

# 8086 Hardware Specifications

Dept. of Computer Science and Engineering BRAC University

**CSE 341 Team** 





## Book:

Microprocessors and Interfacing: Programming and Hardware,

Author: Douglas V. Hall

The 8086/8088 Family: Design, Programming, And Interfacing,

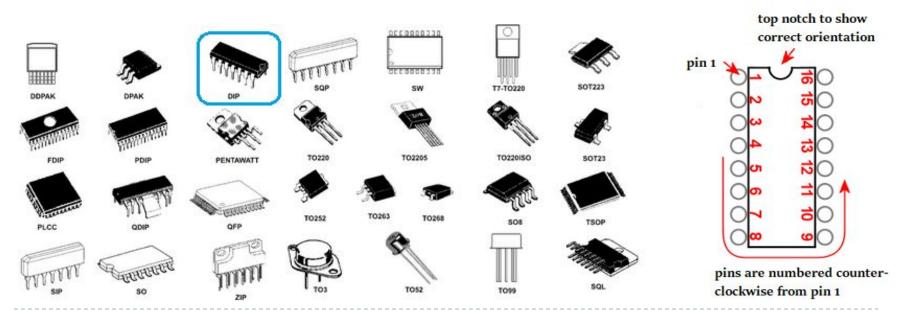
Author: John Uffenbeck.







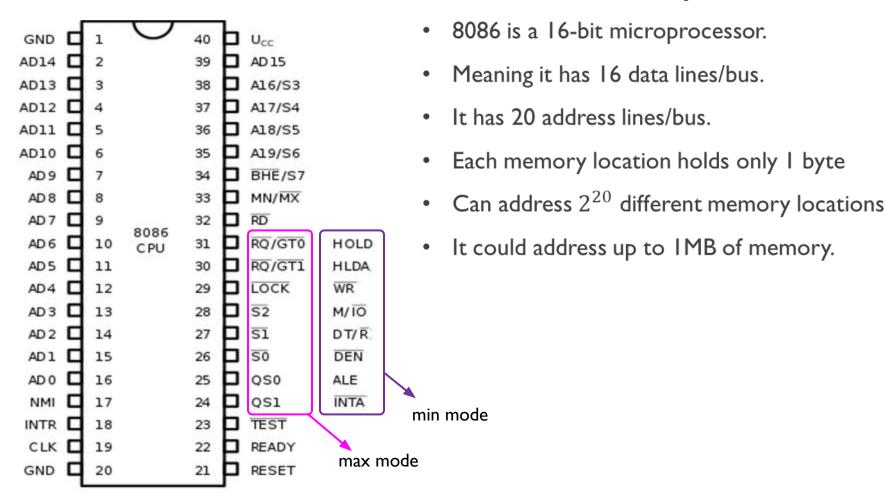
- is a 40-pin DIPs; Dual in-line package
- DIP refers to a rectangular housing with two parallel rows of electrical connection pins.
- DIPs have a notch on one end to show its correct orientation.
- The pins are then numbered as shown in the figure below.







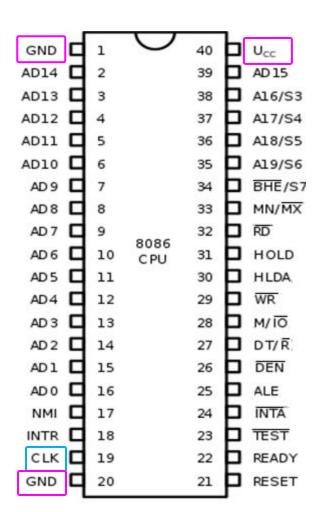
### Recap





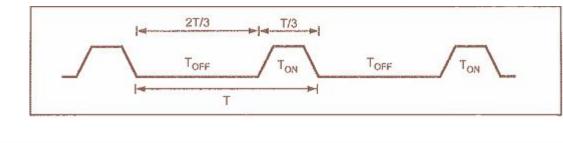
# 8086 Pin Specification

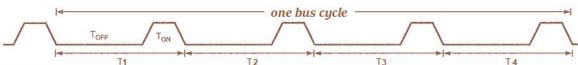




## CLK, input

- provides basic timing to control processor operation
- frequencies of different versions are 5, 8 or 10 MHz
- asymmetric with a 33% duty cycle

