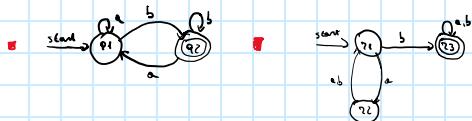


# ESEMPIO - LINGUAGGI REGOLARI CHIUSI PER UNIONE

martedì 28 marzo 2023 09:46

Supponiamo di avere due DFA

- $M_1$  riconosce  $L_1 \Rightarrow L(M_1)$
- $M_2$  riconosce  $L_2 \Rightarrow L(M_2)$



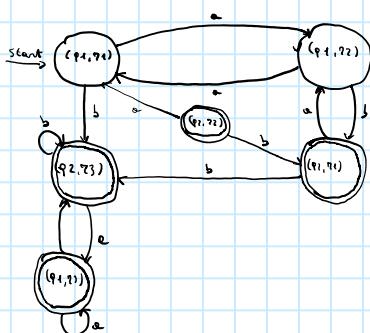
Supponiamo di voler costruire un DFA  $M_3$  che riconosce l'unione  $L(M_1) \cup L(M_2)$

4) La prima cosa da fare è definire l'insieme di stati, dato dal prodotto cartesiano

$$B_3 = \{(q_1, q_2), (q_1, q_3), (q_1, q_1), (q_2, q_2), (q_2, q_3), (q_3, q_3)\}$$

5) Poniamo parola al disegno

$$\begin{aligned} - \delta_3((q_1, q_2), a) &= (\delta_1(q_1, a), \delta_2(q_2, a)) = (q_1, q_2) \\ - \delta_3((q_1, q_2), b) &= (\delta_1(q_1, b), \delta_2(q_2, b)) = (q_2, q_3) \\ - \delta_3((q_1, q_1), a) &= (\delta_1(q_1, a), \delta_2(q_1, a)) = (q_1, q_1) \\ - \delta_3((q_1, q_2), b) &= (\delta_1(q_1, b), \delta_2(q_2, b)) = (q_2, q_1) \\ - \delta_3((q_2, q_2), a) &= (\delta_1(q_2, a), \delta_2(q_2, a)) = (q_2, q_2) \\ - \delta_3((q_2, q_2), b) &= (\delta_1(q_2, b), \delta_2(q_2, b)) = (q_3, q_3) \\ - \delta_3((q_1, q_3), a) &= (\delta_1(q_1, a), \delta_2(q_3, a)) = (q_1, q_3) \\ - \delta_3((q_1, q_3), b) &= (\delta_1(q_1, b), \delta_2(q_3, b)) = (q_2, q_1) \\ - \delta_3((q_1, q_1), b) &= (\delta_1(q_1, b), \delta_2(q_1, b)) = (q_1, q_1) \\ - \delta_3((q_2, q_3), a) &= (\delta_1(q_2, a), \delta_2(q_3, a)) = (q_2, q_3) \\ - \delta_3((q_2, q_3), b) &= (\delta_1(q_2, b), \delta_2(q_3, b)) = (q_3, q_2) \\ - \delta_3((q_1, q_1), a) &= (\delta_1(q_1, a), \delta_2(q_1, a)) = (q_1, q_1) \\ - \delta_3((q_2, q_1), b) &= (\delta_1(q_2, b), \delta_2(q_1, b)) = (q_2, q_1) \end{aligned}$$



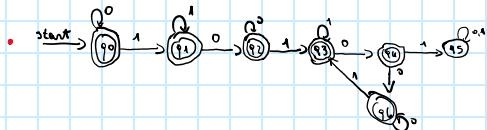
6) Ili stati eccellenti:  $F_3 = \{(q_2, q_3), (q_2, q_1), (q_1, q_3), (q_1, q_1)\}$

# ESERCITAZIONE - DFA

martedì 28 marzo 2023 14:36

#1

Scrivere un DFA che accetta tutte le stringhe su  $\Sigma = \{0, 1\}$  che contengono al più un'occorrenza di "01"



#2

Dire, fornire un DFA che accetta almeno un'occorrenza di "01"



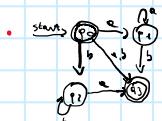
#3

Un DFA che contiene un'occorrenza di "01" come suffisso



#4

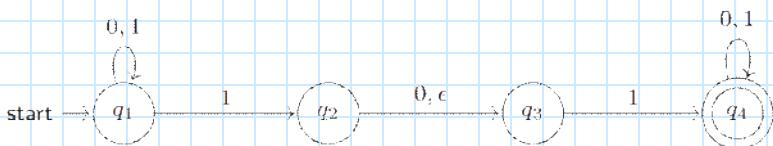
Progettare un DFA che accetta tutte le stringhe che terminano con un simbolo che non compare in nessuna altre posizione di "w"  
 $a^nb - b^na$



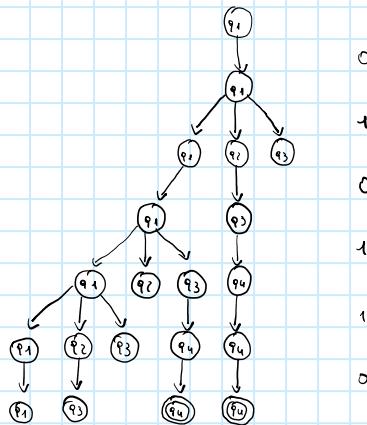
#5

# ESEMPIO - ALBERO DI COMPUTAZIONE DI UN NFA

mercoledì 29 marzo 2023 10:08



Vogliamo controllare se la stringa  $0+0110$  è accettata, quindi procediamo con l'albero delle computazioni.



In questo caso, la stringa è accettata.

