# NYU Tandon Bridge

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# Homework 1

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# Question 1

Α.

1.

$$10011011_2 = 1 \times 2^7 + 0 \times 2^6 + 0 \times 2^5 + 1 \times 2^4 + 1 \times 2^3 + 0 \times 2^2 + 1 \times 2^1 + 1 \times 2^0$$
  
= 128 + 0 + 0 + 16 + 8 + 0 + 2 + 1  
= 155

**2.** 

$$456_7 = 4 \times 7^2 + 5 \times 7^1 + 6 \times 7^0$$
$$= 196 + 35 + 6$$
$$= 237$$

**3.** 

$$38A_{16} = 3 \times 16^{2} + 8 \times 16^{1} + 10 \times 16^{0}$$
$$= 768 + 128 + 10$$
$$= 906$$

$$2214_5 = 2 \times 5^3 + 2 \times 5^2 + 1 \times 5^1 + 4 \times 5^0$$
$$= 250 + 50 + 5 + 4$$
$$= 309$$

# В.

1.

$$69 \div 2 = 34 \mod 1$$
  
 $34 \div 2 = 17 \mod 0$   
 $17 \div 2 = 8 \mod 1$   
 $8 \div 2 = 4 \mod 0$   
 $4 \div 2 = 2 \mod 0$   
 $2 \div 2 = 1 \mod 0$   
 $1 \div 2 = 0 \mod 1$ 

$$69_{10} = 1000101_2$$

**2**.

$$485 \div 2 = 242 \mod 1$$

$$242 \div 2 = 121 \mod 0$$

$$121 \div 2 = 60 \mod 1$$

$$60 \div 2 = 30 \mod 0$$

$$30 \div 2 = 15 \mod 0$$

$$15 \div 2 = 7 \mod 1$$

$$7 \div 2 = 3 \mod 1$$

$$3 \div 2 = 1 \mod 1$$

$$1 \div 2 = 0 \mod 1$$

$$485_{10} = 111100101_2$$

3.

Each hexadecimal digit can be translated into its 4-bit binary equivalent:

$$\begin{aligned} &6_{16} = 0110_2 \\ &D_{16} = 13_{10} = 1101_2 \\ &1_{16} = 0001_2 \\ &A_{16} = 10_{10} = 1010_2 \end{aligned}$$

Therefore,

$$6D1A_{16} = 0110110100011010_2$$

 $\mathbf{C}.$ 

1.

$$1101011_2 = 01101011_2$$

Each 4-bit binary can be translated into its hexadecimal digit equivalent:

$$0110_2 = 6_{16}$$

$$1011_2 = B_{16}$$

Therefore,

$$1101011_2 = 6B_{16}$$

$$895 \div 16 = 55 \mod 15$$

$$55 \div 16 = 3 \qquad \mod 7$$

$$3 \div 16 = 0 \qquad \mod 3$$

$$15_{10} = F_{16}$$

$$895_{10} = 37F_{16}$$

1.

$$\begin{array}{r}
11111 \\
7566_8 \\
+4515_8 \\
\hline
14303_8
\end{array}$$

$$7566_8 + 4515_8 = 14303_8$$

**2.** 

$$\begin{array}{c} \begin{smallmatrix} 1 & 1 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 & 0 & 0 & 1 & 1_2 \\ + & 0 & 0 & 0 & 1 & 1 & 0 & 1_2 \\ \hline & 1 & 1 & 0 & 0 & 0 & 0 & 0_2 \end{smallmatrix}$$

$$10110011_2 + 1101_2 = 11000000_2$$

**3.** 

$$\begin{array}{c} {}^{1} \ {}^{1} \\ {}^{7}A \ 6 \ 6_{16} \\ {}^{+} \ 4 \ 5 \ C \ 5_{16} \\ {}^{C} \ 0 \ 2 \ B_{16} \end{array}$$

$$7A66_{16} + 45C5_{16} = C02B_{16}$$

$$\begin{array}{r} {\overset{2}{\cancel{3}}\,\overset{4}{\cancel{0}}\,\overset{1}{\cancel{2}}\,2_{5}} \\ -\,2\,4\,3\,3_{5} \\ \hline 3\,4_{5} \end{array}$$

$$3022_5 - 2433_5 = 34_5$$

# Α.

1.

