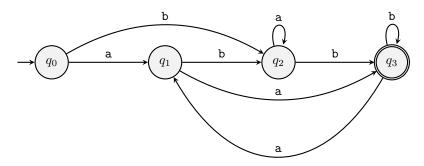
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Solve problems 1 through 6. Problem 7 is optional.

Problem 1: Converting DFA to Regular Expressions (10 points)

Convert the following DFA into an equivalent regular expression using the state elimination method. First eliminate q_1 , then q_2 and so on. You must show work.



Problem 2: Designing Context-Free Grammars (5 points)

Solve the following problems.

(a) Give a context-free grammar for the following language. (2 points)

$$A = \{w \in \{0, 1\}^* : w \text{ contains at least two 0s}\}\$$

(b) Give a context-free grammar for the following language. (3 points)

$$L = \{w \in \{0, 1\}^* : w = 0^{3i}v1^{2i} \text{ where } v \in A \text{ and } i \ge 0\}$$

Problem 3: Derivations, Parse Trees, and Ambiguity (5 points)

$$S \rightarrow 0S1 \mid 1S0 \mid SS \mid 01 \mid 10$$

Take a look at the grammar above and solve the following problems.

- (a) Show that the grammar above is ambiguous by demonstrating two different parse trees for 011010. (4 points)
- (b) Find a string w of length six such that w has exactly one parse tree in the grammar above. (1 point)

Problem 4: Chomsky Normal Form (10 points)

Convert the following grammar into Chomsky Normal Form. You must show work.

$$\begin{split} S &\to \mathtt{a} \mid \mathtt{a}A \mid B \\ A &\to \mathtt{a}BB \mid \varepsilon \mid D \\ B &\to A\mathtt{a} \mid \mathtt{b} \\ D &\to \varepsilon \end{split}$$

Here a, and b are terminals and the rest are variables.

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Problem 5: The CYK Algorithm (10 points)

Use the CYK algorithm to determine whether the string abcccc can be derived in the following grammar. You must show the entire CYK table.

$$\begin{split} S &\to AB \\ A &\to CD \mid CF \\ B &\to \mathbf{c} \mid EB \\ C &\to \mathbf{a} \\ D &\to \mathbf{b} \\ E &\to \mathbf{c} \\ F &\to AD \end{split}$$

Here, a, b, and c are terminals and the rest are variables.

Problem 6: Constructing Pushdown Automata (5 points)

Construct a pushdown automaton for the following language.

$$L = \{ w\overline{w}^{\mathcal{R}} : w \in \{0, 1\}^* \}$$

Here, $\overline{w}^{\mathcal{R}}$ denotes the reverse complement of the string w.

For example, $01001101 \in L$ because the second half of the string, 1101, is the reverse complement of the first half, 0100, i.e. $1101 = \overline{0100}^{\mathcal{R}}$.

Problem 7: (Bonus) Deja Vu (5 points)

(Note that this is a bonus problem. Attempt it only after you are done with everything else. Even if you do not attempt it, you can get a perfect score. So, do not worry if you find it too hard!)

Take a look at the following languages. They should seem familiar.

$$L_1 = \{ w \in \{0,1\}^* : w = 0^m 1^n \text{ where } m \text{ and } n \text{ are both even} \},$$

$$L_2 = \{w \in \{0,1\}^* : w = 0^m 1^n \text{ where } m \text{ and } n \text{ are both odd}\}.$$

We will use these two languages to define the following language.

$$L = \{w \in \{0,1\}^* : w = uv, \text{ where } u \in L_1, v \in L_2, \text{ and } |u| = |v|\}$$

- (a) Write down a string $w \in L$ such that w is of length eight. (0 point)
- (b) Construct a pushdown automaton that recognizes L. Describe what your automaton is doing in two or three sentences. (5 points)