
<Ex14>: Fixed-Point Implementation

Objective: Write a program to implement Fixed-Point Implementation using EPB_C5515 target board

Workflows you learned in the **previous lab**

- Connecting your DSP kit EPB_C5515 to CCS5.3
- Creating a new project or copying an existing project into the workspace
- Configuring the linker options and file-search paths
- Building/Compiling and running/Executing a project on the kit EPB_C5515
- Making use of breakpoints for debugging the code and using watch window to track variable values.
- Storing the data read from PC to the buffer in the CCS5.3 and view buffer data runtime.
- Generating new sine wave signal from lookup table
- Re-generating sine wave using loop-back method
- Controlling Gain of encoder for AIC3204 audio codec
- Profiling for FIR filter using linear buffering by C language and assembly language
- Profiling for FIR filter using Circular buffering by C language and assembly language

After reading **this section** you will be able to,

- Understand Fixed-Point Implementation using EPB_C5515 kit.

Part List:

- PC
- Code Composer Studio
- +5v DC Power supply
- EPB_C5515
- Emulator + Emulator cable (USB A to Mini-A Cable, 14 pin FRC Flat cable)

List of Files Required:

- fixed_float.c (Program application main.c file)
- lnkx.cmd (Command file)
- usbstk5515bsl.lib (Library file)

Lab Goals:

- Understand Fixed-point representation of numbers and Fixed-point computations
- Compare Fixed-point and Floating-point implementations of an FIR Filter in terms of cycle count and accuracy

1. Introduction:

TMS320C5515 is a Fixed-point processor. Fixed-point arithmetic is generally used when hardware resources are limited and we can afford a reduction in accuracy in return for higher execution speed. Fixed-point processors are either 16-bit or 24-bit devices, while floating point processors are usually 32-bit devices. A typical 16-bit processor such as the TMS320C55x, stores data as a 16-bit integer or a fraction format in a Fixed range. Although signals are only stored with 16-bit precision, intermediate values during the arithmetic operations may be kept at 32-bit precision. This is done using the internal 40-bit accumulator, which reduces cumulative

round-off errors. Fixed-point DSP devices are usually cheaper and faster than their Floating-point counterparts because they use less silicon, have lower power consumption, and require fewer external pins. Most high volume low cost embedded applications, such as appliance control, cellular phones, hard disk drives, modems, audio players, and digital cameras use Fixed-point processors.

Floating-point arithmetic greatly expands the dynamic range of numbers. A typical 32-bit Floating point DSP processor, such as the TMS320C67x, represents number with a 24-bit mantissa and an 8-bit exponent. The mantissa represents a fraction in the range -1.0 to +1.0, while the exponent is an integer that represents the number of places that the binary point must be shifted left or right in order to obtain the true value. For example, in decimal number system we have some number like 10.2345. We can write this in the form of 0.102345×10^2 . So in this format, 0.102345 is called mantissa and 2 is called exponent.

A 32-bit Floating-point format covers a large dynamic range, thus the data dynamic range restriction may be virtually ignored in a design using Floating-point DSP processors. But in Fixed-point format, the designer has to apply scaling factors and other techniques to prevent arithmetic overflow. This is usually a difficult and time consuming process. As a result, Floating-point DSP processors are generally easy to program and use, but are more expensive and have higher power consumption. In this session, you will learn the issues related to Fixed-point arithmetic

Two's complement Integer Representation

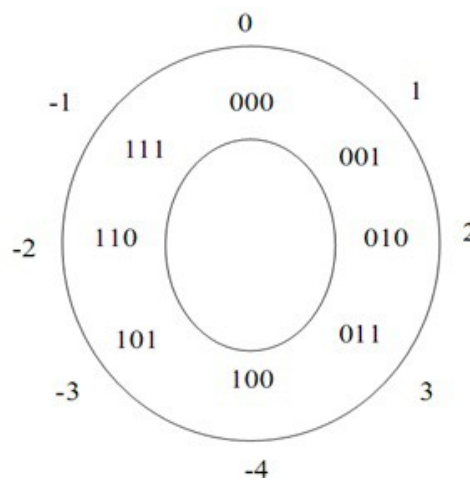
This representation is most common method for representing a signed integer. The two's complement of a binary number is defined as the value obtained by subtracting the number from a large power of two (specifically, from 2^N for an N-bit two's complement). As far as the hardware is concerned, Fixed-point number systems represent data as B-bit integers.

