

CSE321

Take Home Quiz

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sec: 8 (suw)

Spring 24

①

Ans to Q 1

	Allocation R_1, R_2, R_3	max R_1, R_2, R_3	Available R_1, R_2, R_3	need R_1, R_2, R_3
P_0	2 5 0	8 5 9	10 7 7	6 0 9
P_1	1 1 1	5 6 8		4 5 7
P_2	2 5 2	6 10 8		4 5 6
P_3	4 0 3	5 6 10		1 6 7
P_4	0 3 0	9 5 9		9 2 9

Finish :

F	R_T	R_T	R_T	R_T
P_0	P_1	P_2	P_3	P_4

sequence : $P_1 \rightarrow P_2 \rightarrow P_3 \rightarrow P_4 \rightarrow P_0$

① P_0 : No allocation

② P_1 : Allocate

work =

11	8	8
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③ P_2 : Allocate

work =

13	13	10
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④ P_3 : Allocate

work =

17	13	13
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⑤ P_4 : Allocate

work =

17	16	13
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⑥ P_0 : Allocate

work =

19	21	13
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\therefore The state is safe :

$P_1 \rightarrow P_2 \rightarrow P_3 \rightarrow P_4 \rightarrow P_0$

11

Request = [0, 2, 7]

for P_3 :

① Available = [10, 5, 0]

② Allocated = [4, 2, 10]

③ need = [1, 4, 0]

Banker's algorithm:

$F = \begin{bmatrix} P_T & P_T & P_T & P_T & P_T \\ P_0 & P_1 & P_2 & P_3 & P_4 \end{bmatrix}$

seq = $P_3 \rightarrow P_4 \rightarrow P_0 \rightarrow P_1 \rightarrow P_2$

① P_0 : 6, 0, 9 \nless 10, 5, 0

② P_1 : 4, 5, 7 \nless 10, 5, 0

③ P_2 : 4, 5, 6 \nless 10, 5, 0

④ P_3 : 1, 4, 0 \less 10, 5, 0 : Allocate
work = [14, 7, 10]

⑤ P_4 : 3, 2, 9 \less 14, 7, 10 : Allocate
work = [14, 10, 10]

⑥ P_0 : 6, 0, 9 \less 14, 10, 10 : Allocate
w = [16, 15, 10]

⑦ P_1 : 4, 5, 7 \less 14, 10, 10 : Allocate
w = [17, 16, 11]

⑧ P_2 : 4, 5, 6 \less 14, 10, 10 : Allocate
w = [19, 21, 13]

\therefore Yes system remain safe.

①

Ans to Q 2

	Need			Available R_1, R_2, R_3
	R_1	R_2	R_3	
P_0	10	5	3	$2, 5, 7$
P_1	4	3	3	
P_2	2	3	2	
P_3	2	3	2	
P_4	3	8	6	

$$f = \left[\begin{array}{c|c|c|c|c} R_1 & R_2 & R_3 & R_4 & R_5 \\ \hline P_0 & P_1 & P_2 & P_3 & P_4 \end{array} \right] \quad W = [2, 5, 7]$$

$$\text{Seq} = P_2 \rightarrow P_3 \rightarrow P_4 \rightarrow P_0 \rightarrow P_1$$

$$\textcircled{1} P_0 : 10, 5, 3 \not\leq 2, 5, 7$$

$$\textcircled{2} P_1 : 4, 3, 3 \not\leq 2, 5, 7$$

$$\textcircled{3} P_2 : 2, 3, 2 \leq 2, 5, 7 : W = [5+2, 4+5, 5+7] \\ = [7, 9, 12]$$

$$\textcircled{4} P_3 : 2, 3, 2 \leq 7, 9, 12 : \text{Allocate}$$

$$W = [7+3, 9+2, 12+2] = [10, 11, 14]$$

$$\textcircled{5} P_4 : 3, 8, 6 \leq 10, 11, 14 : \text{allocate}$$

$$W = [10+3, 11+8, 14+6] = [13, 19, 20]$$

$$\textcircled{6} P_0 : 10, 5, 3 \leq 13, 19, 20 : \text{allocate}$$

$$W = [13+10, 19+5, 20+3] = [23, 24, 23]$$

$$\textcircled{7} P_1 : 4, 3, 3 \leq 23, 24, 23$$

$$W = [23+4, 24+3, 23+3] = [27, 27, 26]$$

\therefore Safe state : $P_2 \rightarrow P_3 \rightarrow P_4 \rightarrow P_0 \rightarrow P_1$

ii) for P_3 : $P_3 < \text{Need}$, $P_3 < \text{Available}$

$$P_3: \begin{array}{ccc} \text{Allocation} & \text{Available} & \text{need} \\ [0, 2+1, 2+3] & [2, 5-1, 7-2] & [2, 3-1, 2-2] \\ = [0, 3, 5] & [2, 4, 5] & [2, 2, 0] \end{array}$$

$$F = \begin{array}{|c|c|c|c|c|} \hline 0 & 1 & 2 & 3 & 4 \\ \hline F & F & F & F & F \\ \hline \end{array} \quad W = [2, 4, 5]$$

$$S \Rightarrow P_2 \rightarrow P_3 \rightarrow P_4 \rightarrow P_0 \rightarrow P_1$$

