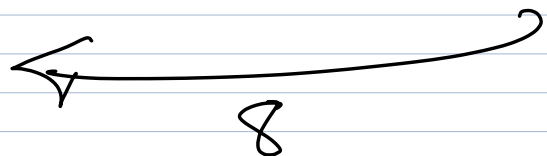


$p =$   
 4 5 6 7 8 9 10 11  
 5ft 5.2ft 4.8ft 6ft 4.3ft 5.1ft 8.2ft 5.0ft

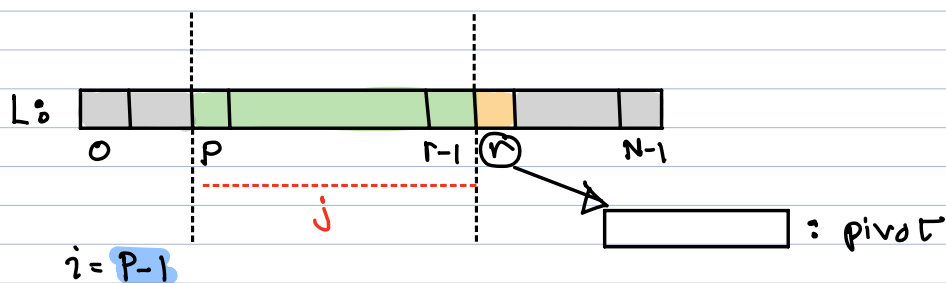
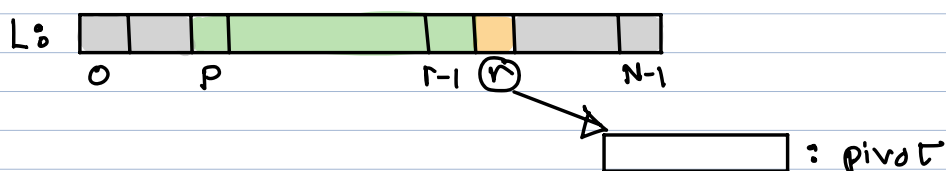
Partition  
 4 5 6 7 8 9 10 11  
 5ft 4.8 4.3 3.2 5ft 5.2 6ft 5.1



Tony Hoare : Computer Scientist : Quick Sort.

1970's : Communicating Seq. Processes (CSP) : Go.  
 goroutine / async def / co-routine

Adversary Input Randomized Quick.



$L[i] \leftrightarrow L[j]$

```

j = p
while j < r:
    L[j]
    j = j + 1

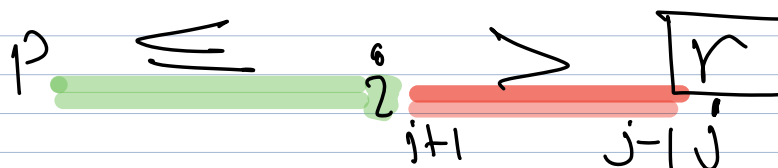
```