The two's complement number system usually used is

$$k = \begin{cases} \text{binary integer representation,} & \text{if } 0 \leq k \leq 2^{B-1} - 1 \\ \text{bitwise complement of } k + 1, & \text{if } -(2^{B-1}) \leq k \leq -1 \end{cases}$$

The most significant bit is known as the sign bit. It is 0 when the number is non-negative and 1 when the number is negative.

Below Figure 1 is an easy way to visualize two's complement representation.



Fractional Fixed-point number representation

For the purposes of signal processing, we often regard the Fixed-point numbers as binary fractions in the range $[-1, 1)$, by implicitly placing a decimal point after the sign bit. Fixed-point representation of a fractional number x is illustrated in Figure 2 below

$$x = b_0 . b_1 b_2 \cdots b_{M-1} b_M$$

The word-length is $B(= M + 1)$ bits, i.e. M magnitude bits and one sign bit. The most significant bit (MSB) is the sign bit, which represents the sign of the number as follows.

$$b_0 = \begin{cases} 0, & \text{if } x \geq 0 \\ 1, & \text{otherwise} \end{cases}$$

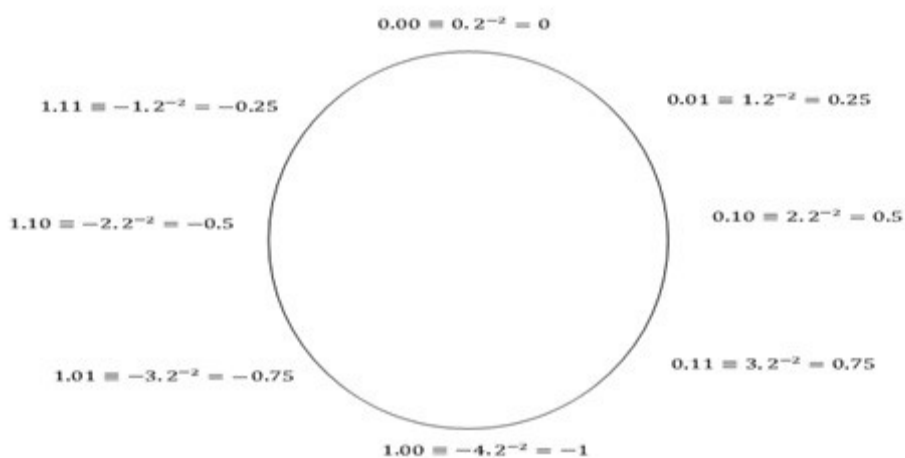
The remaining M bits give the magnitude of the number. The rightmost bit b_M is called the least significant bit (LSB), which represents precision of the number.

The decimal value corresponding to a binary fractional number x can be expressed as,

$$x = b_0 + b_1 2^{-1} + b_2 2^{-2} + \dots + b_M 2^{-M}$$

$$x = b_0 + \sum_{m=1}^M b_m 2^{-m}$$

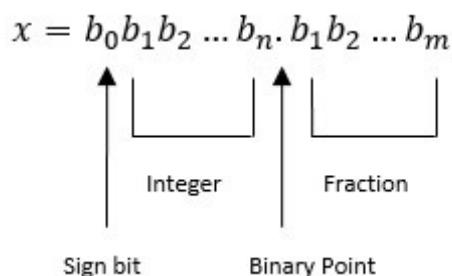
Below figure 3 provides a visual representation for a 3 bit fractional representation.



For example, from Figure 2, the easiest way to convert a normalized 16-bit fractional binary number into an integer that can be used by C55x assembler is to move the binary point to the right by 15 bits (at the right of b_M). Since shifting the binary point 1 bit right is equivalent to multiplying the fractional number by 2, moving the binary point to the right by 15 bits can be done by multiplying the decimal value by $2^{15} = 32768$.

Q-format

Q-format is a formal mechanism to keep track of radix or Fixed point. Format Q_{nm} is illustrated in Figure 4 below, where $m = M = B - 1$. There are n bits to the left of the binary point representing integer portion, and m bits to the right, representing fractional value.



The most popular fractional number format is Q0.15 format ($n = 0$ and $m = 15$), which is simply referred to as Q15 format since there are 15 fractional bits.