① P_0 : $10, 5, 3 \neq 2, 4, 5$

② P_1 : $4, 3, 3 \neq 2, 4, 5$

③ P_2 : $2, 3, 2 < 2, 4, 5$: $W = [2+5, 4+4, 5+5]$
 $= [7, 8, 10]$

④ P_3 : $2, 2, 0 < 7, 8, 10$: Allocate

$$W = [7, 8+3, 10+5] = [7, 11, 15]$$

⑤ P_4 : $3, 8, 6 < 7, 11, 15$: Allocate

$$W = [7+3, 11+2, 15] = [10, 13, 15]$$

⑥ P_0 : $10, 5, 3 \leq 10, 13, 15$:

$$W = [10, 15, 18]$$

⑦ P_1 : $4, 3, 3 \leq 10, 15, 18$:

$$W = [10+2, 15+3, 18+2]$$

$$= [12, 18, 20]$$

Yes, system in safe state

$$P_2 \rightarrow P_3 \rightarrow P_4 \rightarrow P_0 \rightarrow P_1$$

Ans to q 3

	Need		
	R_1	R_2	R_3
P_0	5	6	2
P_1	5	9	0
P_2	5	6	5
P_3	3	9	3
P_4	1	7	0

$$W = [2, 9, 1]$$

$$F = \begin{array}{|c|c|c|c|c|} \hline 0 & 1 & 2 & 3 & 4 \\ \hline \cancel{F} & \cancel{F} & \cancel{F} & \cancel{F} & \cancel{F} \\ \hline \end{array}$$

$$S: P_4 \rightarrow P_0 \rightarrow P_1 \rightarrow P_2 \rightarrow P_3$$

$$① P_0: 5, 6, 2 \not< 2, 9, 1$$

$$② P_1: 5, 9, 0 \not< 2, 9, 1$$

$$③ P_2: 5, 6, 5 \not< 2, 9, 1$$

$$④ P_3: 3, 9, 3 \not< 2, 9, 1$$

$$⑤ P_4: 1, 7, 0 < 2, 9, 1: \text{Allocate}$$

$$W = [4+2, 1+9, 5+1] = [6, 10, 6]$$

$$⑥ P_0: 5, 6, 2 < 6, 10, 6:$$

$$W = [5+6, 1+10, 5+6] = [11, 11, 11]$$

$$⑦ P_1: 5, 9, 0 < 11, 11, 11:$$

$$W = [3+11, 11, 5+11] = [14, 11, 16]$$

$$⑧ P_2: 5, 6, 5 < 14, 11, 16:$$

$$W = [14, 1+11, 4+16] = [14, 12, 20]$$

$$⑨ P_3: 3, 9, 3 < 14, 12, 20:$$

$$W = [18, 13, 22]$$

safe state: $P_4 \rightarrow P_0 \rightarrow P_1 \rightarrow P_2 \rightarrow P_3$

ii) P_0 request = $[2, 3, 0]$ $<$ need
 \leq available

	allocated	available	need
P_0	$[7, 4, 5]$	$[0, 6, 1]$	$[3, 3, 2]$

if $F = \boxed{F/F/F/F/F}$ S:

$W = [0, 6, 1]$

① P_0 : $3, 3, 2 \nless 0, 6, 1$

② P_1 : $5, 9, 0 \nless 0, 6, 1$

③ P_2 : $5, 6, 5 \nless 0, 6, 1$

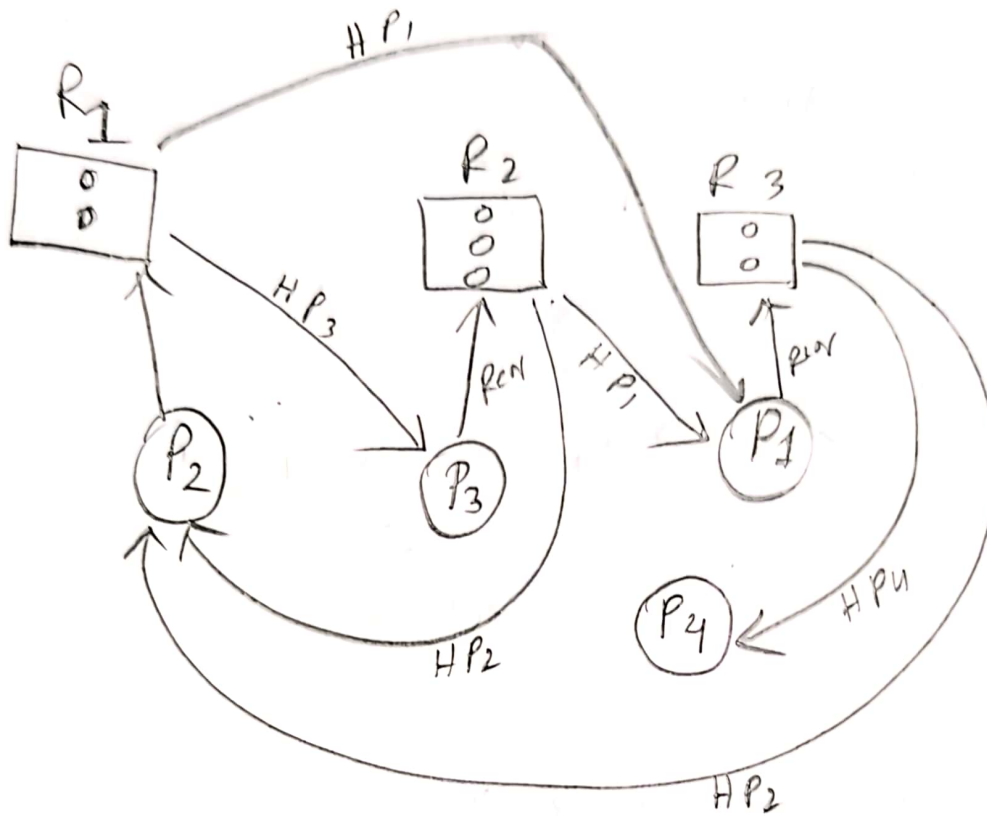
④ P_3 : $3, 9, 3 \nless 0, 6, 1$

⑤ P_4 : $1, 7, 0 \nless 0, 6, 1$

second iteration same result

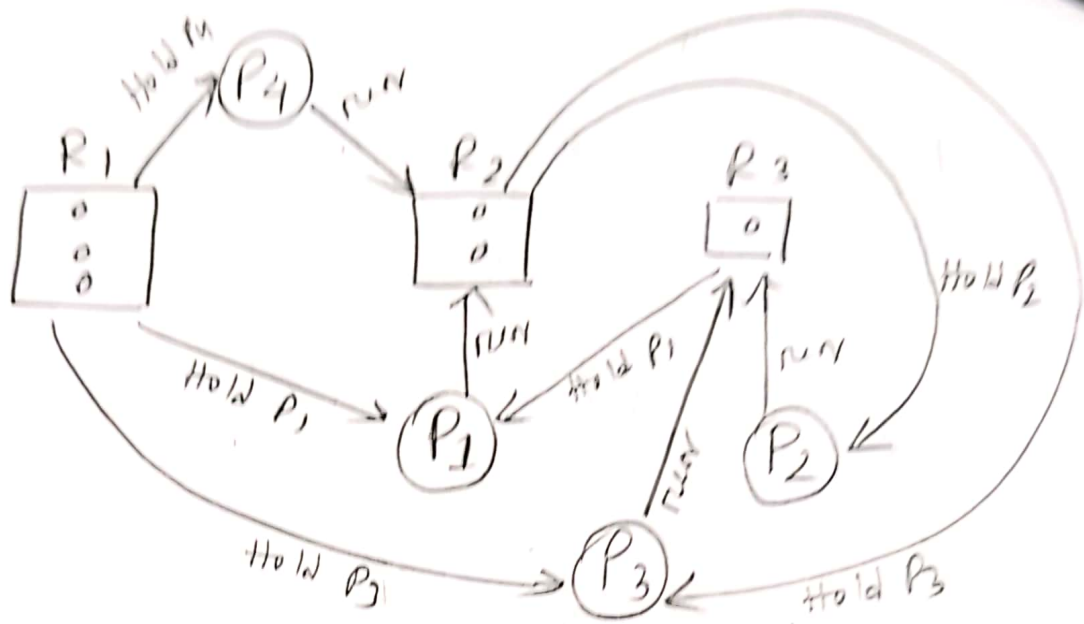
\therefore No, the system doesn't remain in safe state if request approved.