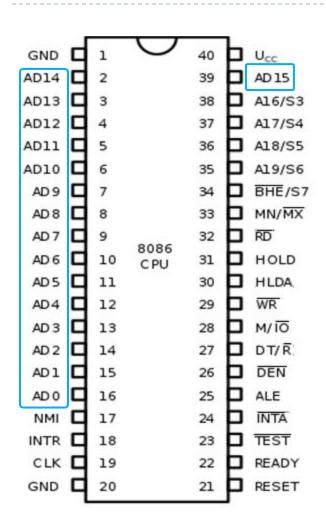










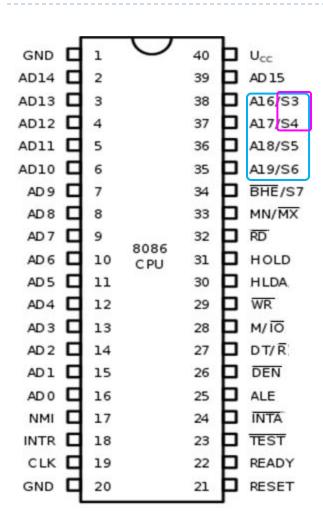


$$AD_0 - AD_{15}$$
,  $bi - directional$ 

- lines are multiplexed bidirectional address/data bus.
- During  $T_1$ , they carry 16-bit address.
- In remaining clock cycles  $T_2$ ,  $T_3$  and  $T_4$ , I 6-bit data.
- $AD_0 AD_7$  carry lower order data byte
- $AD_8 AD_{15}$  carry higher order data byte







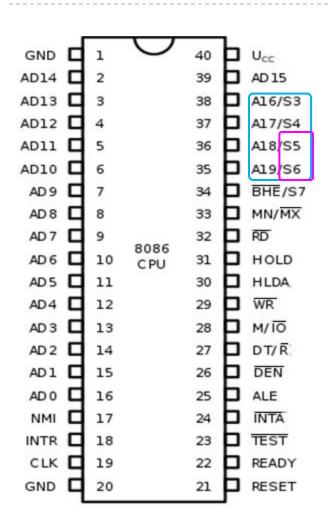
$$A_{19}/S_6$$
,  $A_{18}/S_5$ ,  $A_{17}/S_4$ ,  $A_{16}/S_3$ , **output**

- lines are multiplexed address and status bus.
- During  $T_1$ , they carry the highest order 4-bit address.
- During  $T_2$ ,  $T_3$  and  $T_4$ , status signals.
- $S_3$  and  $S_4$ , segment identifiers as in table below

S4	<i>S3</i>	Function
0	0	Extra segment access
0	1	Stack segment access
1	0	Code segment access
1	1	Data segment access







$$A_{19}/S_6$$
,  $A_{18}/S_5$ ,  $A_{17}/S_4$ ,  $A_{16}/S_3$ , **output**

 $S_5$ : Indicates if interrupt is enabled or disabled.

- If  $S_5 = I$ , then the IF = I, so the interrupt is enabled.
- If  $S_5 = 0$ , then the IF = 0, so the interrupt is disabled.

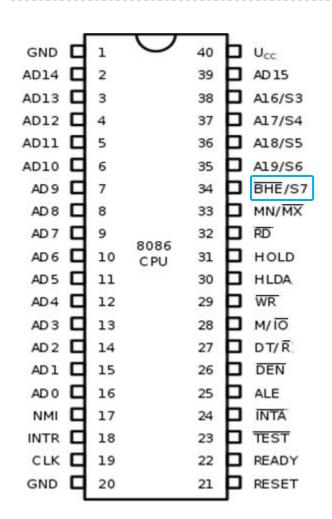
 $S_6$ : Indicates if 8086 is the bus master or not

- If  $S_6 = 0$ , 8086 is the bus master
- If  $S_6 = 1$ , 8086 is not the bus master









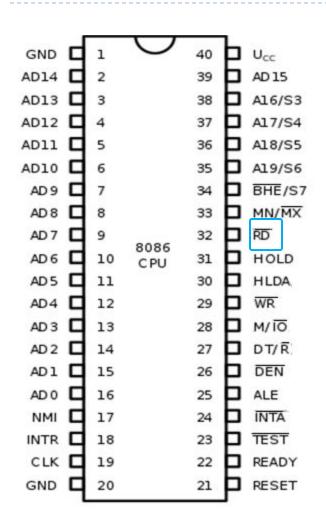
 $\overline{BHE}/S_7$ , output

- Bus High Enable
- $\overline{BHE}$  is active low
- To indicate the transfer of data over  $AD_8 AD_{15}$
- Related to memory bank
- Selects odd/high memory bank when  $\overline{BHE}$  is 0
- $S_7$ : Reserved for further development