# ESERCITAZIONE - CONVERSIONE DA NFA A DFA

giovedì 30 marzo 2023 11:51

Trasformare il seguente NFA in DFA



$\delta$	0	1	$\epsilon$
$q_0$	$\{q_0\}$	$\{q_0\}$	$\{q_1\}$
$q_1$	$\emptyset$	$\emptyset$	$\emptyset$
$q_2$	$\emptyset$	$\emptyset$	$\emptyset$

$$Q_M = \{\emptyset, \{q_0\}, \{q_1\}, \{q_2\}, \{q_0, q_1\}, \{q_0, q_2\}, \{q_1, q_2\}, \{q_0, q_1, q_2\}\}$$

$$\delta_M(\{q_0\}, 0) = E(\delta_N(\{q_0\}, 0)) = \{q_0\}$$

$$\delta_M(\{q_0\}, 1) = E(\delta_N(\{q_0\}, 1)) = \{q_0\}$$

$$\delta_M(\{q_1\}, 0) = E(\delta_N(\{q_1\}, 0)) = \{q_2\}$$

$$\delta_M(\{q_1\}, 1) = E(\delta_N(\{q_1\}, 1)) = \{\emptyset\}$$

$$\delta_M(\{q_2\}, 0) = E(\delta_N(\{q_2\}, 0)) = \{\emptyset\}$$

$$\delta_M(\{q_2\}, 1) = E(\delta_N(\{q_2\}, 1)) = \{q_0\}$$

$$\delta_M(\{q_0, q_1\}, 0) = E(\delta_N(\{q_0, q_1\}, 0)) \cup E(\delta_N(\{q_1, q_0\}, 0)) = \{q_0\} \cup \{q_1\} = \{q_0, q_1\}$$

$$\delta_M(\{q_0, q_1\}, 1) = E(\delta_N(\{q_0, q_1\}, 1)) \cup E(\delta_N(\{q_1, q_0\}, 1)) = \{q_0\} \cup \emptyset = \{q_0\}$$

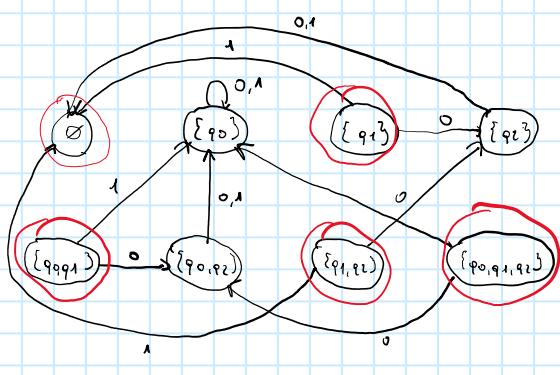
$$\delta_M(\{q_0, q_2\}, 0) = E(\delta_N(\{q_0, q_2\}, 0)) \cup E(\delta_N(\{q_2, q_0\}, 0)) = \{q_0\} \cup \{\emptyset\} = \{q_0\}$$

$$\delta_M(\{q_0, q_2\}, 1) = E(\delta_N(\{q_0, q_2\}, 1)) \cup E(\delta_N(\{q_2, q_0\}, 1)) = \{q_0\} \cup \{q_1\} = \{q_0\}$$

$$\delta_M(\{q_1, q_2\}, 0) = E(\delta_N(\{q_1, q_2\}, 0)) \cup E(\delta_N(\{q_2, q_1\}, 0)) = \{q_1\} \cup \{q_2\} = \{q_1, q_2\}$$

$$\delta_M(\{q_0, q_1, q_2\}, 0) = E(\delta_N(\{q_0, q_1, q_2\}, 0)) \cup E(\delta_N(\{q_1, q_0, q_2\}, 0)) \cup E(\delta_N(\{q_0, q_2, q_1\}, 0)) = \{q_0\} \cup \{q_1\} \cup \{q_2\} = \{q_0, q_1, q_2\}$$

$$\delta_M(\{q_0, q_1, q_2\}, 1) = E(\delta_N(\{q_0, q_1, q_2\}, 1)) \cup E(\delta_N(\{q_1, q_0, q_2\}, 1)) \cup E(\delta_N(\{q_0, q_2, q_1\}, 1)) = \{q_0\} = \{q_0\}$$



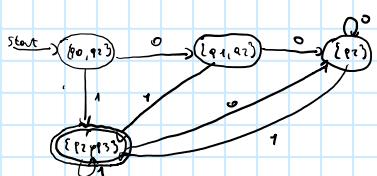
$$\delta_M(\{q_1, q_2\}, 1) = E(\delta_N(q_1, 1)) \cup E(\delta_N(q_2, 1)) = \{q_2, q_3\}$$

$$\delta_M(\{q_2, q_3\}, 0) = E(\delta_N(q_2, 0)) \cup E(\delta_N(q_3, 0)) = \{q_1\}$$

$$\delta_M(\{q_1, q_3\}, 1) = E(\delta_N(q_1, 1)) \cup E(\delta_N(q_3, 1)) = \{q_2, q_3\}$$

$$\delta_M(\{q_2\}, 0) = E(\delta_N(q_2, 0)) = \{q_1\}$$

$$F = \{q_1, q_2\}$$



# ESEMPIO - DA REGEX AD NFA

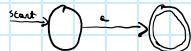
venerdì 31 marzo 2023 15:14

Troviamo un NFA che riconosce il linguaggio descritto dalla REGEX

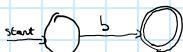
- $(ab \cup a)^*$

Procediamo per passi, individuando di volta in volta un NFA che riconosce un simbolo

- a



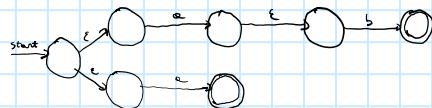
- b



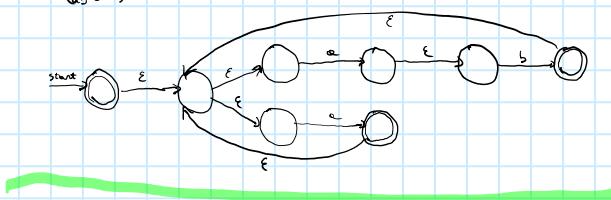
- ab



- $ab \cup a$



- $(ab \cup a)^*$



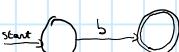
Troviamo un NFA che riconosce il linguaggio descritto dalla REGEX

