



<http://pl.skuniv.ac.kr>

컴파일러 입문

제 8 장 LR 구문 분석



목 차

8.1 LR Parser

8.7 Implementation of an LR Parser



8.1 LR Parser

- an efficient *Bottom-up* parser for a large and useful class of context-free grammars.
- the "L" stands for *left*-to-right scan of the input; the "R" for constructing a *Rightmost* derivation in reverse.
- The attractive reasons of LR parsers
 - (1) LR parsers can be constructed for *most* programming languages.
 - (2) LR parsing method is more *general* than LL parsing method.
 - (3) LR parsers can detect *syntactic errors* as soon as possible.

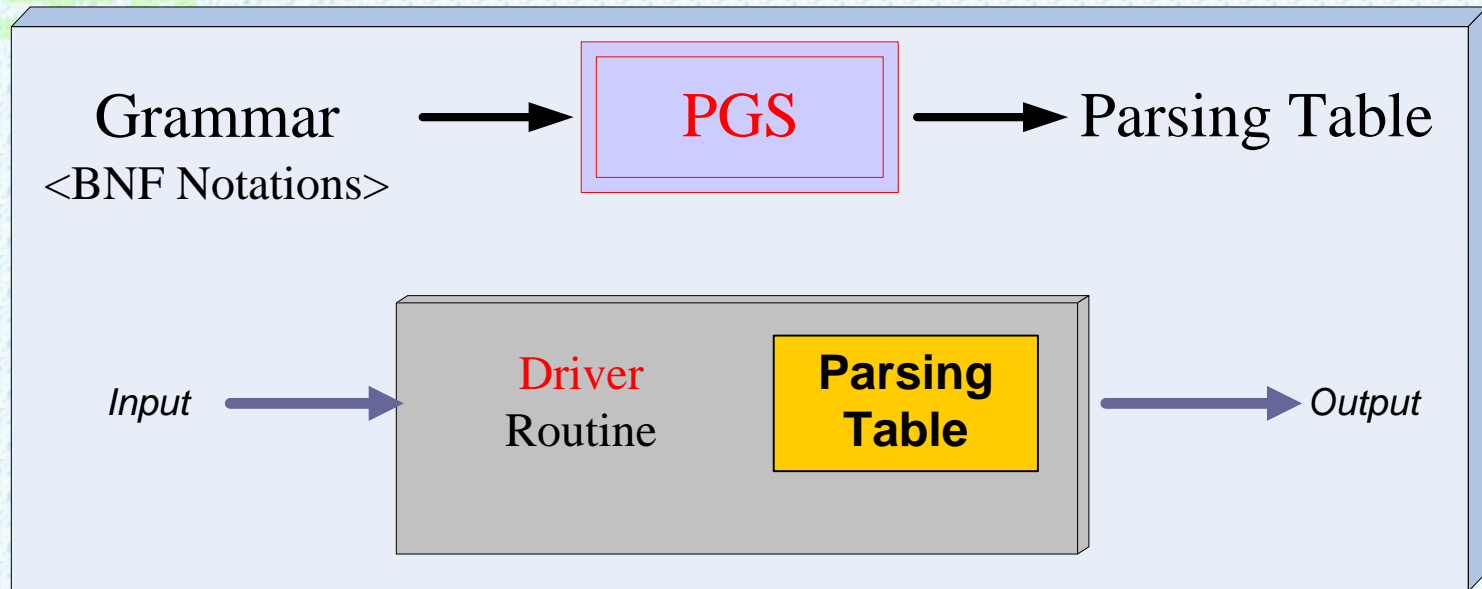
But,

- it is too much work to implement an LR parser *by hand* for a typical programming-language grammar.