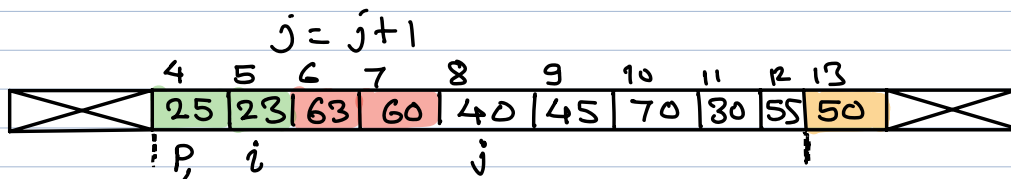
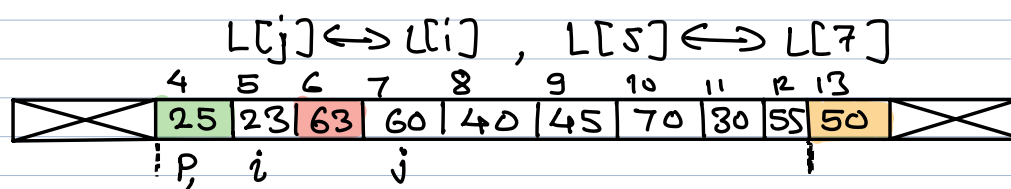
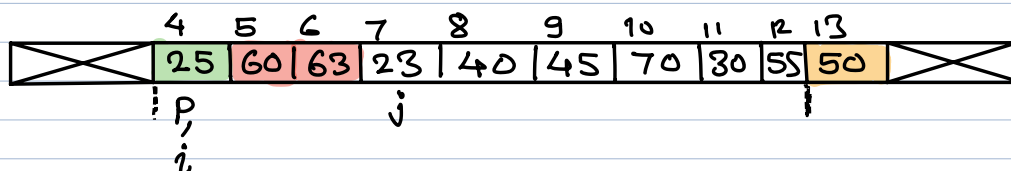
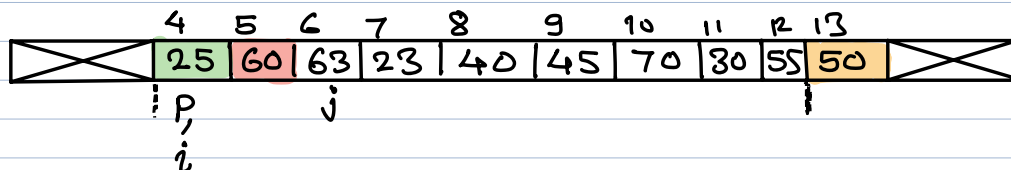
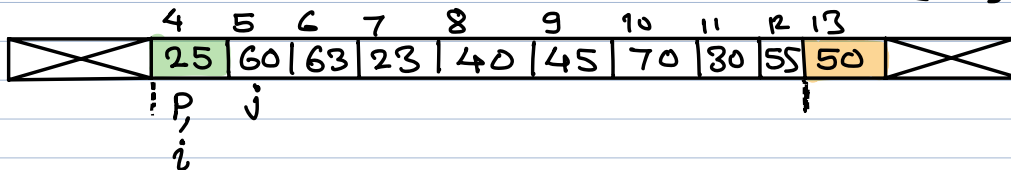
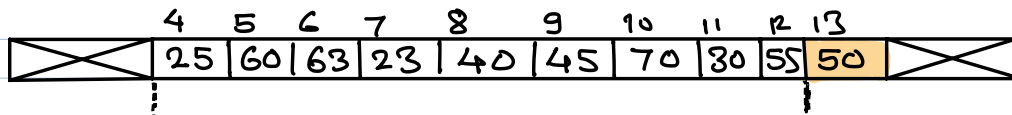
```

j = p
while j < r:
    if L[j] <= pivot:
        j = j + 1

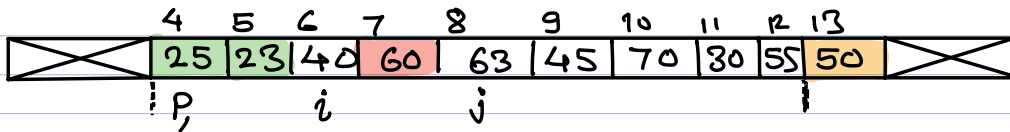
```

- 1) Set index  $j \leftarrow p$ .  
 Set index  $i \leftarrow p-1$   
 Define tmp variable for swapping  
 Set pivot  $\leftarrow L[r]$
- 2) Visit all elements from  $p$  to  $r-1$  using index  $j$ .
  - 2-1) Compare  $L[j] \leq \text{pivot}$
  - 2-2) Condition false
    - $\rightarrow$  Do nothing
    - $\rightarrow$  Goto 2-4)
  - 2-3) Condition True.
    - $i \leftarrow i+1$
    - $L[i] \leftrightarrow L[j]$
    - Goto 2-4)
  - 2-4) Go to the next element.

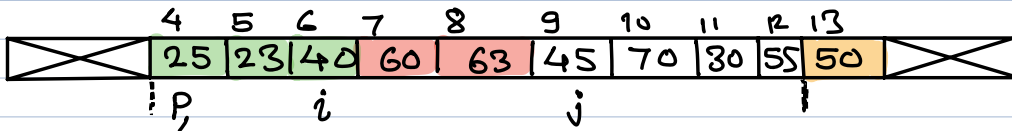




$L[i] \leftrightarrow L[j], L[6] \leftrightarrow L[8]$



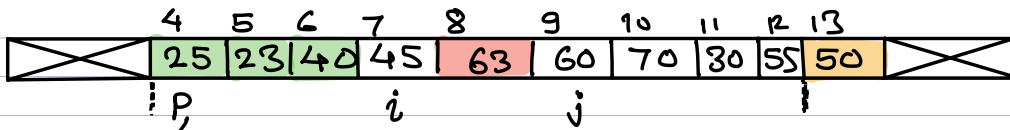
$j = j + 1$



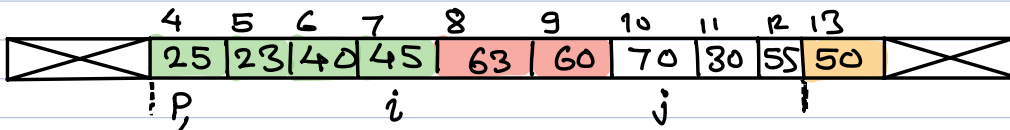
$j = 9, i = 6, \text{pivot} = 50, L[j] \leq \text{pivot}, 45 \leq 50, \text{True.}$