- $(a \cup b)^* aba$

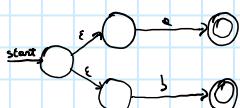
- a



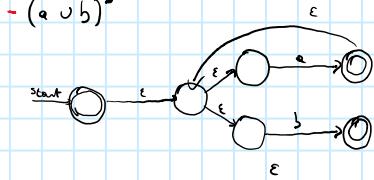
- b



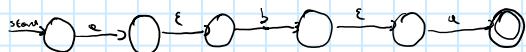
- $(a \cup b)$



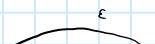
- $(a \cup b)^*$

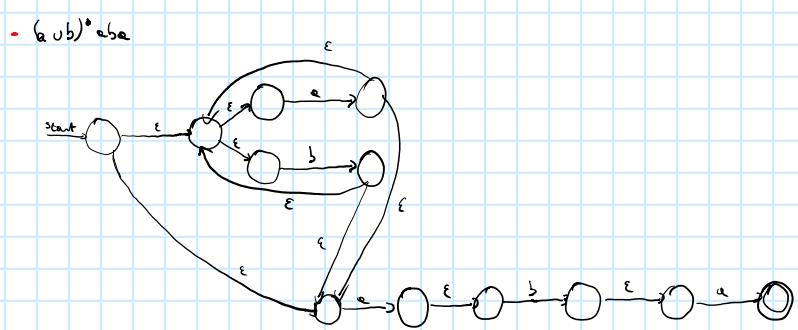


- aba



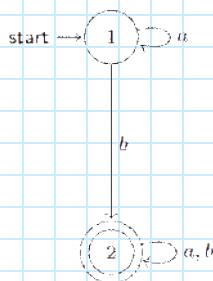
- $(a \cup b)^* aba$



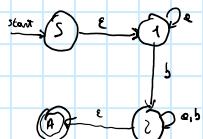


# ESEMPI - DA DFA A REGEX

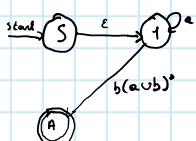
sabato 1 aprile 2023 16:26



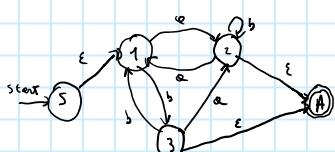
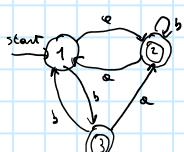
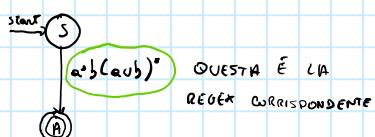
Per prime cose, andiamo ad aggiungere uno stato iniziale e uno finale



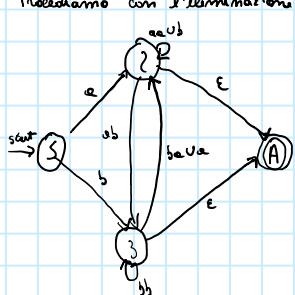
Concentriamoci nell'eliminare 2, conservando tutto le transizioni



Ora eliminiamo anche 1

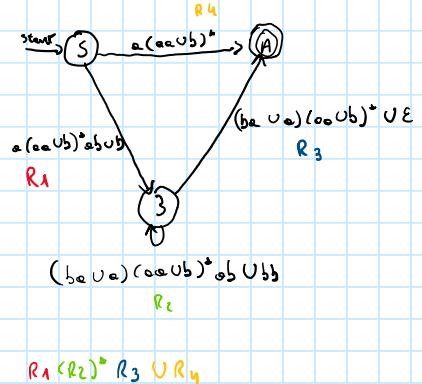


Procediamo con l'eliminazione di 1



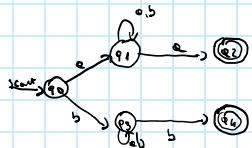
Eliminiamo il 2

Eliminiamo il 2



# ESEMPIO DI PROVA SCRITTA

lunedì 3 aprile 2023 10:17



Dato il seguente NFA, scrivere le definizioni delle quintupla.

$$N = \{Q, \Sigma, \delta, q_0, F\}$$

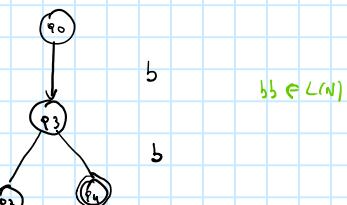
$$Q_N = \{q_0, q_1, q_2, q_3, q_4\}$$

$$\begin{matrix} q_0 = \text{stato iniziale} \\ F = \{q_2, q_3, q_4\} \end{matrix}$$

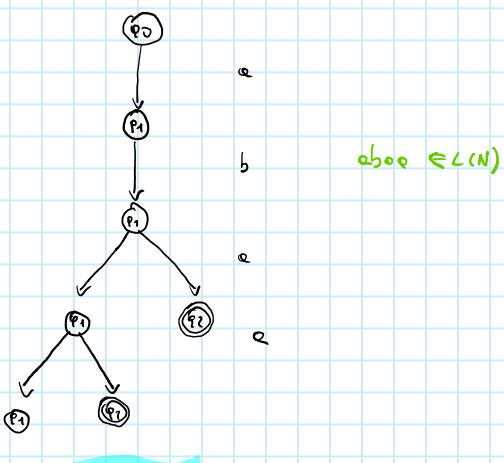
$$\Sigma = \{a, b\}$$

	a	b
$\rightarrow q_0$	{ $q_1, q_3$ }	{ $q_3, q_4$ }
$q_1$	{ $q_1, q_2$ }	{ $q_1$ }
$q_2$	$\emptyset$	$\emptyset$
$q_3$	{ $q_3$ }	{ $q_3, q_4$ }
$q_4$	$\emptyset$	$\emptyset$

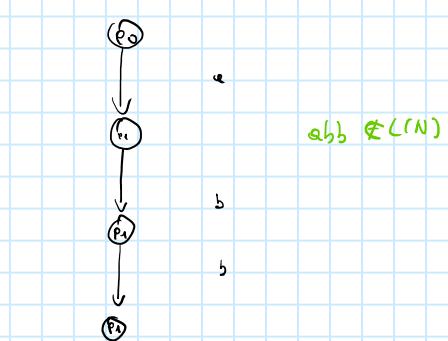
Determinare se le stringhe  $bb$ ,  $abaa$ ,  $abb$  appartengono a  $L(N)$ .