Integer Word-Length (IWL)

The number of bits required to express the integer part of a Floating point number is called as the integer word-length of that number. IWL of a number gives us an idea of the maximum precision with which it can be represented in a finite sized shift register. For example, the IWL of 56:25 is 6 (Why?). Therefore, it can be stored in a 16 bit register just by using the 6 LSBs. But, this results in wastage of the 9 MSBs which all are zeros. Thus, to store 56:25 in a 16 bit register, we will have to multiply it by 2^{15-6} , i.e., $2^9 = 512$, so that it occupies the whole 16 bit register. Note that, the decimal point is implicit while storing the Floating point number. To understand this better, carry out the multiplication yourself and represent the product in signed binary format. Compare this representation with that of 56. This concept of IWL and its use for effective storage is used in section 5.

Round-off Error in Fixed Point Implementation

Fixed-point arithmetic is often used with DSP hardware for real-time processing because it offers fast operation and relatively economical implementation. Its drawbacks include a small dynamic range and low resolution. A Fixed-point representation also leads to arithmetic errors, which if not handled correctly will lead to erroneous operations. These errors have to be handled in such a way as to have some trade-off among accuracy, speed and cost.

Consider the multiplication of two binary fractions, as shown here.

	Fractional Interpretation	Integer Interpretation
0.10	1/2	2
x 0.11	x 3/4	x 3
<u>. 010</u>		
.010		
<u>0.00</u>		
0.0110	3/8	6

From above calculation, we see that full-precision multiplication almost doubles the number of bits. If we wish to return the product to a b-bit representation, we must truncate the (b - 1) least significant bits. However this introduces truncation error. Truncation error is also known as round off error (the number is rounded to the nearest b-bit fractional value rather than truncated). This type of error occurs after the multiplication.

Example code:

Go through below example. Two arrays x floating and y floating of size 100 each, are defined globally.

We want to perform the operation:

y floating[i] = Ax floating[i] + B ;where A = 3:2 and B = 2:7.

The function compute floating() carries out this operation with the help of floating point variables. On the other hand, the compute fixed() deals with only the finite precision data-types such as short int, long int etc. It demonstrates the steps involved in the conversion of floating-point computations to fixed-point. This code is self explanatory.

For better understanding, go to:

https://www.youtube.com/watch?v=om7Bvk7WzJg&list=PLV5Sz2W8vMyagrltexjFRw2X_X99qpX6v

This is another method of looking at the calculations implemented in compute fixed.

Try:

- Function compute_floating() computes y[i] = Ax[i]+B using regular floating-point arithmetic. Firstly, call only this function in main and perform profiling to record the clock-cycle count of its execution.
- Repeat this exercise with compute_fixed() function (Call only one of the two functions at a time in main).

Steps for importing new project:

Open CCS V5.3 from desktop shortcut



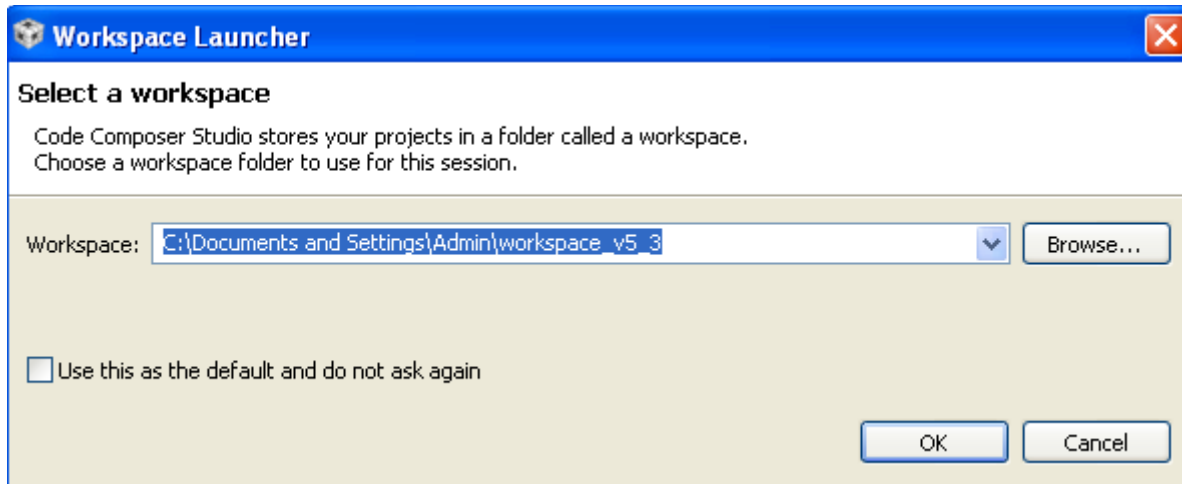
It will open default CCS V5 screen.