$j = j + 1, i = i + 1$

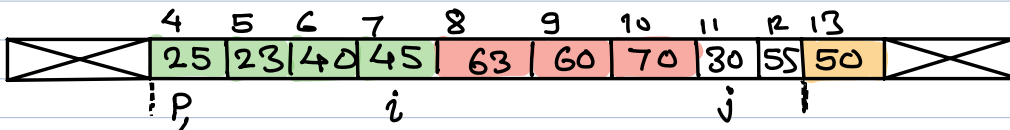
$L[i] \leftrightarrow L[j], L[7] \leftrightarrow L[9]$



$j = j + 1$



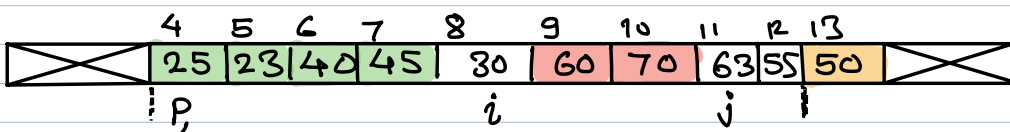
$j = 10, i = 7, \text{pivot} = 50, L[j] \leq \text{pivot}, 70 \leq 50: \text{false}$   
 $j = j + 1, \text{Go to next}$



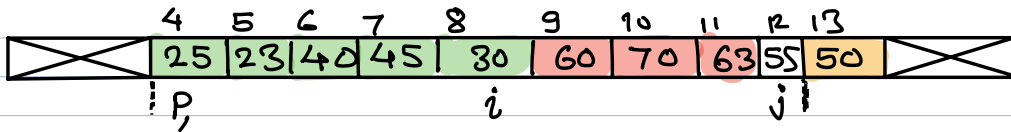
$j = 11, i = 7, \text{pivot} = 50, L[j] \leq \text{pivot}, 30 \leq 50: \text{True.}$

$j = j + 1, i = i + 1$

$L[i] \leftrightarrow L[j], L[8] \leftrightarrow L[11]$

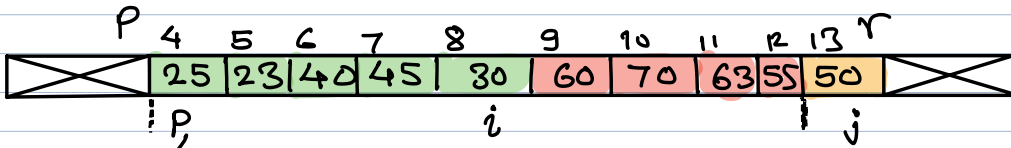


$$j = j + 1$$

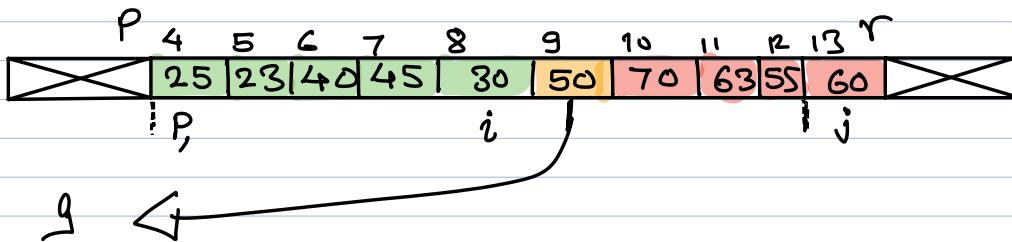


$j = 12$ ,  $i = 8$ ,  $\text{pivot} = 50$ ,  $L[j] \leq \text{pivot}$ ,  $55 \leq 50$ , False

$j = j + 1$ , Goto next  $\rightarrow$  Stop.



$$L[i+1] \leftrightarrow L[j]$$



Loop-Invariant:

$$\forall k \in [P, i] \quad , \quad L[k] \leq \text{pivot}$$

$$\forall k \in (i, j) \quad , \quad L[k] > \text{pivot}$$

Formal verification: David Gries, Dijkstra, Tony Hoare

Prop. logic, Predicate logic, Set theory, Induction.

loop-variant : kernel