 instead of 3.
- “EQU” directive, the program handles Equate directive by directly printing the line containing “EQU” using the address in its operand field, and nothing happens to the Locctr so that it can be used normally for instructions after “EQU”.
- “ORG” directive, when the program reads line containing “ORG” it simply sets Locctr to number in operand field.
- Free format is implemented the user is free to write the four Fields in any place in the file, comments has been handled also, thus there can be two types of comments in the program : a line comment which is the normal case of a line starting with a (.) which means that this line is a comment. or a line which has it's fourth field as a comment. the parser checks to see of the instruction is of which format and thus knows if the word preceding is an operand or a comment with a certain algorithm specifically designed to handle this issue.
- The End directive is handled to be used as a label which breaks the assumption already stated, but we made little trials to test whether it can be done or not, and it worked. we didn't do it for the rest of the code for the sake of time and simplicity. also the end directive is harder than any other directive or instruction because it's the only one which can have an operand or not.
- A silent installer that adds the path of the .exe file of the program in the environment variables of the computer so that the program can run directly from the cmd.
- Two input methods for running the pass1 Program. a Drag and drop of the file to the cmd, which pass the argument to the program, this argument is the path of the file, the output file generated in the same folder from which the original file was dragged.
- expressions are handled
- Literals are also handled

Sample runs :

```
INDIRECTTEST - Notepad
File Edit Format View Help
.23456789012345678
indirect START 0
LDA U
STA UU
LDA V
STA VV
LDA #M
STA MM
JSUB FUN
J *
FUN LDA UU
ADD VV
STA @MM
RSUB
WORD 3
V WORD 4
M RESW 1
UU RESW 1
VV RESW 1
MM RESW 1
END

INDIRECTTESTOUT - Notepad
File Edit Format View Help
OUTPUT
# ADDRESS LABEL MNEMONIC OPERAND COMMENTS
.23456789012345678
1 000000 INDIRECT START 0
2 000000 LDA U
3 000003 STA UU
4 000006 LDA V
5 000009 STA VV
6 00000C LDA #M
7 00000F STA MM
8 000012 JSUB FUN
9 000015 J *
10 000018 FUN LDA UU
11 00001B ADD VV
12 00001E STA @MM
13 000021 RSUB
14 000024 U WORD 3
15 000027 V WORD 4
16 00002A M RESW 1
17 00002D UU RESW 1
18 000030 VV RESW 1
19 000033 MM RESW 1
20 000036 END
>> E N D O F P A S S 1
>>
>>

TOUT - Notepad
File Edit Format View Help
OUTPUT
# ADDRESS LABEL MNEMONIC OPERAND COMMENTS
1 001000 TEST START 1000
2 001000 READ LDX ZERO
3 001003 RLOOP TD INDEV
4 001006 JEQ RLOOP COMMENT
5 001009 RD INDEV
6 00100C STCH RECORD,X
7 00100F TIX K100
8 001012 JLT RLOOP
9 001015 INDEV BYTE X'F3'
10 001016 RECORD RESB 100
11 00107A ZERO WORD 0
12 00107D K100 WORD 100
13 001080 END
>> E N D O F P A S S 1
>>
>>
>> S Y M B O L T A B L E (VALUES IN HEXA.)
NAME VALUE ABSOL/RELOC
INDEV 1015 RELOCATABLE
K100 107D RELOCATABLE
READ 1000 RELOCATABLE
RECORD 1016 RELOCATABLE
RLOOP 1003 RELOCATABLE
ZERO 107A RELOCATABLE

T - Notepad
File Edit Format View Help
TEST START 1000 READ LDX ZERO
RLOOP TD RLOOP INDEV Comment
JEQ RD RECORD,X
STCH TIX K100
JLT RLOOP
RECORD RESB 100 INDEV BYTE X'F3'
K100 WORD 100 ZERO WORD 0
END
```

TOUT - Notepad

File Edit Format View Help

OUTPUT

#	ADDRESS	LABEL	MNEMONIC	OPERAND	COMMENTS
					.23456789012345678901234567890123456
1	001000	TEST	START	1000	
2	001000	READ	LIDX	ZERO	
					ERROR:*** INVALID OPERATOR !!
3	001000	RLOOP	TD	INDEV	
4	001003		JEQ	RLOOP	NDNSJSJNDJSND
5	001006	RDD	INDEV		
					ERROR:*** INVALID OPERATOR !!
6	001006	INDEV	BYTE	X'F3'	
7	001007	RECORD	RESB	1000	
					ERROR:*** INVALID OPERAND IN RESB DIRECTIVE !!
8	001007	ZERO	WORD	0	
9	00100A	K100	WORD	1000	
					ERROR:*** INVALID OPERAND IN WORD DIRECTIVE !!
10	00100A		END		

>> END OF PASS 1

T - Notepad

File Edit Format View Help

.23456789012345678901234567890123456

TEST START 1000

READ LIDX ZERO

RLOOP TD INDEV

JEQ RLOOP

RDD INDEV

INDEV BYTE X'F3'