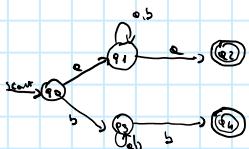
$bb \in L(N)$



$abaa \in L(N)$



$abb \notin L(N)$



Lo stato iniziale è  $q_0$

$$\delta_M = E(\{q_0\}, a) = \delta_N(E(\{q_0\}, a)) = \{q_1\}$$

$$\delta_M = E(\{q_0\}, b) = \delta_N(E(\{q_0\}, b)) = \{q_3\}$$

$$\delta_M = E(\{q_1\}, a) = \delta_N(E(\{q_1\}, a)) = \{q_2\}$$

$$\delta_M = E(\{q_1\}, b) = \delta_N(E(\{q_1\}, b)) = \{q_3\}$$

$$\delta_M = E(\{q_2\}, a) = \delta_N(E(\{q_2\}, a)) = \{q_2\} \cup \{q_3\} = \{q_1, q_2\}$$

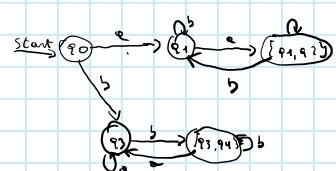
$$\delta_M = E(\{q_2\}, b) = \delta_N(E(\{q_2\}, b)) = \{q_3\} \cup \{q_3\} = \{q_3\}$$

$$\delta_M = E(\{q_3\}, a) = \delta_N(E(\{q_3\}, a)) = \{q_3\}$$

$$\delta_M = E(\{q_3\}, b) = \delta_N(E(\{q_3\}, b)) = \{q_3\}$$

$$\delta_M = E(\{q_3, q_2\}, a) = \delta_N(E(\{q_3, q_2\}, a)) = \{q_3\} \cup \{q_3\} = \{q_3\}$$

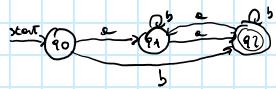
$$\delta_M = E(\{q_3, q_2\}, b) = \delta_N(E(\{q_3, q_2\}, b)) = \{q_3\} \cup \{q_3\} = \{q_3\}$$



# ESERCITAZIONE - DA DFA A REGEX

lunedì 3 aprile 2023 14:07

Convertire il seguente DFA in una REGEX



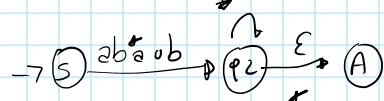
$$(S, q_0, q_2) = a$$

$$(S, q_0, q_1) = b$$



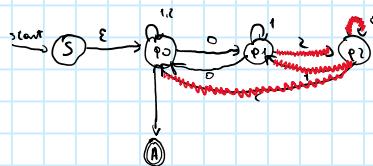
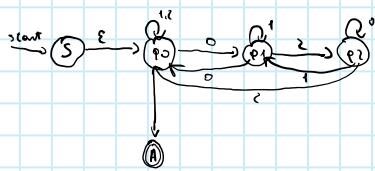
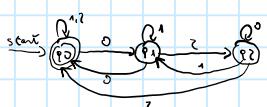
$$(S, q_1, q_2) = ab^*a^*b$$

$$(q_1, q_2, q_3) = a^*b^*a^*b$$



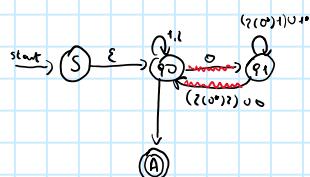
$$(ab^*a^*b)(ab^*a^*b)$$

→



$$q_1 \xrightarrow{a} q_2 \xrightarrow{b} q_1 \equiv q_1 \xrightarrow{a^*b^*} q_1$$

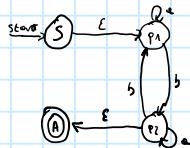
$$q_1 \xrightarrow{a} q_2 \xrightarrow{b} q_0 \equiv q_1 \xrightarrow{a^*b^*} q_0$$



$$(A \cup B)^* = ((B^* \cup A^*)^*)^*$$

$$p_1 \longrightarrow p_0$$

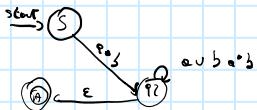
$$(A \cup B)^* = ((B^* \cup A^*)^*)^* = (B^* \cup A^*)^*$$



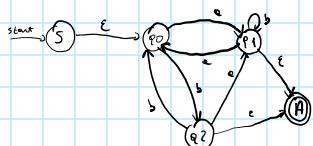
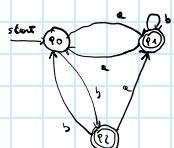
Togliamo  $p_1$

$$(S, p_1, q_2) = a^* b$$

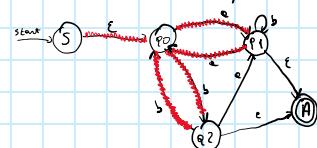
$$(q_2, p_1, q_2) = b a^* b$$



$$(S, q_2, A) = a^* b (a \cup b a^* b)^*$$



Cominciamo eliminando  $p_0$



$$(S, p_0, p_1) = a$$

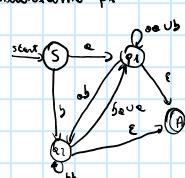
$$(p_1, p_0, p_1) = aa \cup b$$

$$(S, p_0, q_2) = bb$$

$$(q_2, p_0, q_2) = ba \cup a$$

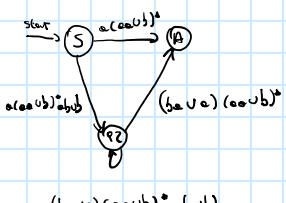
$$(q_2, p_1, q_3) = ab$$

Rimoviamo  $p_1$

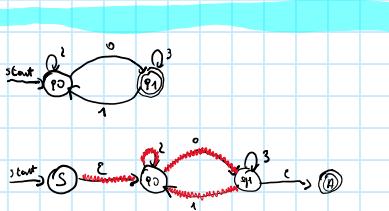


$$(S, q_1, A) = \alpha(\text{aaab})^*$$

$$(S, q_1, q_2) = \alpha(\text{aaab})^* \text{ ab}$$



$\{ba, ab, a\}(aaab)^*$



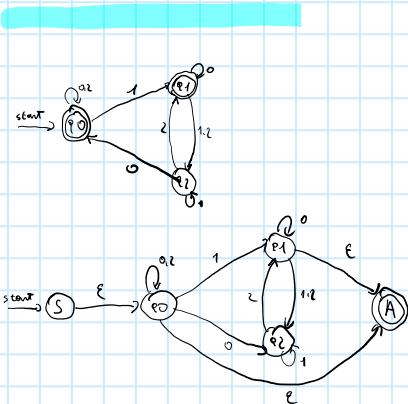
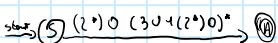
Vogliamo eliminare  $q_0$