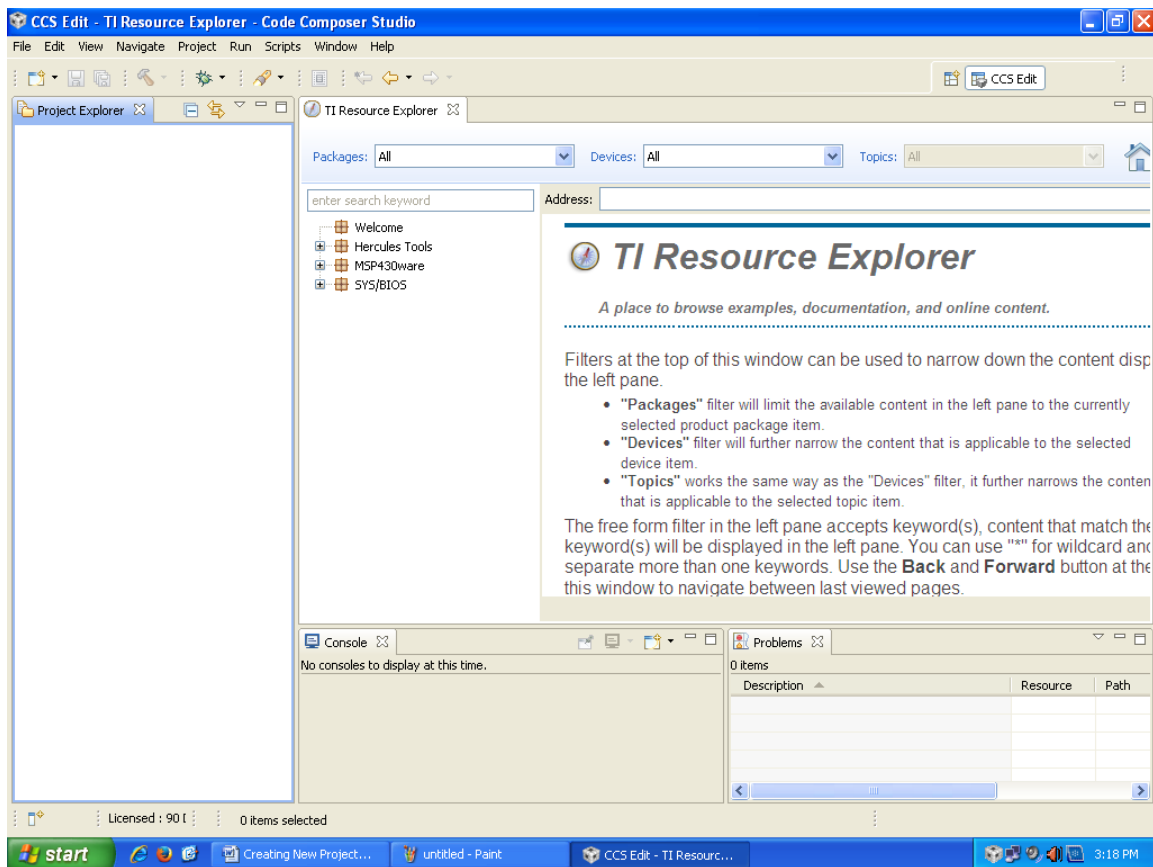
Then it will ask for workspace path

Select path "*C:\Documents and Settings\<User Name>\workspace_v5_3*" for windows XP OS