$$124 \div 2 = 62 \mod 0$$

$$62 \div 2 = 31 \mod 0$$

$$31 \div 2 = 15 \mod 1$$

$$15 \div 2 = 7 \mod 1$$

$$7 \div 2 = 3 \mod 1$$

$$3 \div 2 = 1 \mod 1$$

$$1 \div 2 = 0 \mod 1$$

 $124_{10} = 011111100_{8 \ bit \ 2's \ comp}$ 

## **2.**

Convert 124 to its 8 bits binary expression,

$$\begin{aligned} &124 \div 2 = 62 \mod 0 \\ &62 \div 2 = 31 \mod 0 \\ &31 \div 2 = 15 \mod 1 \\ &15 \div 2 = 7 \mod 1 \\ &7 \div 2 = 3 \mod 1 \\ &3 \div 2 = 1 \mod 1 \\ &1 \div 2 = 0 \mod 1 \end{aligned}$$

$$124_{10} = 011111100_2$$

Invert the bits,

10000011

Add 1,

$$10000011 + 00000001 = 10000100$$

Therefore,

$$-124_{10} = 10000100_{8 \ bit \ 2's \ comp}$$

3.

$$109 \div 2 = 54 \mod 1$$
 $54 \div 2 = 27 \mod 0$ 
 $27 \div 2 = 13 \mod 1$ 
 $13 \div 2 = 6 \mod 1$ 
 $6 \div 2 = 3 \mod 0$ 
 $3 \div 2 = 1 \mod 1$ 
 $1 \div 2 = 0 \mod 1$ 

$$109_{10} = 01101101_{8 \ bit \ 2's \ comp}$$

4.

Convert 79 to its 8 bits binary expression,

$$79 \div 2 = 39 \mod 1$$
  
 $39 \div 2 = 19 \mod 1$   
 $19 \div 2 = 9 \mod 1$   
 $9 \div 2 = 4 \mod 1$   
 $4 \div 2 = 2 \mod 0$   
 $2 \div 2 = 1 \mod 0$   
 $1 \div 2 = 0 \mod 1$ 

$$79_{10} = 01001111_2$$

Invert the bits,

10110000

Add 1,

$$10110000 + 00000001 = 10110001$$

Therefore,

$$-79_{10} = 10110001_{8 \ bit \ 2's \ comp}$$

В

1.

$$00011110_{8\ bit\ 2's\ comp} = -0\times2^7 + 0\times2^6 + 0\times2^5 + 1\times2^4 + 1\times2^3 + 1\times2^2 + 1\times2^1 + 0\times2^0$$
 
$$= -0+0+0+16+8+4+2+0$$
 
$$= 30$$

$$11100110_{8\ bit\ 2's\ comp} = -1 \times 2^7 + 1 \times 2^6 + 1 \times 2^5 + 0 \times 2^4 + 0 \times 2^3 + 1 \times 2^2 + 1 \times 2^1 + 0 \times 2^0$$

$$= -128 + 64 + 32 + 0 + 0 + 4 + 2 + 0$$

$$= -26$$

3.

$$00101101_{8\ bit\ 2's\ comp} = -0\times2^7 + 0\times2^6 + 1\times2^5 + 0\times2^4 + 1\times2^3 + 1\times2^2 + 0\times2^1 + 1\times2^0$$
 
$$= -0+0+32+0+8+4+0+1$$
 
$$= 45$$

$$10011110_{8\ bit\ 2's\ comp} = -1 \times 2^{7} + 0 \times 2^{6} + 0 \times 2^{5} + 1 \times 2^{4} + 1 \times 2^{3} + 1 \times 2^{2} + 1 \times 2^{1} + 0 \times 2^{0}$$

$$= -128 + 0 + 0 + 16 + 8 + 4 + 2 + 0$$

$$= -98$$

# 1. Exercise 1.2.4

# section b

p	q	$p \lor q$	$\neg (p \lor q)$
Т	Т	Т	F
Т	F	Т	F
F	Т	Т	F
F	F	F	Т

# section c

p	q	r	$p \land \neg q$	$r \lor (p \land \neg q)$
Т	Т	Т	F	Т
Т	Т	F	F	F
Т	F	Т	Т	Т
Т	F	F	Т	Т
F	Т	Т	F	Т
F	Т	F	F	F
F	F	Т	F	Т
F	F	F	F	F

