

Experiment No: 12

AIM: Implementation of sum of subset problem using Backtracking

THEORY:

Given n distinct positive numbers (called weights), we need to find all combinations of these numbers whose sums are m . This is called the *sum of subsets problem*.

The problem is efficiently solved in programming using a concept called Backtracking.

Backtracking is an algorithmic technique for solving problems recursively by trying to build a solution incrementally, one piece at a time, removing those solutions that fail to satisfy the constraints of the problem at any point of time.

The following bounding functions are used:

1. $\sum_{i=1}^k w_i x_i + \sum_{i=k+1}^n w_i \geq m$
2. $\sum_{i=1}^k w_i x_i + w_{k+1} \leq m$

ALGORITHM:

Algorithm SumOfSub(s,k,r)

// Find all subsets of $w[1:n]$ that sum to m . The values of $x[j]$,

// $1 \leq j < k$ have already been determined.

// The $w[j]$'s are in nondecreasing order.

// It is assumed that $w[1] \leq m$

```
{  
    // Generate left child. Note:  $s+w[k] < m$  since  $B_{k-1}$  is true.  
     $x[k]:=1$ ;  
    if ( $s+w[k] = m$ ) then write ( $x[1:k]$ ); // Subset found  
    // There is no recursive call here as  $w[j] > 0$ ,  $1 \leq j \leq n$ .  
  
    else if ( $s+w[k] + w[k+1] \leq m$ )  
        then SumOfSub( $s+w[k]$ ,  $k+1$ ,  $r-w[k]$ );  
  
    // Generate right child and evaluate  $B_k$ .
```

```

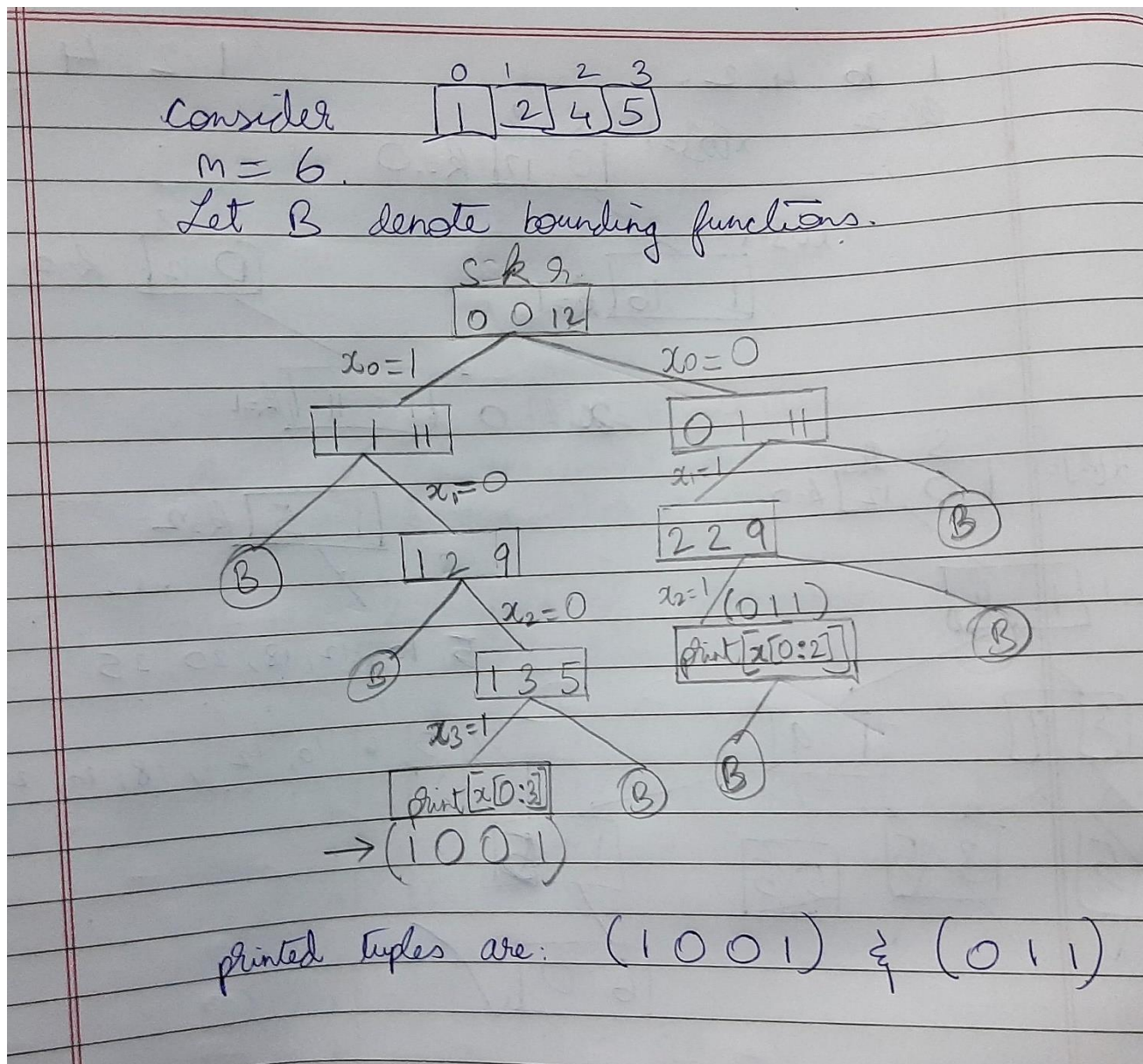
if ((s + r - w[k] ≥ m) and (s + wk+1 ≤ m)) then
{
    x[k]=0;
    SumOfSub(s,k+1,r-w[k]);
}
}

