Generally, When two station transmit data Simultaneounly and collinion occurro then They may restrainmit the data immediately ofter the dollinion. Thin may again lead to the collipion and again retransmit of dota. Station wait bore Some trandon back of time and then Thin waiting time bott which the station waith before retransmitting The data is called an back obt time atter under going the collinion Transmithing station chances a Trandom number in the range To, 2"-1] Stille time praket in unterraine collinion borthe nith time.

To solve thin problem we use binory exponential back of algorithm (BEB).

There the two node have, how much to wait to transmit their data again.

here, A= 40,1,2,3,4,5,6,7,8,9,10,11,12,13,14,15]

B=40,1,2,3,4,5,6,7,8,9,10,11,12,13,14,15]

There means the number of adlinion

there and B in 4.

tore A in 3 and B in 4.

There are and B pick same number

New 4th A and B pick same number

reardomly they will collide again, on both A = 3 and | n= number of Then n 6017 B = 4 So it will be set= for 23-13 So A= h 0,1,2,3,4,5,6,7 B= 201, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14, 15} Sk they Piex some number again pick they will collide. Like it both pick 3×512 Bit time - 1536 (waiting time 3 then They will vait and transmit their again and will collide But at the same point st. They pick

dibbettent number, then the collinion would not occurre.

minimized dollinion in Wlenn

a transminsion medium-a cable, a transminsion medium-a cable, a transminsion medium-a cable, but witteless networks don't use cable, but witteless networks don't use cable, but even WLAN, all device involved send even WLAN, all device involved send and preceived.

The medium accern control (MAC) in the kuredian that arrbitrates the accern to the kuredian that arrbitrates the accern to the channel. In Willelenn metwork, the MAC Channel. In Willelenn metwork, the MAC protocoln play a Kes Male in maximizing protocoln play a Kes Male in maximizing the channel Utilization

Cormier Senese multiple Accen (CSMA) &

Time divinion multiple accen (TDMA)

Time divinion multiple accen (TDMA)

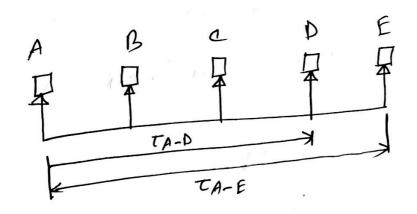
one two know medium accen mechanism

som share medium communication syntems

TDMA treavirsen tight time synchronization among the participating station. Additionally among the participating station. Additionally a prior set-up in treavuirsed to annight a prior station state active station. It the active station. It is the active station to have no data tready to transmit the station has no data tready to transmit the station time to wanted.

There grate an variety is TDMA in that in miniamize collinion and that in miniamize collinion and that in miniamize collinion and achive high channel ethiciency. may achive high channel ethiciency and also Learing - BEB for any number and also Learing a collinion and et station by tredveing a collinion and et station by tredveing a collinion and irenime The number of Sucrembull State.

Ann to the av no 82



By competing transmitted and treceive signal. St there similler then no signal. St there similler the distancent the collinion, and also st distancent the collinion.

Here A in a source and D in the

Dentination.

Dentination period in the time of Arronne

Contention period in the time of Arronne

The sender. When it in a

Theturn to the sender. When it in a

Theturn to the sender.

Collinion for a report the proposion

send brom A-D in JA-P delay to Detacation time - (T-E) + (t-E) 50 · 2 TA-D (2x propa-Contention portiod 2T Sending to D packet dota then A to D WIN also get the dota. E will not get it, become Dwill receive it belonce that A and & one the tronthert hort at this medium. It is dolled wort done And also there A to E in I, the contention penied Deloy

In 2 TA-E.

By some herre. A to E

forther hart and propagation

the forther hart and propagation

telon led other every pother

delay less others Ati E.

x rx lrx