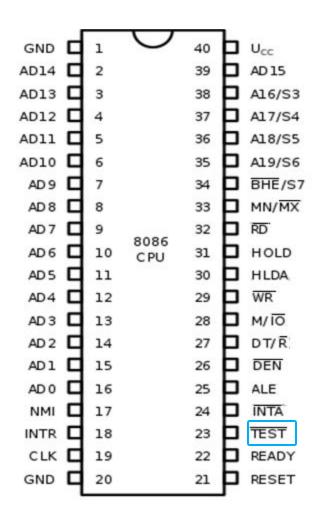
 $\overline{RD}$ , output

- is active low
- Indicates read operation when low
- Processor reading from memory or I/O device
- Is low during  $T_2$ ,  $T_3$  and  $T_w$  states of the read cycle







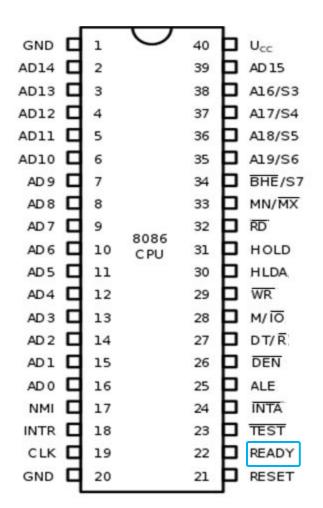


## $\overline{TEST}$ , input

- Is examined by the WAIT instruction.
- If this pin is Low, execution continues.
- Else the processor waits in an idle state.

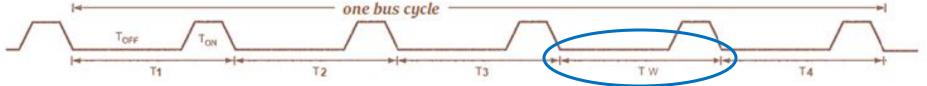
# 8086 Pin Specification





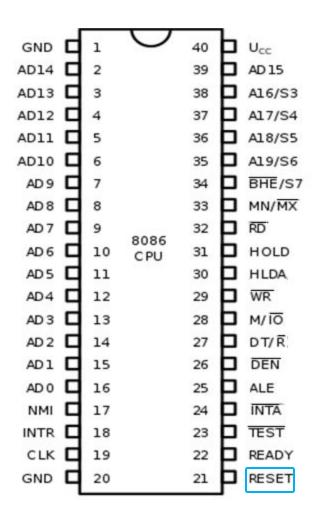
### READY, input

- acknowledgement from a slow I/O device or memory
- To indicate ready/completion of data transfer
- When low, microprocessor enters wait state,  $T_w$ .









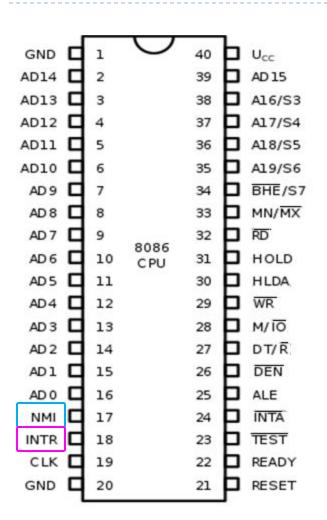
#### RESET, input

- To reset the system reset.
- And terminates the current activity.
- Must be active for at least four clock cycles









## INTR, input

- Interrupt request
- Used to request a hardware interrupt.
- Can be masked.