=====> ∴ **Parser Generator**



Parser Generating Systems

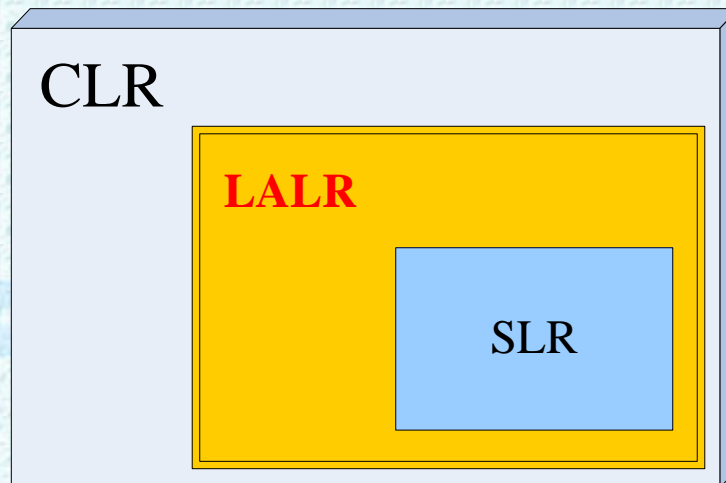


- The driver routine is the **same** for all LR parsers; only the **parsing table** changes from one parser to another.



Three Methods

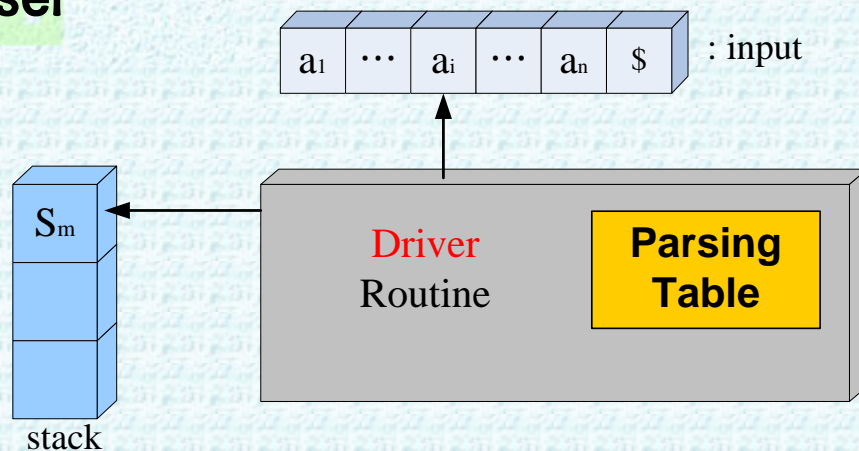
- The *techniques* for producing LR parsing tables
 - Simple LR(*SLR*) - LR(0) items, *FOLLOW*
 - Canonical LR(*CLR*) - LR(1) items
 - Lookahead LR(*LALR*) - ① LR(1) items
② LR(0), *Lookahead*





LR Parser의 구조 [1/3]

■ LR parser



■ **Stack** : $S_0X_1S_1X_2 \dots X_mS_m$, where S_i : state and $X_i \in V$.

■ **Configuration of an LR parser** :

$(\underbrace{S_0X_1S_1 \dots X_mS_m}_{\text{stack contents}}, \underbrace{a_ia_{i+1} \dots a_n\$}_{\text{unscanned input}})$



LR Parser의 구조 [2/3]

- LR Parsing Table (**ACTION** table + **GOTO** table)

		ACTION Table	GOTO Table
states \ symbols		<Terminals>	<Nonterminals>
		:	:

- The LR parsing algorithm

::= same as the **shift-reduce** parsing algorithm.

- Four Actions :

- shift
- reduce
- accept
- error



LR Parser의 구조 [3/3]

1. ACTION[S_m, a_i] = *shift* S

$::= (S_0 X_1 S_1 \dots X_m S_m, a_i a_{i+1} \dots a_n \$)$

$\Rightarrow (S_0 X_1 S_1 \dots X_m S_m \mathbf{a_i S}, a_{i+1} \dots a_n \$)$

2. ACTION[S_m, a_i] = *reduce* $A \rightarrow \alpha$ and $|\alpha| = r$

$::= (S_0 X_1 S_1 \dots X_m S_m, a_i a_{i+1} \dots a_n \$)$

$\Rightarrow (S_0 X_1 S_1 \dots \mathbf{X_{m-r} S_{m-r}}, a_i a_{i+1} \dots a_n \$), \text{GOTO}(S_{m-r}, A) = S$