$$(S, q_0, q_1) = 2^* 0$$

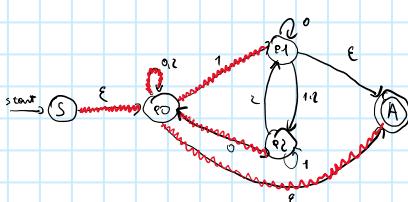
$$(q_1, q_0, q_1) = 1 2^* 0$$



Eliminiamo  $q_1$



Introduciamo  $q_0$

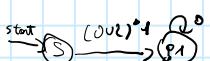


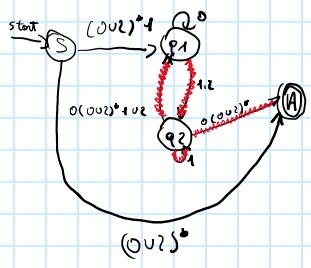
$$(S, q_0, q_1) = S \xrightarrow{\epsilon} q_0 \xrightarrow{1} q_1 = S \xrightarrow{0} q_1$$

$$(S, q_0, A) = S \xrightarrow{\epsilon} q_0 \xrightarrow{1} q_1 = S \xrightarrow{0} A$$

$$(q_2, q_0, q_1) = q_2 \xrightarrow{0} q_0 \xrightarrow{1} q_1 = q_2 \xrightarrow{0} q_1$$

$$(q_2, q_0, A) = q_2 \xrightarrow{0} q_0 \xrightarrow{1} A = q_2 \xrightarrow{0} A$$





Togliiamo  $q_2$

$$(q_1, q_2, q_1) = \underbrace{q_1 \xrightarrow{101} q_1}_{0(0010)^*101} \xrightarrow{1,2} \underbrace{q_1 \xrightarrow{0^*(101)^* (0(0010)^*101)} q_2}$$

# ESERCITAZIONE - DA REGEX A NFA

lunedì 3 aprile 2023 14:40

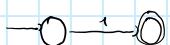
Dare le REGEX

•  $(01 \cup 00)^*$

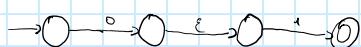
- 0



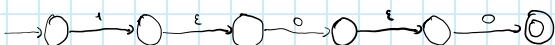
- 1



- 01



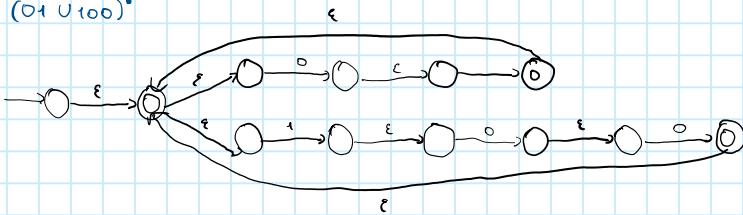
- 100



- 01 ∪ 100

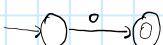


-  $(01 \cup 100)^*$

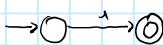


$(0 \cup (11)^*)^* 00$

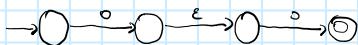
- 0



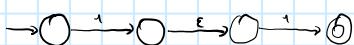
- 1



- 00



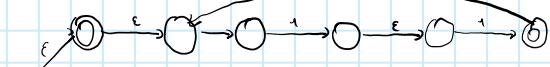
- 11

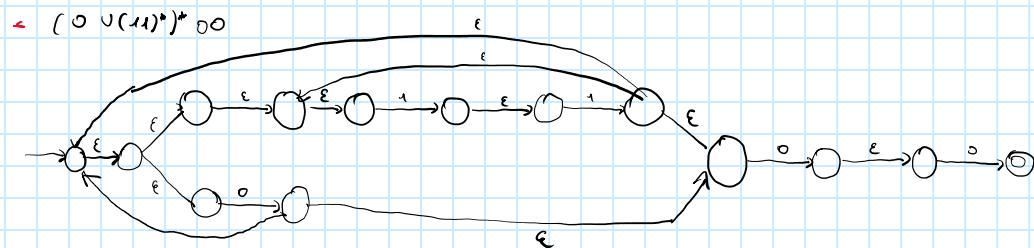
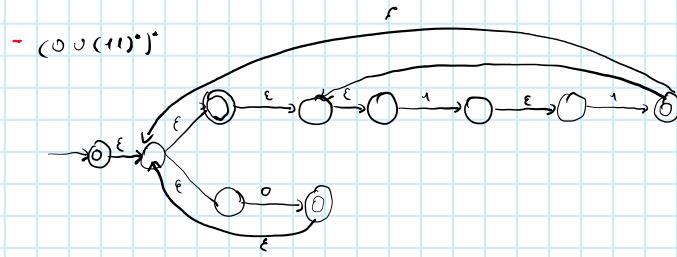
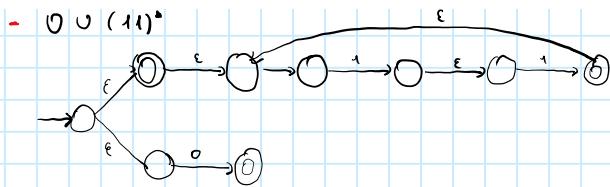