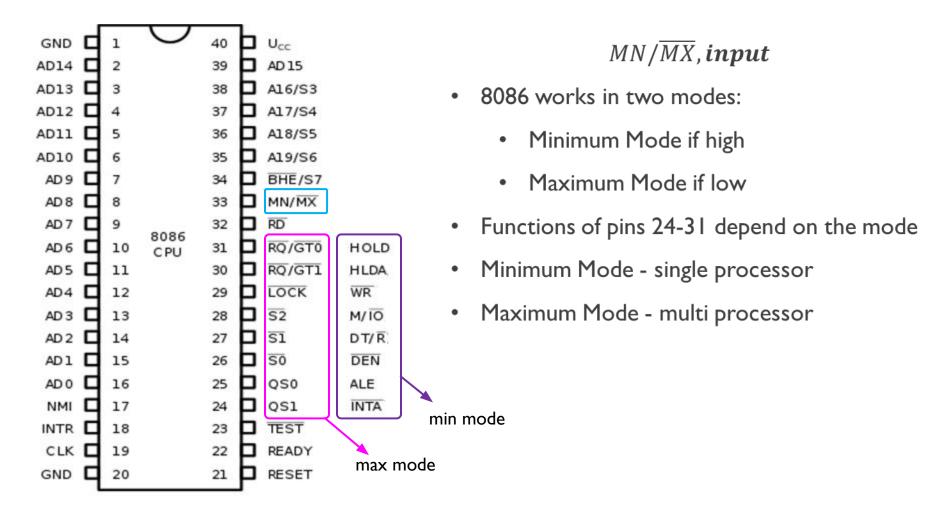
#### NMI, input

- Non-maskable interrupt signal.
- Causes a type-2 interrupt.
- Initiates the interrupt at the end of the current instruction.

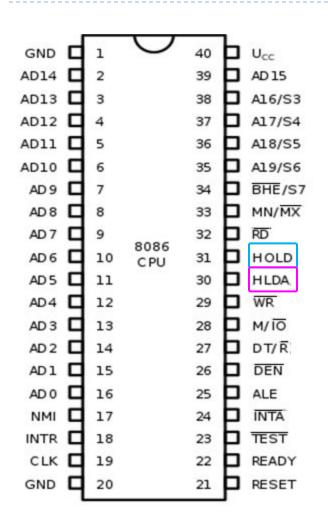












#### HOLD, input

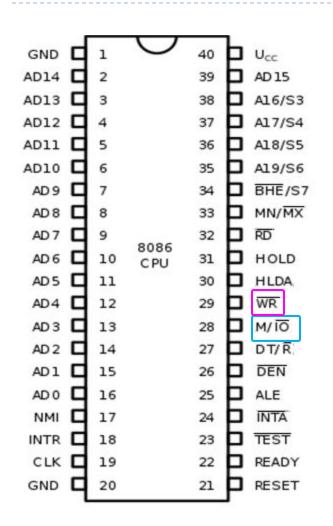
- To request for bus by another device.
- It is an active HIGH signal.

#### HLDA, output

- Hold Acknowledgment.
- When acknowledged, it relinquish the bus to the requesting device







## $\overline{WR}$ , output

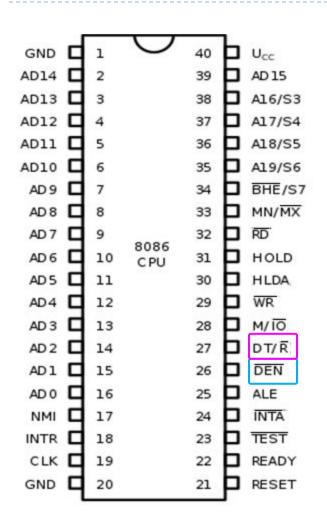
- Active low write signal.
- Writes data to memory or output device depending on  $M/\overline{IO}$  signal.

## $M/\overline{IO}$ , output

- Differentiates memory access from I/O access.
- When high, memory is accessed.
- When low, I/O devices are accessed.







## $DT/\bar{R}$ , output

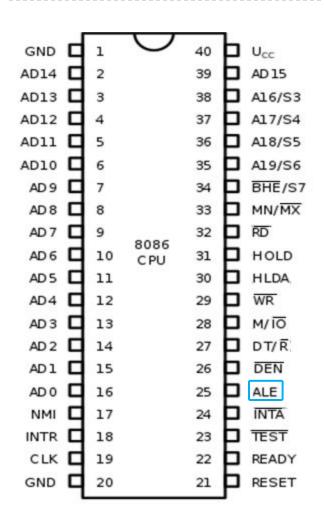
- Data Transmit/Receive signal.
- indicates the direction of flow through the transceiver.
- When high, data is transmitted out i.e. written to.
- When low, data is received in i.e. read in.