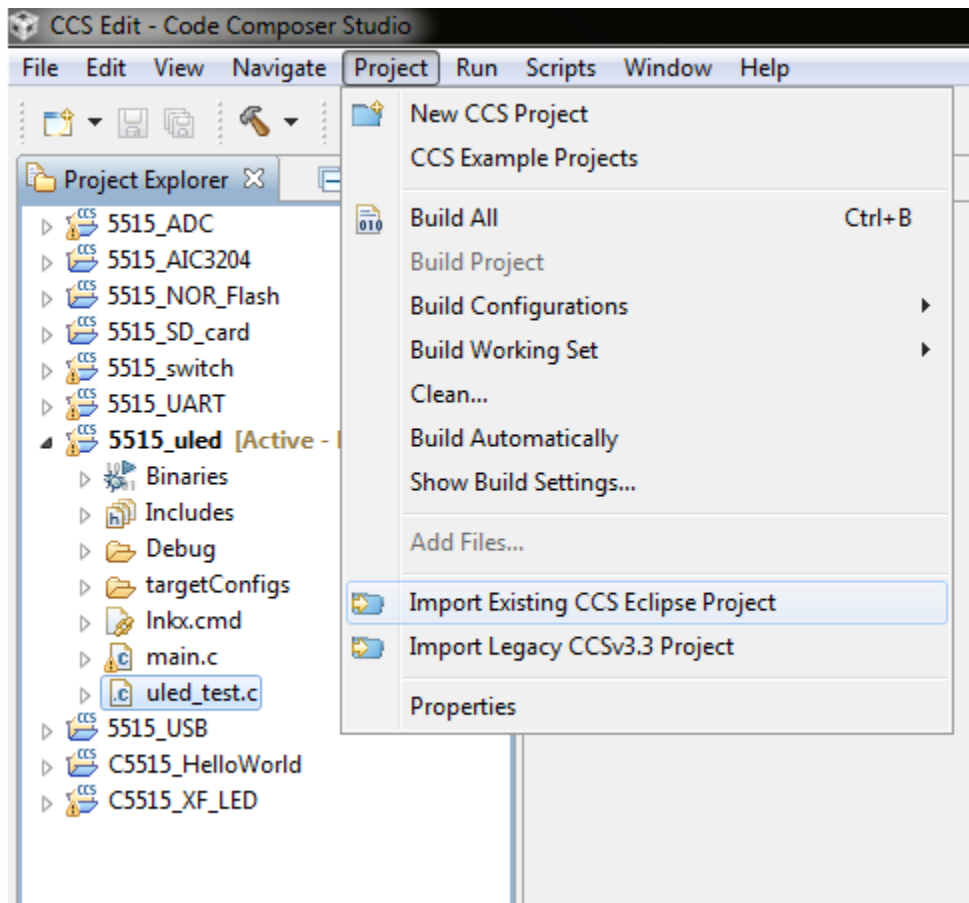
Select path "*C:\Users\<User Name>\workspace_v5_3*" for windows7 OS



Then it will open Default CCS5 screen as shown below



Click “*Project -> Import Existing CCS Eclipse Project*” menu.



-
- Now Copy the project path and paste it to “Select search-directory” location and press enter. Or you can browse the path also by clicking browse.
 - Select project “IIT_LAB5.1” and also select “Copy projects into workspace” and **Finish**
 - Here project is already imported and it can be seen from “**project explorer**”
 - Compile the program by “*right click-> build project*” or “*right click-> rebuild project*” as shown. It will generate “IIT_ LAB5.1.out” file in “*debug*” folder.
 - It will generate IIT_ LAB5.1.out file in “*debug*” folder.

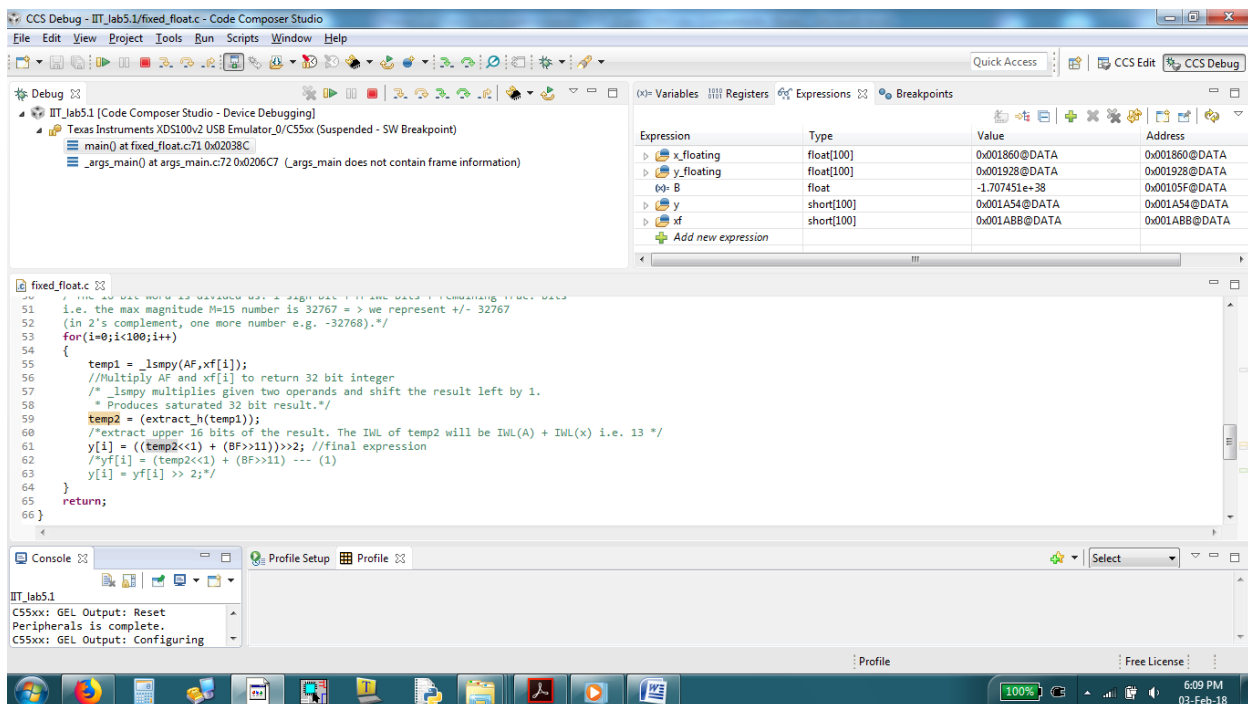
How to Run Program:

Hardware Connection:

- Power on EPB_C5515 hardware using +5V Power supply **or** USB A-to-B cable
- Connect XDS100V2 with EPB_C5515 using USB A-to-miniA cable with CPU
- Reset DSP kit EPB_C5515 by pressing RESET switch.

Steps to Run Program:

- Now to debug the program click “**debug**” icon **OR** from “**run->debug**” menu.
- It will configure/connect EPB_C5515 kit with the CCSV5 using XDS100V2 and download the program in C5515 CPU. It will be done automatically.



Apply profiling for “compute_floating(A,B)” and “compute_fixed()” function.
Apply profiling steps as per the previous lab-4 and check the result for
“compute_floating(A,B)” as shown below

The screenshot shows the CCS Debug interface with the 'Profile Setup' window open. The 'Profile' tab is selected, displaying a table of profiling results for the function 'compute_floating(float, float)'. The table includes columns for Name, Calls, Excl Count Min, Excl Count Max, Excl Count Average, Excl Count Total, Incl Count Min, Incl Count Max, Incl Count Average, Incl Count Total, and Filename.

Name	Calls	Excl Count Min	Excl Count Max	Excl Count Average	Excl Count Total	Incl Count Min	Incl Count Max	Incl Count Average	Incl Count Total	Filename
1 compute_floating(float, float)	1	59126	59126	59126.00	59126	59126	59126	59126.00	59126	C:\User...
2 main()	1	-	-	21.00	21	-	-	59151.00	59151	C:\User...
3 USBSTK3515_init()	1	4	4	4.00	4	4	4	4.00	4	C:\User...

The console window shows the following output:

```

C55xx: GEL Output: Reset
Peripherals is complete.
C55xx: GEL Output: Configuring
PLL (100 MHz).
C55xx: GEL Output: PLL Init
Done.C55xx: GEL Output: Target
Connection Complete.
  
```

Apply profiling steps as per the previous lab-4 and check the result for
“compute_fixed()” as shown below

The screenshot shows the CCS Debug interface with the 'Profile Setup' window open. The 'Profile' tab is selected, displaying a table of profiling results for the function 'compute_fixed()'. The table includes columns for Name, Calls, Excl Count Min, Excl Count Max, Excl Count Average, Excl Count Total, Incl Count Min, Incl Count Max, Incl Count Average, Incl Count Total, and Filename.