```

Time Complexity

- The time complexity of this algorithm is $O(2^n)$.

Problem Tracing



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PROGRAM IMPLEMENTATION:

```
#include<iostream>
```

```
using namespace std;
```

```
int n,m,*w,*x;    //soln vector
```

```
void show_set(int k)
```

```
{  
    cout<<" ";  
    for(int i=0;i<=k;i++)  
        cout<<x[i]<<" ";  
    cout<<" "<<endl;
```

```
}
```

```
void sum_of_sub(int s, int k, int r)
```

```
{  
    x[k] = 1;  
  
    if(s+w[k] == m)  
        show_set(k);  
  
    else if(s + w[k] + w[k+1] <=m)    //current sum is less than or equal to m  
        sum_of_sub(s+w[k],k+1,r-w[k]);  
  
    if( (s + r - w[k] >= m) && (s + w[k+1] <=m))    //bounding function says that
```

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```
        {  
            //w obtained so far + remaining w  
            x[k] = 0;                                //must  
            be >=m  
            sum_of_sub(s,k+1,r-w[k]);  
        }  
        //and the w obtained so far + next  
  
        //w must be <=m  
    }
```

```
int main()  
{  
    int r=0;          //temp sum  
    cout<<"Enter number of elements:\n";  
    cin>>n;  
  
    cout<<"Enter "<<n<<" elements: ";  
  
    w=new int[n];  
  
    for(int i=0;i<n;i++)  
        cin>>w[i],r+=w[i];  
  
    x = new int[n];
```

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```
cout<<"\nEnter the sum: ";
```

```
cin>>m;
```

```
for(int i=0;i<n;i++)
```

```
    x[i]=0;
```

```
int s=0,k=0;
```

```
sum_of_sub(s,k,r);
```

```
return 0;
```

```
}
```

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OUTPUTS:

1. When $n=6$

```
C:\WINDOWS\SYSTEM32\cmd.exe
Enter number of elements: 6
Enter 6 elements: 5 10 12 13 15 18
Enter the sum: 30
( 1 1 0 0 1 )
( 1 0 1 1 )
( 0 0 1 0 0 1 )
```

2. When $n=8$

```
C:\WINDOWS\SYSTEM32\cmd.exe
Enter number of elements: 8
Enter 8 elements: 3 5 6 10 12 18 20 25
Enter the sum: 30
( 1 1 0 1 1 )
( 0 1 0 0 0 0 0 1 )
( 0 0 0 1 0 0 1 )
( 0 0 0 0 1 1 )
```

Conclusion:

- This algorithm works when the elements are arranged in non-decreasing order of their weights.
- Time complexity of the algorithm is $O(2^n)$