RECORD RESB 1000

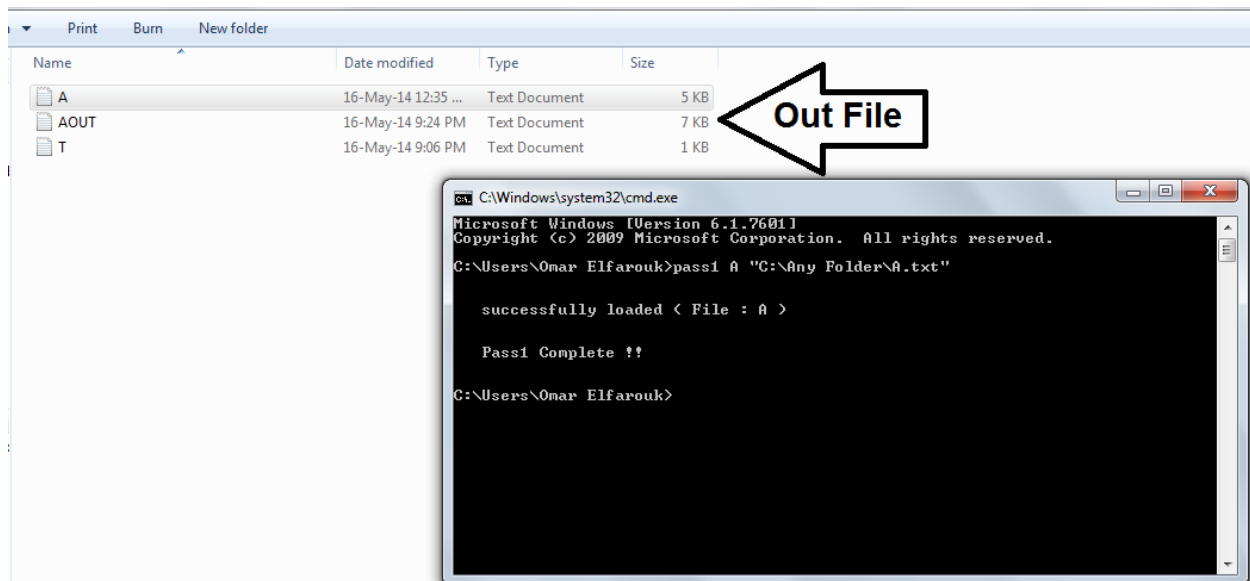
ZERO WORD 0

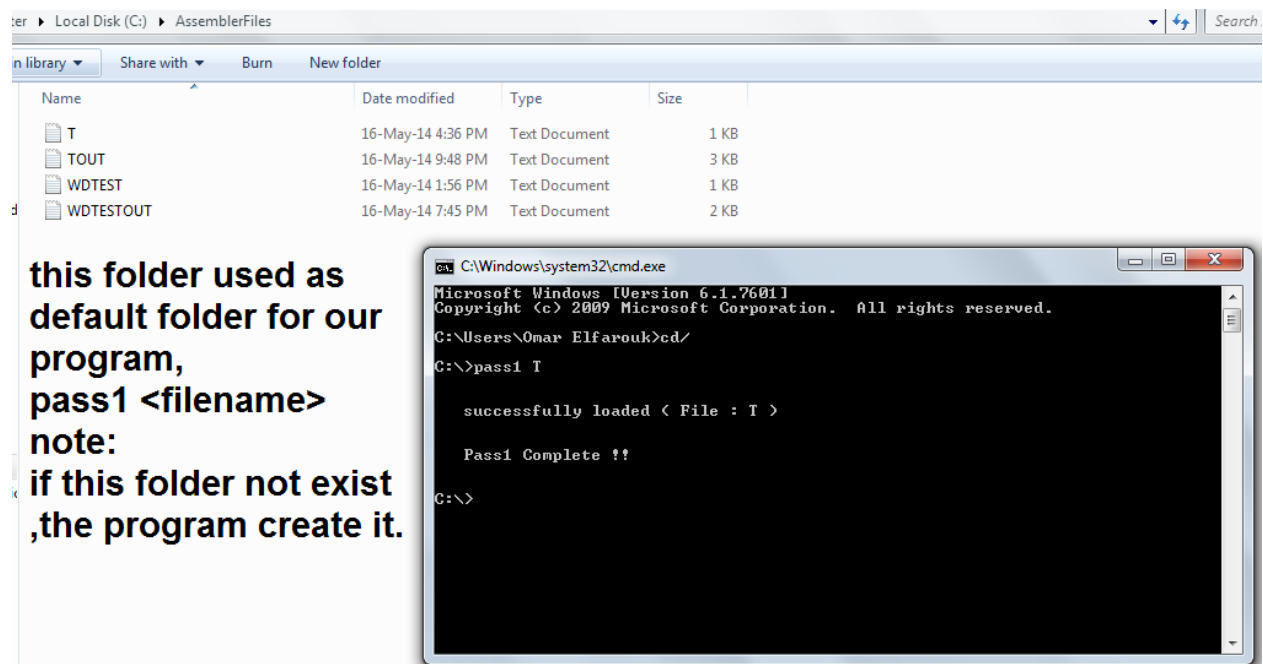
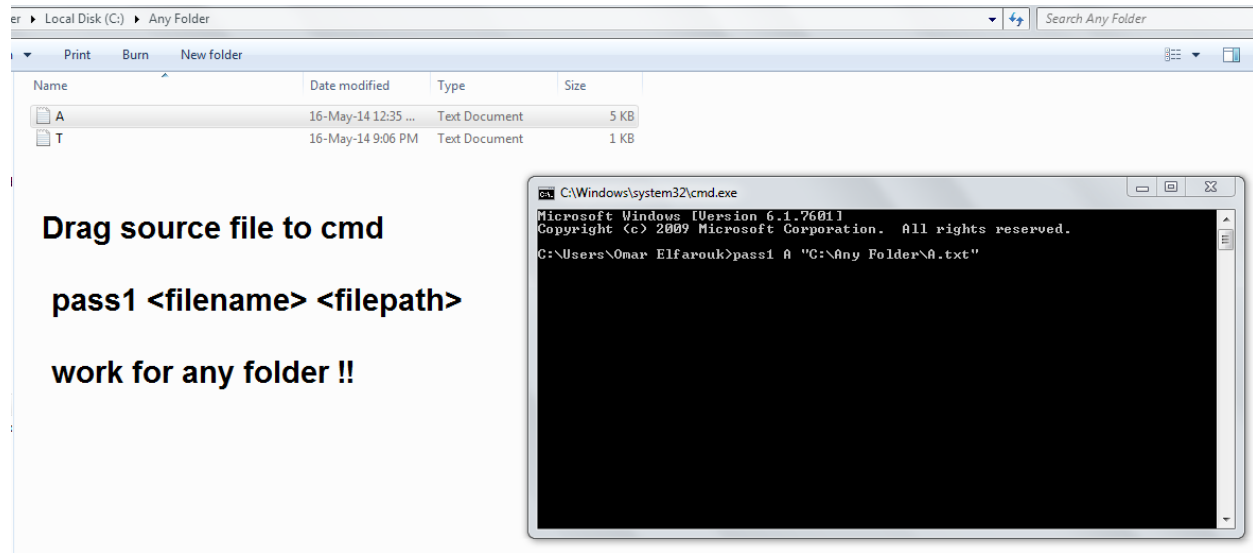
K100 WORD 1000