## $\overline{DEN}$ , output

- Data Enable signal.
- Used to enable a transceiver connected to the μP





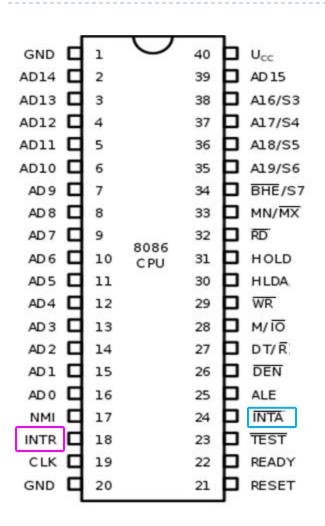


#### ALE, output

- Address Latch Enable
- indicates an address is available on bus  $AD_0 AD_{15}$ .
- active high during  $T_1$  state







## *INTA*, output

- An active low signal.
- An interrupt acknowledge signal.
- When microprocessor receives an INTR signal, it acknowledges the interrupt by generating this signal
- When low it indicates an interrupt is being serviced.



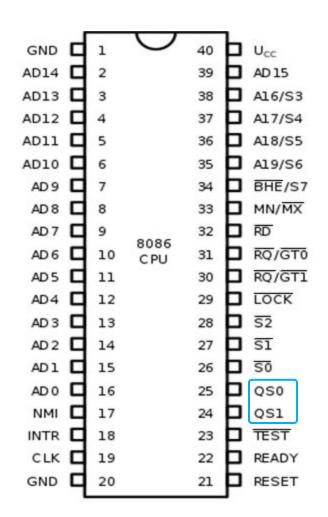


## 8086 Maximum Mode Pins

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## $QS_0$ and $QS_1$ , output

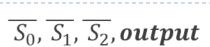
- Instruction queue status.
- Instruction queue is 6 bytes long

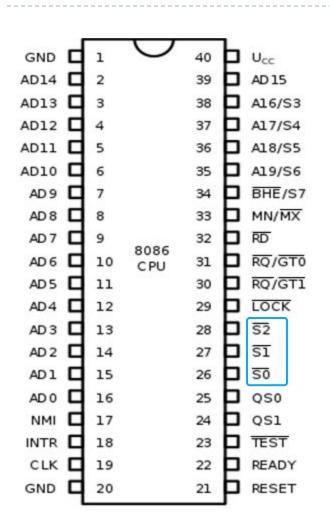
Qs1	Qs0	Function
0	0	No Operation. During the last clock cycle, nothing was taken from the queue.
0	1	First Byte. The byte taken from the queue was the first byte of the instruction.
1	0	Queue Empty. The queue has been reinitialized as a result of the execution of a transfer instruction.
1	1	Fetch subsequent byteSubsequent Byte. The byte taken from the queue was a subsequent byte of the instruction.







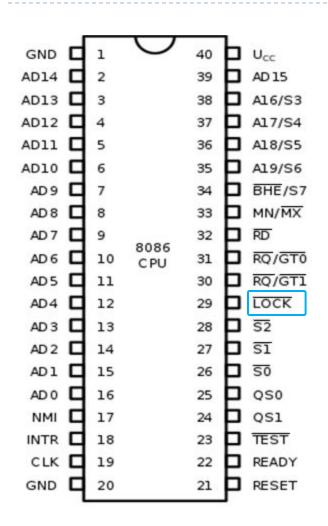




- Status Signals.
- indicate operation done by the microprocessor
- Related to memory and I/O access control signals.

S2	S1	S0	Function
0	0	0	Interrupt acknowledgement
0	0	1	Read data from I/O port
0	1	0	Write data from I/O port
0	1	1	Halt
1	0	0	Opcode fetch
1	0	1	Memory read
1	1	0	Memory write
1	1	1	Passive state





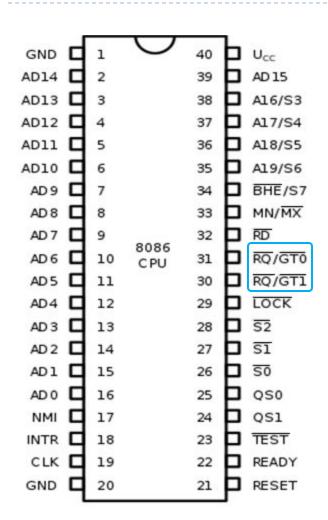
## $\overline{LOCK}$ , output

- When low, all interrupts are masked
- Indicates to other processors to not request for system bus.
- No HOLD request is granted.
- No bus is relinquished to the other processors









$$\overline{RQ/GT_0}$$
 and  $\overline{RQ/GT_1}$ , bi - directional

- Request/Grant pins.
- Other processors request the CPU through these for system bus.
- CPU sends acknowledge signal on the same lines.
- $\overline{RQ/GT_0}$  has higher priority than  $\overline{RQ/GT_1}$ .



