

# Pattern Avoidance in Sequences

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## 1 Introduction

## 2 Saturation for Sequences

## 3 New Results

In the field of **extremal combinatorics**, we ask questions like

## Extremal Question

What is the largest possible size of an object which avoids a given forbidden substructure?

For example:

- What is the largest possible graph on  $n$  vertices that does not contain  $K_3$  as a subgraph? (Turan's Theorem & the forbidden subgraph problem)
- What is the largest possible group of people, such that for any set of  $k$  people, they are not all friends or not all strangers? (Ramsey Theory)
- What is the largest possible subset of  $\{1, \dots, n\}$  that does not contain a  $k$ -term arithmetic progression? (Szemerédi's Theorem)

The **saturation question** is a bit more complicated.

## Saturation Question

What is the SMALLEST possible structure that avoids a given forbidden substructure, BUT making it larger in any way induces a copy of the forbidden structure?

In other words: the *minimum* size of a *maximal* structure, rather than the *maximum* size.

## Definition

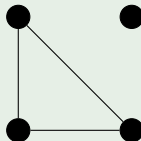
Let  $G$  and  $H$  be graphs. We say  $G$  is  $H$ -saturated if  $G$  avoids  $H$  as a subgraph, but adding any new edge to  $G$  induces a copy of  $H$ .

## Example

Consider



The graph



is  $H$ -saturated.

## Definition

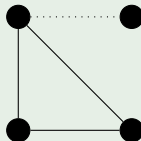
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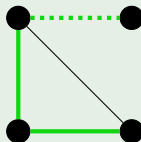
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## Example

Consider



The graph



is  $H$ -saturated.



## Definition

The **saturation function**  $\text{Sat}(n, H)$  is the minimum number of edges in a  $H$ -saturated graph on  $n$  vertices.

$\text{Sat}(n, H)$  exhibits an **dichotomy**:

## Theorem (Kászonyi-Tuza, 1986)

*We have  $\text{Sat}(n, H) = O(1)$  or  $\text{Sat}(n, H) = \Theta(n)$ .*

In many other settings, it has been seen that the saturation function exhibits the same dichotomy.

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# Sequence Saturation

We use “letters” to refer to the terms in a sequence.

## Definition

Let  $s$ ,  $u$  be two sequences. We say that  $s$  contains a **copy** of  $u$  if  $s$  has a subsequence that can be turned into  $u$  by a *one-to-one* renaming of letters.

## Example

The sequence  $s = 1, 2, 3, 2, 3, 1, 2$  contains a copy of  $u = abab$ :

|          |          |   |   |   |          |          |
|----------|----------|---|---|---|----------|----------|
| 1        | 2        | 3 | 2 | 3 | 1        | 2        |
| <i>a</i> | <i>b</i> |   |   |   | <i>a</i> | <i>b</i> |

However,  $s = 1, 2, 3, 2, 1$  does not.

## Definition

A sequence  $s$  is  $r$ -sparse if every consecutive  $r$  letters are pairwise distinct.

## Example

The sequence  $s = 1, 2, 3, 2, 3, 1, 2$  is 2-sparse, but not 3-sparse.

## Definition

Let  $u$  be a sequence with  $r$  distinct letters. A sequence  $s$  is  $u$ -saturated if  $s$  avoids  $u$ ,  $s$  is  $r$ -sparse, and inserting any new letter into  $s$  either induces a copy of  $u$  or violates  $r$ -sparsity.

If we dropped the  $r$ -sparsity condition, we would have arbitrarily long sequences like  $1, 1, \dots$  which avoid  $u$ .

## Example

If  $u = abca$  then the sequence  $s = 1, 2, 3$  is  $u$ -saturated.

First,  $r = 3$  since  $u$  has 3 distinct letters. Now we check:

$s$  **avoids**  $u$ : Evident.

$s$  **is 3-sparse**: Evident.

**Saturation**: Suppose we insert a 1 into  $s$ . The possibilities are:

1 1 2 3 (violates 3-sparsity)

1 1 2 3 (violates 3-sparsity)

1 2 1 3 (violates 3-sparsity)

1 2 3 1  
a b c a (has a copy of  $u$ )

Similar checks for the other letters.

## Definition

The **saturation function**  $\text{Sat}(n, u)$  is the length of the shortest  $u$ -saturated sequence with  $n$  distinct letters.

In 2021, Anand, Geneson, Kaustav, and Tsai conjectured

## Conjecture

*We have  $\text{Sat}(n, u) = O(n)$ .*

This implies the dichotomy  $\text{Sat}(n, u) = O(1)$  or  $\Theta(n)$ . They proved

## Theorem (Anand-Geneson-Kaustav-Tsai, 2021)

*We have  $\text{Sat}(n, u) = O(n)$  for all sequences  $u$  with two distinct letters.*

However, the cases for  $u$  having  $\geq 3$  distinct letters remained completely open.

