

# Theory-of-Languages-and-Machines

Shahla fathollahi

Introduction to Automata Theory, Formal Languages and Computation  
Shyamalendu Kandar

## Example 8.18

*Construct a TM over  $\{a, b\}$  which contains a substring  $abb$ .*

**Solution:** In regular expression, it can be written as  $(a, b)^*abb(a, b)^*$ . In this expression, the substring is important. The string may be  $abb$  only. In that case, the machine gets 'a' as input in the beginning state. Before traversing the substring 'abb', there is a chance to traverse 'a' of  $\{a, b\}^*$ . The transitional functions are

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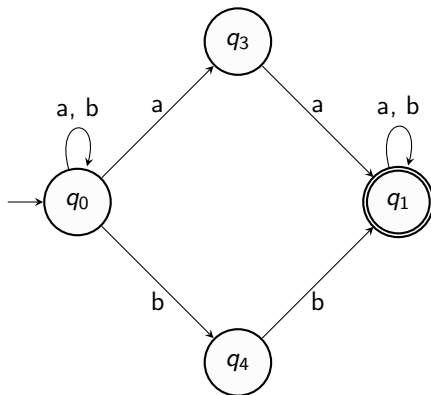
$$\begin{aligned}\sigma(q_1, a) &\rightarrow (q_1, a, R), (q_2, a, R) \\ \sigma(q_1, b) &\rightarrow (q_1, b, R) \\ \sigma(q_2, b) &\rightarrow (q_3, b, R) \\ \sigma(q_3, b) &\rightarrow (q_4, b, R) \\ \sigma(q_4, a) &\rightarrow (q_4, a, R) \\ \sigma(q_4, b) &\rightarrow (q_4, b, R) \\ \sigma(q_4, B) &\rightarrow (q_f, B, H)\end{aligned}$$

In state  $q_1$  for input 'a', there are two transitional functions. So, it is a non-deterministic TM.

## Example 8.19

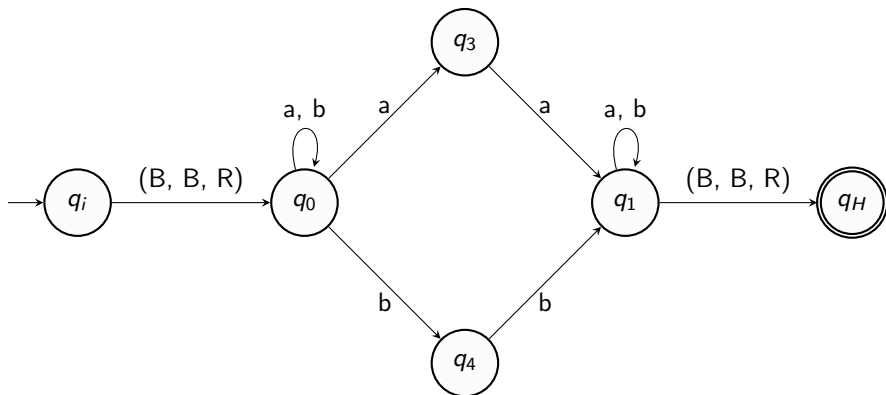
Construct a TM for the regular expression  $(a + b)^*(aa + bb)(a + b)^*$ .

**Solution:** The finite automata accepting the regular expression is



## Example 8.19

**Continue the solution:** Inserting a new initial state  $q_i$  and final state  $q_H$ , the finite automata becomes



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**Continue the solution:** Converting the levels of the inputs, the TM becomes