$\Rightarrow (S_0 X_1 S_1 \dots X_{m-r} S_{m-r} \mathbf{AS}, a_i a_{i+1} \dots a_n \$)$

3. ACTION [S_m, a_i] = *accept*, parsing is completed.

4. ACTION [S_m, a_i] = *error*, the parser has discovered an error and calls an error *recovery* routine.



LR 파싱 예제 [1/2]

- G:
1. LIST \rightarrow LIST , ELEMENT
 2. LIST \rightarrow ELEMENT
 3. ELEMENT \rightarrow a

■ Parsing Table :

states \ symbols	a	,	\$	LIST	ELEMENT
0	s3			1	2
1		s4	acc		
2		r2	r2		
3		r3	r3		
4	s3				5
5		r1	r1		

where, **sj** means *shift* and stack state j,
ri means *reduce* by production numbered i,
acc means *accept*, and blank means *error*.



LR 파싱 예제 [2/2]

- Input : $\omega = a, a$
- Parsing Configuration :

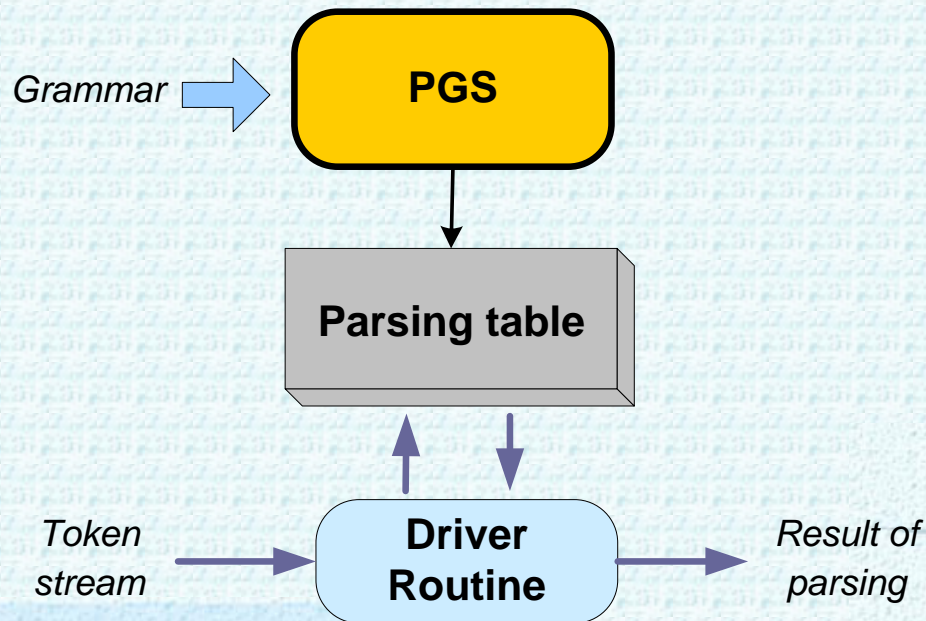
initial
configuration

STACK	INPUT	ACTION
0	a,a\$	s3
0 a 3	,a\$	r3 GOTO 2
0 ELEMENT 2	,a\$	r2 GOTO 1
0 LIST 1	,a\$	s4
0 LIST 1, 4	a\$	s3
0 LIST 1, 4 a 3	\$	r3 GOTO 5
0 LIST 1, 4 ELEMENT 5	\$	r1 GOTO 1
0 LIST 1	\$	accept



8.7 구문 분석기의 작성 [1/3]

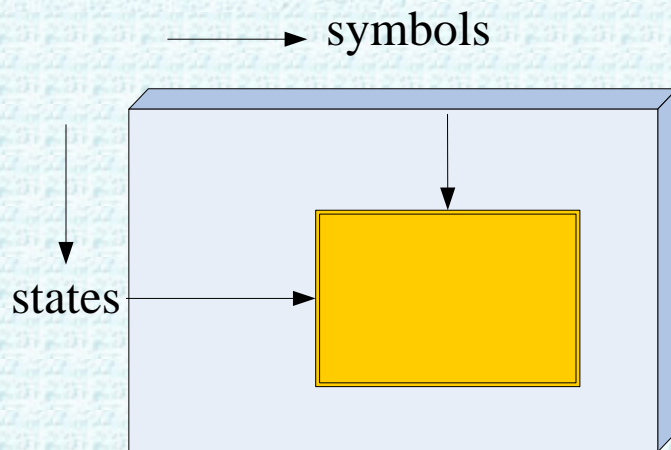
Parser Generating System





구문 분석기의 작성 [2/3]