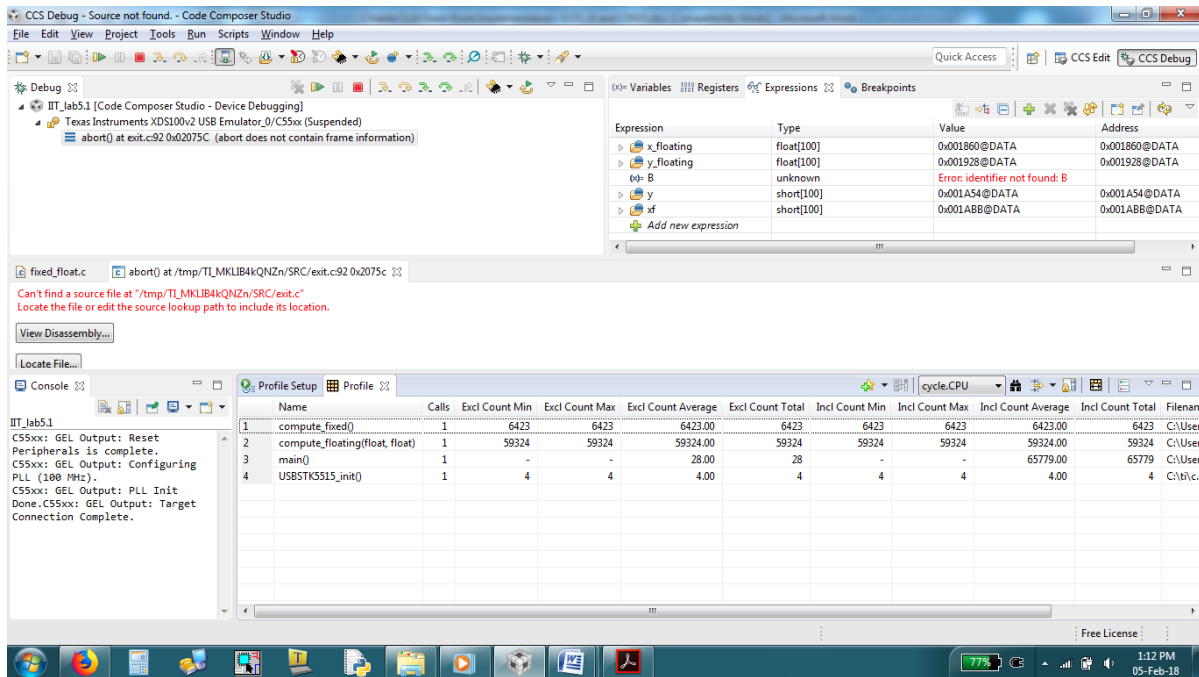
Name	Calls	Excl Count Min	Excl Count Max	Excl Count Average	Excl Count Total	Incl Count Min	Incl Count Max	Incl Count Average	Incl Count Total	Filename
1 compute_fixed()	1	6423	6423	6423.00	6423	6423	6423	6423.00	6423	C:\User...
2 main()	1	-	-	18.00	18	-	-	6445.00	6445	C:\User...
3 USBSTK3515_init()	1	4	4	4.00	4	4	4	4.00	4	C:\User...

The console window shows the following output:

```

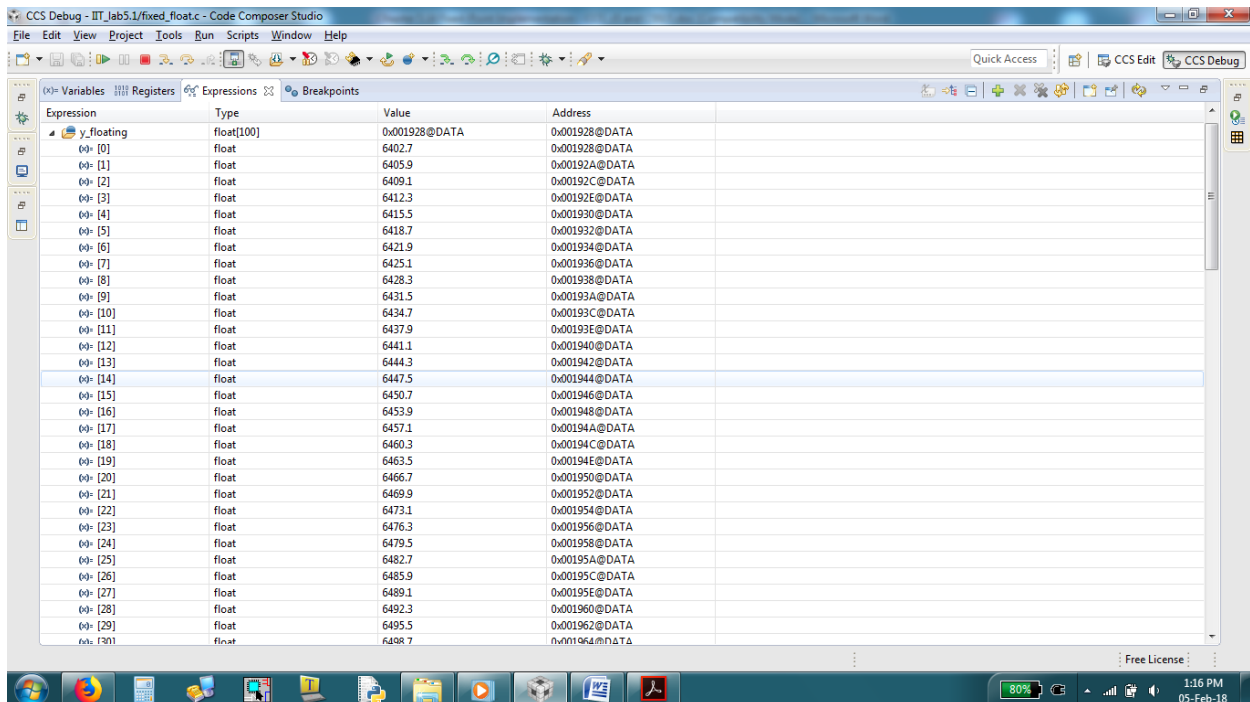
C55xx: GEL Output: Reset
Peripherals is complete.
C55xx: GEL Output: Configuring
PLL (100 MHz).
C55xx: GEL Output: PLL Init
Done.C55xx: GEL Output: Target
Connection Complete.
  