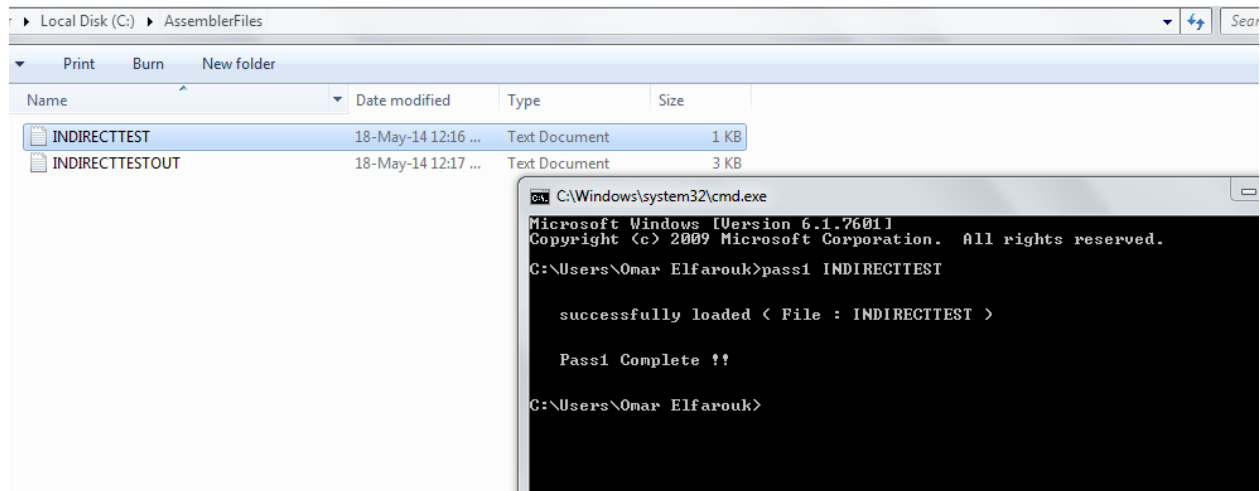
END

ndnsjsjndjsnd

Bonus Screen Shots:







NetBeansProjects	12-May-14 12:03 ...	File folder	
Razer	05-Apr-14 6:34 PM	File folder	
Visual Studio 2010	10-May-14 12:50 P...	File folder	
setup-pass1	18-May-14 12:10 ...	Application	1,655 KB

To initialize the program by add executable file path to environment variable

Parsing pass 2 :

Modify

How it works:

This method takes field 3 and field 4, which are (opcode and operand fields), apply operations on them and sets bits array accordingly, the bits array is the array which contains the “n-i-x-b-p-e” bits, these bits are responsible for telling whether it is format 4 or not, pc relative or base relative, immediate or indirect or neither nor and finally whether there is indexing or not.

Also this method calls method `infixToPostfix`, sending the operand as a parameter such that if the operand was expression then the string returned is the expression in postfix notation, however if the operand wasn't expression then the string returned will be the same string sent.

After that the modify method calls `evaluatePostfix` method sending the string returned from `infixToPostfix` method as a parameter, the `evaluatePostfix` method returns integer representing the address of the operand in both case whether the operand was expression or not.

Finally the modify method returns that address.

Algorithm:

```
Modify (Locctr, Field3, Field4)
    if (field3 is preceded by"+") Then
        eBit = 1
        remove "+" from field3
```

```

end if
else if (field4 is preceded by "@") Then
    nBit = 1
    remove "@" from field4
end if
else If (field4 is preceded by "#") Then
    iBit = 1
    remove "#" from field4
end if
else
    nBit = 1
    iBit = 1
end if
if(field4 contains ",X"at the last two indexes)Then
    xBit = 1
    remove",X" from field4
end if
String postfix = infixToPostfix(field 4)
// get field 4 in postfix notation
int address = evaluatePostfix(postfix)
// get address of operand
return address

```

Evaluate postfix :

This method takes a string converted from infix to postfix from another method

it iterates on the string concatenating each character if it was a Label or a number and adds it to the value stack

this algorithm was made to handle all kinds of expressions and know at the end of the expression if that expression was an Absolute expression, a Relative Expression or a "Neither-Nor" Expression.

it does so by using certain rules of the written algorithm

Data Structures for this algorithm : Two stacks of Integers : Value, Relativity

the value stack is the normal stack of the values of labels or numbers being pushed in the stack, like a normal postfix evaluation algorithm

the relativity stack is used whenever something is inserted in the value stack, the corresponding value is inserted in it

assume that anything absolute - i.e 1000 , 500 , Buffend-buffer - has a relative value of zero, so when anything absolute is pushed in the value stack a zero is pushed in the relativity stack

but if a label was to be inserted we get its value from the symtab and insert it into values, and get if it was relocatable or not, if it was we insert 1 in the relativity stack, meaning it has a relative rank of 1, if it was absolute then we put 0, meaning it has relativity rank of 0