# 2. Exercise 1.3.4

# section b

p	q	$p \rightarrow q$	$q \rightarrow p$	$(p \to q) \to (q \to p)$
Т	T	Т	Т	T
Т	F	F	Т	Т
F	Т	Т	F	F
F	F	Т	Т	Т

# $\mathbf{section}\ \mathbf{d}$

p	q	$p \leftrightarrow q$	$p \leftrightarrow \neg q$	$(p \leftrightarrow q) \oplus (p \leftrightarrow \neg q)$
Т	Т	Т	F	T
Т	F	F	Т	Т
F	Т	F	Т	Т
F	F	Т	F	Т

# 1. Exercise 1.2.7

section b

$$(B \wedge D) \vee (B \wedge M) \vee (D \wedge M)$$

section c

$$B \vee (D \wedge M)$$

# 2. Exercise 1.3.7

section b

$$(s\vee y)\to p$$

section c

$$p \rightarrow y$$

section d

$$p \leftrightarrow (s \land y)$$

section e

$$p \to (s \vee y)$$

# 3. Exercise 1.3.9

section c

$$c \to p$$

 $\mathbf{section}\ \mathbf{d}$ 

$$c \to p$$

## 1. Exercise 1.3.6

### section b

If Joe is eligible for the honors program, then he has maintained a B average.

### section c

If Rajiv can go on the roller coaster, then he is at least four feet tall.

### section d

If Rajiv is at least four feet tall, then he can go on the roller coaster.

## 2. Exercise 1.3.10

## section c

p	q	r	$p \vee r$	$q \wedge r$	$(p \lor r) \leftrightarrow (q \land r)$
Τ	F	Т	Т	F	F
Т	F	F	Т	F	F

According to the truth table, the truth value of the expression  $(p \lor r) \leftrightarrow (q \land r)$  is false.

### section d

p	q	r	$p \wedge r$	$q \wedge r$	$(p \land r) \leftrightarrow (q \land r)$
Т	F	Т	Τ	F	F
Т	F	F	F	F	T

According to the truth table, the truth value of the expression  $(p \land r) \leftrightarrow (q \land r)$  is unknown.

### section e

p	q	r	$r \lor q$	$p \to (r \lor q)$
Т	F	Т	Т	Т
$\Gamma$	F	F	F	F

According to the truth table, the truth value of the expression  $p \to (r \lor q)$  is unknown.

## section f

p	q	r	$p \wedge q$	$(p \land q) \to r$
Т	F	Т	F	T
Т	F	F	F	T

According to the truth table, the truth value of the expression  $(p \land q) \rightarrow r$  is true.

## Exercise 1.4.5

### section b

 $1^{\rm st}$  sentence can be expressed in symbolic logic as,

$$\neg j \rightarrow (l \lor \neg r)$$

2<sup>nd</sup> sentence can be expressed in symbolic logic as,

$$(r \land \neg l) \to j$$

# Proof:

j	l	r	$\neg j \rightarrow (l \lor \neg r)$	$(r \land \neg l) \to j$
Т	Т	Т	T	Т
Т	Т	F	Т	Т
Т	F	Т	Т	Т
Т	F	F	Т	Т
F	Т	Т	Т	Т
F	Т	F	Т	Т
F	F	Т	F	F
F	F	F	Т	Т

According to the truth table, the two expressions have the same truth value in all situations, therefore they are logically equivalent.

## section c

1<sup>st</sup> sentence can be expressed in symbolic logic as,

$$j \rightarrow \neg l$$

2<sup>nd</sup> sentence can be expressed in symbolic logic as,

$$\neg j \rightarrow l$$

## Proof:

j	l	$j \rightarrow \neg l$	$\neg j \rightarrow l$
Т	Т	F	Т
Т	F	Т	Т
F	Т	Т	Т
F	F	Т	F

According to the truth table, the two expressions have different truth values, therefore they are not logically equivalent.