```

Combine results are as shown here.

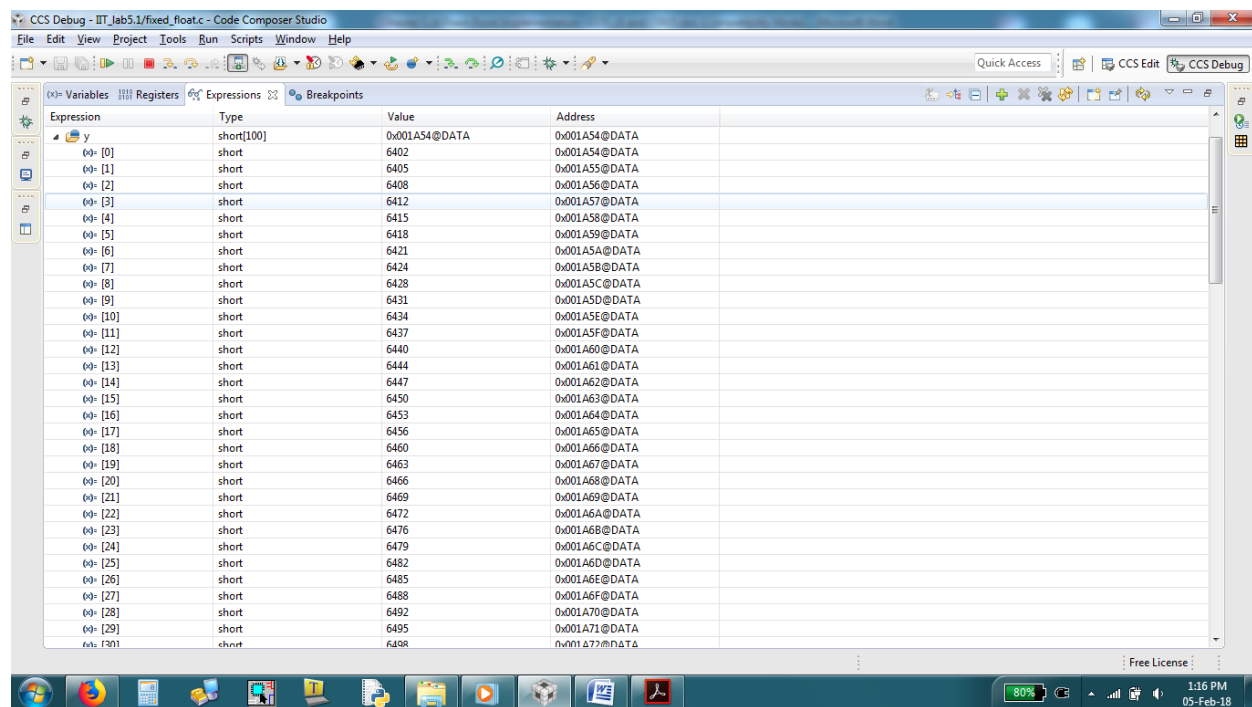


At this time check the variable “y_floating” and “y” in watch window and compare results

- y_floating



- y

**Exercise:**

Modify the given C code (compute floating()) with the values of x floating[i] taking floating point values between 4000 and 4100 (as specified below) and convert this code to fixed point inside the function compute fixed(). The floating point values may be assigned to x floating[] as shown in the following code snippet.

```
for (i=4000; i<4100; i++)
{
    x_floating[i-4000] = i + 0.875;
}
```

So, instead of [2000; 2099], x[i] will take values from the range [4000:875; 4099:875] and you will have to calculate $y[i] = Ax[i] + B$. Note that, since x floating[i] itself contains floating point values, we cannot use the trick of calculating xf[i] by just bitwise right shifting the values in x[i]. You will have to actually multiply the array x floating[i] by correct value to get xf[i].