-if there were a '+', the normal protocol is done on the value stack, pop the last two add them and insert them back

but for the relativity stack we pop the last two and add them back as well

if the last two popped were 0 (corresponding to an absolute value) and 1 (a relative value) then their addition will make also a relative term (push 1 back in the relativity stack)

the same is done with '-' but subtracting the values and the relative terms

multiplication is done differently, if we found two relative terms multiplied by each other we send back an error to the calling line

if an absolute term is multiplied to a absolute term, the values are popped multiplied and pushed back, and relativity is still zero

the main case is when a relative term is multiplied by an absolute term , the value of the absolute term is multiplied by the relativity number of the relative term and pushed into the relativity term, the values is treated as basic operation of multiplication

the same is done with "/" division

at the end of the method basic checks are done like if more than element remains in one of the two stacks then this is an invalid postfix expression

the last element in the value stack is returned to the calling line

the last element in the relativity stack, if it was one, then this expression has a relativity order of 1, meaning it's a relative valid term

if it was zero, then this expression has a relativity order of 0, meaning it's an absolute term

if it was any other value then it's an invalid expression when talking in terms of relativity, meaning it's a neither nor expression an error is returned to the calling line

Processing pass 2 algorithms:

```
Pass2:
begin
  read first input line (from intermediate file)
  if OPCODE = 'START' then
    begin
      write listing line
      read next input line
    end (if START)
  write Header record to object program
  initialise first Text record
  while OPCODE ≠ 'END' do
    begin
      if this is not a comment line then
        begin
          search OPTAB for OPCODE
          if found then
            begin
              if there is a symbol in OPERAND field then
                begin
                  search SYMTAB for OPERAND
                  if found then
                    store symbol value as operand address
                  else
                    begin
                      store 0 as operand address
                      set error flag (undefined symbol)
                    end
                  end (if symbol)
                else
                  store 0 as operand address
                  assemble the object code instruction
                end (if opcode found)
              else if OPCODE = 'BYTE' or 'WORD' then
                convert constant to object code
              if object code will not fit into the current Text record then
                begin
                  write Text record to object program
                  initialise new Text record
                end
              end
              add object code to Text record
            end (if not comment)
            write listing line
            read next input line
          end (while not END)
        write last Text record to object program
        write End record to object program
        write last listing line
      end (Pass 2)
```

Figure 2.4(b) Algorithm for Pass 2 of assembler.

pass2Processing :

this method is responsible for looping on line by line coming from the parser which has parsed the intermediate file, and starts making checks on it so if the opcode is (start) to start writing the header record to the object code file by calling method (formatObjectCode), and then starts writing the first text record by while looping over the lines coming from the parser till operator(End) is found , checks if the line is comment then it call method (printOneLine) then checks if the opcode is found in the optab then starts to check the operand if it is in the symtab or if they are registers in case the opcode is format 2 or if it is immediate operand because all of this changes in the value of displacement which will be sent to method (formatObjectCode) which will combine all the bits into hexacharacter in string and send it to method print . if the operator was a directive like (byte) or (word) or , then it converts the value of the constant to object code, if the length of the line is finished a new line is started and the object code is written in this new line.

convertBitstoint :

this method takes string of bits , changes it to decimal number by checking 2 characters from the string .

findRegisterNumber :

this method takes a character which is a register and find its number ranging from 0 to 6

CheckOneRegister :

this method takes an operand and boolean called svc , if this boolean is true then the operator was svc and this means that the operand was a number . the target of this method is to help in the check done by method (checkRegisterOperand), so that when the operator is format 2 then the operand must be a register , so this method is specified for the operators of format 2 which takes only one character as an operand and returns an array of integer of size 3 where the first slot is a boolean (0 -> false statement , 1 -> true statement) the second and the third slot are the values of the register or (n) in case of operator (svc).

CheckTwoRegister :

this method exactly perform the same functionality as the last described method but for the operators with operands of 2 registers and it returns array of 3 slots where the first slot is a boolean (0 -> false statement , 1 -> true statement) the second and the third slot are the values of the register or (n) in case of operator (shiftr / shiftrl).

CheckTwoRegister :

this method is the father method of the last 2 methods , it takes 2 strings of the operation and the operand . checks if the operation is of operation type that takes 2 registers or only one register and according to this call on of the two previous methods. it returns array of 3 slots as described

formatObjectCode :

this method is responsible for binding the bits to form the object code , this method returns string which is the object code as hexadecimal characters. at the beginning we start to check that the number of each input is correct, where the inputs are : int opcode which is the opcode as decimal and string bits which is nixbpe bits and string displacement and integer format of the opcode and int register 1 and inint register 2. this method. if the format is 1 or 2 then the object code is easy, just the object code or the object code + the number of the registers. if the operation is format 3 or format 4,first we add the decimal value of the object code + the decimal value of the 2 bits (ni) and set the value of the first 2 hexabytes , after this i add the hexadecimal character of the 4 bits (xbpe) and add the displacement to the output string and returns it .