# QUIZ



- Assuming you want to type a secret message using a keypad connected to an 8086 microprocessor, deduce the values of the following pins during that time. Justify your anwers too.
  - a)  $\overline{RD}$  e.g. mention if low (0) / high (1) and why.
  - b)  $\overline{WR}$
  - c)  $M/\overline{IO}$
- 2) Do you think there may be other pins involved? If so, justify your answer.



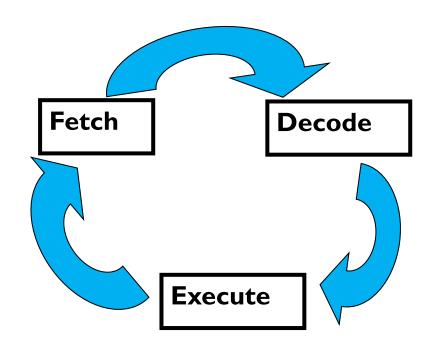
# 8086 Clocks & Timing Diagrams

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# Microprocessor Operation



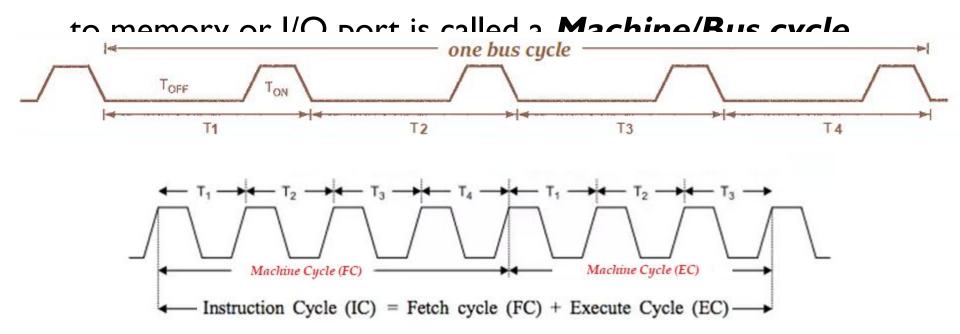
- An instruction e.g. MOV [7531h], AX; SUB CH, [0ABCh] etc
- The time a μP requires to complete fetch-(decode)-execute operation of a single instruction is known as *Instruction*





# Microprocessor Operation

- Instruction Cycle consists of one or more Machine
  Cycles
- $\square$  A basic  $\mu P$  operation such as reading/writing a byte from or



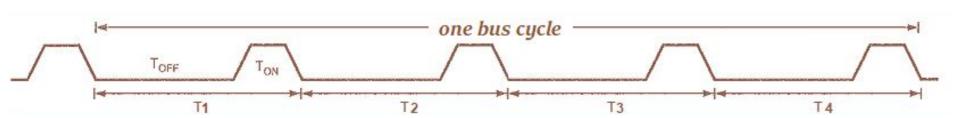






# Microprocessor Operation

- A Machine (bus) cycle consists of at least four clock
   cycles, called T states.
- One cycle of a clock is called a **State**
- Each read or write operation takes I bus cycle.

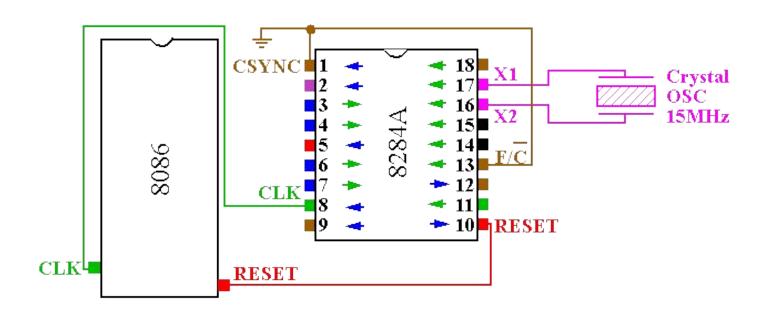








Clock generator circuit is 8284A and connected to pin19 (CLK) of 8086.