Ans to or 4



In this graph no mutual exclusion condition holds which is why there is no possibility of deadlock.

Ans to Q 5



conditions exist:

- ① Mutual exclusion : R3 has one resource
- ② Hold and wait : ex: P4 holds R1, req for R2
- ③ No preemption.
- ④ Circular wait :
ex: $P_1 \rightarrow R_2 \rightarrow P_2 \rightarrow R_3 \rightarrow P_1$

∴ Hence, possibility of deadlock exist.

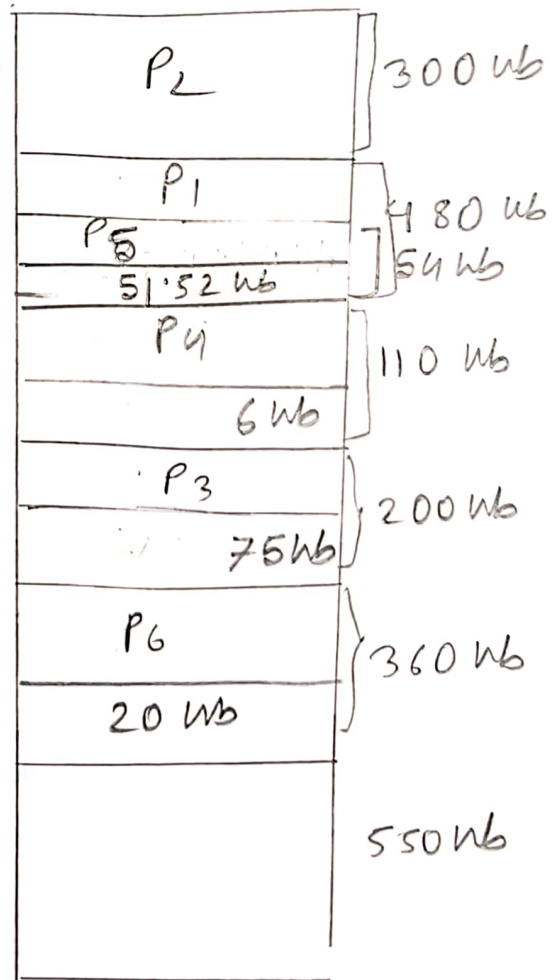
Main Memory Allocation

(1)

(i) First fit :

Process	space	(wb) Allocated to	(wb) waste/hole
P ₁	426 wb	480 wb	54 wb
P ₂	300 wb	300 wb	0 wb
P ₃	125 wb	200 wb	75 wb
P ₄	104 wb	110 wb	6 wb
P ₅	475 b	54 wb	51.52 wb
P ₆	340 wb	360 wb	20 wb

∴ There exists external fragmentation for allocation of future processes as there are holes in between allocated processes.



(ii) Best fit:

P	is	Allocated to	hole
P ₁	426 kb	480 kb	54 kb
P ₂	300 kb	300	0
P ₃	125 kb	200	75
P ₄	104 kb	110	6
P ₅	475 B	6 kb	5.52
P ₆	340	360	20

∴ As holes between process exists, external fragmentation exists.

Compare:

① (1) First fit biggest continuous space
= 550 + 20 = 570 kb

(ii) Best fit = 570 kb

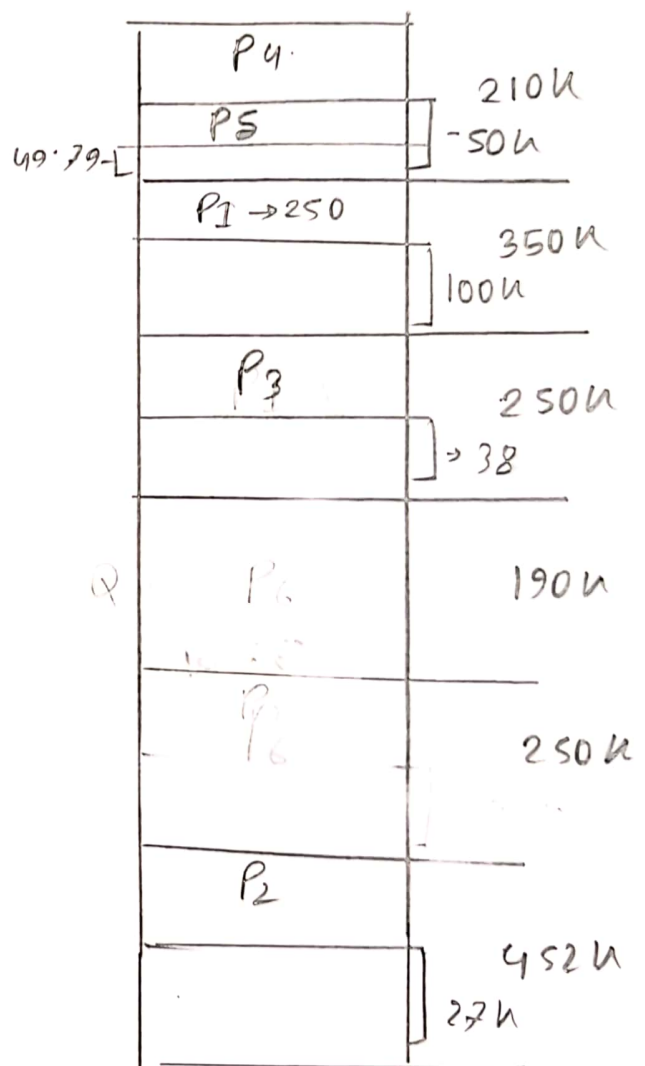
② Comparing other aspects in first fit -
allocating P₅ created 5.52 kb waste and
in best fit it created 5.1 kb waste.
Hence, best fit is more effective.