Initialize Files :

This method to initialize files to make to objects file to read and write from same file. And to reach the first line in code.

```
void initPassTwo()
{
    readFile = fopen ("LISAFILE.txt","r");
    writeFile = fopen("LISAFILE.txt","a");
    objFile = fopen("OBJFILE.txt","w");
    char f[100];
    fgets(f,100,readFile);
    string g(f);
    while( g.at(0) != '#' ) // REACH FIRST LINE FOR THE CODE
    {
        g = fgetline(readFile);
    }
    print(writeFile,">>
*****
*****");
    print(writeFile,">>  S t a r t   o f   P a s s   I I");
    print(writeFile,">>  A s s e m b l e d   p r o g r a m   l i s t i n
g");
    current = fgetline(readFile);
    while ( current.at(0) == ' ' ){
        print(writeFile,current);
        current = fgetline(readFile);
    }
}
```

Parse line (from lisa File):

Return string array with three parameter (next location counter , operator and operand)
the comment line print direct to file and get next line.

```
void getParsedLine(string out[])
{
    out[3] = current;
    if ( out[3].at(0) == ' ' )
    {
        current = fgetline(readFile);
        return;
    }
    if(out[3].at(0) == '>')
    {
        eFile = -1;
    }
}
```

```

        return ;
    }
    int length=0;
    for ( int i = 26 ; i < current.size() ; i++)
    {
        if ( current.at(i) == ' ' )
        {
            break;
        }
        length++;
    }
    out[1] = current.substr(26,length);
    length=0;
    for (int i = 38 ; i < current.size(); i++)
    {
        if ( current.at(i) == ' ' )
        {
            break;
        }
        length++;
    }
    out[2] = current.substr(38,length);
    currentLoc = current.substr(6,6);;
    current = fgetline(readFile);
    while ( current.at(0)==' ' ){
        print(writeFile,current);
        current = fgetline(readFile);
    }
    out[0] = current.substr(6,6);
}

```

Split by specific string (for every char for string):
Return vector of string.

```

vector<string> split(string str1,string sf)
{
    vector<string> f;
    char * pch;
    for ( int i = 0 ; i < str1.length() ; i++ )
    {
        if ( str1.at(i) > 41 && str1.at(i) < 48 || str1.at(i)=='(' ||
str1.at(i)==')' )
        {
            str1 = str1.substr(0,i)+' '+str1.at(i)+'
'+str1.substr(i+1,str1.length());
            i++;
        }
    }
    char str[100];
    strcpy(str,str1.c_str());
    pch = strtok (str,sf.c_str());
    while (pch != NULL)
    {
        f.push_back(pch);
        pch = strtok (NULL, sf.c_str());
    }
}

```

```
    return f;
}
```

Convert infix expression to postfix:

```
string infixToPostfix(vector<string> str)
{
    stack<string> stack;
    string postFix="";
    int comingPr=0;
    int topPr=0;
    for ( int i = 0 ; i < str.size() ; i++ )
    {
        if( str.at(i).at(0)=='(' )
        {
            stack.push(str.at(i));
        }
        else if(str.at(i).at(0) ==')')
        {
            while( stack.top().at(0) !='(' )
            {
                postFix += stack.top()+" ";
                stack.pop();
            }
            stack.pop();
        }
        else if ( str.at(i).at(0) > 41 && str.at(i).at(0) < 48 )
        {
            if( str.at(i).at(0)=='(' )
            {
                stack.push(str.at(i));
            }
            else if(str.at(i).at(0) ==')')
            {
                while( stack.top().at(0) !='(' )
                {
                    postFix += stack.top()+" ";
                    stack.pop();
                }
                stack.pop();
            }
            else if ( str.at(i).at(0) == '+' || str.at(i).at(0) == '-' )
            {
                comingPr=1;
            }
            else
            {
                comingPr=2;
            }
            if ( !stack.empty() && (stack.top().at(0) == '+' ||
stack.top().at(0) == '-' ) )
            {
                postFix += stack.top()+" ";
                stack.pop();
                topPr =1;
            }
        }
    }
}
```

```

        else if(!stack.empty() && (stack.top().at(0) == '*' ||
stack.top().at(0) == '/') )
        {
            topPr =2;
        }
        if ( topPr < comingPr || stack.top().at(0) == '(' )
        {
            stack.push(str.at(i));
        }
        else
        {
            while( !stack.empty() && topPr >= comingPr )
            {
                postFix += stack.top()+" ";
                stack.pop();
                if ( !stack.empty() && ( stack.top().at(0) == '+' ||
stack.top().at(0) == '-' ) )
                {
                    topPr =1;
                }
                else if ( !stack.empty() && (stack.top().at(0) == '*' ||
stack.top().at(0) == '/') )
                {
                    topPr =2;
                }
            }
            stack.push(str.at(i));
        }
    }
    else
    {
        postFix += str.at(i)+" ";
    }
}
while( !stack.empty() )
{
    postFix += stack.top()+" ";
    stack.pop();
}
return postFix;
}

```

Get char in hex base from binary string:

```

char binaryToHexChar(string str)
{
    char * ptr;
    long parsed = strtol(str.c_str(), & ptr, 2);
    return hexaChange(parsed).at(0);
}

```

Return object code for WORD Directive:

```

string wordObj(string str)
{
    string temp = hexaChange(atoi( str.c_str()));
}

```

```

string zero = "";
for ( int i = 0 ; i < 6-temp.size(); i++ )
{
    zero += '0';
}
return zero+temp;
}

```

Return object code for BYTE Directive:

```

string byteObj( string str )
{
    if( str.at(0) == 'X' )
    {
        vector<string> s = split(str,"X");
        return s[0];
    }
    else if( str.at(0) == 'C' )
    {
        vector<string> s = split(str,"C");
        string out = "";
        for ( int i = 0 ; i < s[0].size() ; i++ )
        {
            int c = (int) (s[0].at(i));
            cout << c << endl;
            out += hexaChange(c);
        }
        return out;
    }
}

```

assumptions :

- 1. the user is not allowed to enter the expression with spaces*
- 2. the user is free to enter whatever expression S/he would like*
- 3. the user is free to enter comments where ever S/he would like*
- 4. &&& cannot be used as a label.*