# section d

1<sup>st</sup> sentence can be expressed in logic as,

$$(r \vee \neg l) \rightarrow j$$

 $2^{\rm nd}$  sentence can be expressed in logic as,

$$j \to (r \land \neg l)$$

# Proof:

j	l	r	$(r \lor \neg l) \to j$	$j \to (r \land \neg l)$
Т	Т	Т	Т	F
Т	Т	F	Т	F
Т	F	Т	Т	Т
Т	F	F	Т	F
F	Т	Т	F	Т
F	Т	F	Т	Т
F	F	Т	F	Т
F	F	F	F	Т

According to the truth table, the two expressions have different truth values, therefore they are not logically equivalent.

### 1. Exercise 1.5.2

section c

$$(p \to q) \land (p \to r) \equiv (\neg p \lor q) \land (p \to r)$$
 (Conditional identity) 
$$\equiv (\neg p \lor q) \land (\neg p \lor r)$$
 (Conditional identity) 
$$\equiv \neg p \land (q \lor r)$$
 (Distributive law) 
$$\equiv p \to (q \lor r)$$
 (Conditional identity)

section f

$$\neg(p \lor (\neg p \land q)) \equiv \neg((p \lor \neg p) \land (p \lor q)) \qquad \qquad \text{(Distributive law)} \\
\equiv \neg(T \land (p \lor q)) \qquad \qquad \text{(Complement law)} \\
\equiv \neg(p \lor q) \qquad \qquad \text{(Identity law)} \\
\equiv \neg p \land \neg q \qquad \qquad \text{(De Morgan's law)}$$

section i

$$(p \wedge q) \rightarrow r \equiv \neg (p \wedge q) \vee r$$
 (Conditional identity) 
$$\equiv (\neg p \vee \neg q) \vee r$$
 (De Morgan's law) 
$$\equiv r \vee (\neg p \vee \neg q)$$
 (Commutative law) 
$$\equiv (r \vee \neg p) \vee \neg q$$
 (Associative law) 
$$\equiv (\neg p \vee r) \vee \neg q$$
 (Commutative law) 
$$\equiv (\neg p \vee \neg \neg r) \vee \neg q$$
 (Double negation law) 
$$\equiv \neg (p \wedge \neg r) \vee \neg q$$
 (De Morgan's law) 
$$\equiv (p \wedge \neg r) \rightarrow \neg q$$
 (Conditional identity)

# 2. Exercise 1.5.3

section c

$$\neg r \lor (\neg r \to p) \equiv \neg r \lor (\neg \neg r \lor p) \qquad \qquad \text{(Conditional identity)} \\
\equiv \neg r \lor (r \lor p) \qquad \qquad \text{(Double negation law)} \\
\equiv (\neg r \lor r) \lor p \qquad \qquad \text{(Associative law)} \\
\equiv T \lor p \qquad \qquad \text{(Complement law)} \\
\equiv T \qquad \qquad \text{(Domination law)}$$

# section d

# 1. Exercise 1.6.3

section c

$$\exists x \ (x = x^2)$$

 ${\bf section}\ {\bf d}$ 

$$\forall x \ (x \le x^2 + 1)$$

2. Exercise 1.7.4

section b

$$\forall x \ (\neg S(x) \land W(x))$$

 $\mathbf{section}\ \mathbf{c}$ 

$$\forall x \ (S(x) \to \neg W(x))$$

 $\mathbf{section}\ \mathbf{d}$ 

$$\exists x \ (S(x) \land W(x))$$

### 1. Exercise 1.7.9

#### section c

True.

This expression means that there exists an x such that either x = c is false or P(x) is true. According to the truth table, if x = a, x = c is false, so the expression is true.

### section d

True.

According to the truth table, if x = e, then Q(x) is true and R(x) is true. It follows  $Q(x) \wedge R(x)$  is true. Therefore, there exists an x such that  $Q(x) \wedge R(x)$  is true.

### section e

True.

According to the truth table, Q(a) is true and P(d) is true. Therefore,  $Q(a) \wedge P(d)$  is true.