- 11\*



-  $0 \cup (11)^*$





$(0 \cup 1)^* 000 (0 \cup 1)^*$

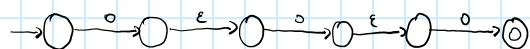
- 0



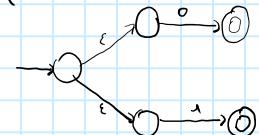
- 1



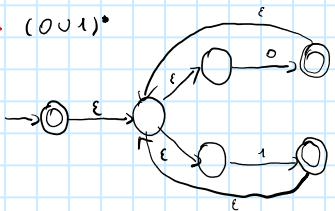
- 000



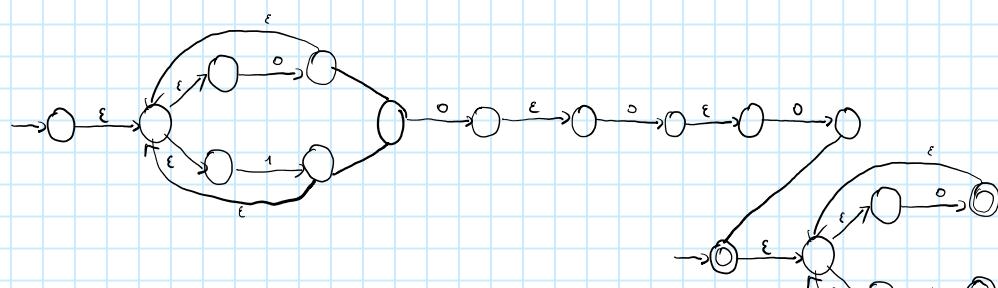
-  $(0 \cup 1)$



-  $(0 \cup 1)^*$



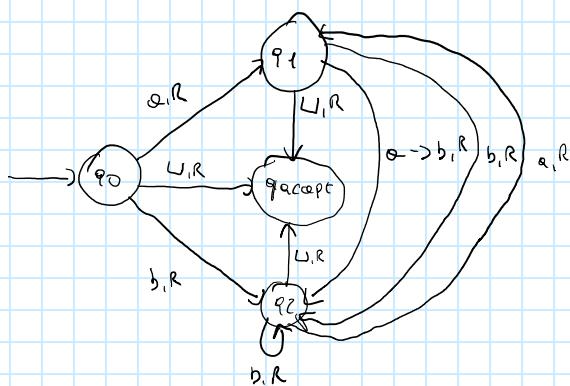
-  $(0 \cup 1)^* 000 (0 \cup 1)^*$





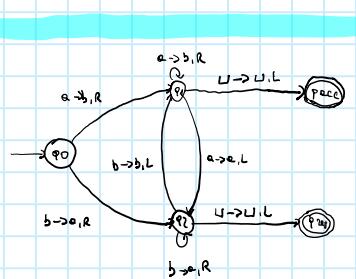
# ESERCITAZIONE - MACCHINE DI TURING

lunedì 3 aprile 2023 14:54



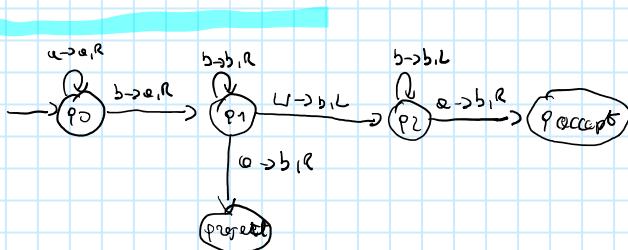
abbbae

$C_1 = q_0 a b b b a e a$ ,  $C_2 = a q_1 b b b a e a$ ,  $C_3 = a b q_2 b b a e a e$   
 $C_4 = a b b q_2 b b a e a e$ ,  $C_5 = a b b b q_2 a e a e$ ,  $C_6 = a b b b a q_2 a e a$   
 $C_7 = a b b b b q_2 a e a$ ,  $C_8 = a b b b a b q_2 a e a$ ,  $C_9 = a b b b a b a q_2$ ,  $C_{10} = a b b b a b a q_{\text{accept}}$   
 La stringa è accettata



$(Q, \Sigma, \Gamma, \delta, q_0, q_{\text{accept}}, q_{\text{reject}})$

$Q = \{q_0, q_1, q_2, q_{\text{accept}}, q_{\text{reject}}\}$   
 $\Sigma = \{a, b\}$   
 $\Gamma = \Sigma \cup \{\#\}$   
 $\delta = \delta(Q, \Sigma) = \{(q_0, a) = (q_1, b, R)$   
 $\delta(q_0, b) = (q_1, a, R)$   
 $\delta(q_0, \#) = (q_{\text{accept}}, L, R)$   
 $\delta(q_1, a) = (q_1, b, R)$   
 $\delta(q_1, b) =$



$(Q, \Sigma, \Gamma, \delta, q_0, q_{\text{accept}}, q_{\text{reject}})$

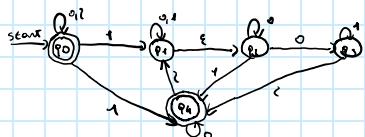
$Q = \{q_0, q_1, q_2, q_{\text{accept}}, q_{\text{reject}}\}$   
 $\Sigma = \{a, b\}$   
 $\Gamma = \Sigma \cup \{\#\}$   
 $\delta = Q \times \Gamma \rightarrow Q \times \Gamma \times \{\text{L}, \text{R}\}$   
 $\delta(q_0, a) = (q_1, a, R)$   
 $\delta(q_0, b) = (q_1, a, R)$   
 $\delta(q_0, \#) = (\text{reject}, L, R)$   
 $\delta(q_1, a) = (\text{reject}, b, R)$

$\delta(q_0, b) = (q_1, a, R)$   
 $\delta(q_0, L) = (q_{accept}, L, R)$   
 $\delta(q_1, a) = (q_{reject}, b, R)$   
 $\delta(q_1, b) = (q_2, b, R)$   
 $\delta(q_2, L) = (q_2, b, L)$   
 $\delta(q_2, a) = (q_{accept}, b, R)$   
 $\delta(q_2, b) = (q_2, b, L)$   
 $\delta(q_2, L) = (q_{reject}, L, R)$

# ESERCITAZIONE - DEFINIZIONE DI QUINTUPLA E ALBERO DI COMPUTAZIONE

martedì 4 aprile 2023 09:21

Determinare la quintupla dell'automa più definito



$$M = (Q, \Sigma, \delta, q_0, F)$$

$$Q = \{q_0, q_1, q_2, q_3, q_4\}$$

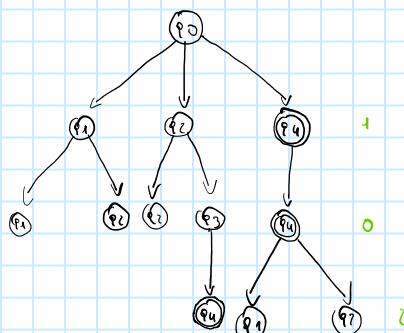
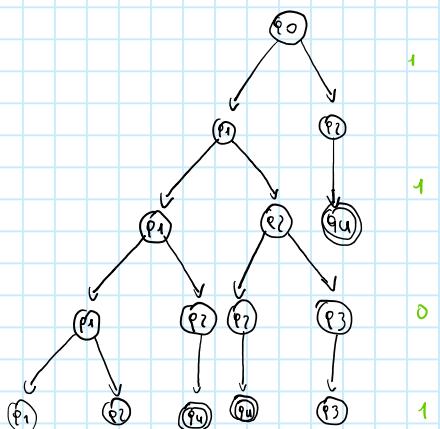
$$\Sigma = \{0, 1, \epsilon\}$$

	0	1	2	$\epsilon$
$q_0$	$\{q_0\}$	$\{q_1\}$	$\{q_3\}$	$\emptyset$
$q_1$	$\{q_0\}$	$\{q_2\}$	$\emptyset$	$\{q_4\}$
$q_2$	$\{q_1, q_3\}$	$\{q_3\}$	$\emptyset$	$\emptyset$
$q_3$	$\emptyset$	$\{q_3\}$	$\{q_4\}$	$\emptyset$
$q_4$	$\{q_4\}$	$\emptyset$	$\{q_1\}$	$\emptyset$

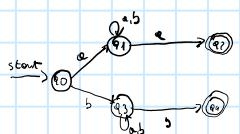
$q_0$  stato iniziale

$$F = \{q_0, q_4\}$$

1001 1024 è numero accettato?