Screen shots :

```
.bubble sort
.temp = t
.2345678901234567890123
bubble      START      1000
            .LTORG
            base        temp
            LDA          temp
            LDA          #0
            LDT          #0
            LDX          #0
            LDS          #0
one         RESW        10
two         RESB        100
iLOOP      LDX          #0
            STX          =c'asda'
jLOOP      LDX          3**+iLoop+*
            LDCH         STR+1,X
            .LTORG
            RMO          A,S
THREE      RESW        10
            LDCH         STR,X
            COMPR        A,S
tar        resb        4000
            equ          tar
            equ          tar-jloop
            equ *
            equ          tar-tar
            equ          3*tar-3*tar+100-40
yy         JGT          JSWAP
            J            BREAK
JSWAP      JSUB          SWAP
            .
BREAK     LDX          LNTH
            LDT          =X'1234'
            SUBR         X,T
            LDX          jjj
            TIXR         T
            STX          jjj
            JLT          jLOOP
            .
            LDX          iii
            TIX          LNTH
            STX          iii
            JLT          iLOOP
            J            HALT
            .
SWAP      LDX          jjj
            LDCH         STR,X
            STCH         TEMP
            LDCH         STR+1,X
```


	STCH	STR,X
	LDCH	TEMP
	STCH	STR+1,X
	RSUB	
HALT	J	*
STR	BYTE	C'3142'
iii	WORD	0
jjj	WORD	0
LNPTH	WORD	4
LNPTH2	WORD	3
TEMP	RESB	1
	END	

Output #	Address	Label	Mnemonic	Operand	Comments
		.BUBBLE	SORT		
		.TEMP	= T		
		.2345678901234567890123			
1	001000	BUBBLE	START	1000	
		.LTORG			
2	001000		BASE	TEMP	
3	001000		LDA	TEMP	
4	001003		LDA	#0	
5	001006		LDT	#0	
6	001009		LDX	#0	
7	00100c		LDS	#0	
8	00100f	ONE	RESW	10	
9	00102d	TWO	RESB	100	
		.			
10	001091	ILOOP	LDX	#0	
11	001094		STX	=c'asda'	
		.			
12	001097	JLOOP	LDX	3**+ILOOP+*	
13	00109a		LDCH	STR+1,X	
		.LTORG			
14	00109d		RMO	A,S	
15	00109f	THREE	RESW	10	
16	0010bd		LDCH	STR,X	
17	0010c0		COMPR	A,S	
18	0010c2	TAR	RESB	4000	
		.SHE	EQU	TAR	
19	002062	DAS	EQU	TAR-JLOOP	
		.ASDA	EQU *		
20	002062	D	EQU	TAR-TAR	
21	002062	YY	EQU	3*TAR-3*TAR+100-40	
22	002062		JGT	JSWAP	
23	002065		J	BREAK	
24	002068	JSWAP	JSUB	SWAP	
		.			
25	00206b	BREAK	LDX	LNPTH	
26	00206e		LDT	=X'1234'	
27	002071		SUBR	X,T	
28	002073		LDX	JJJ	
29	002076		TIXR	T	
30	002078		STX	JJJ	
31	00207b		JLT	JLOOP	
		.			
32	00207e		LDX	III	
33	002081		TIX	LNPTH	
34	002084		STX	III	
35	002087		JLT	ILOOP	
36	00208a		J	HALT	
		.			
37	00208d	SWAP	LDX	JJJ	
38	002090		LDCH	STR,X	

```

39      002093      STCH      TEMP
40      002096      LDCH      STR+1,X
41      002099      STCH      STR,X
42      00209c      LDCH      TEMP
43      00209f      STCH      STR+1,X
44      0020a2      RSUB
45      0020a5      HALT      J      *
46      0020a8      STR      BYTE      C'3142'
47      0020ac      III      WORD      0
48      0020af      JJJ      WORD      0
49      0020b2      LENGTH   WORD      4
50      0020b5      LENGTH2  WORD      3
51      0020b8      TEMP     RESB      1
52      0020b9      END
      0020B9      =C'ASDA'
      0020BD      =X'1234'
>>  e n d   o f   p a s s   1
>>  ****
>>  ****
>>  s y m b o l   t a b l e   (values in Hexa.)
      Name      value      Absol/Reloc
      -----
      BREAK      206b      Relocatable
      D           0         Absolute
      DAS         2b        Absolute
      HALT        20a5      Relocatable
      III         20ac      Relocatable
      ILOOP       1091      Relocatable
      JJJ         20af      Relocatable
      JLOOP       1097      Relocatable
      JSWAP       2068      Relocatable
      LENGTH      20b2      Relocatable
      LENGTH2     20b5      Relocatable
      ONE         100f      Relocatable
      STR         20a8      Relocatable
      SWAP        208d      Relocatable
      TAR         10c2      Relocatable
      TEMP        20b8      Relocatable
      THREE       109f      Relocatable
      TWO         102d      Relocatable
      YY          3c        Absolute
>>  ****
>>  s t a r t   o f   P a s s   I I
>>  A s s e m b l e d   p r o g r a m   l i s t i n g
      .BUBBLE SORT
      .TEMP = T
      .2345678901234567890123
      .LTOrg
1      001000      BUBBLE      START      1000
2      001000      BASE      RELC N = 1 I = 1 X = 0 B = 0 P = 0 E = 0
                        TEMP