### section f

True.

According to the truth table, for all x, if  $x \neq b$ , Q(x) is always true.

## section g

False.

According to the truth table, if x = c, P(x) is false and R(x) is false. It follows  $P(x) \vee R(x)$  is false. Therefore, it is false that for all x,  $P(x) \vee R(x)$  is true.

### section h

True.

According to the truth table, for all x, either P(x) is true or R(x) is false. If P(x) is true, then  $R(x) \to P(x)$  is true. If R(x) is false, then  $R(x) \to P(x)$  is also true. Therefore, for all x,  $R(x) \to P(x)$  is always true.

### section i

True.

According to the truth table, when x = a, Q(x) is true and R(x) is false. It follows  $Q(x) \vee R(x)$  is true. Therefore, there exists an x such that  $Q(x) \vee R(x)$  is true.

### 2. Exercise 1.9.2

### section b

True.

According to the truth table, when x = 2, Q(x, y) is true for all y.

### section c

True.

According to the truth table, when y = 1, P(x, y) is true for all x.

### section d

False.

According to the truth table, S(x,y) is always false. Therefore, there does not exist an x and a y such that S(x,y) is true.

### section e

False.

According to the truth table, when x = 1, there does not exist a y such that Q(x, y) is true. Therefore, it false that, for all x, there exist a y such that Q(x, y) is true.

### section f

True.

According to the truth table, for all x, there exists a y such that P(x, y) is true. When x = 1, if y = 1, then P(x, y) is true. When x = 2, if y = 1, then P(x, y) is true. When x = 3, if y = 1, then P(x, y) is true.

## section g

False.

According to the truth table, when x = 1 and y = 2, P(x, y) is false. So, it is false that, for all x and y, P(x, y) is true.

### section h

True.

According to the truth table, when x = 2 and y = 1, Q(x, y) is true. So, there exists an x and a y such that Q(x, y) is true.

# section i

True.

According to the truth table, for all x and y, S(x,y) is always false. Therefore,  $\neg Q(x,y)$  is always true.

### 1. Exercise 1.10.4

section c

$$\exists x \exists y \ (x + y = xy)$$

section d

$$\forall x \forall y \ (((x>0) \land (y>0)) \rightarrow \left(\frac{x}{y}>0\right))$$

section e

$$\forall x \; (((x>0) \land (x<1)) \rightarrow \left(\frac{1}{x} > 1\right))$$

section f

$$\forall x \exists y \ (y < x)$$

section g

$$\forall x \exists y \ ((x \neq 0) \to \left(y = \frac{1}{x}\right))$$

## 2. Exercise 1.10.7

section c

$$\exists x \ (N(x) \land D(x))$$

section d

$$\forall y \ (D(y) \to P(\operatorname{Sam}, y))$$

section e

$$\exists x \forall y \ (N(x) \land P(x,y))$$

section f

$$\exists x \forall y \ ((N(x) \land D(x)) \land ((N(y) \land (y \neq x)) \rightarrow \neg D(y)))$$

### 3. Exercise 1.10.10

 $\mathbf{section}\ \mathbf{c}$ 

$$\forall x \exists y \ (T(x,y) \land (y \neq \text{Math } 101))$$

 $\mathbf{section}\ \mathbf{d}$ 

$$\exists x \forall y \ ((y \neq \text{Math } 101) \rightarrow T(x,y))$$

section e

$$\forall x \exists y_1 \exists y_2 \ ((x \neq \operatorname{Sam}) \to (T(x, y_1) \land T(x, y_2) \land (y_1 \neq y_2)))$$

# section f

 $\exists y_1 \exists y_2 \forall y_3 \ (T(\operatorname{Sam}, y_1) \land T(\operatorname{Sam}, y_2) \land (y_1 \neq y_2) \land (y_3 \neq y_1) \land (y_3 \neq y_2)) \rightarrow \neg T(\operatorname{Sam}, y_3))$ 

### 1. Exercise 1.8.2

### section b

Statement:

 $\forall x \ (D(x) \lor P(x))$ 

Negation:

 $\neg \forall x \ (D(x) \lor P(x))$ 

Applying De Morgan's law:

 $\exists x \neg (D(x) \lor P(x))$ 

Applying De Morgan's law:

 $\exists x \ (\neg D(x) \land \neg P(x))$ 

English: there is a patient who was neither given the medication nor the placebo.