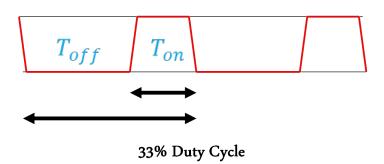


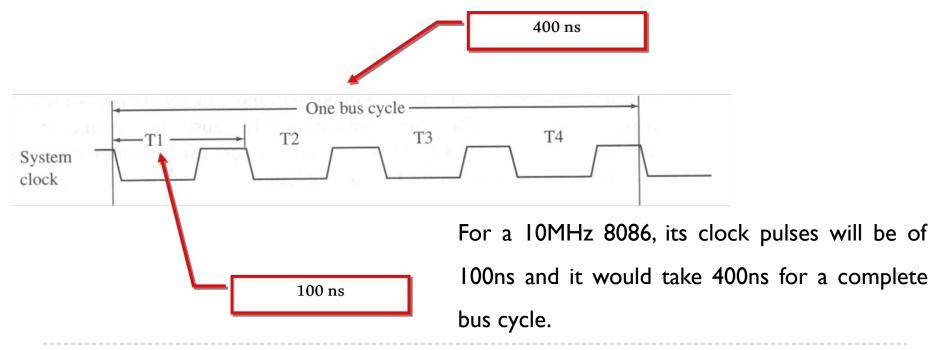




# System Clock Concept

- 8086 is found to operate in between 5 to 10 Mhz.
- An 8086 running at 5MHz, its clock pulses will be of 200ns and it would take 800ns for a complete bus cycle.





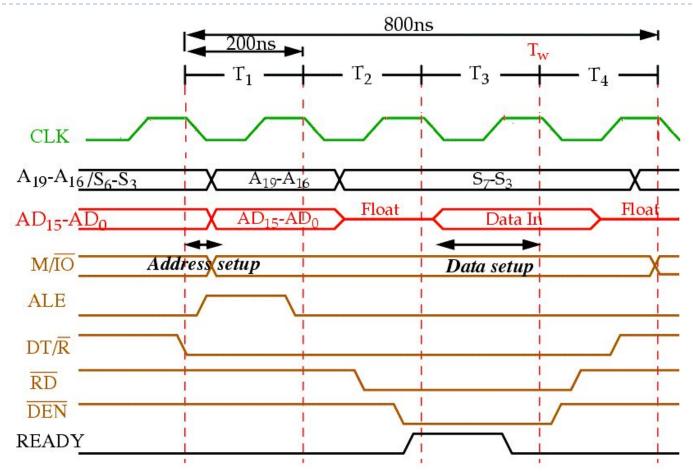


## Clock States - Why are there T states?

- In 8086, address and data lines are multiplexed to reduce number of pins e.g.  $AD_{0-15}$  else 32 pins would have been needed instead of 16
  - The μp needs time to change the signals during each bus cycle.
  - Memory devices need time to interpret the address value and then read/write the data (access time)
  - A specific defined action occurs during each T state  $(T_1 T_4)$ 
    - $\vdash T_1$ : Address is output
    - T<sub>2</sub>: Bus cycle type (Mem/IO, read/write)
    - ► T<sub>3</sub>: Data is supplied / Data is received
    - $T_4$ : Data latched by CPU, control signals removed



T<sub>1</sub>: Address is output

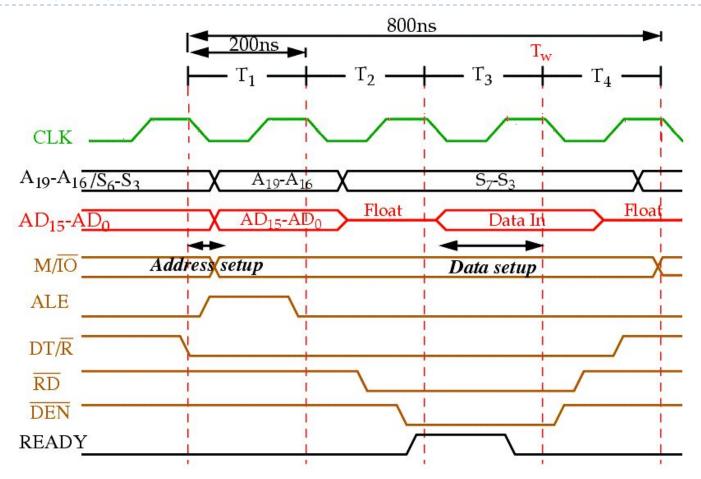


Address of memory is sent out by 8086 via address bus Used Control signals: ALE, DT/R', M/IO' shows some output





T<sub>2</sub>: Bus cycle type (MEMORY/IO, READ/WRITE)

