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Ans: to the qu: NO: 2

B, C, D, G, H are interested.

	0	1	2	3	4	5	
B (101110)	1	0	1	1	1	0	
C (101010)	1	0	1	0	X	X	
D (100010)	1	0	0	X	X	X	
G (101100)	1	0	1	1	0	X	
H (011011)	0	X	X	X	X	X	
OR,	1	0	1	1	1	0	→ B

so, the winner will be "B"

Ans: to the qu: NO: 3

for A :

$$\text{set } A = \{0, 1, 2, 3\}$$

$$\rightarrow 0 \times 512 \text{ bit time} = 0 \text{ bit waiting time}$$

$$\rightarrow 1 \times 512 \text{ bit time} = 512 \text{ bit waiting time}$$

$$\rightarrow 2 \times 512 \text{ bit time} = 1024 \text{ bit waiting time}$$

$$\rightarrow 3 \times 512 \text{ bit time} = 1536 \text{ bit waiting time}$$

Here, B is similar as A.

$$\text{and, } C = \{0, 1, 2, 3, 4, 5, 6\}$$

first 3 bit is similar, but for 4, 5, 6

we will get 2048, 2560, 3072 bit waiting time.

If we compare between A and C,
if we choose any random number
from A and C they will wait
before transmitting,

Here,

$\{0, 1, 2, 3\}$ is common to all

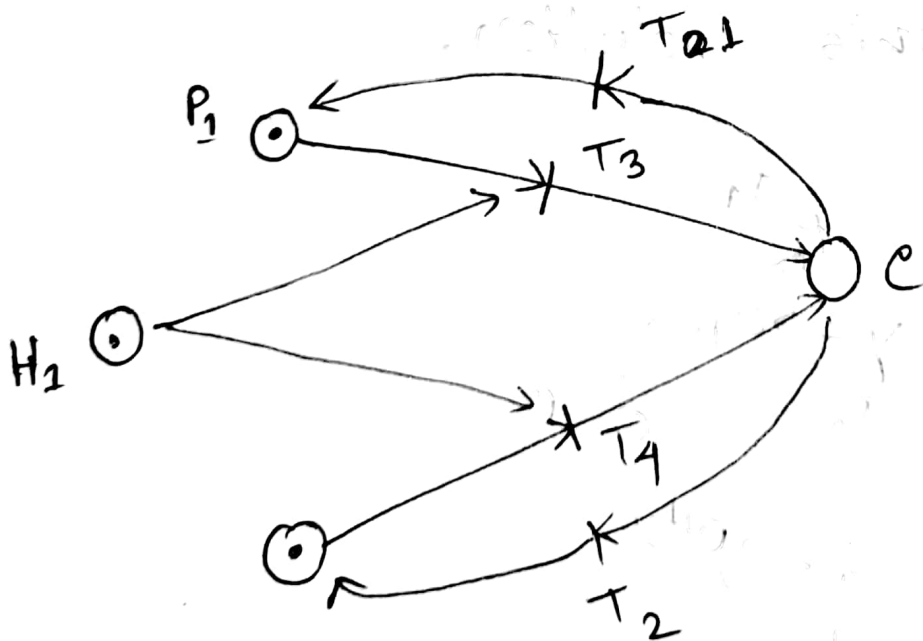
set (A, B, C). so, they will wait 1024
bit time to transmit and they will
occur a collision for same time.

C can't communicate with other nodes
because if A tries to communicate
with C for picked number it
can occur collision in sending.
Same goes for b and C.