### section c

Statement:

 $\exists x \ (D(x) \land M(x))$ 

Negation:

 $\neg \exists x \ (D(x) \land M(x))$ 

Applying De Morgan's law:

 $\forall x \ \neg (D(x) \land M(x))$ 

Applying De Morgan's law:

 $\forall x \ (\neg D(x) \lor \neg M(x))$ 

English: every patient was not given the medication or did not have migraines or both.

### section d

Statement:

 $\forall x (P(x) \to M(x))$ 

Negation:

 $\neg \forall x (P(x) \to M(x))$ 

Applying De Morgan's law:

 $\exists x \neg (P(x) \to M(x))$ 

Applying conditional identity:

 $\exists x \neg (\neg P(x) \lor M(x))$ 

Applying De Morgan's law:

 $\exists x (\neg \neg P(x) \land \neg M(x))$ 

Applying double negation law:

 $\exists x (P(x) \land \neg M(x))$ 

English: there is a patient who was given the placebo and did not have migraines.

### section e

Statement:

$$\exists x (M(x) \land P(x))$$

Negation:

$$\neg \exists x \ (M(x) \land P(x))$$

Applying De Morgan's law:

$$\forall x \neg (M(x) \land P(x))$$

Applying De Morgan's law:

$$\forall x \ (\neg M(x) \lor \neg P(x))$$

English: every patient did not have migraines or was not given not have placebo.

### 2. Exercise 1.9.4

#### section c

$$\neg \exists x \forall y (P(x,y) \rightarrow Q(x,y)) \equiv \forall x \neg \forall y (P(x,y) \rightarrow Q(x,y)) \qquad \qquad \text{(De Morgan's law)}$$
 
$$\equiv \forall x \exists y \neg (P(x,y) \rightarrow Q(x,y)) \qquad \qquad \text{(Conditional identity)}$$
 
$$\equiv \forall x \exists y (\neg P(x,y) \land \neg Q(x,y)) \qquad \qquad \text{(De Morgan's law)}$$
 
$$\equiv \forall x \exists y (P(x,y) \land \neg Q(x,y)) \qquad \qquad \text{(De Morgan's law)}$$
 
$$\equiv \forall x \exists y (P(x,y) \land \neg Q(x,y)) \qquad \qquad \text{(Double negation law)}$$

## section d

$$\neg \exists x \forall y (P(x,y) \leftrightarrow P(y,x)) \equiv \forall x \neg \forall y (P(x,y) \leftrightarrow P(y,x)) \qquad \qquad \text{(De Morgan's law)}$$
 
$$\equiv \forall x \exists y \neg (P(x,y) \leftrightarrow P(y,x)) \qquad \qquad \text{(De Morgan's law)}$$
 
$$\equiv \forall x \exists y \neg ((P(x,y) \rightarrow P(y,x)) \land (P(y,x) \rightarrow P(x,y))) \qquad \text{(Conditional identity)}$$
 
$$\equiv \forall x \exists y \left( \neg (P(x,y) \rightarrow P(y,x)) \lor \neg (P(y,x) \rightarrow P(x,y)) \right) \qquad \text{(De Morgan's law)}$$
 
$$\equiv \forall x \exists y \left( (\neg \neg P(x,y) \lor P(y,x)) \lor \neg (\neg P(y,x) \lor P(x,y)) \right) \qquad \text{(Conditional identity)}$$
 
$$\equiv \forall x \exists y \left( (\neg \neg P(x,y) \land \neg P(y,x)) \lor (\neg \neg P(y,x) \land \neg P(x,y)) \right) \qquad \text{(De Morgan's law)}$$
 
$$\equiv \forall x \exists y \left( (\neg \neg P(x,y) \land \neg P(y,x)) \lor (\neg \neg P(y,x) \land \neg P(x,y)) \right) \qquad \text{(Double negation law)}$$

### section e