P ₂ → 300	300 kb
P ₁ → 426 kb	480 kb
	54 kb
P ₄ → 104 kb	110 kb
P ₅ → 475 B	6 kb
	→ 5.52 kb
P ₃ → 125	200 kb
	75 kb
P ₆ → 340	360 kb
	20 kb
	550 kb

② (i) First fit :

P	space	space allocated	waste
P ₁	250K	—	—
P ₂	425K	—	—
P ₃	212K	—	—
P ₄	160K	—	—
P ₅	210K	—	—
P ₆	440K	—	—
		10X	0X

∴ External fragmentation exists as holes exist in between allocated processes.



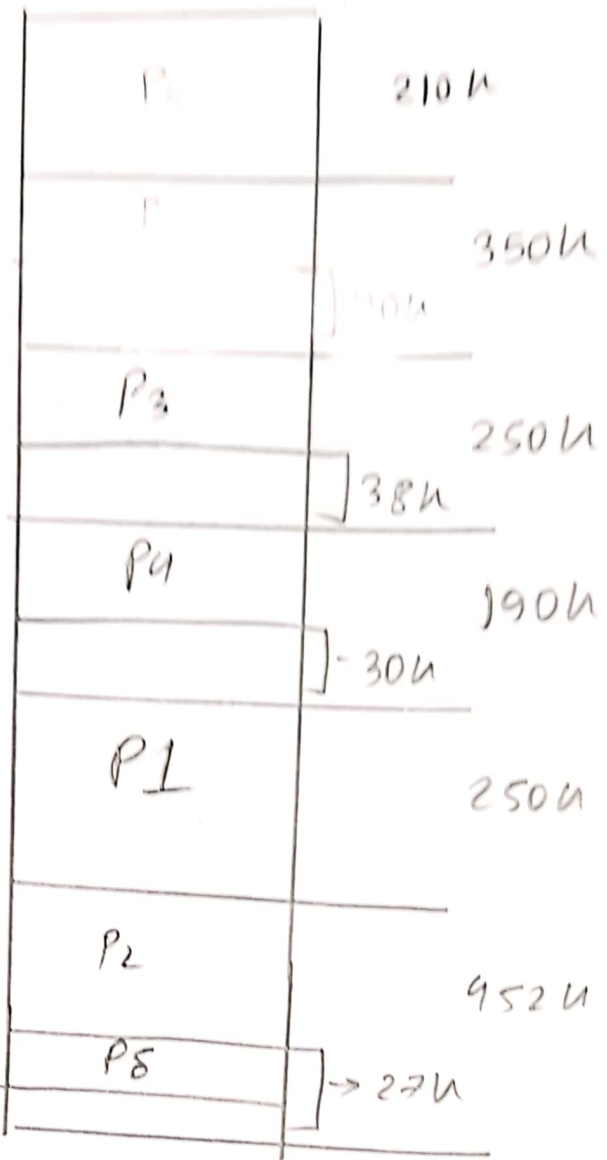
(ii) Best fit

P	Allocated to	req
P ₁	250	0
P ₂	452	22
P ₃	250	38
P ₄	190	30
P ₅	27	26.29
P ₆	210 X 0	1 X

External fragmentation exists.

Compare

26.29



① First fit biggest continuous.

$$\text{space} = 138 + 190 + 250 = 478 \text{ Kb}$$

② Best fit biggest continuous space = 210 + 350 Kb

∴ Best fit is more efficient.