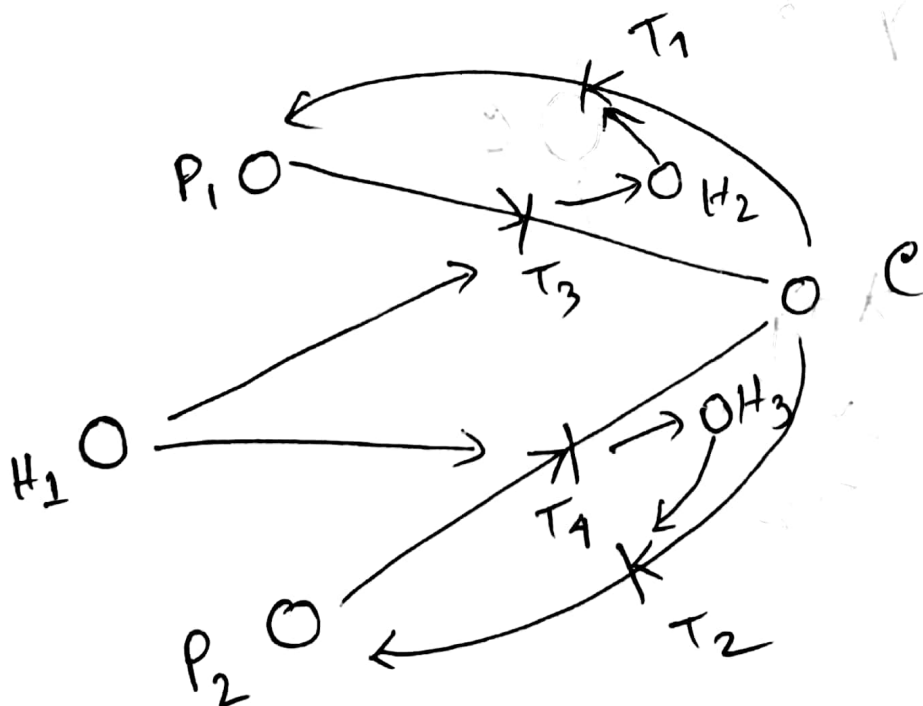
Ans: to the qu: NO: 5

There is a problem of returning the token and mutual exclusion.

Here, P_1 and H_1 sends token to C_1 but it can't complete process for because of return function.



In the diagram if P_1 and H_1 send two token it will go to the T_3 through t_3 and while returning, it will return two token to P_1 which is irrelevant. and same for P_2 , so, we should use two more (h_2, h_3) to solve this solution.



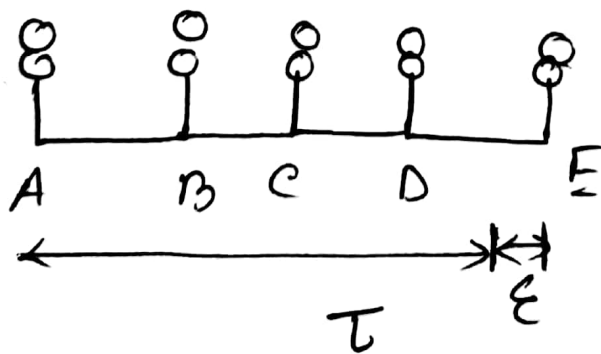
Ans: to the qu: No: 4

CSMA/CD is inherently half duplex because it cannot send and receive data at the same time. It can send data when it is not receiving and it will wait to transmit.

If may occur collision if they do send and receive at a same time.