```

3	001000	LDA	RELC N = 1 I = 1 X = 0 B = 1 P = 0 E = 0 TEMP	034000
4	001003	LDA	ABS N = 0 I = 1 X = 0 B = 0 P = 0 E = 0 #0	010000
5	001006	LDT	ABS N = 0 I = 1 X = 0 B = 0 P = 0 E = 0 #0	750000
6	001009	LDX	ABS N = 0 I = 1 X = 0 B = 0 P = 0 E = 0 #0	050000
7	00100c	LDS	ABS N = 0 I = 1 X = 0 B = 0 P = 0 E = 0 #0	6D0000
8	00100f	ONE .	RESW 10	
9	00102d	TWO	RESB 100 ABS N = 0 I = 1 X = 0 B = 0 P = 0 E = 0	
10	001091	ILOOP .	LDX #0	050000
11	001094	STX	RELC N = 1 I = 1 X = 0 B = 1 P = 0 E = 0 =c'asda'	134001
12	001097	JLOOP Error:*** undefined symbol in operand!!	LDX 3**+ILOOP+*	
	.LTORG			
13	00109a	LDCH Error:*** out of boundry!!	RELC N = 1 I = 1 X = 1 B = 0 P = 1 E = 0 STR+1,X	
14	00109d	RMO	RELC N = 1 I = 1 X = 0 B = 0 P = 1 E = 0 A,S	AC04
15	00109f	THREE	RESW 10 RELC N = 1 I = 1 X = 1 B = 0 P = 1 E = 0	
16	0010bd	LDCH Error:*** out of boundry!!	STR,X RELC N = 1 I = 1 X = 0 B = 0 P = 1 E = 0	
17	0010c0	.SHE EQU	COMPR A,S	A004
		TAR		
18	0010c2	TAR .ASDA EQU *	RESB 4000	
19	002062	DAS	EQU ABS N = 1 I = 1 X = 0 B = 0 P = 0 E = 0 TAR-JLOOP	
20	002062	D	EQU ABS N = 1 I = 1 X = 0 B = 0 P = 0 E = 0 TAR-TAR	
21	002062	YY	EQU ABS N = 1 I = 1 X = 0 B = 0 P = 0 E = 0 3*TAR-3*TAR+100-40	
22	002062	JGT	RELC N = 1 I = 1 X = 0 B = 0 P = 1 E = 0 JSWAP	372003
23	002065	J	RELC N = 1 I = 1 X = 0 B = 0 P = 1 E = 0 BREAK	3F2003
	.			

24	002068	JSWAP	JSUB	RELC N = 1 I = 1 X = 0 B = 0 P = 1 E = 0 SWAP RELC N = 1 I = 1 X = 0 B = 0 P = 1 E = 0	4B2022
25	00206b	BREAK	LDX	LNGLTH RELC N = 1 I = 1 X = 0 B = 0 P = 1 E = 0	072044
26	00206e		LDT	=X'1234' RELC N = 1 I = 1 X = 0 B = 0 P = 1 E = 0	77204C
27	002071		SUBR	X,T RELC N = 1 I = 1 X = 0 B = 0 P = 1 E = 0	9415
28	002073		LDX	JJJ RELC N = 1 I = 1 X = 0 B = 0 P = 1 E = 0	072039
29	002076		TIXR	T RELC N = 1 I = 1 X = 0 B = 0 P = 1 E = 0	B850
30	002078		STX	JJJ RELC N = 1 I = 1 X = 0 B = 0 P = 1 E = 0	132034
31	00207b		JLT	JLOOP Error:*** out of boundry!! RELC N = 1 I = 1 X = 0 B = 0 P = 1 E = 0	
32	00207e		LDX	III RELC N = 1 I = 1 X = 0 B = 0 P = 1 E = 0	07202B
33	002081		TIX	LNGLTH RELC N = 1 I = 1 X = 0 B = 0 P = 1 E = 0	2F202E
34	002084		STX	III RELC N = 1 I = 1 X = 0 B = 0 P = 1 E = 0	132025
35	002087		JLT	ILOOP Error:*** out of boundry!!	
36	00208a		J	RELC N = 1 I = 1 X = 0 B = 0 P = 1 E = 0 HALT	3F2018
37	00208d	SWAP	LDX	RELC N = 1 I = 1 X = 0 B = 0 P = 1 E = 0 JJJ	07201F
38	002090		LDCH	RELC N = 1 I = 1 X = 1 B = 0 P = 1 E = 0 STR,X	53A015
39	002093		STCH	TEMP RELC N = 1 I = 1 X = 1 B = 0 P = 1 E = 0	572022
40	002096		LDCH	STR+1,X RELC N = 1 I = 1 X = 1 B = 0 P = 1 E = 0	53A010
41	002099		STCH	STR,X RELC N = 1 I = 1 X = 0 B = 0 P = 1 E = 0	57A00C
42	00209c		LDCH	TEMP RELC N = 1 I = 1 X = 1 B = 0 P = 1 E = 0	532019
43	00209f		STCH	STR+1,X RELC N = 1 I = 1 X = 0 B = 0 P = 0 E = 0	57A007
44	0020a2		RSUB	RELC N = 1 I = 1 X = 0 B = 0 P = 1 E = 0	4F0000
45	0020a5	HALT	J	*	3F2017
46	0020a8	STR	BYTE	C'3142'	33313432
46	0020a8	STR	BYTE	C'3142'	33313432
47	0020ac	III	WORD	0	000000
48	0020af	JJJ	WORD	0	000000
49	0020b2	LNGLTH	WORD	4	000004
50	0020b5	LNGLTH2	WORD	3	000003
51	0020b8	TEMP	RESB	1	
	0020b9		=C'ASDA'		
	0020bd		=X'1234'		
52	0020b9		END		