▣ Parsing Table 구조



❖ $ptbl[S, X] =$

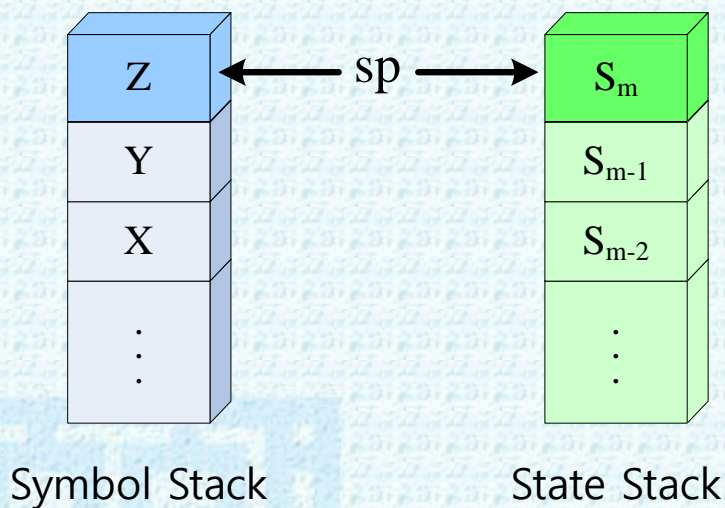
shift	: > 0
reduce	: < 0
accept	: $NO_RULES + 1$
error	: 0



구문 분석기의 작성 [3/3]

■ Parsing stack

- Parsing stack은 병렬로 운행되는 Symbol stack과 State stack으로 구성
 - Symbol stack : 문법 심벌 저장
 - State stack : 상태 저장





LR parser for Mini C

▣ Text pp.361-365 에 수록

▣ Mini C Grammar(Text. pp. 619-622)

- | | |
|-----------------------|-------|
| (1) number of rules | : 97 |
| (2) number of symbols | : 85 |
| (3) number of states | : 153 |

AST를 위한 문법



Mini C grammar

1. 부록 A : pp.619-622

- ▣ LALR(1) 문법
- ▣ AST를 구성하기에 적당한 형태
- ▣ **miniC.tbl**
 - ▣ #define NO_RULES 97
 - ▣ #define GOAL_RULE (NO_RULES + 1)
 - ▣ #define NO_SYMBOLS 85
 - ▣ #define NO_STATES 153

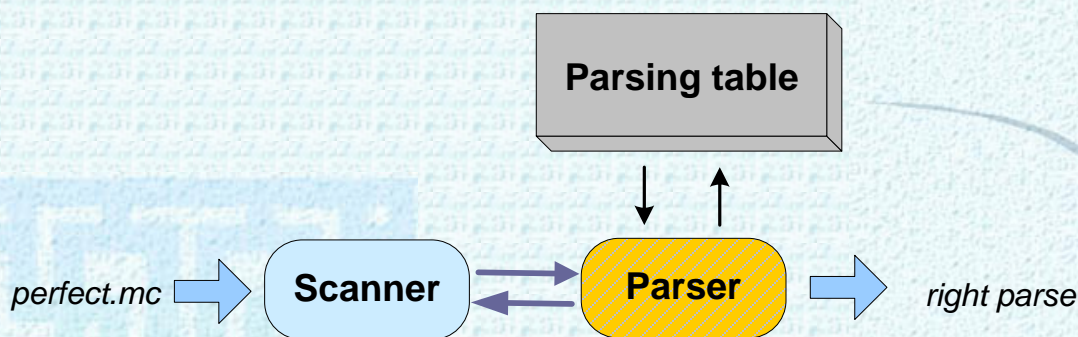
2. 10.3 절 : pp.434-437

- ▣ 부록 A에 있는 문법에 AST 정보를 추가
- ▣ 스탠포드 PGS의 입력 형태



Programming Assignment #2

- Implement an LR parser for Mini C grammar(pp.371-372 => 8.20)
- Problem Specifications
 - input : perfect.mc(Text p. 616-617)
 - output : **right parse**
 - methods :
 - use the existing scanner for Mini C.
 - get the parsing table for Mini C at the web site.
<http://pl.skuniv.ac.kr>
 - program an LR parser(Text pp.361-365)