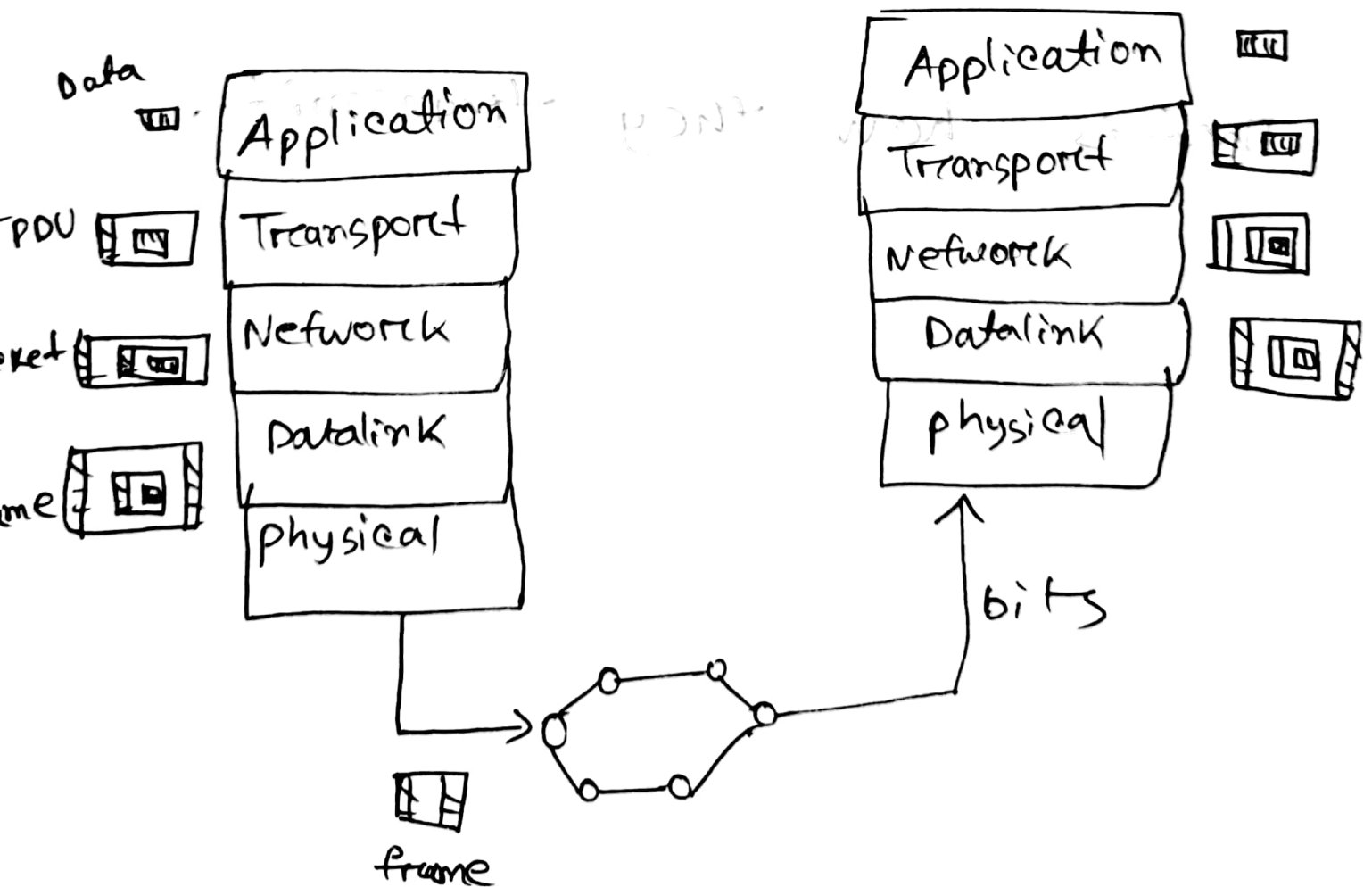
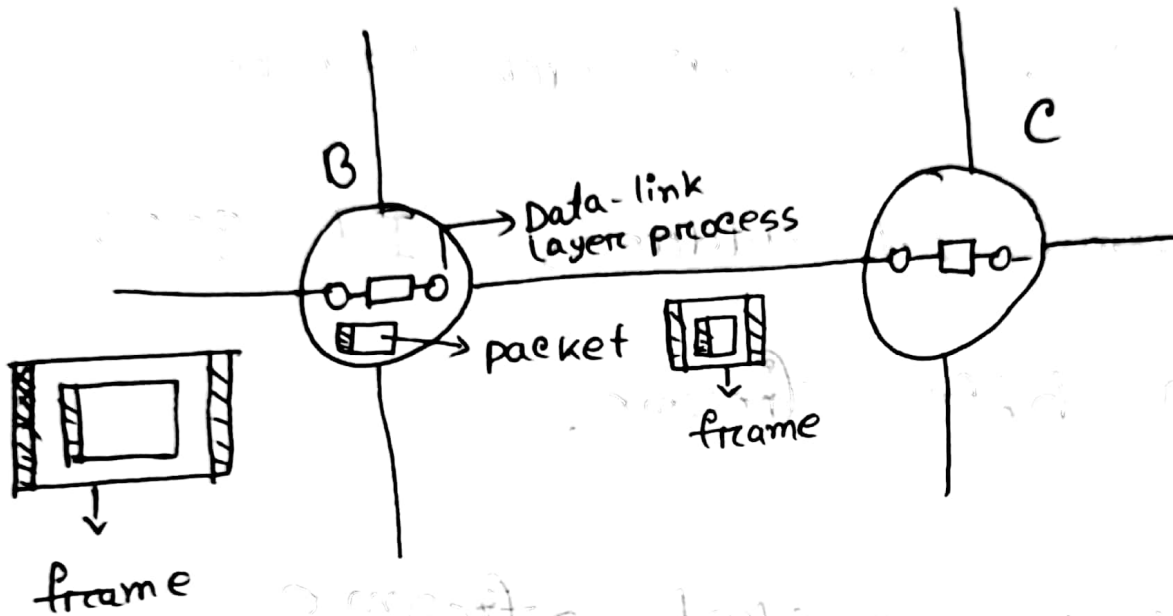
contention period always counted between the two farthest node.

Here, we can see the farthest node is E. So, even C is destination, it will propagate to the E.



∴ contention period = 2τ

Ans: to the qu: NO: 1



Here, is the encapsulation. Here packet holds id address, that's why it is in network layer.

~~Destination~~ Id address It goes as a bits frame.

Data \rightarrow tpdu \rightarrow packet \rightarrow frame.

That is how they transmit.