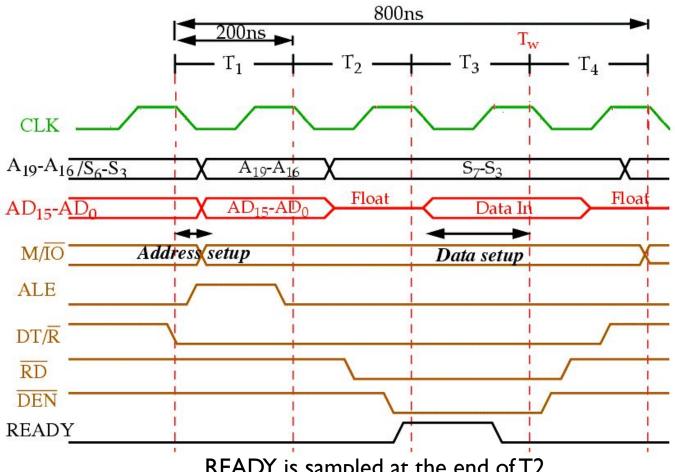


8086 issues either RD' or WR' and DEN' In case of WRITE (WR) operation, data to be written appear on data bus





T<sub>2</sub>: Data is supplied



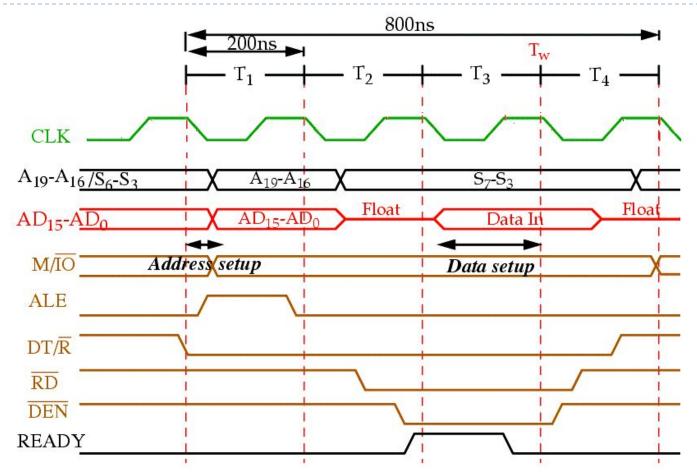
READY is sampled at the end of T2

If READY is low, T3 becomes a wait state (Tw), means no operation (NOP). In READ bus cycle data bus is sampled at end of T3





T<sub>2</sub>: Data latched by μP, control signals removed



All bus signals deactivated in preparation for next bus cycle  $\mu P$  sampled data bus for data that read from M or I/O



## **Clock States**



- $\Box$  A specific, defined action occurs during each T states  $(T_1 T_4)$
- T<sub>1</sub>: Address is output
  - Address of memory is sent out by 8086 via address bus
  - Used Control signals: ALE, DT/R', M/IO' shows some output
- $\square$   $\mathbf{T}_2$ : Bus cycle type (MEMORY/IO, READ/WRITE)
  - 8086 issues either RD' or WR' and DEN'
  - ☐ In case of **WRITE (WR)** operation, data to be written appear on data bus

## **Clock States**



- $\Box$   $T_3$ : Data is supplied
  - READY is sampled at the end of T<sub>2</sub>
    - If READY is low,  $T_3$  becomes a wait state  $(T_w)$ , means no operation (NOP).
    - In READ bus cycle data bus is sampled at end of T<sub>3</sub>
- $\Box$   $T_{a}$ : Data latched by  $\mu$ P, control signals removed
  - All bus signals deactivated in preparation for next bus cycle
  - □ µP sampled data bus for data that read from M or I/O
  - At trailing edge of WR', transfer data to M or I/O





- The READY input is controlled to insert "Wait states" into the timing of the microprocessor for slower memory and I/O components..
- If the READY pin is at a logic 0 level, the micro-processor enters into wait states and remains idle.
- When it is high (logic I), it indicates that the device is ready to transfer data.
- A wait state is a situation in which a computer processor is waiting for the completion of some event before resuming activity.
- A program or process in a wait state is inactive for the duration of the wait state.







- When a computer processor works at a faster <u>clock speed</u> than the random access memory (RAM) that sends it instructions, it is set to go into a wait state for one or more clock cycles so that it is synchronized with RAM speed. In general, the more time a processor spends in wait states, the slower the performance of that processor.
- Wait states are a pure waste for a processor's performance. Modern designs try to eliminate or hide them using a variety of techniques: <u>CPU caches</u>, <u>instruction pipelines</u>, <u>instruction prefetch</u>, <u>simultaneous multithreading</u> and others.





# Thank You Questions are welcome in the discussion class