Determinare la quintupla del seguente automa



$$Q = \{q_0, q_1, q_2, q_3, q_4\}$$

$$\Sigma = \{a, b\}$$

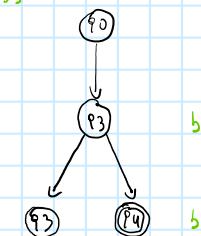
$\delta =$

	a	b	
$q_0$	$\{q_1\}$	$\{q_3\}$	
$q_1$	$\{q_0, q_2\}$	$\{q_2\}$	
$q_2$	$\emptyset$	$\emptyset$	
$q_3$	$\{q_3\}$	$\{q_1, q_2\}$	
$q_4$	$\emptyset$	$\emptyset$	

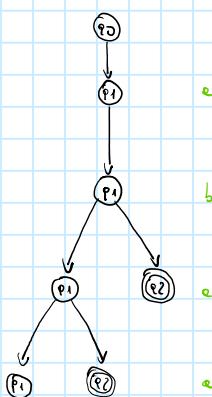
$q_0$  stato iniziale

$$F = \{q_2, q_4\}$$

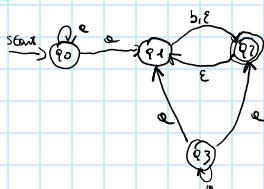
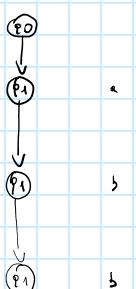
bb



abaae



ahj



$$M = \{Q, \Sigma, \delta, p_0, F\}$$

$$Q = \{q_0, q_1, q_2, q_3\}$$

$$\Sigma = \{a, b\}$$

$\delta =$

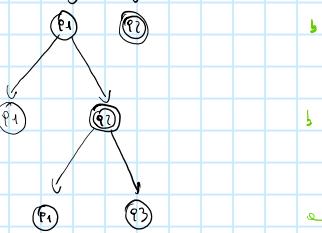
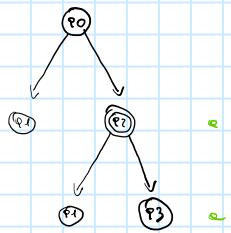
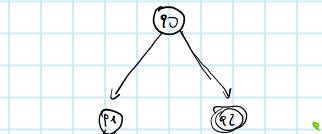
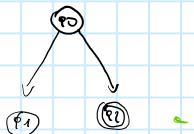
	a	b	$\epsilon$
$q_0$	$\{q_0, q_2\}$	$\emptyset$	$\emptyset$
$q_1$	$\emptyset$	$\{q_3\}$	$\{q_2\}$
$q_2$	$\{q_3\}$	$\emptyset$	$\{q_1\}$

$q_0$	$\{q_0, q_1\}$	$\emptyset$	$\emptyset$
$q_1$	$\emptyset$	$\{q_1\}$	$\{q_2\}$
$q_2$	$\{q_2\}$	$\emptyset$	$\{q_3\}$
$q_3$	$\{q_1, q_2\}$	$\emptyset$	$\emptyset$

$q_0$  è stato inserito

$F = \{q_2\}$

e se abba



# ESERCITAZIONE - INTERSEZIONE

mercoledì 5 aprile 2023 14:36



$$Q_3 = \{ (q_0, 20), (q_0, 21), (q_0, 22), (q_0, 23), (q_1, 20), (q_1, 21), (q_1, 22), (q_1, 23), (q_2, 20), (q_2, 21), (q_2, 22), (q_2, 23), (q_3, 20), (q_3, 21), (q_3, 22), (q_3, 23) \}$$