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# Algorithm

Consider the following algorithm:

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```
1: Input: Alphabet  $A = \{1, \dots, n\}$ , forbidden sequence  $u$ 
2: Output:  $u$ -saturated sequence
3: Initialize the sequence:  $s \leftarrow 1, 2, \dots, r - 1$            ▷ Initial sequence avoids  $u$ 
4: while it is possible to extend the sequence do
5:   for each letter  $x \in A$  do
6:     if  $x$  can be properly inserted into  $s$  then
7:       Insert  $x$  appropriately into  $s$  to form  $s'$            ▷ Smallest  $x$ , leftmost position
8:       Update  $s \leftarrow s'$                                ▷ New sequence
9:       break                                               ▷ Exit loop after the first valid insertion
10:    end if
11:  end for
12: end while
13: Return  $s$                                                ▷ Final sequence is  $u$ -saturated
```

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The output of the algorithm:

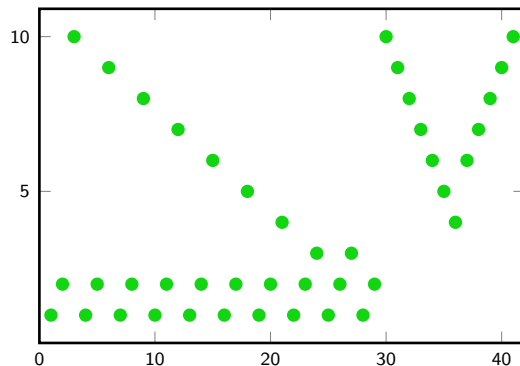


Figure: Algorithm on  $u = abcacbc$ .

Here, we represent  $s = s_1 \cdots s_\ell$  by plotting the points  $(i, s_i)$ . Using this pattern, we get  $\text{Sat}(n, abcacbc) = O(n)!$

Some more pictures:

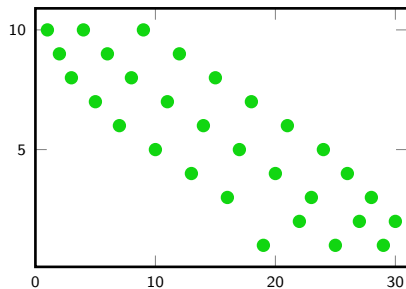
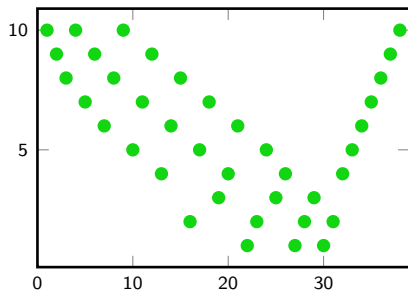


Figure: Algorithm on  $u = abbacac$  (left) and  $abcacba$  (right).

And even more:

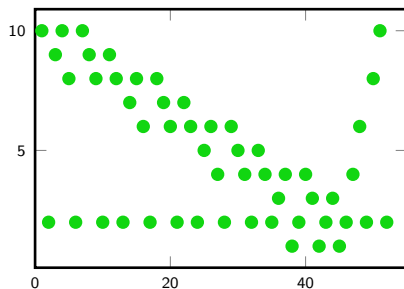
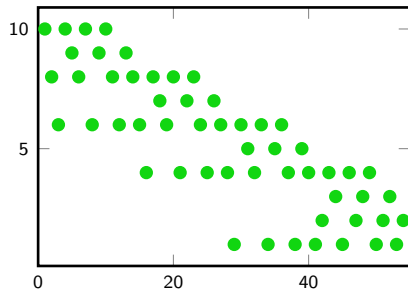


Figure: Algorithm on  $u = abcacbacb$  (left) and  $abcbacbac$  (right).

This lets us resolve the conjecture for many specific sequences  $u$ .

Say  $u$  is **irreducible** if  $u$  cannot be decomposed into sequences  $u = u_1 u_2$  such that  $u_1$  and  $u_2$  have no letters in common.

**Theorem (Kanungo, 2025+)**

*If  $u$  is irreducible and of the form  $aa \cdots bb$ , then  $\text{Sat}(n, u) = O(n)$ .*

Suppose  $u$  is a sequence on 3 letters, and  $u = abc \cdots xyz$  where  $a, b, c$  are distinct. Define

$$f_0(u) = \#\{\text{consecutive pairs of the form } ab, bc, ca\},$$

$$f_1(u) = \#\{\text{consecutive pairs of the form } ac, ba, cb\}.$$

## Theorem (Kanungo, 2025+)

Let  $u = abc \cdots xyz$  be a three-letter sequence with  $a, b, c$  distinct. Suppose

$$xyz \in \{abc, bca, cab\}, \quad f_0(u) \geq f_1(u) + 5.$$

Then  $\text{Sat}(n, u) = O(n)$ .

## Corollary

For any sequence  $u$  on 3 letters,  $\text{Sat}(n, (abc)u(abc)^t) = O(n)$  for large enough  $t$ .

- My mentor, Prof. Jesse Geneson, for his guidance and support.
- The PRIMES organizers, for making this experience possible.
- My parents, without whom none of this would be possible.

Questions?

Thank You!

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