LR Parsing



Parser Main Routine(pp.361-362)

```
#include <stdio.h>
#include "scanner.c "
#include "miniC.tbl "           // Parsing Table

#define NO_RULES                70           // number of rules
#define GOAL_RULE               (NO_RULES+1) // accept rule
#define NO_SYMBOLS 74          // number of grammar symbols
#define NO_STATES              131          // number of states
#define PS_SIZE    100         // size of parsing stack

void dpush(int,int);
int sp,pstk[PS_SIZE];
void semantic(int);
void parser();

void main()
{
    cout<<" start of parser"<<endl;
    parser();
    cout<<" end of parser"<<endl;
}
```

Parser Main Routine

```
void parser()
{
    int entry, ruleNumber, lhs;
    int current_state;
    struct tokentype token ;

    sp = 0; pstk[sp] = 0;
    token = scanner();
    while (1) {
        current_state = pstk[sp];
        entry = ptbl[current_state][token.number];
        if (entry > 0)          // shift action
        { dpush(token.number, entry);
          token = scanner();
        }
    }
}
```

Parser Main Routine

```

else if (entry < 0)           // reduce action
{
    ruleNumber = -entry;
    if (ruleNumber == GOAL_RULE)    // accept action
    {
        printf(" *** valid source ***\n");
        return;
    }
    semantic(ruleNumber);
    sp = sp - right[ruleNumber]*2;
    lhs = left[ruleNumber];
    current_state = ptbl[pstk[sp]][lhs];
    dpush(lhs, current_state);
}
else {
    printf(" === error in source ===\n");
    exit(1);
}
}
}

```


Parser Main Routine

```
void semantic(int n)
{
    cout<<"reduced rule number = "<< n <<endl;
}
```

```
void dpush(int a, int b)
{
    pstk[++sp] = a;
    pstk[++sp] = b;
}
```

Mini C Program(Perfect Number)(p.616)

```
const int max = 500;
void main()
{   int i, j, k;
    int rem, sum ;

    i = 2;
    while (i<= max) {
        sum = 0;
        k = i / 2;
        j = 1;
        while (j <= k) {
            rem = i % j;
            if (rem == 0) sum += j;
            ++j;
        }
        if (i == sum) write(i);
        ++i;
    }
}
```

- LALR(1) Grammar의 예제 – pp.345-348,
pp.372-373

0. $S' \rightarrow S \$$

1. $S \rightarrow L = R$

2. $S \rightarrow R$

3. $R \rightarrow L$

4. $L \rightarrow * R$

5. $L \rightarrow id$

Token : $= * id$

Input Data : $id=id \$$

LR Parser

```
#include <iostream.h>
#include <stdio.h>
#include <string.h>
#include <stdlib.h>
#include <ctype.h>
```

```
#define NO_RULES          5           // number of rules
#define GOAL_RULE         (NO_RULES+1) // accept rule
#define NO_SYMBOLS        7           // number of grammar symbols
#define NO_STATES         10          // number of states
#define PS_SIZE           100         // size of parsing stack
```

```
void dpush(int,int);
int sp, pstk[PS_SIZE];
void semantic(int);
void parser();
```

Grammer

- | | |
|---------------|---------------|
| 1. S -> L = R | 2. S -> R |
| 3. R -> L | 4. L -> * R |
| 5. L -> id | 0. S' -> S \$ |

Grammar Symbols

id * = \$ L R S S'