$$\delta_3((q_0, 20), 0) = (\delta_1(q_0, 0), \delta_2(20, 0)) = (q_1, 20)$$

$$\delta_3((q_0, 20), 1) = (\delta_1(q_0, 1), \delta_2(20, 1)) = (q_0, 21)$$

$$\delta_3((q_1, 20), 0) = (\delta_1(q_1, 0), \delta_2(20, 0)) = (q_1, 20)$$

$$\delta_3((q_1, 20), 1) = (\delta_1(q_1, 1), \delta_2(20, 1)) = (q_0, 22)$$

$$\delta_3((q_0, 22), 0) = (\delta_1(q_0, 0), \delta_2(22, 0)) = (q_1, 20)$$

$$\delta_3((q_0, 22), 1) = (\delta_1(q_0, 1), \delta_2(22, 1)) = (q_0, 23)$$

$$\delta_3((q_0, 23), 0) = (\delta_1(q_0, 0), \delta_2(23, 0)) = (q_1, 23)$$

$$\delta_3((q_0, 23), 1) = (\delta_1(q_0, 1), \delta_2(23, 1)) = (q_0, 23)$$

$$\delta_3((q_1, 23), 0) = (\delta_1(q_1, 0), \delta_2(23, 0)) = (q_1, 23)$$

$$\delta_3((q_1, 23), 1) = (\delta_1(q_1, 1), \delta_2(23, 1)) = (q_0, 20)$$

$$\delta_3((q_2, 20), 0) = (\delta_1(q_2, 0), \delta_2(20, 0)) = (q_3, 20)$$

$$\delta_3((q_2, 20), 1) = (\delta_1(q_2, 1), \delta_2(20, 1)) = (q_2, 21)$$

$$\delta_3((q_2, 21), 0) = (\delta_1(q_2, 0), \delta_2(21, 0)) = (q_3, 21)$$

$$\delta_3((q_2, 21), 1) = (\delta_1(q_2, 1), \delta_2(21, 1)) = (q_2, 22)$$

$$\delta_3((q_2, 22), 0) = (\delta_1(q_2, 0), \delta_2(22, 0)) = (q_3, 22)$$

$$\delta_3((q_2, 22), 1) = (\delta_1(q_2, 1), \delta_2(22, 1)) = (q_2, 23)$$

$$\delta_3((q_2, 23), 0) = (\delta_1(q_2, 0), \delta_2(23, 0)) = (q_3, 23)$$

$$\delta_3((q_2, 23), 1) = (\delta_1(q_2, 1), \delta_2(23, 1)) = (q_2, 23)$$

$$\delta_3((q_3, 20), 0) = (\delta_1(q_3, 0), \delta_2(20, 0)) = (q_3, 20)$$

$$\delta_3((q_3, 20), 1) = (\delta_1(q_3, 1), \delta_2(20, 1)) = (q_3, 21)$$

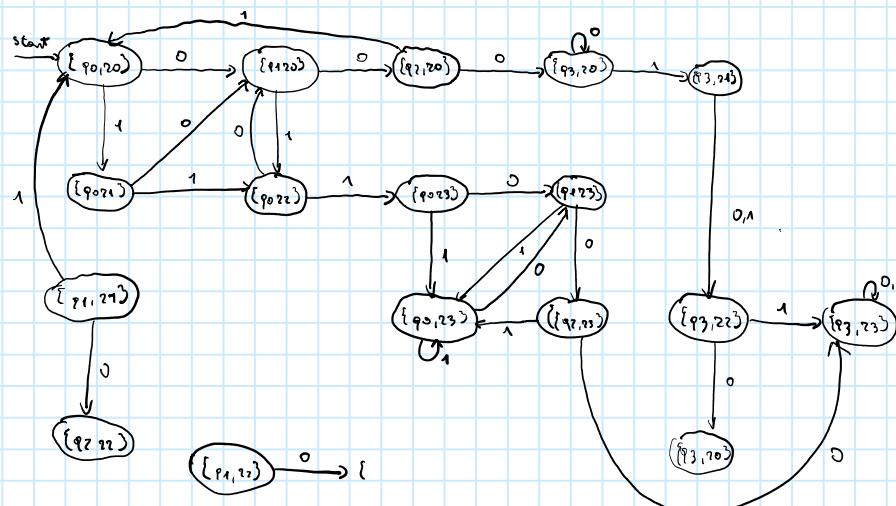
$$\delta_3((q_3, 21), 0) = (\delta_1(q_3, 0), \delta_2(21, 0)) = (q_3, 21)$$

$$\delta_3((q_3, 21), 1) = (\delta_1(q_3, 1), \delta_2(21, 1)) = (q_3, 22)$$

$$\delta_3((q_3, 22), 0) = (\delta_1(q_3, 0), \delta_2(22, 0)) = (q_3, 22)$$

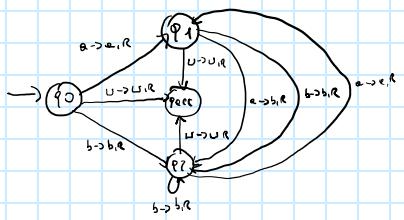
$$\delta_3((q_3, 22), 1) = (\delta_1(q_3, 1), \delta_2(22, 1)) = (q_3, 23)$$

$$\delta_3((q_3, 23), 0) = (\delta_1(q_3, 0), \delta_2(23, 0)) = (q_3, 23)$$



# ESEMPIO - CONFIGURAZIONE DI UNA MDT

martedì 18 aprile 2023 15:08



Formare le sequenze di configurazione per la seguente struttura: *aabbacaa*

- *aaaaabbaaa*
- *aaabbbaaaa*
- *abababaaaa*
- *abbbabaaaa*
- *abbbbabaaaa*
- *abbbbabbaaa*
- *abbbbabbaaa*
- *abbbbabbaaa*
- *abbbbabbaaa*
- *abbbbabbaaa*
- *abbbbabbaaa*

# RIDUZIONE 3-SAT < INDEPENDENT-SET

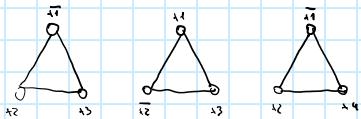
lunedì 12 giugno 2023 11:11

date un'istanza  $\phi$  di 3-SAT, costruiamo un'istanza  $(U, E)$  di INDEPENDENT-SET, tale che  $G$  ha un insieme indipendente di almeno  $k$ , se e solo se  $\phi$  è soddisfacibile.

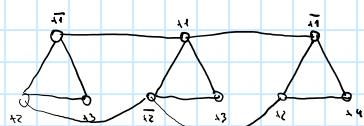
Cioè facciamo nel fatto che in 3-SAT abbiamo delle variabili che possono avere uno o zero, e anche in INDEPENDENT-SET, facciamo una regola di includere un modo nello insieme indipendente.

- $\phi = (x_1 \vee \bar{x}_2 \vee x_3) \wedge (\bar{x}_1 \vee x_2 \vee x_3) \wedge (\bar{x}_1 \vee \bar{x}_2 \vee x_4)$

Per ogni clause creiamo un triangolo e collegiamo i vertici:



Prestiamo un solo vertice e collegiamo qui l'eterno al suo negato.



Per soddisfare le clausole, restiamo  $k$  al numero di clausole.

# RIDUZIONE 3-SAT < SUBSET-SUM

lunedì 12 giugno 2023 12:07

$$\emptyset = (\bar{x}_1 \vee x_2 \vee \bar{x}_3) \wedge (x_1 \vee x_2 \vee x_3) \wedge (x_1 \vee \bar{x}_2 \vee x_3) \wedge (\bar{x}_1 \vee \bar{x}_2 \vee \bar{x}_3) \quad 6 + 8 = 14$$

$2m+2k$

	$x_1$	$x_2$	$x_3$	$c_1$	$c_2$	$c_3$	$c_4$
$x_1$	1	0	0	1	1	0	
$\bar{x}_1$	0	1	0	0	0	1	
$x_2$	0	1	0	1	1	0	
$\bar{x}_2$	0	0	0	0	1	1	
$x_3$	0	0	1	0	1	1	0
$\bar{x}_3$	0	0	1	1	0	0	1
	0	0	0	1	0	0	1
	0	0	0	0	1	0	0
	0	0	0	0	0	1	0
	0	0	0	0	0	0	1
	0	0	0	0	0	0	0
	1	1	1	1	1	1	1

$$x_1 = \text{nero}$$

$$x_2 = \text{giallo}$$

$$x_3 = \text{nero}$$