H^BUBBLE^001000^0010bf
T^001000^0f^0340000100007500000500006D0000
T^001091^0b^05000013400153A000AC04
T^0010bd^05^53A000A004
T^002062^1c^3720033F20034B202207204477204C9415072039B8501320343B2000
T^00207e^1e^07202B2F202E1320253B20003F201807201F53A01557202253A01057A00C
T^00209c^1c^53201957A0074F00003F201733313432000000000000000004000003
E^001000

 . LTOrg
 base temp
 LDA temp
 LDAf #0
 LDT #0
 LDX #0
 LDS #0

Output #	Address	Label	Mnemonic	Operand	Comments
		.BUBBLE SORT			
		.TEMP = T			
		.2345678901234567890123			
1	001000	BUBBLE	START	1000	
		.LTORG			
2	001000		BASE	TEMP	
3	001000		LDA	TEMP	
4	001003	LDAF	#0		
5	001003		Error:*** invalid operator !!		
6	001006		LDT	#0	
7	001009		LDX	#0	
8	00100c	ONE	RESW	10	
9	00102a	TWO	RESB	100	
		.			
10	00108e	ILOOP	LDX	#0	
11	001091		STX	=c'asda'	
		.			
12	001094	JLOOP	LDX	3**+ILOOP+*	
13	001097		LDCH	STR+1,X	
		.LTORG			
14	00109a		RMO	A,S	
15	00109c	THREE	RESW	10	
16	0010ba		LDCH	STR,X	
17	0010bd		COMPR	A,S	
18	0010bf	TAR	RESB	4000	
		.SHE	EQU	TAR	
19	00205f	DAS	EQU	TAR-JLOOP	
		.ASDA	EQU *		
20	00205f	D	EQU	TAR-TAR	
21	00205f	YY	EQU	3*TAR-3*TAR+100-40	
22	00205f		JGT	JSWAP	
23	002062		J	BREAK	
24	002065	JSWAP	JSUB	SWAP	
		.			
25	002068	BREAK	LDX	LNKTH	
26	00206b		LDT	=X'1234'	
27	00206e		SUBR	X,T	
28	002070		LDX	JJJ	
29	002073		TIXR	T	
30	002075		STX	JJJ	
31	002078		JLT	JLOOP	
		.			
32	00207b		LDX	III	
33	00207e		TIX	LNKTH	
34	002081		STX	III	
35	002084		JLT	ILOOP	
36	002087		J	HALT	
		.			
37	00208a	SWAP	LDX	JJJ	

```

37 00208a SWAP LDX JJJ
38 00208d LDCH STR,X
39 002090 STCH TEMP
40 002093 LDCH STR+1,X
41 002096 STCH STR,X
42 002099 LDCH TEMP
43 00209c STCH STR+1,X
44 00209f RSUB
45 0020a2 J *
46 0020a5 STR BYTE C'3142'
47 0020a9 III WORD 0
48 0020ac JJJ WORD 0
49 0020af LENGTH WORD 4
50 0020b2 LENGTH2 WORD 3
51 0020b5 TEMP RESB 1
52 0020b6 END
0020B6 =C'ASDA'
0020BA =X'1234'
>> Incomplete p a s s 1
>> ****
>> ****
>> s y m b o l t a b l e (values in Hexa.)
Name value Absol/Reloc
-----
BREAK 2068 Relocatable
D 0 Absolute
DAS 2b Absolute
HALT 20a2 Relocatable
III 20a9 Relocatable
ILOOP 108e Relocatable
JJJ 20ac Relocatable
JLOOP 1094 Relocatable
JSWAP 2065 Relocatable
LDAF 1003 Relocatable
LENGTH 20af Relocatable
LENGTH2 20b2 Relocatable
ONE 100c Relocatable
STR 20a5 Relocatable
SWAP 208a Relocatable
TAR 10bf Relocatable
TEMP 20b5 Relocatable
THREE 109c Relocatable
TWO 102a Relocatable
YY 3c Absolute

J HALT
SWAP LDX j j j
LDCH STR,X
STCH TEMhP
LDCH STR+1,X
STCH STR,X
LDCH TEMP
STCH STR+1,X
RSUB

REL C N = 1 I = 1 X = 0 B = 0 P = 0 E = 0
002093 STCH TEMHP
Error:*** undefined symbol in operand!!
REL C N = 1 I = 1 X = 1 B = 0 P = 1 E = 0
002096 LDCH STR+1,X ..... 53A010
REL C N = 1 I = 1 X = 1 B = 0 P = 1 E = 0
002099 STCH STR,X ..... 57A00C
REL C N = 1 I = 1 X = 0 B = 0 P = 1 E = 0

```

track x5 int x5 sant xInve xPuzz xInter xinter xinter xS Subs xNew Ta x[C++ x[C++ x[C++ xDiff x

www.diffchecker.com/diff

The two files are identical

1 HBUBBLE^001000^000076
2 T001000^1D^0100007500000500006D000005000013205A07205753A04EAC0453A048
3 T00101D^1E^A0043720033F20034B20220720447500019415072039B850132034382FD7
4 T00103B^1E^07202B2F202E1320253B2F0C3F201807201F53A01557202253A01057A00C
5 T001059^1C^53201957A0074F00003F2FFD333134320000000000000004000003
6 E001000
7

ORIGINAL TEXT

1 HBUBBLE^001000^000076
2 T001000^1D^0100007500000500006D000005000013205A07205753A04EAC0453A048
3 T00101D^1E^A0043720033F20034B20220720447500019415072039B850132034382FD7
4 T00103B^1E^07202B2F202E1320253B2F0C3F201807201F53A01557202253A01057A00C
5 T001059^1C^53201957A0074F00003F2FFD333134320000000000000004000003
6 E001000
7

CHANGED TEXT

Dont store diff

Find Difference!

slkjnushi
Cancelled

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11:15 PM
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