LR Parser

```
int leftSymbol[NO_RULES + 1] = { 7, 6, 6, 5, 4, 4};
int rightLength[NO_RULES+1] = { 2, 3, 1, 1, 2, 1};
int parsingTable[NO_STATES][NO_SYMBOLS+1] = {
    /* terminal + nonterminal   :
    /* terminal + nonterminal #:
    /* state 0 */
    /* state 1 */
    /* state 2 */
    /* state 3 */
    /* state 4 */
    /* state 5 */
    /* state 6 */
    /* state 7 */
    /* state 8 */
    /* state 9 */
};
```

Grammar

- | | |
|---------------|---------------|
| 1. S -> L = R | 2. S -> R |
| 3. R -> L | 4. L -> * R |
| 5. L -> id | 0. S' -> S \$ |

	id	*	=	\$	L	R	S	S'
	0	1	2	3	4	5	6	7
/* state 0 */	{ 5,	4,	0,	0,	2,	3,	1,	0},
/* state 1 */	{ 0,	0,	0,	-6,	0,	0,	0,	0},
/* state 2 */	{ 0,	0,	6,	-3,	0,	0,	0,	0},
/* state 3 */	{ 0,	0,	0,	-2,	0,	0,	0,	0},
/* state 4 */	{ 5,	4,	0,	0,	8,	7,	0,	0},
/* state 5 */	{ 0,	0,	-5,	-5,	0,	0,	0,	0},
/* state 6 */	{ 5,	4,	0,	0,	8,	9,	0,	0},
/* state 7 */	{ 0,	0,	-4,	-4,	0,	0,	0,	0},
/* state 8 */	{ 0,	0,	-3,	-3,	0,	0,	0,	0},
/* state 9 */	{ 0,	0,	0,	-1,	0,	0,	0,	0}

LR Parser

```
enum tsymbol{ tnull = -1 , tident, ttimes, tequal, teof};
```

```
struct tokentype{  
    int number;          /* token number */  
    char var;            /* token value */  
};
```

```
struct tokentype token;
```

```
void scanner();    void parser();
```

```
void main()  
{  
    printf(" start of parser\n");  
    parser();  
    printf(" end of parser\n");  
}
```


LR Parser

```
void scanner()
{
    char ch;
    token.number = tnull;
    do {
        while (isspace(ch = getchar())) ;
        if (ch == 'a')
        {
            token.number = tident;
            token.var = ch; }
        else switch(ch)
        {
            case '*': token.number = ttimes;
                    break;
            case '=': token.number = tequal;
                    break;
            case '$': token.number = teof;
                    break;
            default: return;
        }
    } while (token.number == tnull);
}
```

LR Parser

```
void parser()
{
    int entry, ruleNumber, lhs;
    int current_state;
    struct tokentype token ;

    sp = 0; pstk[sp] = 0;
    token = scanner();
    while (1) {
        current_state = pstk[sp];
        entry = parsingTable[current_state][token.number];
        if (entry > 0)           // shift action
        { dpush(token.number, entry);
          token = scanner();
        }
    }
}
```

LR Parser

```
else if (entry < 0 )           // reduce action
{
    ruleNumber = -entry;
    if (ruleNumber == GOAL_RULE)    // accept action
    {
        printf(" *** valid source ***\n");
        return;
    }
    semantic(ruleNumber);
    sp = sp - rightLength[ruleNumber]*2;
    lhs = leftSymbol[ruleNumber];
    current_state = parsingTable[pstk[sp]][lhs];
    dpush(lhs, current_state);
}
else {
    printf(" === error in source ===\n");
    exit(1);
}
}
}
```


LR Parser

```
void semantic(int n)
{
    printf("reduced rule number = %d\n", n);
}
```

```
void dpush(int a, int b)
{
    pstk[++sp] = a;
    pstk[++sp] = b;
}
```

Data : $a=a\$$

```
C:\ 명령 프롬프트

C:\WUNIUWLEC\CDTWCompiler>ExParser < input.txt
*** Start of parser ***
reduced rule number = 5
reduced rule number = 5
reduced rule number = 3
reduced rule number = 1
*** Valid Source ***
*** End of parser ***

C:\WUNIUWLEC\CDTWCompiler>
```